

Functional Programming – Laboratory 6

Iterative functions

Isabela Drămnesc

March 30, 2012

1 Concepts

- `rplaca`, `rplacd`
- `nconc`, `nreverse`
- simulations of `do`

2 Questions from Laboratory 5

- How many types of blocks do we know?
- How can we exit from each of the blocks? (What instruction do we use?)

3 Exercises

3.1 `let`, `let*` (a test)

;; attention at the binding of the variables when they have the same name!

```
> (setq a 10 b 20)
```

```
> (let ((a 1)
        (b 2))
    (+ a b))
```

```
> a
```

```
> b
```

```
> (let* ((a 1)
         (b a))
    (+ a b))
```

```

    )
> a
> b

(let ((a 5))
  (let ((a 1)
    (b (+ a 1))
  )
    (print (list a b))
  )
)
```

;; test the difference!

```

(let ((a 5))
  (let* ((a 1)
    (b (+ a 1))
  )
    (print (list a b))
  )
)
```

3.2 **rplaca**, **rplacd**

*; **rplaca**: replace the contents of car*
*; **rplacd**: replace the contents of cdr*

```
(setq a '(x y))
```

a

```
(rplaca a 'f)
```

a

```
(rplaca a '(q f))
```

a

```
(rplacd a 'g)
```

a

```
(rplacd a '(h))
```

a

```

(rplacd a '(u v))

a

(setq x '(r s t))

(setq y '(u v w))

(setq a (append x y))

(setq b (append x y))

x
y
a
b

(rplaca y 'new)
y
a
b

(rplacd (cddddr a) 'end)
a
b
y

(rplaca a 'start)
a
b

```

3.3 do - a special iteration form

```

(do ((var-1 init-1 stepper-1)
      (var-2 init-2 stepper-2)
      ...
      (var-n init-n stepper-n))
  (end-test
   end-form-1
   ...
   end-form-k
   return-value)
  body-1
  ...
  (return value)           ; optional
  ...
  body-m)

```

*; Simulate the behavior of the Lisp interpreter for
; the following programs:*

; Try to understand the following examples and explain them

```
1)
(do ((v1 1 (+ 1 v1))
      (v2 () (cons 'a v2))
      )
      ((> v1 5) (return v2) (print 'end))
      (print v1)
    )
```

```
2)
(do ((v1 1 (+ 1 v1))
      (v2 () (cons 'a v2))
      )
      ((> v1 3) (if (> v1 4) (return v2))
        v1
      )
      (print v1)
    )
```

```
3)
(do ((v1 1 (+ 1 v1))
      (v2 () (cons 'a v2)))
      ((> v1 3)
        (if (> v1 3) (return v2))
        v1
      )
      (print v1))
```

```
4)
(do ((v1 1 (+ 1 v1))
      (v2 () (cons 'a v2)))
      ((> v1 5)
        v1
      )
      (if (> v1 4) (return v2))
      (print v1))
```

```
5)
(do ((v1 1 (+ 1 v1))
      (v2 () (cons 'a v2))
      (v3 9))          ; no stepper: value
                        ; remains unchanged —> could be in a LET
      ((> v1 5)
        (return v2) ; uses return in end-form!
      )
    )
```

```

(car v2))    ; return value is ignored!
(print v1))  ; side effect

```

```

6)
(do ((v1 1 (if (< v1 3)
               (+ 1 v1)
               (return 'done))))
    (v2 () (cons 'a v2)))
    ((> v1 5)
     v1)
    (print v1))

```

```

7)
(do ((v1 (if nil
             1
             (return 'done)))
    (+ 1 v1))
    (v2 () (cons 'a v2)))
    ((> v1 5)
     v1)
    (print v1))

```

```

8)
(do ((v1 1 (+ 1 v1))
    (v2 () (cons 'a v2))
    (v3 9)) ; no stepper: value
        ; remains unchanged --> could be in a LET
    ((> v1 5)
     (print v3) ; side effect
     v2)        ; return value
    (print v1)) ; side effect

```

```

9)
(do ((v1 1 (+ 1 v1))
    (v2 () (cons 'a v2))
    (v3 9)) ; no stepper: value
        ; remains unchanged --> should be in a LET
    ((> v1 5)
     ) ; no return value
    (print v1)) ; side effect

```

```

10)
(do ((v1 1 (+ 1 v1))
      (v2 ()) (cons 'a v2))
      (v3 9))      ; no stepper: value
                  ; remains unchanged —> should be in a LET
((print v3) ; funny end cond
)          ; no return value
(print v1)) ; side effect

```

```

11)
(do ((v1 1 (+ 1 v1))
      (v2 ()) (cons 'a v2))
      (v3 9))      ; no stepper: value
                  ; remains unchanged —> should be in a LET
((print v3) ; funny end cond
 33)        ; return value 33
(print v1)) ; side effect

```

```

12)
(do ((v1 1 (+ 1 v1))
      (v2 ()) (cons 'a v2))
      (v3 9))
((print nil) ; end cond is never satisfied
)          ; no return value
(print v1))

```

```

13)
(do ((v1 1 (+ 1 v1))
      (v2 ()) (cons 'a v2))
      (v3 9))
()          ; no end cond
(print v1))

```

```

14)
(do ((v1 1 (+ 1 v1))
      (v2 ()) (cons 'a v2))
      (v3 9))
()          ; no end cond
)          ; no body

```

```

15)
(do ((v1 1 (+ 1 v1))
      (v2 ()) (cons 'a v2))

```

```

      (v3 9))
    )

```

3.4 Destructive append and reverse

```

(setq x '(r s t))

```

```

(setq y '(u v w))

```

```

x

```

```

y

```

```

(append x y)

```

```

(setq a (append x y))

```

```

(eq a x)

```

```

(eq a y)

```

```

(eq (cddddr a) y)

```

; the destructive version of append is nconc

```

(setq a (nconc x y))

```

```

a

```

```

x

```

```

y

```

```

(eq x a)

```

```

(defun my-nconc (x y)
  (if (null x) y
      (progn (nconc-aux x y) x)))

```

```

(defun nconc-aux (x y)
  (if (null (cdr x)) (rplacd x y)
      (nconc-aux (cdr x) y)))

```

```

(progn (+ 5 3) 11)

```

```

(setq x '(r s t))

```

```

(setq y '(u v w))

(setq a (my-nconc x y))

a

x

y

(eq x a)

; One more experiment.

(setq x ())

(setq y '(u v w))

(setq a (my-nconc x y))

a

x

y

(eq x a)

; Not actually!

;-----
; the destructive version of reverse is nreverse

(nreverse '(1 2 3))

(setq a '(1 2 3))

(nreverse a)

a

(setq a '(1 2 3))

(setq a (nreverse a))

a

(setq a '(1 2 3))

(setq b (append a '(3 4 5)))

```


b

```
(setq c (append '(3 4 5) a))
```

c

```
(nreverse a)
```

a

b

c

```
(setq a '(1 2 3))
```

```
(setq a (nreverse a))
```

```
(setq b '(1 2 3 4 5 6))
```

```
(setq c (cdddr b))
```

c

```
(nreverse b)
```

c

; Why? Draw the box notations for b and c (the cons cells).

```
(defun step (rem sofar)  
  (list (cdr rem) (rplacd rem sofar)))
```

```
(setq r '(1 2 3))
```

```
(setq s ())
```

```
(step r s)
```

```
(step '(2 3) '(1))
```

```
(step '(3) '(2 1))
```

```
(defun nreverse-aux (rem sofar)  
  (if (null rem) sofar  
      (nreverse-aux (cdr rem) (rplacd rem sofar))))
```

```
(nreverse-aux '(1 2 3) nil)
```

```

(defun my-nreverse (x)
  (nreverse-aux x nil))

(setq a '(1 2 3))

(my-nreverse a)

(nreverse '(1 2 3 . 4))

(my-nreverse '(1 2 3 . 4))

(defun nreverse-aux (rem sofar)
  (if (atom rem) sofar
      (nreverse-aux (cdr rem) (rplacd rem sofar))))

(setq a '(1 2 3))

(my-nreverse a)

(nreverse '(1 2 3 . 4))

(my-nreverse '(1 2 3 . 4))

```

4 Homework - Deadline: next lab

4.1 Simulate the behavior of the Lisp interpreter for the following programs:

```

1)
(do ((x 0 (+ 1 x))
      (y 0 (+ x y))
      (z 0 (+ y z)))
  ((>= x 10)
   (list x y z))
  (print (list x y z)))

2)
(do ((v1 1 (+ 1 v1))
      (v2 ()) (cons 'a v2)))
  ((> v1 5)
   v2)
  (print v1))

3)
(do ((v1 1 (print (+ 1 v1)))
      (v2 (print ())) (cons 'a v2)))

```

```

(> v1 0)
v2)
(print (list v1 v2)))

```

```

4)
(do ((v1 1 (+ 1 v1))
      (v2 'nothing (cons 'a v2)))
    (> v1 0)
    v2)
(print v1))

```

```

5)
(do ((v1 1 (+ 1 v1))
      (v2 () ))
    (> v1 3)
    v2)
(setq v2 (cons 'a v2)))

```

```

6)
(do ((v1 1 (+ 1 v1))
      (v2 () (cons 'a v2)))
    (> v1 3)
    v2)
(setq v2 (cons 'b v2)))

```

```

7)
(do ((v1 1 (+ 1 v1))
      (v2 '(b) (cons 'a v2))
      (v3 1 (list v1 v2 v3)))
    (> v1 3)
    v2)
(print (list v1 v2 v3)))

```

```

8)
(do ((v1 1 (+ 1 v1))
      (v2 () (cons 'a v2))
      (v3 9))          ; no stepper: value
                                ; remains unchanged --> should be in a LET
    ((print v1)         ; end cond
     )                  ; no return value
    (print v3))

```

4.2 Define a function for: the length of a list, the reverse of a list, the greatest common divisor of two natural numbers in 3 ways:

1. the recursive version;
2. the tail recursive version;
3. the iterative version.