

Functional Programming – Laboratory 5

Tail recursion, Control

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March 28, 2012

1 Concepts

- Tail recursion
- when, unless, let, let*
- block, tagbody, loop
- progn, prog1, prog2
- prog, prog*
- do, do*, dolist, dotimes

2 Questions from Laboratory 4

- How many types of recursion do you know? Which one is the most efficient?
- How do we recognize a tail recursive function? What is the method used for writing tail recursive functions? Give one example of a tail recursive function (write the definition of the function in LISP).
- Simulate (trace functionname), where functionname is the name of your own function.

3 Control

3.1 when, unless

```
> (when (= 8 9) 4 9)
> (when (< 8 9) (print 'azi) 4 9)
> (when (> 10 9))
> (unless (> 10 9) 100 'nothing)
```

```

> (unless (> 10 9) 100 (print 'nothing))

> (unless (< 10 9) 100 (print 'nothing))

> (unless (< 10 9) (print 'nothing) 100)

```

when and *unless* are equivalent with:

```

(unless p a b c) == (cond ((not p) a b c))

(unless p a b c) == (when (not p) a b c)

(when p a b c) == (cond (p a b c))

(when p a b c) == (unless (not p) a b c)

```

3.2 block, tagbody, loop

; construction of a block — the structure

```
(block block_name <form1> <form2> .... <formn>)
```

*; <form1> <form2> <formn> are optional
; in the case of missing it will return NIL*

```

> (block block_name (print 'expr1) (print 'expr2) (print 'expr3))

> (block block_name 1 2 3 4)

> (block block_name 2)

```

*; Blocks with abrupt exit
; for this we use
; return-form or return.*

*; When return-form block_name is evaluated the program exits
; the block block_name with the value resulted from the
; evaluation of the second argument (which is optional
; <form1> <form2> <formn>),
; if this is missing, then we get NIL.*

```

> (block block_name3
    (setq x 1)
    (print (1+ x))
    (return-from block_name3 (1+ x))
    (print 4))

> (block block_name33
    (setq x 1)
    (print (1+ x))
    (setq x (1+ x))

```

```

        (return-from block_name33 (1+ x))
        (print 4))

; return is used in order to exit blocks with the name NIL
; do, dolist, dotimes and loop include implicitly a
; block with the name nil):

> (dolist (x '(1 2 3)))

> (dolist (x '(1 2 3)) (print x))

> (dolist (x '(2 3 -7 5))
      (print x)
      (if (< x 0) (return 'done)))

; Any function definition is a block with the name the name of the function
; from the body of the function we exit using
; return-from

> (defun f()
      (print 'a)
      (return-from f 10)
      (print 'b))

> (f)

; Blocks within which you can use gotos

> (tagbody again
      (setq x (1+ x))
      (if (< x 5) (go again)
          (go end)
      )
      end
      (print x)
  )

> (tagbody lala (setq v 9) (print v))

; read a number until that number is greater than 0

> (tagbody retake
      (print 'Introduce>)
      (if (plusp (read))
          (go retake)
          'done
      )
  )
INTRODUCE>
23

```

INTRODUCE➤

45

INTRODUCE➤

12

INTRODUCE➤

-90

NIL

```
> (block some
    (setq d (1+ 3))
    (print d)
    (if (< d 4) (go some)
        (return-from some 100))
    )
    (print 'somethingStrage)
)
```

```
> (tagbody some
    (setq d (1+ 3))
    (print d)
    (if (< d 4) (go some)
        (return 100))
    )
    (print 'somethingStrage)
)
```

```
> (tagbody some
    (setq d (1+ 3))
    (print d)
    (if (< d 4) (go some)
        (go lala)
    )
    lala
    (print 'somethingStrage)
)
```

Conclusions:

- (block block-name expression-1 expression-2 ... expression-n)
- (return [result])
- (return-from name [result])
- (return value) == (return-from nil value)
- tagbody accepts symbols - they are not evaluated.
- If we reach the end of a tagbody nil is returned.
- loop repeatedly evaluates his parameters.
- We exit a loop using return or throw.

- Advise: try to not use go!

3.3 progn, prog1, prog2

```
> (setq w 11 xx 22)
> (values w xx)
> (values 1 2 3 4)
> (values '(1 2 3 4))
> (progn 10 (print 20) 30)
> (progn 10 (print 20) 30 (values 10 20 30))
> (progn 100 200 300 (values 10 20 30))
> (progn 1 2 3 4 5 6 7)
> (prog1 1 2 3 4 5 6 7)
> (prog2 1 2 3 4 5 6 7)
> (prog2 'la (values 2 3 4) 9)
> (prog2 'la (values '(2 3 4) 90) 9)
> (progn (+ 2 3) (+ 3 5) (+ 11 22))
> (prog1 (+ 2 3) (+ 3 5) (+ 11 22))
> (prog2 (+ 2 3) (+ 3 5) (+ 11 22))
```

Conclusions:

- *progn* returns the last evaluation form (if the last evaluation returns multiple values, then progn will return them all);
- *prog1* returns the first form evaluation (only its first value).
- $(prog2abc...z) == (progna(prog1bc...z))$

3.4 let, let*

let is used to group expressions and performs parallel binding of local variables to values.

let* performs sequential binding of local variables to values.

```
(let [*] (
  (var-1 value-1)
```

```

        (var-2 value-2)
        ...
        (var-m value-m)
    )
    expression-1
    expression-2
    ...
    expression-n)

> (let ((x 1) (y 2) (z 3)) (setq w (+ x y z)) (list x y z w))

> (let* ((x 1) (y (+ x 1)) (z (+ y 1))) (list x y z))

```

3.5 prog, prog*

```

(prog (var-1 var-2 (var-3 init-3) var-4 (var-5 init-5))
      expression-1
      value-1
      expression-2
      expression-3
      expression-4
      value-2
      expression-5
      ...
    )

;    prog is a combination between block, tagbody and let*
;    prog opens implicitly a nil block, therefore
;    in order to exit we use return [result].
; Example:

> (prog (i (sum 0))
      retake
      (print 'Introduce>)
      (setq i (read))
      (if (> i 0)
          (progn (setq sum (+ sum i)) (go retake))
          (return sum)
      )
    )
INTRODUCE>
34
INTRODUCE>
23
INTRODUCE>
11
INTRODUCE>
4.5
INTRODUCE>

```

-3
72.5

; use catch and throw

*; Example: if x is greater than 0 returns ok,
; otherwise the result is the value of x*

```
> (defun f1 (x)
    (catch 'ex1 (f2 x)))
```

```
> (defun f2 (x)
    (if (minusp x)
        (throw 'ex1 x)
        )
    'ok
)
```

```
> (f1 2)
```

```
> (f1 -2)
```

*;; the mechanism of catch and throw is:
; - when throw is evaluated, it does not evaluate the
; forms after throw, but it takes up the forms
; (located in the interpreter stack)
; until it reaches catch with the same label as
; throw (the second argument). In this moment the
; result of throw is the value returned by catch.*

Concluzii:

- In blocks like *block* we use *return - from* or *return* in order to force the exit;
- In blocks like *tagbody* we have *go*;
- The definition of a function is a block with the name of the function;
- For non-local gotos use *catch* and *throw*
- sequential blocks are: *progn*, *prog1*, *prog2*
- *prog* is a combination between *block*, *tagbody*, *let**

3.6 loop,do,do*

Iteration statements in Lisp can be done in multiple ways. The most commonly used is *do*. The general form is:

```
(do (((<var> [<init> [<step>]]))*)
    (<test-end> {<result>}*)
    <body>)
```

```
)
```

```
; or
```

```
(do ( (var1 init1 step1)
      (var2 init2 step2)
      ...
      (varn initn stepn) )
    (test-end result)
  body)
```

The difference between *do* and *do** is that *do* is used for binding in parallel and *do** is used for binding sequential (similar with *psetq* and *let*, or *setq* or *let**

```
; prints the natural numbers between two values:
```

```
> (defun printing (start ending)
    (do ((i start (1+ i)))
        ((> i ending))
      (print i)
    )
  )
```

```
> (printing 2 10)
```

```
; if we want to return some values
; then we write the function:
```

```
> (defun printing2 (start ending)
    (do ((i start (1+ i)))
        ((> i ending) 'done)
      (print i)
    )
  )
```

```
> (printing2 2 8)
```

```
> (defun factorial (n)
    (do* ((i 1 (1+ i))
          (result 1 (* result i)))
        ((= i n) result)
    )
  )
```

```
> (factorial 8)
```

```
; using loop
```

```
> (loop (print 10) (print 20) (print 30) (return))
```



```
> (loop (print '>')
      (if (eq (read) 'stop)
          (return 'exit)
        )
    )
```

```
>
256
>
7896540
>
9
>
-980
>
lalala
>
stop
EXIT
```

```
>(loop for i from 1 to 6 do (print i))
```

```
> (do ((x 1 (1+ x)) (y 1 (* x y))) ((> x 5) y))
```

DOTIMES is a version of the function DO and is equivalent with FOR. The syntax is:

```
(DOTIMES (variable counter result) body)
```

Equivalent with:

```
FOR variable=0 until counter-1 DO
    body
RETURN result
```

Examples:

```
> (dotimes (i 10 (1+ i)) (print 'today))
```

```
> (dotimes (i 10 (1+ i)) (prin1 i) (princ " "))
```

```
> (dotimes (i 4 (* i 2)))
```

```
> (dotimes (i 4 (* i 2)) (print 'one))
```

```
> (dotimes (i 4 (* i 2)) (1+ 89))
```

```
> (dotimes (i 4 (* i 2)) (print (1+ 89)))
```

```
> (dotimes (i 5 (* i i)) (prin1 i))
```

```
> (dotimes (i 5 (* i i)) (prin1 i) (princ " "))
```

```
> (dotimes (i 10 (* i i)) (prin1 i) (princ " ")))
```

DOLIST is similar to DOTIMES but when using DOTIMES the variable receives values between 0 to counter-1, and when using DOLIST the variable receives all the values of the elements from the list.

The syntax is:

```
(DOLIST (variable my-list result) body)
```

Equivalent to:

```
FOR variable=first elem from the list
UNTIL the last element from the list DO
    body
RETUR result
```

Examples:

```
> (dolist (var '(1 2 3)) (print var))
```

```
> (dolist (x '(a b c)) (prin1 x) (princ " ")))
```

4 Homework - deadline: next lab

1. Write a recursive function which returns a list with all the atoms from a list given as parameter:

```
(squash '(a b c (d e) ((f) g)))
=> (a b c d e f g)
```

```
(squash '(a b))
=> (a b)
```

```
(squash '(() (((a)))) ()))
=> (a)
```

2. Write a non-recursive function which calculates the sum of squares of the elements from a list.

```
(square-sum '(1 2 3))
=> 14
```