

Where is Aristotle in our brain? On biologically plausible reasoning embedded in neuronal computation[†]

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Abstract

Human cognition involves tightly interleaved numerical but also symbolic[†] (including logical)[†] computations. ~~How—can~~ Determining how the brain can implement such processing is an important issue, and we would like to address some aspects of this issue here[†] by combining two approaches. On the one hand, ontology-oriented languages allow us to describe symbolic structured knowledge and perform logical inference using entailment rules. To what extent this could provide a rather natural representation of usual human reasoning is an open question[†] that we are going to discuss here, considering a generalization to modal logic[†]; ~~and we will~~ also go beyond deductive reasoning, discussing whether this could also apply to inductive and abductive reasoning. On the other hand, spiking neuronal networks are biologically plausible implementations of brain circuit computations, ~~which—meaning~~ that they can provide a way to manipulate symbols embedded as numeric vectors that carry semantic information. In the present work, we consider such an architecture with the ~~Vector—Symbolic~~ Architecture—vector symbolic architecture (VSA) formalism[†], allowing us to describe neuronal implementations at an algebraic level. This development illustrates how the former cognitive mechanisms can naturally emerge from ~~usual—~~distributed calculus, yielding neuro-symbolic computations. Our ~~argument—aim~~ is to show

that it can be implemented, considering a VSA approach, into biological neuronal computations. **Author query: In the previous sentence, it is not completely clear what the word “it” is referring to.** Such neuro-symbolic deductive mechanisms are especially useful in complex problem-solving.

Keywords: Ontology, Modal Logic, Semantic Pointer Architecture, Abstract Thought, Neuro-symbolism, Problem Solving.

1 Introduction

1.1 From sensorimotor processing to logical reasoning

It is generally admitted (as reviewed, e.g., in [1]) that human logical reasoning emerges progressively from the sensorimotor association, with the formation of stable concepts, even at a symbolic level, before ~~being able it is possible~~ to manipulate them at a concrete level, ~~by~~ performing inductive reasoning, ~~and perform up to~~ more formal deductive reasoning. Furthermore, as pointed out in [2], such a mechanism includes a cultural bias, since not all cultures feel the need to develop formal logical operation competencies, and obviously most people do not use such formal operations in all aspects of their lives. However, as discussed in, e.g., [3], deductive reasoning, especially for goal-driven behavior, is deeply interleaved with heuristic deduction, ~~involving which involves~~ analogy and metaphor. Furthermore, as thoroughly studied by, e.g., [4], the experience of conscious or subconscious emotion has a powerful influence on rational decisions, including the choice of alternatives in deductive reasoning. Does this mean that the human brain does *not* need to develop deductive reasoning, except for singular cultural needs (e.g., for scholarship ~~constraints~~ or to practice formal science)? Our understanding is that the situation is not binary. We need to make deductions to solve ~~problems in~~ everyday life ~~problems~~, while the elements we briefly reviewed here, demonstrate that such deductions are not Boolean (either true or false) but related to a given context, weighted by a certain level of belief, and biased by motivational elements in ~~the a~~ wide sense.

In this study, we attempt to reconcile deductive reasoning with such cognitive mechanisms of inference, ~~up to~~ ~~reach~~ a biologically plausible neuronal implementation and show to what extent this could be extended to approximate deductive reasoning, in addition to inductive and abductive reasoning mechanisms¹.

¹Here we make the distinction between:
- *deductive reasoning*, ~~which is the process of determining for~~ the formal logical consequence of some assumptions, considered as true, or approximately true;
- *inductive reasoning*, which is the process of inferring some general principle from a set of knowledge and plausible induction rules; ~~and~~
- *abductive reasoning*, ~~which~~ is the process of inferring an explanation of some assertions, i.e., hypothesizing the precondition of a consequence.

~~A-Taking this discussion~~ a step further, how does the brain give meaning to the symbols considered here? We will not address this issue here but would like to clarify some points. First of all, in neuro-symbolic studies, as reviewed in [5], grounding is understood as the process ~~to embed of embedding~~ symbolic computations onto real-valued features [6] because it ~~allows providing provides~~ a semantic interpretation or model (in the sense of a model of a set of logical assertions) of the symbolic system. This means that it is no ~~more longer~~ an abstract set of assessments (potentially without any concrete implementation) but something that corresponds to a real formal object. Our approach thus does not solve the grounding problem, but in some sense it does solve what we can call the “anchoring” problem of relating symbols ~~with to~~ the neural substrate. This will be further discussed in the last section after ~~detailling the~~ present approach is described in detail.

1.2 Representing neuronal activity at a symbolic level

As a possible entry point to considering a biologically plausible implementation at a symbolic level, ~~Vector Symbolic Architectures~~ vector symbolic architectures (VSAs) ~~was were~~ introduced as a way to manipulate symbolic information represented as numeric vectors (see, e.g., [7] for an introduction). VSAs have been proven to be helpful ~~to model in modeling~~ high-level cognition and accounting for multiple biological features [8, 9]. More specifically, the ~~Semantic Pointer Architecture~~ semantic pointer architecture (SPA) [9] instantiates so-called semantic pointers (i.e., vectors that carry semantic information) and ~~their manipulation makes it possible to manipulate them~~ in networks of spiking neurons. This approach ~~mtakes~~ a significant step towards the unification of symbolic and sub-symbolic processing in that it provides a way to translate the former into the latter. Consequently, complex knowledge representations in the form of compositional structures that are traditionally restricted to symbolic approaches can now be distilled in numerical and even neural² systems [10].

How can we represent a symbol in a neuronal assembly? A localist representation (one neuron or neuron group represented by a symbol) does not correspond to what is observed in the brain, and the basic idea is that a symbol corresponds to a pattern of activity of the whole assembly. Let us consider a spiking neuron network and quantify its activity ~~by using~~ some statistics, e.g., the neuron rates, or higher-order statistics (see, e.g., [11] for a discussion). As developed in [12], this includes timing codes and population codes (i.e., relative timing codes between neurons), and the authors of [12] show how, with their developed ~~Neural Engineering Framework~~ neural engineering framework (NEF), we can collect this high-dimensional set of bounded quantitative values, which can be normalized, as a unitary stochastic vector in a ~~huge~~ high-dimensional space (~~of with~~ a few thousands dimensions for a biological neuronal map, ~~and~~ often a few hundred dimensions at the simulation level),

²The term “neural” refers to any type of nerve cell, whereas “neuronal” is specifically related to neurons.

defining an ~~SPA~~-~~Semantic Pointer Architecture~~. This includes a time representation in spiking neuron systems, not ~~only-just~~ a rate representation. The key point is that compared to other representations, e.g., based on synchrony within the neural assembly, the NEF alternative is much more scalable.

In the present study, we consider these developments as prerequisites and will simply consider that neural assembly activity is represented by a high-dimensional unary stochastic vector. We also need to specify transformations and will define them at this abstract algebraic level. Mainly, following [13], we will consider the auto-association mechanism, as developed in [14], and functional transformations, as detailed in [12]; their development is based, in a nutshell, on parameterized kernel-based approaches.

1.3 What is ~~the this~~ paper about?

We first revisit how to encode symbols, within the VSA approach based on the framework introduced in [9], analyzing in more detail than ~~previously-done~~ previous works the numerical approximation statistic, because this is needed to model the VSA mechanism at a macroscopic level. We also describe how to generalize symbol encoding considering a related degree of belief, beyond binary information, and ~~after-following~~ [13], we explain the semantic interest of such a generalization.

We then consider knowledge hierarchical structure encoding, in the sense of, e.g., [15], ~~in-as-a~~ complement to associative and sequential memorization, and we discuss how to implement such a memory structure, using the VSA ~~architecture~~. To this end, we have to review ~~usual~~-VSA data structures and ~~showing-their-link-with-both-demonstrate-that-they-are-linked-to~~ cognitive memory classification according to [15]. To better understand their computational properties, we also illustrate how such existing VSA data structures ~~can~~ compare to ~~with-usual~~ programming language containers.

This new data structure is the basic tool that is used to then study to what extent symbolic inference could be implemented by specific connectivity feedback, and, as a second contribution, we show that the vanilla fixed-point algorithm used for a forward deduction on a decidable set of entailment rules can be adapted to this biologically plausible framework, ~~allowing-making-it-possible~~ to perform deductively reasoning, ~~but and~~ also, to some extent, using inductive and abductive ~~deductions-reasoning~~.

We finally illustrate the proposed mechanism ~~by-using~~ both a simulation at the mesoscopic level, ~~considering-utilizing~~ the well-established Nengo simulator, ~~but-and~~ also a simulation at ~~a-the~~ macroscopic scale, ~~showing~~ We show, as the third contribution of this paper, that such computations may be, up to a certain point, approximated without explicitly ~~computing-performing~~ computations at the vector component level; ~~but-using~~ instead, an algorithmic ersatz can be used.

2 Symbolic information encoding

2.1 Symbol encoding

At the numerical level, each symbol is implemented as a randomly drawn fixed unit d -dimensional vector $\mathbf{x} \in \mathcal{R}^d$. Typically, $d \simeq 100 \cdots 1000$, and we expect to manipulate $k \simeq 100 \cdots 10000$ symbols at the simulation level. A dimension typical dimension of 256 has been considered in related studies, such as [13]. Here, our macroscopic implementation uses the same dimension. In a cortical or brain map, the order of magnitude is higher since the vector corresponds to the neuronal map activity (thus it is closer to $10^{5 \cdots 6}$) and the number of encoded symbols depends on which map is considered.

The vector components are drawn from a normal distribution $\mathcal{N}(0, \sigma)$, i.e., this distribution has of zero mean and a standard-deviation $\sigma \stackrel{\text{def}}{=} 1/\sqrt{d}$, in order to have an $O(1)$ magnitude.

A similarity measure is now introduced in order to semantically compare two vectors. Classically, the cosine similarity (i.e., normalized dot product, denoted \cdot) is used to compute the semantic similarity between two unit vectors:

$$\mathbf{x} \cdot \mathbf{y} \stackrel{\text{def}}{=} \mathbf{x}^\top \mathbf{y} = \cos(\widehat{\mathbf{x}, \mathbf{y}}),$$

where \mathbf{x}^\top denotes the transpose of \mathbf{x} . It is in obvious correspondence with This measure obviously corresponds to the angular distance between both the vectors.

The key property is that, provided that the space dimension d is large enough, two randomly chosen different vectors will be approximately orthogonal. More precisely,

$$\mathbf{x} \cdot \mathbf{y} \sim \delta_{\mathbf{x}=\mathbf{y}} + \mathcal{N}(0, 1/\sqrt{d}),$$

i.e., it is almost 1 if equal and 0 otherwise, plus an centered normal noise [16]. **Author query: In the previous sentence, the meaning of the phrase “if equal” is not completely clear. Do you mean “if the vectors are equal”?** At the numerical level³, we have verified for $d \simeq 100 \cdots 1000$ that we generate unary vectors with a bias below 0.3%, while orthogonality is verified with a bias below 0.4%; and the noise standard deviation prediction bias is below 0.3%.

This allows us to define a hypothesis to decide whether the \mathcal{H}_0 hypothesis $\mathbf{x} \cdot \mathbf{y} = 0$ can be rejected. We are in the situation of can consider a two-tailed “z-test” with the alternative hypothesis \mathcal{H}_0 , which states that $\mathbf{x} \cdot \mathbf{y} \neq 0$. Here, the z-score⁴ is obviously, with d samples and a known standard deviation of

³See https://raw.githubusercontent.com/vthierry/onto2spa/main/figures/z_score.mpl for the open-source code used in this subsection.

⁴Given a distribution, the z-score for d samples is defined as

$$z \stackrel{\text{def}}{=} \frac{\bar{X} - \mu}{\sigma/\sqrt{d}},$$

with here-where the expected means is $\mu = 0$, the a priori standard-deviation is $\sigma = O(1/d)$, and the experimental mean $\bar{X} = (\mathbf{x} \cdot \mathbf{y})$ value is obtained from the dot-product.

with an order of magnitude $O(1/d)$, is the following:

$$z \equiv \sqrt{d} (\mathbf{x} \cdot \mathbf{y}).$$

It follows an almost distribution, which-as-it can be easily verified numerically, as shown in Fig. 1₇ (left column). **Author query: In the previous sentence, is “almost distribution” meant to say “almost normal distribution”? A step further considering** For two vectors that are not independent but have an angular dependency, we can numerically observe, in Fig. 1₇ (right column), the similarity dependency as a function of the vector’s relative orientation. This obvious fact is quite important, in the sequel allowing us to develop a macroscopic simulation of VSA operations.

This allows-on-makes it possible, on the one hand, to consider, for instance, a ± 2 threshold on-for the standard deviation, in-link-along with considering this z-score to have a confidence interval better than 99%, and to relate the similarity estimation to an angular dependence between two vectors, as detailed in Fig. 1. To our-best-the best of our knowledge, this obvious implementation recipe has not yet been made explicit, and it is used for-in subsection 5.1, as detailed in Appendix D, allowing us to propose to simulate the different operations defined later in the sequel this paper at a macroscopic scale.

2.2 Modality encoding

Based on this, mMost VSA approaches consider that 2-two vectors \mathbf{x} and \mathbf{y} are semantically equivalent when this-the similarity τ equals 1, but with-there are different ways to interpret the-this result. Here, we enrich the notion of something being either false or true, by using a numeric representation of, e.g., partial knowledge, as illustrated in Fig. 2. The true value corresponds to 1 (fully possible and fully necessary), the false value to -1 (neither possible nor necessary, i.e., impossible), and the unknown value to 0, which corresponds to a fully possible but absolutely not necessary value.

Our representation is in one-to-one correspondence with the dual notions of necessity and possibility representation of-in the standard possibility theory. Information is always related to a certain degree of belief. While almost all partially known information is related to probability, the human “level of truth” is more subtle and related to possibility and necessity, as formalized in the possibility theory discussed in [17] and [18]. This theory stems from modal logic, i.e., something being true in a “given context” [19], which is also considered representative of what is modeled in educational science and philosophy [20]; namely, it corresponds to common-sense reasoning in the sense of Piaget’s-sense [21], which involves taking exceptions into account, i.e., considering non-monotonic reasoning. In other words, possibility theory is devoted to the modeling of incomplete information, which is related to an observer’s belief regarding a potential event and surprise after the event’s occurrence. Furthermore, in symbolic artificial intelligence, i.e., knowledge representation and logical inference, a link has been drawn between this necessity/possibility dual representation and ontology [22]. This must be understood as a deterministic

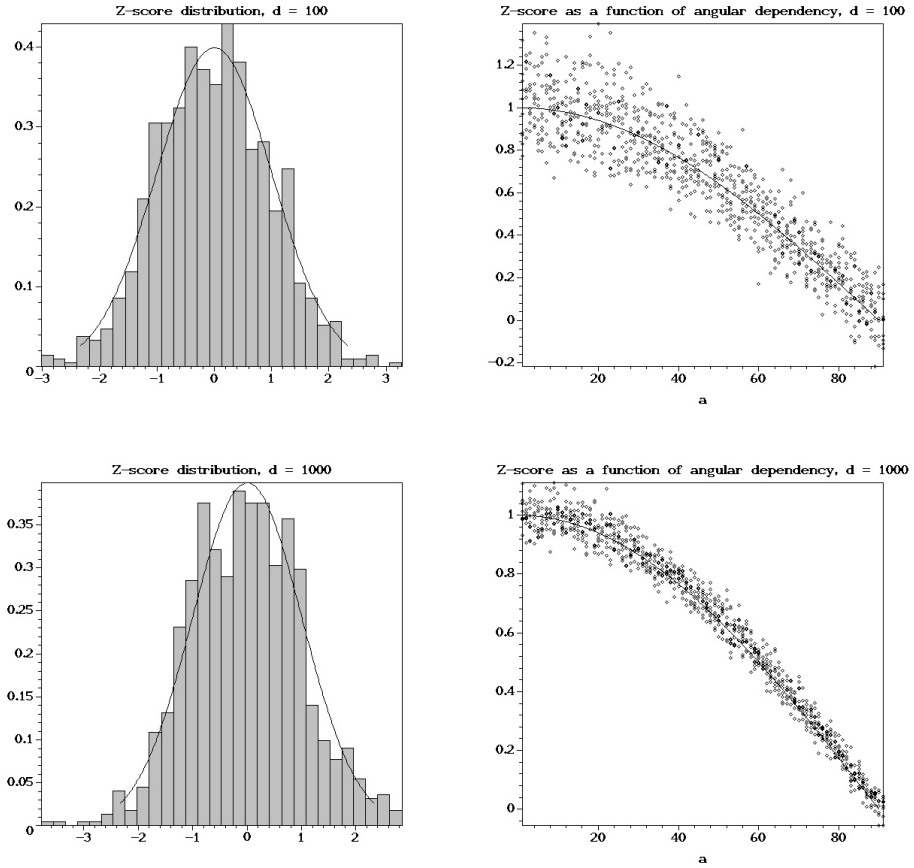


Fig. 1 Numerical observations of the similarity defined by the dot product of two random vectors for $d = 100$ in the upper row and $d = 1000$ in the lower row. The left columns shows the histogram of the z-score ($\sqrt{d}(\mathbf{x} \cdot \mathbf{y})$) for two normal vectors, in comparison with a normal distribution. These experimental distributions have a kurtosis of about 10; thus this is lower than the kurtosis of for a normal distribution, which is expected to have a kurtosis of 3. The right columns shows the z-score as a function of the angle $a \stackrel{\text{dec}}{=} \widehat{\mathbf{x}, \mathbf{y}} = \arccos(\mathbf{x} \cdot \mathbf{y})$, allowing making it possible to visualize the dispersion with respect to the expected cosine profile.

theory, in the sense that partial knowledge is not represented by randomness⁵. Given these elements, it is thus interesting to represent approximate knowledge at the cognitive level by a degree of belief in relation to the notions of

⁵This deterministic representation of partial knowledge can be generalized in order to also include a representation of the randomness belief. In the vanilla possibility theory, the possibility can be seen as an upper probability: Any possibility distribution defines a set of admissible probability distributions, i.e., a consonant plausibility measure in the Dempster-Shafer theory of evidence [23]. In [24, 25], it is proposed to bound the approximate probability, reconsidering the original notion of necessity, in order to also consider a lower bound of probability. This could be an interesting extension for of the present work.

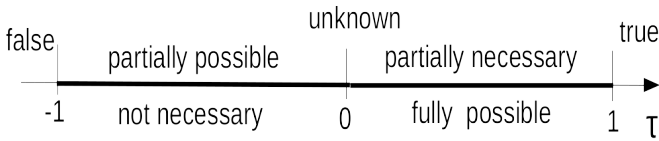


Fig. 2 Representation of the partial truth $\tau \in [-1, 1]$ with regard to necessity and possibility, as defined in the possibility theory. The interpretation is that what is “not so false” is partially possible but not necessary and what is “partially true” is entirely possible but partially necessary. Such a formulation corresponds qualitatively to the human appreciation of the degree of belief in a fact.

possibility and necessity. This will be further developed at a technical level ~~in the sequel~~ later in this paper. **Author query: The phrase “in the sequel” appears in several places in this paper, and its meaning is not clear. I have been changing it to “later in this paper”, but if this is not what you mean, you should reword it accordingly.**

This representation has also been ~~also~~ designed to be compatible with the ternary Kleene logic, ~~besides in addition to~~ being ~~also~~ coherent with respect to the possibility theory, as discussed in detail in [25], where this deterministic representation of partial knowledge is generalized in order to ~~also~~ include a probabilistic representation (using a 2D representation).

This modal notion of partial knowledge or belief is not only epistemic or doxastic but also deontic and so on, i.e., it has several semantic interpretations, depending on the concept feature. Our strong hypothesis here is that all modalities can be encoded using the proposed numerical grounding.

We can thus represent a quantity with a partial degree of belief $\tau \in [-1, 1]$ and use the notation:

$$\hat{\mathbf{x}} \stackrel{\text{def}}{=} \tau \mathbf{x},$$

where \mathbf{x} corresponds to the numerical grounding of a symbol, and $\hat{\mathbf{x}}$ corresponds to the numerical grounding of a symbol weighted by a modality quantification τ .

Interestingly enough, this representation is coherent with the semantic similarity in the following sense: Are two vectors similar? Considering $\mathbf{x} \cdot \mathbf{y}$, if it this value is close to 1, then it is considered true, and ~~both the~~ modal representation and semantic similarity are coherent. If it is almost equal to 0, then the modal representation is *unknown*, and in order for it to be coherent with the semantic similarity, we must consider being in an open world, in which all that is not true is not necessarily false. ~~A~~ To take this a step further, if ~~it this value~~ is negative (down to -1), the modal representation considers that it is false, which is coherent with the semantic similarity, although negative values are not explicitly used, ~~up~~ to the best of our knowledge, ~~of in~~ the literature quoted in this paper.

Given these atomic ingredients, let us now study how they can be used in different cognitive data structures.

3 Knowledge structure encoding

3.1 Knowledge representation

From early artificial intelligence knowledge representations such as ~~say, frames~~ previously avowed ~~frameworks~~ previously, to modern web semantic data structures, the basic idea of symbolic representation is to consider symbols that representing objects and express knowledge through relationships, i.e., triple statements of the form ($\$subject$, $\$predicate$, $\$object$), as schematized in Fig. 3.

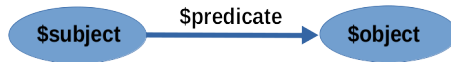


Fig. 3 Atomic representation of knowledge: To express some knowledge regarding a symbol, the subject, we define a feature, with a predicate, ~~with~~ that has an object as an attribute (i.e., a quantitative data value or a qualitative symbol).

Following, e.g., [26], we start by introducing the notion of a concept, with the simple common idea that a concept can be defined by “feature dimensions,” with i.e., attributes with some typed value. The object can be either qualitative or quantitative “data,” or another object, ~~stating that describes~~ relations between objects. This is also the basic syntax of ontology languages. In the brain, such feature dimensions are usually anchored in sensorimotor feature spaces [27], in coherence with the present representation. Furthermore, given a concept, as developed in [26], this choice of representation induces the notion of prototypes, ~~that allows representing which makes it possible to represent~~ the state space region corresponding to the concept.

~~A~~ Taking this a step further, the ~~meaning~~ definition of a concept is completed by relations between concepts, i.e., predicates. Predicates can be generic, in the sense of, e.g., [28], defining a hierarchical taxonomy using the “is-a” predicate, as in almost any such language, but they also have capability qualities (“can”), extrinsic qualities (“has”), and intrinsic qualities (“is”); thus, there are four general predicates. They can also be unconstrained, as in the RDFS framework (see, e.g., [29] for an introduction), describing any property, and in that case, properties also form a taxonomy⁶. **Author query: You may want to write out what “RDFS” stands for since this is the first time it appears in the paper.** This is further illustrated in subsection 3.2. Conversely, as a limit case, we could consider relationships between subjects and objects only, without taking into account the nature of ~~that a~~ relationship (predicate) or its direction.

An important point is that features can be hierarchical because the value itself may have some features: ~~F~~For instance, a quantitative physical value is

⁶For instance, stating that Amid is-a-descendant-of Yang-Li ~~implied~~ implies that Amid is-a-relative-of Yang-Li, the former property being a sub-property of the latter. This example also illustrates that such properties have ~~meta-property~~ meta-properties, such as being transitive or symmetric.

not ~~only-just~~ a “number” stating the current or default value. ~~It~~ may also be specified by a unit and a precision value or some bounds. In terms of ~~the~~ data structure, this forms a tree, and the whole data structure is a set of trees, i.e., a forest, ~~and~~ thus ~~it is~~ a graph.

We are going to specify how we can represent such symbolic information at a biologically plausible level in the present study, ~~thanks to using vector symbolic architecture VSAs.~~

3.2 Hierarchical schematic organization of information

To make things more concrete, we aim at manipulating the symbolic representation of knowledge of the form ~~of~~ ~~shown in~~ Fig. 4.

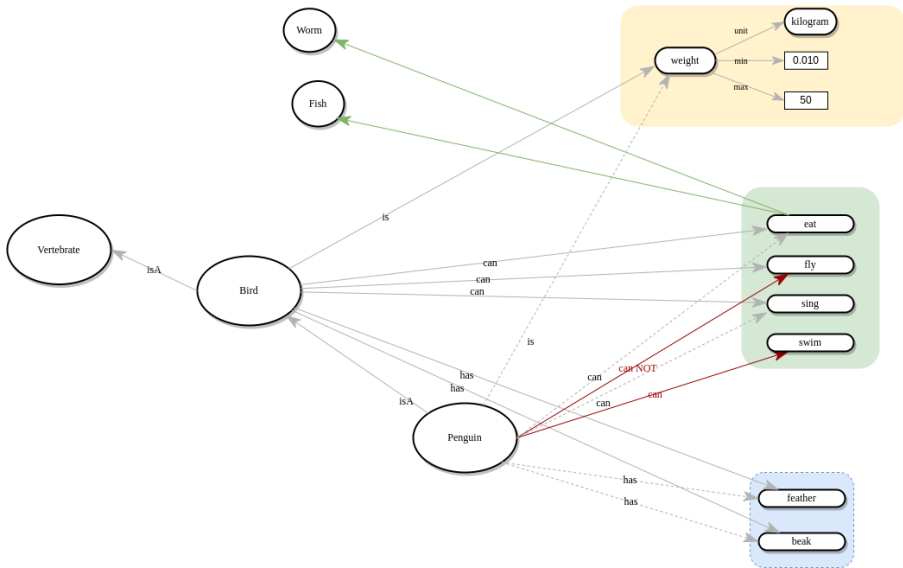


Fig. 4 In our context, we represent concepts as a hierarchical data structure. Concepts are anchored in ~~an~~ input/output, i.e., stimulus/response, ~~framework, say, which might consist of~~ sensorimotor feature spaces (colored regions) corresponding, for example, to different sensor modalities. Inherited features (e.g., the penguin “is-a” bird and thus inherits the features of a bird) are ~~in-shown with~~ dotted lines, while red lines represent overwritten values (e.g., a penguin can also swim but cannot fly). Green arrows point toward concepts that are themselves attributes of other concept features, accounting for inter-concept relationships. Values are completed by meta-information, ~~that is~~ not explicitly manipulated by the agent, ~~that is~~ used for ~~the~~-process specification or interpretation (e.g., the weight unit and bounds).

In other words, we follow [26], with the simple idea that an individual resource can be defined by “feature dimensions,” ~~i.e.,~~ attributes with some typed value. For instance, a bird could be ~~the following~~⁷:

⁷The ~~used~~-syntax ~~used~~ is a weak form of the JSON syntax: <https://line.gitlabpages.inria.fr/aide-group/wjson>.

```
bird: {
  is_a: vertebrate
  can: { sing fly eat: { worm fish } }
  has: { feather beak }
  is: { weight : { min: 0.010 max: 50 unit: kilogram } }
},
```

with some exceptions like penguins:

```
penguin: {
  is_a: bird
  can: { fly: false walk }
}.
```

and equivalently, this could be decomposed⁸ in terms of triples such as `bird is_a vertebrate` and so on.

Here, we choose the general approach of semantic knowledge representation as using a hierarchical taxonomy (`is-a`), with capability features (`can`), including in relation with those related to other resources, extrinsic features (`has`), and intrinsic features (`is`) [28]. This illustrative example is sufficient to allow us to detail the main characteristics of our representation. Some features are properties, and others are relations. A property can be qualitative, e.g., the is-covered-by property takes a value in an enumeration (e.g., {`sing`, `fly`}), or quantitative (e.g., the weight). The features can be hierarchical, either because the value is an enumeration (e.g., `can`) or because the value has some features (e.g., `weight`).

Such a data structure defines a “concept” in the sense of [26] (e.g., “a bird”), which is both a convex region of the state space (e.g., the region of all birds) and a prototype: Each feature has a default value, and this also defines a prototype (e.g., a typical, i.e., prototypical, bird). It corresponds to the third cognitive memory architecture, as proposed by [15], and we are now going to discuss how to implement it in a VSA framework.

Let us now review how such a cognitive symbolic data structure can be implemented, in a biologically plausible way, ~~thanks to~~ using the proposed VSA framework.

3.3 Biologically plausible implementation

To proceed, we thus have to consider a set of items of knowledge of the form of shown in Fig. 3, bounded to a triple of vectors $\{(\mathbf{s}_1, \mathbf{p}_1, \mathbf{o}_1), \dots (\mathbf{s}_N, \mathbf{p}_N, \mathbf{o}_N)\}$. This corresponds to relations, which are the basis of semantic information.

⁸Conceptually, considering a distributed version of this hierarchical specification by enumerating each node and directed edges is rather obvious; at the concrete implementation level, some details have to be carefully taken into account and the so-called turtle specification takes care of proposing a well-defined one-to-one correspondence (with an open-source implementation). ~~this being-It is~~ not necessary to describe this in detail here.

Container	VSA mechanism	Cognitive usage	Main available o perations
Set	Bundling or superposition		+ Element insertion/modification + Check membership - No enumeration
Map or D ictionary	Binding superposition	Associative memory	+ Element insertion/modification + Value-s retrieval from key + Key-s retrieval from value + Exact symbol recovery from approximate input - No enumeration
Indexed and C hained list	Ordinal binding superposition	Sequential memory	+ Element insertion/modification + Value-s enumeration
Relational map	(see next subsection)	Hierarchical memory	+ Element insertion/modification - No E numeration

Table 1 Biologically plausible data containers; see Appendix B for details.

3.3.1 Using bundling and binding to store information

As developed in detail in Appendix B and summarized in Table 1, the VSA formalism ~~allows~~ makes it possible to implement different data structures corresponding to different memory architectures, as defined by [15], and to different programming containers. The key point of the present work is, ~~one~~ the one hand, to verify that these VSA mechanisms generalize to modal symbol encoding, which is technically obvious (see Appendix B) but worthwhile to mention, and, on the other hand, to show that ~~it-they~~ also very easily generalizes to the so-called “relational map,” as developed in the next subsection. Further details on a scalable biologically plausible knowledge representation, can be found in [10].

3.3.2 Relational maps

How should we ~~to~~ store hierarchical information in a biological data structure? We propose ~~to~~ considering a distributed representation, a kind of “triple store” using superposition, i.e., bundling, and a few binding operations. The most natural choice might be to consider the triple set \mathcal{T} grounded to a vector \mathbf{t} :

$$\mathbf{t}_{\text{pso}} \stackrel{\text{def}}{=} \sum_i \mathbf{B}_{\mathbf{p}_i} \mathbf{B}_{\mathbf{s}_i} \mathbf{o}_i.$$

Here,

- ~~quantities~~ \mathbf{o}_i are represents vectors encoding symbols;

- quantities of the form \mathbf{B}_y are binding matrices ~~allowing that make it possible~~ to create a $\mathbf{B}_y \mathbf{x}$ key-value pair (~~the~~ key is y and ~~the~~ value is x), as defined in equation (A1) of Appendix A; ~~while~~ ~~+ their combination creating~~ ~~combining these quantities to create~~ a vector $\mathbf{B}_{p_i} \mathbf{B}_{s_i} \mathbf{o}_i$ ~~allows makes it possible~~ to encode a triple as a new random vector, i.e., as a neuronal ensemble activity; ~~and~~ ~~and~~ finally, the sum ~~allows makes it possible~~ to superpose a different triple. A reader not familiar with ~~such this~~ VSA formalism will find in Appendix B a didactic introduction, while the choice of the binding operator ~~from~~ among several available ~~binding operators~~ [16] is discussed in Appendix A. For this section to be self-contained, one just has to consider that binding ~~allows makes it possible~~ to create a key-value symbol pair, while ~~the~~ unbinding operation, which ~~writes is written as~~ $\mathbf{B}_y \sim$, ~~makes it possible allows~~ to retrieve the value from the key, as discussed ~~now below~~.

This design corresponds to a nested associative map: Each property p_j is defined through a mapping between subjects and objects, ~~verifying such that~~ ~~verifies this~~ property:

$$\mathbf{t}_{p_j} \stackrel{\text{def}}{=} \sum_{i, p_j = p_i} \mathbf{B}_{s_i} \mathbf{o}_i,$$

with the property⁹ that: $\mathbf{B}_{p_j} \sim \mathbf{s} = \mathbf{t}_{p_j}$.

Given a triple (s_0, p_0, o_0) , it is ~~obvious straightforward~~ to verify to what extent it is stored in the relational map ~~by through~~ unbinding:

$$(B_{s_0} \sim B_{p_0} \sim \mathbf{s} \cdot \mathbf{o}_0),$$

and this obviously generalizes to a triple multiplied by a modal τ value.

We can also further obtain all objects of a given subject for a given property:

$$\mathbf{t}_{p_j, s_j} \stackrel{\text{def}}{=} \sum_{p_j = p_i, s_j = s_i} \mathbf{o}_i \simeq B_{s_j} \sim \mathbf{o}_{p_j} \mathbf{t}_{p_{so}},$$

using the notation of Appendix A. ~~We~~ ~~can also easily define~~:

$$\mathbf{t}_{p_j, o_j} \stackrel{\text{def}}{=} \sum_{p_j = p_i, o_j = o_i} \mathbf{s}_i \simeq B_{s_j} \sim \mathbf{B} \leftrightarrow B_{p_j} \sim \mathbf{t}_{p_{so}}.$$

A dual construction, $\mathbf{t}_{s_{po}} \stackrel{\text{def}}{=} \sum_i \mathbf{B}_{s_i} \mathbf{B}_{p_i} \mathbf{o}_i$, with similar decoding formulae, ~~allows makes it possible~~ to further access the properties of a given subject $\mathbf{t}_{s_j} \stackrel{\text{def}}{=} \sum_{i, s_j = s_i} \mathbf{B}_{p_i} \mathbf{o}_i$ ~~or the properties of a given subject-object couple~~ $\mathbf{t}_{s_j, o_j} \stackrel{\text{def}}{=} \sum_{i, s_j = s_i, o_j = o_i} \mathbf{p}_i$ ~~using similar formulae~~. A key point is that, ~~up to our best to the best of our~~ knowledge, there is no operation to recover \mathbf{t}_{s_j} or \mathbf{t}_{s_j, o_j} from $\mathbf{t}_{p_{so}}$, and no operation to recover \mathbf{t}_{p_j} or \mathbf{t}_{p_j, o_j} from $\mathbf{t}_{s_{po}}$. This is an important constraint, and it would be interesting to verify if such a constraint is observed at the level of the brain's semantic memory.

⁹Obviously:

$$\mathbf{B}_{p_j} \sim \mathbf{s} = \sum_{i, p_j \neq p_i} B_{p_j} \sim \mathbf{t}_{p_i} + \sum_{i, p_j = p_i} B_{p_j} \sim \mathbf{B}_{p_i} \mathbf{B}_{s_i} \mathbf{o}_i \simeq \mathbf{t}_{p_j}.$$

The first term ~~vanishing vanishes~~ because $\mathbf{t}_{p_i}, p_j \neq p_i$, is generically orthogonal to \mathbf{p}_j , while the second reduces to \mathbf{t}_{p_j} thanks to the approximate inverse property.

Moreover, in order to enumerate the different elements of these maps \mathbf{t}_{\bullet} , we need the corresponding indexing mechanisms discussed ~~before~~ previously. If the basic operation is to enumerate all triples, with order constraints, then the choice of the ~~store-storage~~ architecture is not crucial; ~~this is going to be the case in the sequel~~ later in this paper.

~~A~~ To take this a step further, we can also consider an additional symbol “something,” and each time a triplet $(s_i, \mathbf{p}_i, \mathbf{o}_i)$ is added, we can also adding $(\sigma, \mathbf{p}_I, \mathbf{o}_i)$, $(s_i, \sigma, \mathbf{o}_i)$, and $(s_i, \mathbf{p}_I, \sigma)$. This ~~allows retrieving~~ makes it possible to retrieve the fact that there is a link between the predicate and object, subject and object, and subject and predicate, without requiring ~~enumerating~~ the enumeration of the different elements, as previously discussed.

At the cognitive level, this corresponds to cognitive maps ~~in interaction~~ interacting with each other and is a proposal to formalize the notion of hierarchical memory organization, as discussed ~~by~~ in, e.g., [15].

At the computer programming level, this corresponds to a “triple store” used in ontology reasoners and is in fact a distributed representation of an oriented graph, in the form of an adjacency set for ~~the~~ \mathbf{t}_{spo} construction and a hierarchical edge set for ~~the~~ \mathbf{t}_{pso} construction.

At this stage, we have reviewed and ~~completed~~ developed the ~~Vector Symbolic Architecture (VSA)~~ elements needed for the next step, in particular the notion of a relational map, which is a combination of associative maps, enumerated by an indexed list ~~and that~~ builds on the bundling structure reviewed at the beginning of ~~the~~ this section. We are now going to explain how symbolic computation can be performed, by making ~~interacting~~ such memory structures interact.

4 Knowledge transformation encoding

Let us now consider how to define operators ~~allowing that make it possible~~ to enrich the memorized information, ~~using~~ biologically plausible transformations.

4.1 Considering symbolic operations

We are going to consider symbolic operations as entailment rules of deductive, inductive, or abductive inference. It is worthwhile to note that both induction, as, for instance, in [30], and abduction could be formalized by inference rules (see, e.g., [31] for a computational example and [32] for a contribution regarding temporal reasoning). Our position~~ing~~ is that reasoning in the brain is mainly related to the construction of mental models, as discussed in, e.g., [33], considering common-sense reasoning. Such a framework better corresponds to human reasoning than pure logic reasoning, including modal logic reasoning [34]. In any case, a mechanism to construct, enrich, and modify such a mental model must be described.

Here, we focus on rule-based reasoning. Regarding an ontology such as OWL, the RDFS layer is purely rule-based, and several property inferences

such as the inverse, symmetry, transitivity, and rules of equating objects are rule-based. **Author query: You may want to write out what “OWL” stands for since this is the first time it appears in the paper.** ~~†~~This corresponds to the so-called OWL 2 RL profile. **Author query: You may want to write out what “RL” stands for since this is the first time it appears in the paper.** It allows scalable reasoning without sacrificing too much expressive power and can be implemented using rule-based reasoning mechanisms¹⁰. ~~Carefully-By~~ carefully considering the related entailment rules of this specification, we have observed that they require either one, two, or at most three premises, which is an important aspect in our context. Beyond this, other rule-based mechanisms can complete the description of rule-based reasoning, such as by adding function-free Horn rules (e.g., using the SWRL¹¹ language with suitable conditions [35]). **Author query: You may want to write out what “SWRL” stands for since this is the first time it appears in the paper.** This is what we target here.

Beyond deductive reasoning, following [30], for instance, or using inductive logic programming¹², inductive reasoning is easily formulated using the rule-based mechanism provided; ~~7~~ we can add counting operations in order to perform enumerative induction. This means that we will have rules ~~which-that~~ are going to scan the whole relational map, but we are going to observe that this is also required for other rules. This being stated, we will not further develop this aspect, and it is only mentioned as an interesting perspective of the present work.

~~A-Taking this a~~ A-Taking this a step further, rule-based abduction has been formalized, following, e.g., [31], as in fact a set of possible causes of a given observation, ~~7~~ that is deduced (to make short a more sophisticated story). In this rather restricted but still powerful context, the rule-based mechanism does not “invent” the cause, ~~7~~ but does infer ~~s~~ parameterized predefined causes, performing what could be called model-based abduction.

To summarize, our design choice is to propose a rather generic mechanism, specific enough to ~~enjoy-produce~~ enjoy-produce a robust implementation, ~~7~~ and ~~allowing making it possible~~ allowing making it possible to specify what has been reviewed here.

4.2 A generic mechanism

In order to proceed, let us write:

$$(\$s_0, \tau_0 \$p_0, \$o_0.),$$

which represents the fact that a variable subject $\$s_0$ is associated ~~to-with~~ to a variable object $\$o_0$ in an associative table related to the predicate variable $\$p_0$ with a level of belief τ_0 , i.e., ~~that~~ in the triple store, ~~there is:~~

$$t_{pso} = \tau_0 B_{\$p_0} B_{\$s_0} \$o_0 + \dots,$$

¹⁰See, for instance, https://www.w3.org/TR/owl2-profiles/#OWL_2_RL for details.

¹¹https://en.wikipedia.org/wiki/Semantic_Web_Rule_Language

¹²https://en.wikipedia.org/wiki/Inductive_logic_programming

where we use the \$ prefix to make explicit the fact that it-a given symbol is a variable symbol.

We are going to consider ~~the~~ entailment rules of the form:

$$\bigwedge_{i=1}^{i=I} (\$s_i, \$p_i, \$o_i) \rightarrow \bigoplus_{j=1}^{j=J} (\$s_j, \tau_j(\cdot) \$p_j, \$o_j),$$

where:

- The left-hand-side expression corresponds to a conjunction of premises, with-the $\$x_i$ ($\$x$ stands for the subject, predicate, or object) receiving the corresponding item as input.
- The weight $\tau_j(\cdot)$ is a function of left-hand-side elements; ~~equaling it is~~ equal to 0 if the rule does not apply to the element, it is negative if the consequence is to partially delete an element, and it is positive otherwise.
- Each right-hand-side value $\$x_j$ ($\$x$ stands for the subject, predicate, or object) is either a constant or equals a left-hand-side element $\$x_i$, or it could be a more complex expression of left-hand-side elements.
- The right-hand-side expression corresponds to $t_{ps_o} += \sum_j \tau_j(\cdot) B_{\$p_j} B_{\$s_j} B_{\$o}$, i.e., ~~to update it updates~~ the related structure. In almost all the cases that we have considered, only one triple is generated; this part of the setup is thus stated for future use.

This includes modifying, (including deleting), existing triplets since they are weighted by a τ value that can be modified, and thus set to 0. The key point is that the calculation of $\tau_j(\cdot)$ allows the integration of not only ~~of~~ modal logical formulae but also threshold mechanisms or counting operations that are required for induction.

Such a general setting clearly corresponds to the usual production rules. In the binary mode, e.g., if $\tau \in \{0, 1\}$, this is nothing but a specification equivalent to Horn clauses. Introducing the calculation of $\tau_j(\cdot)$ allows us to go beyond this, but at the ~~price-cost~~ of lowering the decidability and implementation tractability, as detailed ~~in the sequel later in this paper~~, so we will need to consider suitable convergence conditions.

At the implementation level, the application of such rules can simply be implemented by-using feedback connections between an iterator over the triples of the relational map, as discussed in detail for a typical example, before ~~describing~~ a general implementation is described. **Author query: The meaning of the phrase “...feedback connections between an iterator over the triples of the relational map,” is not clear. Between the iterator and what?**

4.3 Class inheritance entailment rule

In order to better understand what is proposed here, let us start by detailing an illustrative and quite universal example.

4.3.1 Class inheritance as a major deductive mechanism

The most common entailment rule is likely the class inheritance rule, which states that “if a subject belongs to a class, and if this class is a subset of a

superclass, then the subject belongs to the superclass,” e.g., “if Tom is a cat, and cats are animals, then Tom is an animal.” This is a deductive rule, a particular syllogism, which belongs to common-sense reasoning and is well understood as soon as formal reasoning emerges in children [21].

The notion of “class” corresponds to Boolean properties (e.g., if you are alive, you belong to the class of living organisms) and capabilities (e.g., if you can fly, you belong to the class of flying organisms). It is associated with a given individual feature. In addition, such Boolean features are organized in a hierarchy, leading to a taxonomy that describes the information about the considered individuals.

At the syntax level, it is based on two predicates, say namely, is a, to state which states that an individual belongs to a class, and say are, to state which states that all individuals of a class are also individuals of a more general class. This corresponds They correspond, respectively, to **rdf:type** and **rdfs:subClassOf**, when using the RDFS vocabulary is used, as discussed in Appendix C.

Implementing The implementation of the class inheritance rule using the VSA has already been successfully developed and numerically experimented tested in by [13] using the Nengo simulator [36], while we propose here to rely on this work to propose a more general mechanism. In Appendix C, we make explicit the fact that the mechanisms of a commonly used Semantic Web semantic web modeling language, the RDFS (Resource Description Framework Schema) (RDFS), can have its main mechanisms be implemented in such a numerical framework.

4.3.2 Design of the inference rule

The class inheritance rule writes can be written in its “binary equality” form as follows:

$$(\$s, \text{is_a}, \$c_1.) \wedge (\$c_1, \text{are}, \$c_2.) \Rightarrow (\$s, \text{is_a}, \$c_2.),$$

that which states that if any subject $\$s$ belongs to the class $\$c_1$, i.e., is of “type” $\$c_1$, and this class $\$c_1$ is a subclass of $\$c_2$, then $\$s$ also belongs to the class $\$c_2$. Here, the equality is either verified or not; we are in the exact reasoning case.

When considering an approximate form, as schematized in Fig. 5, the rule writes is written as follows:

$$(\$s_1, \$p_1, \$o_1.) \wedge (\$s_2, \$p_2, \$o_2.) \Rightarrow (\$s_1, \tau \text{ is_a}, \$o_2.),$$

with:

$$\tau \stackrel{\text{def}}{=} (\$p_1 \cdot \text{is_a}) \ \& \ (\$o_1 \cdot \$s_2) \ \& \ (\$p_2 \cdot \text{are}),$$

so that τ equals 0 unless the three equalities are at least partially verified. If τ equals to 0 nothing is inferred, whereas if higher it is greater than 0, a new triple is output. It is a forward schema in the sense that given some input data, a new result is inferred. This second form allows approximate correspondences, e.g., $\$p_1$ to can be approximately is a.

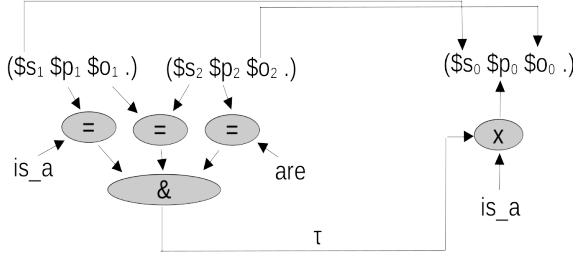


Fig. 5 The forward implementation of the class inheritance entailment rule.

H-The rule requires two numerical operators:

- The similarity operator $(\mathbf{x} \cdot \mathbf{y})$, and
- a numerical implementation of the conjunction $(\mathbf{x} \ \& \ \cdots \ \& \ \mathbf{y})$ operator, which is equals to:

- 0 as soon if one operand is equal to or less than 0,
- 1 if all operands are equal to 1,

with intermediate values if the operands are between 0 and 1. This could be a min operator in coherence accordance with the numeric modal logic implementation, as discussed in [25], and this is our proposal what we propose here;; which it is easily implementable¹³ at a neuronal level [37], while any other T-norm¹⁴ would be suitable.

4.3.3 Forward versus backward inference

This previous setup corresponds to a forward application of the rule, i.e., given two left-hand-side triples, it describes how to numerically calculate the right-hand side.

In its backward form, given the right-hand side, we calculate:

$$\tau \stackrel{\text{def}}{=} (\$s_0 \cdot \$s_1) \ \& \ (\$o_0 \cdot \$o_2) \ \& \ (\$p_0 \cdot \text{is_a}) \\ \& \ (\$p_1 \cdot \text{is_a}) \ \& \ (\$o_1 \cdot \$s_2) \ \& \ (\$p_2 \cdot \text{are}),$$

i.e., given a right-hand-side input, we evaluate to what extent two left-hand-side triples could correspond to what is expected. See, e.g., [38] for an overview of this forward versus backward duality.

In the brain, as discussed, for instance, in [39] at a computational level and [40] at a more conceptual level, both forward and backward mechanisms are mixed in cognitive behaviors, as studied, for instance, in [41] at a more experimental level. In a nutshell, entailment rules are applied both “on query” when, given a goal-directed behavior, some information is required and at a

¹³To be very precise, the max operator biological implementation is often discussed, which it is obviously equivalent to a max up to linear transform: **Author query: The meaning of “obviously equivalent to a max up to linear transform” is not clear. Do you mean “obviously equivalent to a linear transformation of a min operator”?**

$\min(x_1, \dots, x_n) = 1 - \max(1 - x_1, \dots, 1 - x_n)$, with $x_i \in [0, 1]$.

¹⁴<https://en.wikipedia.org/wiki/T-norm>

more data-driven level, which here is the stimulus-driven level. ~~In the sequel~~ In this paper, we are going to focus on forward inference, while in [13], a backward mechanism based on the VSA, for a specific couple of inference rules, ~~has been~~ was proposed and validated.

4.3.4 Exact versus approximate inference

This rule thus considers approximate similarity: If the input triple states that $\$s_1$ is of class $\$o_1$, only approximately, with $\$p_1 = \tau$ is a, ~~with where~~ $\tau \in [0, 1]$, and similar approximate equalities, while the similarity operator outputs a value between 0 and 1 and the conjunction operators interpolate values between 0 and 1 given the input, we will obtain as the output an appropriate predicate value τ' is a, with $\tau' \in [0, 1]$, which corresponds to the fact that the predicate is only approximately true.

It is interesting to note that in the case of exact inference (i.e., $\$p_1 = \text{is_a}$, $\$o_1 = \s_2 , $\$p_2 = \text{are}$), we obtain $\tau = 1$ as expected.

4.3.5 Implementation of negative inference

~~A step further, t~~This, however, does not immediately generalize to negative inference (i.e., the fact that it approximately does *not* belong to a given class, which is beyond the RDFS specification but present at the OWL level). ~~, but still:~~ In ~~that this~~ case, a symmetric rule:

$$(\$s, \neg \text{is_a}, \$c_1.) \wedge (\$c_2, \text{are}, \$c_2.) \Rightarrow (\$s, \neg \text{is_a}, \$c_2.)$$

has to be considered, with ~~totally a~~ similar development. The key point, ~~here,~~ is that we must only consider the positive part of the similarity, i.e., ~~combine~~ it is combined with a rectification operator.

4.3.6 Implementation of the inference rule

At the implementation level, what is schematized in Fig. 5, is a simple analog calculation of a τ value given some left-hand-side ~~size~~ input, while the backward implementation is a calculation that also involves ~~also~~ the right-hand-side values: It is no more than a weighted connection between triples, i.e., between memory elements, with the generation of a new triple in the forward case and an evaluation of the queried triple in the backward case.

Since we consider approximate inference, and contrary to the binary equality case¹⁵, we cannot select triples that exactly match a given pattern. In the

¹⁵ In the binary equality case, let us illustrate our purpose by explicitizing, how to generate the entailment rule closure, i.e., all deducible triples, given this rule and a set of input triples. An efficient mechanism must be able to select the pertinent triples, i.e., use a select operator, e.g., ~~thanks to~~ based on associative tables, indexed by property and subject. Given two constant values c_i , and c_j ,

```
for all ( $\$s_i$   $\$p_i$   $\$o_i.$ ),  $\$p_i = c_i$  &  $\$s_i = c_j$  do
  ..../..
end for
```

can be enumerated directly without scanning all triples ~~but by scanning~~ only the one to be selected. This will allow us to incrementally calculate the closure ~~as follows~~, given a “closed” set of triples

forward mode, we must iteratively calculate ~~the~~ $\tau(\cdot)$ on the whole conjunction of left-hand-side triples, generating new triples when $\tau(\cdot)$ is above a given threshold, as discussed in the general case ~~now~~ below.

4.4 Entailment rules in the general case

Back-Going back to the general case, let us now describe an iterative closure algorithm, as schematized in Fig. 6, and discuss its properties.

The key idea is that the relational map is a two-level hierarchy with:

- a known triples map storing all existing facts and all their derivations given a fixed set of entailment rules, and
- a new triples buffer that is input or feedback ~~by-for~~ entailment rules and for which derivations must still be done.

This is a very common setup used to calculate a fixed point or a closure.

For each new triple, we have to:

- check ~~that~~ if it is already a known triple; ~~;~~ in ~~that~~ this case, the job is done;
- otherwise, ~~;~~
- for each entailment rule,
 - for each left-hand-side premise instantiated by the new triple,
 - for each known triplet set matching the other premises,
 - we have to calculate the related $\tau(\cdot)$ value and generate the corresponding new triple or triples; and
 - move this new triple from the new triple buffer to the known fact map.

At the initialization stage, the known triple map is empty, and the new triple buffer waits for an input.

This basic mechanism raises several issues.

4.4.1 Comparing a new triple with a known existing triple

Given the relation map data structure, it is obvious ~~by-that~~ unbinding can be used to verify if the new triple has a similarity with a known triple of the relation map, as made explicit in subsection 3.3.2. This similarity corresponds to the product of the degree of belief of each triple ~~multiplied-by-and~~ the cosine similarity.

$\{(\$s_i \$p_i \$o_i.) \dots\}$ with all possible triplets generated and a new triple $(\$s_0 \$p_0 \$o_0.)$ to be added. ~~;~~ The standard algorithm writes-is written as follows (see, e.g., [38] for an introduction):

```

input A new triple  $(\$s_0 \$p_0 \$o_0.)$  and a closed set of triples  $\{(\$s_i \$p_i \$o_i.) \dots\}$ .
let  $\{(\$s_0 \$p_0 \$o_0.)\}$  be an "open" triple set, initialized with the new triple inside.
repeat
  pull a triple  $(\$s_0 \$p_0 \$o_0.)$  from the open triple set.
  if  $\$p_0 = \text{is\_a}$  then
    for all  $(\$s_i \$p_i \$o_i.), \$p_i = \text{are} \ \& \ \$s_i = \$o_0$ , in the closed triple set do
      add  $(\$s_0 \text{ is\_a } \$o_i.)$  to the open triple set.
    end for
  else if  $\$p_0 = \text{are}$  then
    for all  $(\$s_i \$p_i \$o_i.), \$p_i = \text{is\_a} \ \& \ \$o_i = \$s_0$ , in the closed triple set do
      add  $(\$s_i \text{ is\_a } \$o_0.)$  to the open triple set.
    end for
  end if
  add the triple  $(\$s_0 \$p_0 \$o_0.)$  to the closed triple set.
until the open triple set is empty.
```

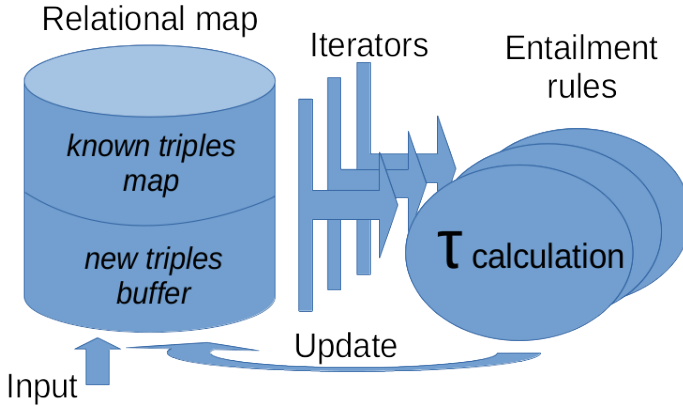


Fig. 6 The entailment rules feedback mechanism in the general case; see [the](#) text for details.

Thanks to the development of a statistical test in section 2 and [the](#) easy first-order estimation of the noise level, [which is as](#) made explicit in Appendix D at the implementation level, we have available, given a probability threshold, a mechanism to decide if the statistical difference [of between the](#) degrees of belief [between of](#) the new and known triples is negligible or not.

If it is not negligible, [this it](#) means that both triples are similar and:

- either the new triple enforces the degree of belief of the known one, up [to](#) the upper bound, which is 1, corresponding to the “true” value,
- or it decreases, potentially [down to contradicting to the point that it contradicts](#) the previous degree of belief (setting the value [down to -1, which corresponds corresponding](#) to the “false” value, or setting it to 0, [which corresponds corresponding](#) to an “unknown” state).

This is very interesting but it has a drawback: [it](#) may generate an unstable state. [Let us explain:](#) In the binary case, the only degree of belief value is 1, [and](#) thus the only possibility is that the new triple equals the known one, [and meaning that](#) we can avoid redoing the inference. At a more general level, if the only possibility is that the new triple enforces the degree of belief of the known one, we still have to redo inferences, but the system is monotonic, in the sense that [the](#) value can only increase and is bounded, [and](#) thus [the values](#) must converge. A simple way to implement a monotonic mechanism is to take into account the new triple; only if its level of belief value is higher than the previously known triple’s belief value.

At this stage, it is really a semantic alternative, to be chosen at the application level, depending on whether we design a cumulative knowledge mechanism in a stable universe or an adaptive mechanism in a changing knowledge environment. Both are possible at this implementation level, while developing this point is beyond the scope of our purpose.

4.4.2 Algorithmic complexity for a given new triple

For one- step, given T inputs, R entailment rules of arity I (i.e., number of left-hand-side premises) $-I$, and S known facts, these nested loops generate $O(T R I S^{I-1})$ calculations of $\tau(\cdot)$. Fortunately, for most of the rules, we have $I = 2$ or even $I = 1$, except for some fragments of OWL-RL entailment rules¹⁶, which-with often have $I = 3$, so that the complexity is mainly linear with respect to the relational map, and seldom quadratic. Furthermore, beyond a sequential system, in our case, we have a completely distributed setup, so that such an operation set is obvious-simple to implement in parallel at a certain stage:

- Each new triple must be treated in sequence, in order to work with a stable known triples map;
- However, given a new triple, the rule enumeration, the left-hand-side assignation, and the known triple enumeration on other premises' locations can all be performed in parallel, generating new triples in any order.

4.4.3 Global algorithmic time and space complexity size.

The global algorithmic complexity-complexities in time and space are linked, since each step generates one new triple, while these operations are easily implemented in a distributed framework.

This complexity is highly dependent on each rule; for instance, it is easy to verify that a reflexivity property generates a closure whose order of magnitude is linear with respect to the input triple, a symmetry property generates a closure whose order of magnitude is quadratic with respect to the input triple, and transitivity generates a closure equivalent to the number of the-paths in an oriented graph, which-thus meaning that the complexity can be exponential for a fully connected graph.

Regarding the space complexity, an important aspect is the-“catastrophic forgetting,” i.e., what happens if too-too many random vectors are superposed in the relation map, meaning-so that approximate orthogonality cannot be guaranteed. This has been numerically studied in, e.g., [16], showing on the one hand that the VT mechanism we have chosen outperforms other representations based on the approximate inverse, and it outperforms other dense methods (but not the sparse one) in terms of the superposition capability. **Author query: You may want to write out what “VT” stands for since this is the first time it appears in the paper.** The order of magnitude is rather low, with about 30 symbols reaching 99% accuracy in a space of dimension 1000, while the best sparse methods have still-limited performances of about 50 symbols, for the same dimension.

As it-is, the proposed mechanism is really-limited to toy applications; we have two tracks to improve. On the one hand, when randomly drawing

¹⁶The situation is even more complex because some entailment rules require managing-the management of a variable arity depending, for instance, on the data structure length. See [42] for a detailed discussion. What we proposed here is thus directly applicable to the RDFS language layer and opportunistically to some OWL 2 reasoning mechanism.

almost orthogonal symbols, it is always possible to use, ~~say~~ for instance, the modified Gram-Schmidt methods to improve the stability and obtain a precise orthogonality, with the benefit ~~to improve of improving the~~ overall stability. On N symbols, the complexity is quadratic in the number of symbols, i.e., ~~in it is~~ $O(dN^2)$. On a sphere of dimension d , we can obviously ~~can~~ draw d orthogonal symbols.

On the other hand, at the microscopic level, ~~on in~~ a real spiking neuronal network, the dimension is several orders higher; ~~a~~ neuronal map typically corresponds to a state space of dimension 10^5 , which is not an order of magnitude higher. **Author query: The meaning of the previous sentence is not completely clear. An order of magnitude higher than what?** Extrapolating the linear approximation ~~for the of~~ [16], which is shown in Fig. 4 of ~~their that~~ paper, we obtain a bit more than 3000 symbols for $d = 10^5$ and a bit more than 30000 symbols for $d = 10^6$.

This is of course impossible to calculate ~~at on~~ a mesoscopic scale by manipulating symbols of such a dimension, but we are going to introduce the idea ~~to simulate of simulating~~ the VSA mechanism at a macroscopic level, i.e., directly manipulating the algebraic symbols and predicting the cumulative noise of the result. We make explicit this alternative method in Appendix D and share a preliminary implementation of such a mechanism, ~~which that~~ is no ~~more longer~~ limited by the vector space dimension.

5 A preliminary experiment~~ation~~

To illustrate the use of this mechanism, we reconsider the example proposed in [13], and show a minimal ontology in Fig. 7, limiting this preliminary experiment~~ation~~ to deductive rules, and using the RDFs vocabulary. **Author query: In the previous sentence, is “RDFs” meant to be “RDFS”?** This allows us to better compare the two approaches. Here, the class inheritance *rdfs9* ~~detailed rules described~~ in this paper, and the *rdfs2* subject domain inference (*rdfs2*), ~~the~~ *rdfs3*-object range inference (*rdfs3*), and ~~the~~ property inheritance (*rdfs7*) entailment rules detailed in Appendix C are implemented (refer to this appendix explaining for an explanation of why this choice makes sense).

5.1 Macroscopic implementation of the VSA system

As discussed throughout this paper, we are in a position to propose an algorithmic ersatz of the usual VSA mesoscopic linear algebra calculations ~~between~~ involving high-dimensional random vectors, and this has been implemented and made available as ~~a~~ public documented open-source code¹⁷.

The For the generation of a symbol, at a given level of belief τ and for a given level of first-order random normal noise with a standard-deviation σ , ~~with~~ the usual similarity, bundling, and binding operations are made available,

¹⁷<https://line.gitlabpages.inria.fr/aide-group/macrovsa/index.html>

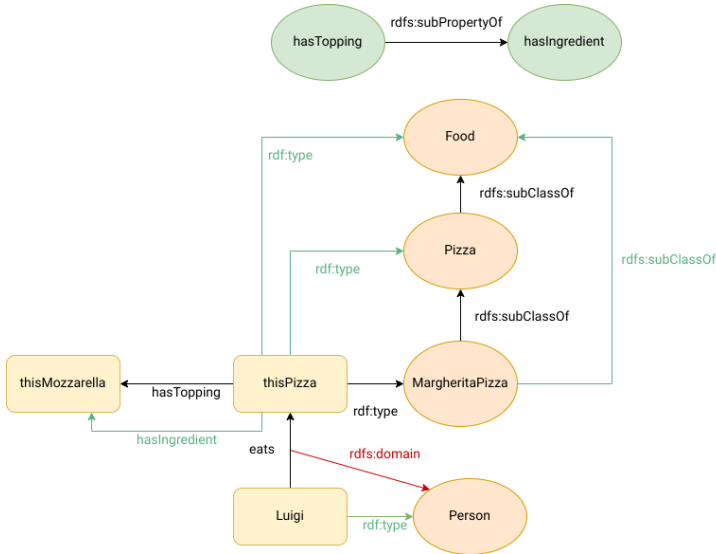


Fig. 7 An example of a simple ontology with three individuals. The black arrows correspond to factual statements input into the database and the green arrows correspond to inferred statements. Rectangular boxes stand for individuals, round boxes stand for classes, and properties are used to label arrows. Here, from the fact that a subject eats an object, we deduce that this subject is a Person, and the object is Food. This is enlightened illustrated by a red arrow. From the fact that the object is a Margherita pizza, which is a Pizza, which is a Food, from according to the class hierarchy, we deduce that the object is a Pizza, and re-deduce that it is a Food. Furthermore, because Luigi (among other activities, since it is an open world) eats a pizza, we deduce it that Luigi is a Person. Because of property heritage, meaning that here a Topping is an Ingredient, we also deduce from the fact that this pizza has mozzarella as a topping, and that it also has it-mozzarella as an ingredient. In the macroscopic implementation, this property is deduced from the fact that Margherita pizza has, say, always has mozzarella as a topping, allowing us to generate compounded inferences.

with some technical details about the noise calculation and software architecture given in Appendix D. We have considered a symbol encoding dimensions of $d = 256$ to fit-be consistent with usual-previous mesoscopic experimentation-experiments, such as those in [13]. This is tested with an associative map and relational map, as described in the previous section, and we have implemented the tiny Pizza experiment¹⁸, obtaining, in the simplest case, the expected closure, as given in Fig. 8.

More interesting-in this-result is what happens considering-when modality is considered, e.g.,

(Luigi 0.5 eats thisPizza).

¹⁸Source. The source code is available here-at https://gitlab.inria.fr/line/aide-group/macrovsa/-/blob/master/src/pizza_experiments.cpp and it is noticeable that the C/C++ implementation of such rules is really-straightforward to write, as documented in the source code. This piece of code output is available here-at https://gitlab.inria.fr/line/aide-group/macrovsa/-/raw/master/src/pizza_experiments.out.txt, in coherence-accordance with Fig. 8, showing-also-and it also shows the intermediate inference steps.

<i>Input triples</i>	<i>Inferred triples</i>
(Luigi eats thisPizza)	(Luigi rdf:type Person)
(thisPizza rdf:type MargheritaPizza)	(thisPizza rdf:type Pizza)
(MargheritaPizza rdfs:subClassOf Pizza)	(thisPizza rdf:type Food)
(Pizza rdfs:subClassOf Food)	(MargheritaPizza rdfs:subClassOf Food)
(eats rdfs:domain Person)	(thisPizza hasIngredient thisMozzarella)
(eats rdfs:range Food)	
(thisPizza hasTopping thisMozzarella)	
(hasTopping rdfs:subPropertyOf hasIngredient)	

Fig. 8 The expected inferences using the proposed RDFS subset of entailment rules—as obtained by the macroscopic algorithmic ersatz of [the](#) VSA implementation.

~~in~~ [In other](#) words, it is possible but not completely necessary that Luigi eats the given pizza. In that case,÷

- it is still possible but no ~~more~~ [longer](#) entirely true that Luigi is a person;
- it is still entirely true that this pizza is some food, even if Luigi did not eat it, because it is true that it is a pizza, which is food.÷

~~which~~ [This](#) is what is obtained by the implementation, as ~~readable on~~ [shown](#) by the open-source tiny experiment output.

Another interesting aspect is the calculation of the standard deviation of the level of noise modeling ~~what that~~ happens at the mesoscopic level for either [the](#) calculation of [the](#) similarity between two random vectors or unbinding operations. On [the](#) one hand, our macroscopic model is ~~in coherence~~ [consistent](#) with [that of](#) [16], which obtains [the following](#) (from ~~the Fig. 4 of that~~ paper^{Fig.4}) for the VTB representation: **Author query: You may want to write out what “VTB” stands for since this is the first time it appears in the paper.**

$$d \gtrsim 32(s + 0.575)$$

~~as a~~ [represents the](#) minimal dimension d [needed](#) to obtain [a](#) 99% accuracy with ~~the a~~ bundling of size s , using [a](#) similarity calculation to extract vectors from the bundling. Our model does not take the negligible bias 0.575 into account but allows us to calibrate [the level of noise](#) to $\sigma \simeq \frac{0.016}{d}$ ~~the level of noise~~, in order to perform a simple ~~test of~~ z-score [test](#) under the normal hypothesis:÷

$$\tau > 2\sigma$$

to decide if the related τ value of the similarity is distinguishable from the noise. On the ~~reverse other~~ [hand](#), we ~~do not~~ consider ~~that~~ two vectors ~~are not to be~~ similar if the similarity is below the ~~noise~~ standard deviation [of the noise](#), preferring a more conservative threshold.

~~A~~ To take this a step further, we also implemented the $(1/d^{1/4})$ noise dependency for unbinding, with the same calibration, and it would be interesting to further investigate, at the mesoscopic level, the numerical precision of unbinding on associative maps; ~~whereas up to our best~~ to the best of our knowledge, ~~it is this was~~ not studied in the papers quoted here.

~~Far~~ Although it is far from being complete, this macroscopic implementation of an algorithmic ersatz of VSA mesoscopic operations seems sound and ~~coherent~~ it is consistent with previous results. It has a final non-negligible advantage: It is quite “simple” in the sense that it did not require very complicated or twisted mechanisms. ~~It requires~~ a bit more than 500 lines of formatted C++ code, including formal symbolic operations on the algebraic operators.

5.2 Comparison with a mesoscopic implementation

Let us now discuss how, considering the NEF methodology [9], as implemented in the Nengo platform [36], the previous mechanisms can be implemented at the mesoscopic level. Such an implementation has already been proposed in [13] for the class inheritance *rdfs9* rules detailed in this paper, as a question-answering system, and thus ~~working it works~~ in backward mode. In order ~~not to lengthen~~ to avoid lengthening the present paper, we do not ~~detail describe~~ again all the implementation mechanisms and expected results, but simply explicitize ~~the how to~~ how to implement these mechanisms for this new approach.

The present implementation works in forward mode, computing the fixed point of the inference loops. To this end, it is possible to create a Nengo vocabulary containing all the resources of our ontology encoded as vectors, and we stored asserted memberships and relationships between these resources into associative memories¹⁹, one for each predicate. For the sake of simplicity, we consider a simplified setup with $\tau = 1$ for each triple, while the use of approximate inference is also discussed in [13] ~~with for~~ a Nengo implementation, ~~also~~ using the VTB operator²⁰.

The RDFS entailment rules are implemented as feedback computations, ~~with the~~ retro-actions enriching the associative memory’s contents, as schematized in Fig. 9. We illustrate here the network connections for *rdfs9* entailments (in forward mode), but similar networks can be implemented for the other rules:-

- Two input cues are fed into the network (in this case, ~~they are~~ initialized as TYPE and SUB_CLASS_OF because these are the predicates involved ~~to~~ apply in applying the rule).
- Then, each associative memory stores a chained list of all known triples for a given predicate. A feedback connection ~~allows~~ makes it possible to use the last retrieved value as ~~a the~~ key ~~to for~~ the next value, ~~allowing making it possible~~ to enumerate the triples just like a relational map, as discussed previously.

¹⁹<https://www.nengo.ai/nengo-spa/examples/associative-memory.html>

²⁰~~Or~~ its transpose (TVTB) ~~can also be used~~ to manage left and right binding: https://www.nengo.ai/nengo-spa/modules/nengo-spa.algebras.html#nengo-spa.algebras.tvtb_algebra.TvtbAlgebra.

Retrieved triples are successively stored into a state²¹ which allows that makes it possible to pass data, in that this case, to enumerate each input triple, and thus acting it acts as the input triple buffer of Fig. 6.

- An action selection module²² (corresponding roughly to a conditional test) allows makes it possible to nest the enumeration of the second associative memory, by detecting the end of the enumeration of the first associative memory.
- Then, the rule itself is implemented by an interconnection similar to the one made explicit in Fig. 5, and another selection module acting as a threshold gate triggers the addition of a new triple if the rule is applicable.

Other entailment rules are implemented using similar networks networks, noticing that the associative memories are shared between them, while the enumeration mechanisms are generic, and only the inferred statement is specific to a given rules. The key point is that the whole computing process is distributed and “programmed” only by spatial inter-connections, as expected.

This shows that our macroscopic implementation can translate to what is proposed to represent biologically plausible processes at a level of representation that can be “compiled” as a set of interacting spiking neural networks [9].

6 Discussion and conclusion

6.0.1 Contributions

In this paper, we have been able to propose, up to the implementation level, a reformulation of the powerful VSA approach with a few add-ons additions:

- Explicitizing We explicitized a degree of belief for each knowledge item in link with that is linked to the possibility theory related to modal logic, and revisiting we revisited the main proposed abstraction of biologically plausible data structures to verify the their compatibility with this generalization, while comparing them with usual programming data structures, and discussing discussing how to efficiently scan (i.e., enumerate) such data structures.
- Proposing We proposed an implementation of hierarchical or relational semantic data structures within the VSA formalism, in link with relation to hierarchical cognitive memory, and allowing us to introduce also symbolic derivations.
- Suggesting We suggested a design of for a forward reasoning mechanism implemented via connectivity feedback on relational data structures, allowing making it possible to perform deductively, but also to some extent inductive and abductive inductions reasoning.
- Introduce We introduced the idea of simulating such a mechanism at a macroscopic, more symbolic level in order to have obtain computations independent of the VSA dimension space, thus allowing sealing making it possible to scale

²¹https://www.nengo.ai/nengo-spa/modules/nengo_spa.modules.html?highlight=state#nengo_spa.modules.State

²²https://www.nengo.ai/nengo-spa/modules/nengo_spa.html?highlight=actionselection#nengo_spa.ActionSelection

up such mechanisms. This idea has been applied to VTB- algebra but is also obviously reusable with other VSA ~~algebra~~-~~algebras~~.

6.0.2 On biological plausibility

We thus propose an anchoring, i.e., a numerical grounding of semantic information. This, indeed, does not mean that the brain-~~exactly~~ performs such operations ~~exactly~~, but this anchoring is biologically plausible in the sense that ontology rule-based algorithms can be implemented in a VSA-~~architecture~~, as developed in this paper, which itself is a model of spiking neuron assembly activity, as developed by [9], with the NEF approach reviewed ~~here in this paper~~.

More technically, we hypothesize that such a neuron assembly stores the ontological assertions (i.e., A-box elements) as a compounded semantic pointer (SPA), while derivations of new assertions (i.e., T-box elements) are implemented as parameterized feedback on this SPA. What we show here is that a general feedback mechanism implements such derivations, while a specific derivation rule corresponds to a specific connectivity of the feedback input layer. This model implies that new rules are coded by new synapses' long-term connections of feedback connections, thanks to repetitive learning, for instance, due to prefrontal-hippocampus interactions, as modeled by, e.g., [43]. **Author query: In the previous sentence, is “long-term connections of feedback connections” meant to be just “long-term feedback connections”?**

This model also implies that assertions are somehow enumerated off-line (in the sense of being scanned and replayed), to infer their consequences, corresponding to forward inference mechanisms (while we discussed a backward mechanism in [13], such a mechanism is also ~~in-coherence-consistent~~ with prefrontal-hippocampus interactions, as discussed in, e.g., [44]).

6.0.3 On numerical versus semantic grounding

~~On~~Numerical grounding, or anchoring, fundamentally differs from the semantic symbol grounding problem, as reviewed and discussed in [45], where symbols are linked to their meanings, and anchored in sensorimotor features, which involves the capacity to pick referents of concepts and a notion of consciousness. In a nutshell, this is still an open problem that we are not going to address here. However, proposals of methods to link abstract symbols to the neuronal reality, enrich the issue of how mental states can be meaningful. Furthermore, the fact that our abstract representation is anchored in sensorimotor features, means that it is also a link between symbols and their potential referents. ~~A~~To take this a step further, when we represent concepts, the chosen design choice associates prototypes, allowing us to anchor an abstract element to a concrete example.

Another aspect not targeted by the present study, is the emergence of symbols, i.e., the fact that a symbolic representation emerges from a biological or any physical system in interaction with its environment. This issue corresponds

to the ungrounding of concrete signs²³, as discussed in, e.g., [46], in link-with-relation-to the emergence of symbolic thinking (see, e.g., [47] for a detailed discussion). At the computational neuroscience level, the issue is addressed in [48] for a toy experiment; emphasizing that paper emphasizes that to address such an issue, we must avoid explicitly embedding any symbol anywhere in the model, a priori or a posteriori. Here, we do not address the emergence issue, but in a sense we do address a *feasibility* issue: †To what extent event-can sophisticated symbolic processing can be anchored in numerical processing, not only-just rudimentary operators? ‡ We also address an *interpretation* issue, i.e., we consider to what extent sub-symbolic sensorimotor anchored processing corresponds to symbolic processing, as discussed in the sequel later in this paper.

6.1 Approach limitations and perspectives

Although these contributions are of some theoretical and practical interest, it is clear that this is only a preliminary work, far from covering all the modeling and representative power of VSA approaches. It is rather limited regarding several aspects: Although implementing the implementation of inductive and abductive rule-based mechanisms has been evoked and argued-considered, the work is still to be done this implementation has not yet been developed. While the RDFS deductive mechanism can clearly be implemented in the present framework, more sophisticated segments of knowledge representation languages such as the OWL 2 family are must still to be explored to evaluate what can be implemented without extending the present mechanism. Last but not least, another limitation is the fact that, depending on the choice of the entailment rules, an unstable or unbounded number of triples could be generated. This latter caveat is can be easily avoided by using our knowledge of the different levels of specifications of semantic inferences, along with their related decidability and algorithm complexity properties.

The algorithm itself is a simple closure mechanism, whereas-although indeed, much more efficient reasoning algorithms are available, (such as, e.g., the so-called “tableau methods”); however, to the best of our knowledge, but up to our best understanding their link to biologically plausible mechanisms is really-improbable, whereas we here show that a simple set of feedback on iterators makes-does the job, which and is much closer to what is expected to happen in neural assemblies (see [9] for a general discussion).

The macroscopic simulator is operational but actually-only limited to the monotonic entailment rules used in forward mode. We did-have not yet implemented ed the capability to delete some rules that could be contradicted by new incoming facts, ‡ or backward reasoning. Both could be easily implemented as a direct extension of the proposed software package, as evoked-described in this

²³In the semiotic hierarchical meaning of an “icon” only-built only from sensori-motor features, structures at an “index” level built by concrete relationships between given objects, giving give rise to a “symbol” in the semiotic sense, corresponding-which corresponds to abstract general relationships between concrete concepts or sensori-motor features.

paper, but both also require ~~making one to make~~ the proper design choices in order for ~~such this~~ more general behavior to be coherent.

Following the usual VSA approaches, the symbolic information is embedded in a compact Riemannian manifold with a very simple topology, a hypersphere, and we have made explicit the fact that finally, the number of encodable symbols is rather limited. Other geometries may offer better performances, and ~~in~~ the particular hyperbolic embedding of hierarchical representations ~~makes a profit of benefits from~~ the fact that due to the hyperbolic negative curvature of the space, even ~~an~~ exponentially growing data structure can be parsimoniously represented [49], because of the expanding geometry (to make ~~short~~ a long story ~~short~~). The idea to embed ~~the~~ data representation in non-Euclidean spaces and especially hyperbolic spaces ~~have has~~ already been explored in detail, for instance, in [50], showing that ~~the~~ satisfiability and algorithmic complexity can be drastically different²⁴. This might be an interesting extension of ~~usual typical~~ VSA approaches ~~that makes it possible~~ to consider the symbol's numerical embedding in such a Riemannian differential manifold. This could be a fruitful perspective of such work.

Acknowledgments. Terrence C. Stewart is gratefully acknowledged for inspiring advice that helped us with some aspects of this work. Gabriel Doriath Döhler is ~~highly~~ thanked for his work on clarifying the use of VTB- algebra and introducing interesting ideas during his undergraduate internship. Frédéric Alexandre and Hugo Chateau-Laurent are ~~deeply gratefully~~ acknowledged for their ~~precious valuable~~ advices and their contributions to previous works on this subject. This work is supported by the ~~Inria, AExAIDE~~ exploratory action.

Appendix A Using ~~the~~ VTB algebra

Author query: Throughout the appendices, sometimes subsections are numbered and sometimes they are not. For consistency, you should probably either number all of the subsections or number none of them. At the mesoscopic level, symbols ~~are represent a~~ numerical grounding to real or complex vectors of dimension d , ~~with~~ each numerical grounding corresponding to some distributed activity of a spiking neuronal ~~assemble assembly~~ and each algebraic operation ~~corresponding~~ to some transformation of this activity.

Let us review and further develop ~~here~~ one of the algebras ~~used~~ to manipulate such symbols at an abstract level ~~use~~ in this paper: ~~the Vector-Derived Transformation Binding~~ ~~vector-derived transformation binding~~ (VTB) ~~algebra~~, following We follow [51] and ~~completing their complete~~ the developments, ~~in that paper~~ by deriving the different operations at the component

²⁴A ~~large audience~~ version of these elements ~~intended for a wider audience~~ is available in ~~this a science~~ popularization ~~scientifique~~ journal: <https://interstices.info/calculer-dans-un-monde-hyperbolique>.

level, yielding an optimal implementation and making explicit the computational complexity and ~~1st-related first-order-related~~ noise. This is in ~~partieuler~~ particular used in Appendix D to derive the macroscopic computations.

We consider, that $d \stackrel{\text{def}}{=} (d')^2$ for some integer d' ; thus, it is a quadratic number and we start from the standard definition of the VTB binding operation:

$$\mathbf{z} \stackrel{\text{def}}{=} \mathbf{B}_{\mathbf{y}} \mathbf{x},$$

where $\mathbf{B}_{\mathbf{y}}$ is a block-diagonal matrix defined as follows:

$$\mathbf{B}_{\mathbf{y}} \stackrel{\text{def}}{=} \begin{bmatrix} \mathbf{B}'_{\mathbf{y}} & 0 & \dots & 0 \\ 0 & \mathbf{B}'_{\mathbf{y}} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \mathbf{B}'_{\mathbf{y}} \end{bmatrix}, \text{ with } \mathbf{B}'_{\mathbf{y}} \stackrel{\text{def}}{=} \sqrt{d'} \begin{bmatrix} y_1 & y_2 & \dots & y_{d'} \\ y_{d'+1} & y_{d'+2} & \dots & y_{2d'} \\ \vdots & \vdots & \ddots & \vdots \\ y_{d-d'+1} & y_{d-d'+2} & \dots & y_d \end{bmatrix},$$

or equivalently²⁵, for $i = 1 \dots d$:

$$\begin{cases} [\mathbf{z}]_i \stackrel{\text{def}}{=} \mathbf{B}_{\mathbf{y}} \mathbf{x} = \sqrt{d'} \sum_{k=1}^{k=d'} [\mathbf{y}]_{k+\beta(i)} [\mathbf{x}]_{k+\alpha(i)}, \\ [\mathbf{B}_{\mathbf{y}}]_{ij} = \sqrt{d'} \delta_{i \leq d'} \text{ and } j \leq d' [\mathbf{y}]_{k+\beta(i)-\alpha(i)}, \end{cases} \text{ written } \begin{cases} \alpha(i) \stackrel{\text{def}}{=} d' ((i-1) \text{ div } d'), \\ \beta(i) \stackrel{\text{def}}{=} d' ((i-1) \text{ mod } d'), \end{cases} \quad (\text{A1})$$

with the matrix multiplication explicitized as a sum, ~~as it which~~ can be easily verified. Here, $[\mathbf{z}]_k$ stands for the k -th coordinate of the vector \mathbf{z} , and $\delta_{\mathcal{P}}$ is 1 if \mathcal{P} is true; ~~else~~ otherwise, it is 0. This is our basic definition, and reformulating the VTB operation using (A1) will ~~be useful~~ allow us to better understand its properties.

This operation is bi-linear in \mathbf{x} and \mathbf{y} , and thus it is distributive with respect to addition and the scalar product.

The $\sqrt{d'}$ renormalization factor allows \mathbf{z} to ~~be have a~~ unary magnitude²⁶.

At the algorithmic implementation level, the calculation of \mathbf{z} is performed in²⁷ $O(d^{\frac{3}{2}})$ operations, and the $\mathbf{y}_{k+\beta[i]}$ and $\mathbf{x}_{k+\alpha[i]}$ indexing can be tabulated in two fixed look-up tables $\beta[i]$ and $\alpha[i]$, avoiding any additional calculations. Furthermore, the fact that $\sqrt{d'}$ is an integer ~~allows limiting~~ makes it possible

²⁵ All algebraic derivations reported here are straightforward and were verified using a piece of symbolic algebra code available at <https://raw.githubusercontent.com/vthierry/onto2spa/main/figures/VTB-algebra.mpl>.

²⁶ More precisely, two random normalized vectors of dimension d drawn from a random normal distribution of independent samples verify that $\mathbf{x} \cdot \mathbf{y} \sim \mathcal{N}(0, 1/d')$, as ~~already detailed described~~ in subsection 2.1. Then, applying a permutation on all indices on a random vector \mathbf{x} yields another random vector, which is not correlated with any vector \mathbf{y} , if \mathbf{x} is not. Thus, when computing the components $[\mathbf{z}]_i$ in (A1) for two general random vectors \mathbf{x} and \mathbf{y} , we compute the dot product of two random vectors of dimension d' renormalized by $\sqrt{d'}$, and thus ~~of this dot product comes from the distribution~~ $\mathcal{N}(0, 1)$; ~~which this~~ corresponds to drawing a random vector unary ~~in-on~~ average.

Author query: Is “random vector unary” meant to be “random unary vector”?

²⁷ Each of the d components $[\mathbf{z}]_i$ requires a dot product of size $d' = \sqrt{d}$ ~~which that~~ is not factorizable in the general case, since involving different elements of the vectors as readable on the matrix form. **Author query:** The meaning of the phrase “since involving different elements of the vectors as readable on the matrix form” is not clear. Do you mean something like “since it involves different elements of the vectors that are readable in matrix form”?

to limit numerical approximations in order to improve numerical conditioning. This will be verified for all other explicit formulae ~~in the sequel~~ later in this paper. We make explicit these formulae in detail not to re-implement these operations, which are already available in the Nengo simulator, but to study in detail their complexity, and also their precision, ~~targeting to propose with the goal of proposing~~ a macroscopic algorithmic ersatz of these operations.

This can be compared to the fastest binding operation, which is convolution implemented via the fast Fourier transform [16], and thus ~~in it has a complexity of~~ $O(d \log(d))$:

	d = 10	d = 100	d = 500	d = 1000
VTB	$10^{1.5}$	10^3	10^4	$10^{4.5}$
Convolution	$10^{1.4}$	$10^{2.6}$	$10^{3.5}$	$10^{3.8}$
Ratio = $\frac{\sqrt{d}}{\log(d)}$	$\simeq 1$	$\simeq 2$	$\simeq 3.5$	$\simeq 4.5$

As stated in [51] and reviewed in [13], the key point is that this binding operation generates a new vector \mathbf{z} that is almost orthogonal to \mathbf{x} and \mathbf{y} :

$$(\mathbf{B}_y \mathbf{x}) \cdot \mathbf{x} \simeq 0,$$

and this operation is neither commutative~~:~~,

$$(\mathbf{B}_x \mathbf{y}) \cdot (\mathbf{B}_y \mathbf{x}) \simeq 0,$$

nor associative²⁸, in the following sense:

$$(\mathbf{B}_{(\mathbf{B}_z \mathbf{y})} \mathbf{x}) \cdot ((\mathbf{B}_z \mathbf{B}_y) \mathbf{x}) \simeq 0.$$

~~which are precious~~ These properties ensure that we do in order not ~~to~~ infer spurious derivations.

A To take this a step further, in the real case, the random matrix is almost orthogonal, i.e.,~~:~~

$$\mathbf{B}_y^\top \mathbf{B}_y \simeq \mathbf{I},$$

for the same reasons evoked just before above. It is straightforward to evaluate the precision of this approximation, which is again in $\mathcal{N}(0, O(1/\sqrt{d'}))$ for each matrix element²⁹. This also means that the level of noise is relatively high,

²⁸Of course, as a product of ~~matrix matrices~~, the combination of three binding or two binding operations and a vector is associative, but ~~not~~ the operator \mathbf{B} itself is not, as made explicit in the formula.

²⁹From (A1), we derive

$$\begin{aligned} [\mathbf{B}_y^\top \mathbf{B}_y]_{ij} &= \sum_{k=1}^{d'} [\mathbf{B}_y]_{ki} [\mathbf{B}_y]_{kj} \\ &= d' \sum_{k=1}^{d'} [\mathbf{y}]_{k+\beta(i)-\alpha(i)} [\mathbf{y}]_{k+\beta(j)-\alpha(j)} \\ &= d' \sum_{l=1}^{d'} [\mathbf{y}]_l [\mathbf{y}]_{k+(\beta(j)-\beta(i))-(\alpha(j)-\alpha(i))} \\ &= d' \sum_{l=1}^{d'} [\mathbf{y}]_l [\mathbf{y}]_{k+d'((j-i) \operatorname{div} d') - ((j-i) \operatorname{div} d')}, \end{aligned}$$

so that~~:~~

- $[\mathbf{B}_y^\top \mathbf{B}_y]_{ii} = \sum_{l=1}^{d'} [\mathbf{y}]_l^2 = 1 + \mathcal{N}(0, O(1/\sqrt{d'}))$, as derived in subsection 2.1;~~:~~ and
- $[\mathbf{B}_y^\top \mathbf{B}_y]_{ij, i \neq j} = \mathcal{N}(0, O(1/\sqrt{d'}))$, because it is easy to verify that $((j-i) \operatorname{div} d') - ((j-i) \operatorname{div} d') \neq 0$ as soon as when $i \neq j$, so that the dot product of d' random components of $[\mathbf{y}]_l$ with d' other random components, yielding a yields normal noise, as again made explicit in subsection 2.1.

in $O\left(\frac{1}{d^{\frac{1}{4}}}\right)$ instead of $O\left(\frac{1}{d^{\frac{1}{2}}}\right)$, which explains the relatively limited numerical performances of simulations with $d < 10^3$ as reported, for instance, in [16].

We thus define:

$$\mathbf{B}_{\mathbf{y}^\sim} \stackrel{\text{def}}{=} \mathbf{B}_{\mathbf{y}}^\top \text{ with } [\mathbf{y}^\sim]_i \stackrel{\text{def}}{=} [\mathbf{y}]_{\sigma(i)},$$

with:

$$\sigma(i) \stackrel{\text{def}}{=} 1 + d'((i-1) \bmod d') + (i-1) \operatorname{div} d'.$$

In other words, $\mathbf{B}_{\mathbf{y}}^\top$ has the same structure as $\mathbf{B}_{\mathbf{y}}$, except that the vector coordinates are subject to a permutation $\sigma(i)$, which is idempotent ($\sigma(\sigma(i)) = i$), and thus its own inverse, so that if $\mathbf{z}' \stackrel{\text{def}}{=} \mathbf{B}_{\mathbf{y}^\sim} \mathbf{x}$ we obtain:

$$[\mathbf{z}']_i = \sqrt{d'} \sum_{k=1}^{k=d'} [\mathbf{y}]_{\sigma(k+\beta(i))} [\mathbf{x}]_{(k+\alpha(i))}$$

(where $\beta(i)$ and $\alpha(i)$ are the indexing defined to calculate $\mathbf{B}_{\mathbf{y}} \mathbf{x}$ explicitly), and it allows this makes it possible to define a left unbinding operation:

$$\mathbf{B}_{\mathbf{y}^\sim} (\mathbf{B}_{\mathbf{y}} \mathbf{x}) = \mathbf{B}_{\mathbf{y}}^\top \mathbf{B}_{\mathbf{y}} \mathbf{x} = \mathbf{x} + \mathcal{N}(\mathbf{0}, O(1/d)) \simeq \mathbf{x}.$$

Thus, again, up to an additive noise of $O(1/d^{1/4})$.

The right identity vector \mathbf{i} such that $\mathbf{B}_i = \mathbf{I}$, writes can be written explicitly as follows:

$$[\mathbf{i}]_i = \frac{1}{\sqrt{d'}} \delta_{i=\sigma(i)}.$$

In other words, we get i_B by “unfolding” the identity matrix I'_d line by line, writing a 1, then d times 0, then another 1, and so on.

Considering the mirroring matrix $\mathbf{B}_{\leftrightarrow}$, which is defined as:

$$[\mathbf{B}_{\leftrightarrow}]_{ij} \stackrel{\text{def}}{=} \delta_{j=\sigma(i)}$$

(with which is thus not block-diagonal as in the way that a matrix of the form $\mathbf{B}_{\mathbf{y}}$ are is), so that $\mathbf{B}_{\leftrightarrow} \mathbf{x} = \mathbf{x}^\sim$, we obtain:

$$\mathbf{B}_{\leftrightarrow} \mathbf{B}_{\mathbf{y}} \mathbf{x} = \mathbf{B}_{\mathbf{x}} \mathbf{y}, \text{ while } \mathbf{B}_{\leftrightarrow} \mathbf{B}_{\leftrightarrow} = \mathbf{I} \text{ and } \mathbf{B}_{\leftrightarrow}^\top = \mathbf{B}_{\leftrightarrow},$$

which allows defining makes it possible to define a right unbinding operation:

$$(\mathbf{B}_{\mathbf{x}^\sim} \mathbf{B}_{\leftrightarrow}) (\mathbf{B}_{\mathbf{y}} \mathbf{x}) = \mathbf{B}_{\mathbf{x}^\sim} \mathbf{B}_{\mathbf{x}} \mathbf{y} \simeq \mathbf{y}.$$

Unfortunately, $\mathbf{B}_{\leftrightarrow}$ is not a binding matrix, i.e., it is not of the form $\mathbf{B}_{\mathbf{z}}$ for some vector \mathbf{z} , which is easily verified by the fact that some components that must be equal to 0 for a binding matrix are equal to 1 in $\mathbf{B}_{\leftrightarrow}$. Furthermore, the left or right multiplication of a binding matrix by this mirroring matrix does not yield a binding matrix, because of the same observation; that components that must be equal to 0 for a binding matrix are equal to 1 in $\mathbf{B}_{\leftrightarrow}$.

Beyond [51], the authors of [13] have introduced a vector composition operator \odot making to make explicit the composition of two binding operations, namely:

$$\mathbf{B}_{\mathbf{v}} = \mathbf{B}_{\mathbf{y}} \mathbf{B}_{\mathbf{x}} \Leftrightarrow \mathbf{v} \stackrel{\text{def}}{=} \mathbf{y} \odot \mathbf{x},$$

which can be explicitly ~~writes-written~~ as follows³⁰:

$$[\mathbf{v}]_i = \sqrt{d'} \sum_{k=1}^{k=d'} [\mathbf{y}]_{(i-1)d'+k} [\mathbf{x}]_{1+d'(k-1)+(i-1) \bmod d'}.$$

At the algebraic level, the key point is that the product of two binding matrices is still a binding matrix. As a consequence, this composition operator is bi-linear, and thus it is distributive with respect to ~~the~~ addition; it is not commutative, but it is associative, and commutes with the inversion as follows:

$$(\mathbf{y} \odot \mathbf{x})^\sim = \mathbf{x}^\sim \odot \mathbf{y}^\sim,$$

while $\mathbf{x}^\sim \odot \mathbf{x} \simeq \mathbf{i}$; all these results ~~are-can be~~ easily derived by considering usual matrix properties. This allows us to combine two binding matrices without an explicit matrix product, ~~but-also~~ in $O(d^{\frac{3}{2}})$ operations only. At the numeric level, since \mathbf{v} is up to a $\sqrt{d'}$ factor, the dot product of segments of random vectors of dimension d' follows a normal distribution of standard-deviation $O(1/d')$. ~~w~~We also ~~obtain-find~~ that:

$$[\mathbf{v}]_i = \mathcal{N}(\mathbf{0}, O(1/\sqrt{d'})),$$

which is a rather high level of noise, when ~~applying~~ an unbinding operation is applied, for instance, to an associative map.

Using ~~the~~ VTB algebra in the complex case.

All of the developments ~~of-described in~~ this section generalize to complex numbers. Although it is not directly used here, such a generalization is of general interest, because complex implementations of VSA frameworks have also been considered [16]. Furthermore, it is also of interest to see ~~that-if~~ our macroscopic implementation could be easily adapted to the complex case.

Stating that two resources are semantically equivalent if the unary vectors are aligned ~~writes-can be written~~ in the complex case as follows³¹:

³⁰Since $\mathbf{B}_\mathbf{y}$ and $\mathbf{B}_\mathbf{x}$ are block-diagonal matrices, it is easy to verify that $\mathbf{B}_\mathbf{v}$ is a block-diagonal matrix with a $d' \times d'$ block $\mathbf{B}_\mathbf{v}' = \mathbf{B}_\mathbf{y}' \mathbf{B}_\mathbf{x}'$ using the notations ~~of-from~~ the beginning of this section, and we can explicitly write that:

$$[\mathbf{B}_\mathbf{v}']_{ij} = \sqrt{d'} \sqrt{d'} \sum_{k=1}^{d'} [\mathbf{y}]_{k+(i-1)d'} \text{div } d' [\mathbf{x}]_{(k-1)d'+j},$$

from which we obtain the desired formula.

³¹If we are in the real case \mathbf{x} and $\mathbf{y} \in \mathcal{R}^d$, with $\|\mathbf{x}\|_2 = \|\mathbf{y}\|_2 = 1$, then the equality ~~writes-is~~ written as

$$\mathbf{x} = \mathbf{y} \Leftrightarrow \mathbf{x} \cdot \mathbf{y} = \sum_i x_i y_i = \cos(\widehat{\overrightarrow{\mathbf{x}}} \widehat{\overrightarrow{\mathbf{y}}}) = 1 \Leftrightarrow \widehat{\overrightarrow{\mathbf{x}}} \widehat{\overrightarrow{\mathbf{y}}} = 0 \pmod{2\pi},$$

i.e., both unary vectors have the same direction; ~~i.e.,~~ in other words, they are aligned. If we are in the complex case \mathbf{x} and $\mathbf{y} \in \mathcal{C}^d$, let us consider the canonical embedding in \mathcal{R}^{2d} , i.e., we consider the real (*Re*) and imaginary (*Im*) parts as two real coordinates, ~~writing-denoting by~~ $\overrightarrow{\mathbf{x}}$ the corresponding vector:

$$\mathbf{x} \stackrel{\text{def}}{=} (x_1, x_2, \dots)^T \Leftrightarrow \overrightarrow{\mathbf{x}} \stackrel{\text{def}}{=} (Re(x_1), Im(x_1), Re(x_2), Im(x_2), \dots)^T,$$

~~for which, writing-where~~ z^* is the conjugate of a complex number z , while $\langle \mathbf{x} | \mathbf{y} \rangle$ stands for the complex inner product:

$$\begin{aligned} \langle \mathbf{x} | \mathbf{y} \rangle &\stackrel{\text{def}}{=} \sum_i x_i y_i^* \\ &= \sum_i (Re(x_i) Re(y_i) + Im(x_i) Im(y_i)) + I (Re(x_i) Im(y_i) - Im(x_i) Re(y_i)) \\ &= \overrightarrow{\mathbf{x}} \cdot \overrightarrow{\mathbf{y}} + I \overrightarrow{\mathbf{x}}^* \cdot \overrightarrow{\mathbf{y}}, \end{aligned}$$

$$\mathbf{x} \simeq \mathbf{y} \Leftrightarrow \langle \mathbf{x} | \mathbf{y} \rangle \simeq 1,$$

while the orientation is usually defined as[‡]

$$\widehat{\mathbf{x} | \mathbf{y}} \stackrel{\text{def}}{=} \arccos(\text{Re}(\langle \mathbf{x} | \mathbf{y} \rangle)),$$

as ~~detailed~~ explained in the previous footnote.

Provided that the space dimension d is large enough, two randomly chosen different complex vectors \mathbf{x} and \mathbf{y} , ~~will be also~~³² will also be approximately orthogonal in the sense that[‡]

$$\mathbf{x} \neq \mathbf{y} \Leftrightarrow \langle \mathbf{x} | \mathbf{y} \rangle \simeq 0.$$

As a consequence, the VTB matrix is almost a unitary matrix, i.e.,

$$\mathbf{B}_{\mathbf{y}}^* \mathbf{B}_{\mathbf{y}} \simeq \mathbf{I},$$

considering ~~now~~ the conjugate transpose.

All other algebraic operations are common to both real ~~or~~ and complex linear algebra, and this ~~to~~ is also the case for other VSA binding operators.

More than just a confirmation, these derivations allow us to observe that using a complex representation would be interesting if the conjugate of a vector could have a semantic interpretation. In that case, if, say, \mathbf{x} and \mathbf{y}^* are similar, then $\langle \mathbf{x} | \mathbf{y} \rangle \simeq I$, as easily verified from the previous ~~development~~ derivations.

Appendix B On VSA data structures

This section revisits the literature, emphasizing the properties of the data structures ~~properties~~, ~~discussing it~~ discusses in more detail their computational properties and limitations, and linking links them ~~with to~~ usual programming language data structures.

B.1 Unordered set or bundling

We first consider an unordered set \mathcal{S} of N symbols grounded to values $\{\mathbf{s}_1, \dots, \mathbf{s}_i, \dots, \mathbf{s}_N\}$, and we would like to be able to store them, ~~in~~ such a way that we can check if ~~the a given~~ symbol is in the set. Very simply, we ground \mathcal{S} to the vector \mathbf{s} :

$$\mathbf{s} \stackrel{\text{def}}{=} \sum_i \mathbf{s}_i,$$

so that $\text{Re}(\langle \mathbf{x} | \mathbf{y} \rangle) = \vec{\mathbf{x}} \cdot \vec{\mathbf{y}}$, and $\|\mathbf{x}\| = \sqrt{\langle \mathbf{x} | \mathbf{x} \rangle} = \|\vec{\mathbf{x}}\| - 2 = \sqrt{\vec{\mathbf{x}} \cdot \vec{\mathbf{x}}}$, and since vectors are unary,[‡]

$$\langle \mathbf{x} | \mathbf{y} \rangle = 1 \Leftrightarrow \vec{\mathbf{x}} \cdot \vec{\mathbf{y}} = 1 \Leftrightarrow \vec{\mathbf{x}} = \vec{\mathbf{y}} \Leftrightarrow \mathbf{x} = \mathbf{y},$$

making explicit the obvious fact that unary real or complex vectors are equal if and only if their inner product equals one, while we consider the “angle” of two complex vectors as the angle of their $2d$ real embedding, i.e.,[‡]

$$\widehat{\mathbf{x} | \mathbf{y}} \stackrel{\text{def}}{=} \arccos(\text{Re}(\langle \mathbf{x} | \mathbf{y} \rangle)).$$

³²Considering again the canonical embedding in \mathcal{R}^{2d} and the fact that

$$\langle \mathbf{x} | \mathbf{y} \rangle = \vec{\mathbf{x}} \cdot \vec{\mathbf{y}} + I \vec{\mathbf{x}} \cdot \vec{\mathbf{y}},$$

because $\vec{\mathbf{x}}$ and thus $\vec{\mathbf{x}}^*$ and $\vec{\mathbf{y}}$ are random vectors, their dot product almost vanishes[‡]; thus, the real and imaginary parts of $\langle \mathbf{x} | \mathbf{y} \rangle$ also almost vanish.

which provides a solution, because given a symbol \mathbf{s}_\bullet , we obviously observe that $\mathbf{s}_\bullet \cdot \mathbf{s} \simeq 1$ if it corresponds to ~~one-a-certain-symbol~~ \mathbf{s}_i , and it is almost 0 otherwise, because random vectors are almost orthogonal, as previously explained. This is called bundling [16] or superposition.

Furthermore, the representation intrinsically includes a notion of transitivity: If a set includes another subset, by construction, it includes the subset elements. ~~More~~ precisely,

$$\mathbf{s} \stackrel{\text{def}}{=} \sum_i \mathbf{s}_i \text{ and } \mathbf{s}_i \stackrel{\text{def}}{=} \sum_j \mathbf{s}_{ij} \Rightarrow \mathbf{s} \stackrel{\text{def}}{=} \sum_{ij} \mathbf{s}_{ij},$$

and thus, $\mathbf{s}_{ij} \cdot \mathbf{s} > 0$ for all subsets ~~elements~~.

This obviously generalizes to weighted symbols $\hat{\mathbf{s}}_i$, i.e., symbols with modal-ity weighting. In that case, $\mathbf{s}_\bullet \cdot \mathbf{s} \simeq \tau$ ~~allows-retrieving-makes-it-possible-to-retrieve~~ the belief weight. This is equivalent to inputting a symbol \mathbf{s}_\bullet ~~which-that~~ is approximately similar to a given symbol \mathbf{s}_i , thus indicating an approximate similarity; ~~but~~ however, it neither ~~allowing-allows-us~~ to retrieve the exact value of \mathbf{s}_i , nor ~~indicating-indicates~~ if a positive value below 1 corresponds to a weighted symbol that has been exactly retrieved or to a symbol approximation.

This has an interesting biological interpretation: \mathcal{S} has ~~common~~-features in common with a Hopfield network or other related attractor networks, where information has been stored in a distributed way, while activating the map with an input ~~allows-retrieving-if-makes-it-possible-to-determine-whether~~ the symbol is stored, or not. The main difference is that attractor networks converge to the exact stored value, providing a mechanism of associative memory, which is now developed while introducing superposition, ~~first~~ allowing us to better understand the need for a more sophisticated mechanism.

Let us provide an analogy with programming data structures, explicitizing the similarities and differences between what is proposed here and what is available in common programming languages³³. This unordered set representation corresponds to a “set” container (e.g., a `std::unordered_set` in C++ or a `set()` in Python), ~~however-with~~ that has only an insertion method and a membership test function, without the capability to intrinsically enumerate the elements, as formerly discussed.

B.1.1 Symbol enumeration

At this stage, this structure does not allow us to directly enumerate all symbols \mathbf{s}_i , because from \mathbf{s} , it is not possible to decode the superposed vectors. In [10], for instance, where data structures are defined using superposition, the intrinsic memory enumeration of the stored information is not addressed. We thus need an external mechanism to select all elements and perform an operation on each one. However, at the implementation level, in Nengo [9], an explicit list of the defined vocabulary $\{\dots \mathbf{s}_i \dots\}$ is maintained, and the way to select the elements is to test $(\mathbf{s}^T \mathbf{s}_i)$ for each element of the vocabulary. ~~Such-This~~ select

³³We will do the same for other cognitive structures because we think that it ~~is-illustrative-of~~ illustrates the computing capability of the cognitive object.

operator has a complexity ~~is~~ of $O(K)$, where K is the size of the vocabulary. ~~In the sequel~~ Later in this paper, we will also propose a biologically plausible indexing mechanism, in order, for instance, to manipulate sequences.

B.2 Associative map

We now consider an unordered associative memory³⁴, or “map”, of N correspondences $\{\mathbf{s}_1 \rightarrow \mathbf{o}_1, \dots, \mathbf{s}_i \rightarrow \mathbf{o}_i \dots \mathbf{s}_N \rightarrow \mathbf{o}_N\}$ between subjects and objects. To this end, we use the binding operation $B_{\mathbf{s}_i}$, defined in Appendix A, with a pseudo-inverse, i.e., an unbinding operator, $B_{\mathbf{s}_i}^\sim$:

$$\mathbf{m} \stackrel{\text{def}}{=} \sum_i B_{\mathbf{s}_i} \mathbf{o}_i,$$

so that:

$$B_{\mathbf{s}_k}^\sim \mathbf{m} \simeq \sum_{i, \mathbf{s}_i = \mathbf{s}_k} \mathbf{o}_i.$$

In other words, the unbinding operation ~~allows retrieving~~ makes it possible to retrieve the set, i.e., the additive superposition of all objects \mathbf{o}_i associated ~~to~~ with a given subject \mathbf{s}_k , while $B_{\mathbf{s}_k}^\sim \mathbf{m} \simeq 0$ if none. **Author query: The meaning of the phrase “if none” is not clear. Could you reword it?** This is done up to a level of noise of $O(1/d^{1/4})$, as derived in Appendix A, which is rather high with respect to the similarity precision, which ~~precision~~ is ~~of~~ $O(1/d^{1/2})$, as observed numerically [16]; ~~whereas~~ however, in biological neuronal networks, where the dimension is an order of magnitude higher, this is no ~~more~~ longer a limitation because d is high.

This allows us to both detect if the information is in the table and retrieve this information in one step if ~~it~~ this is the case. However, as in the previous case, no mechanism allows the enumeration of the map subjects or objects.

This algebraic construction also ~~allows retrieving of~~ makes it possible to retrieve the subjects associated with a given object, because of the commutator $\mathbf{B}_{\leftrightarrow}$, such that:

$$\mathbf{B}_{\leftrightarrow} \mathbf{B}_{\mathbf{o}_i} \mathbf{s}_i = \mathbf{B}_{\mathbf{s}_i} \mathbf{o}_i,$$

yielding:

$$\mathbf{m}_{\leftrightarrow} \stackrel{\text{def}}{=} \mathbf{B}_{\leftrightarrow} \mathbf{m} = \sum_i B_{\mathbf{o}_i} \mathbf{s}_i,$$

which is now the numerical grounding of the reciprocal map $\{\mathbf{o}_1 \rightarrow \mathbf{s}_1, \dots, \mathbf{o}_i \rightarrow \mathbf{s}_i \dots \mathbf{o}_N \rightarrow \mathbf{s}_N\}$.

The algebraic construction also offers the notion of the identity vector \mathbf{i} , with $\mathbf{B}_i = \mathbf{I}$, so that:

$$\mathbf{s}_i = \mathbf{i} \rightarrow B_{\mathbf{s}_i} \mathbf{o}_i = \mathbf{o}_i.$$

~~In~~ other words, the binding reduces to a superposition. Theoretical details

³⁴This is not the only way to implement such an associative memory: In [52], binding/unbinding is not explicitly used, ~~but~~ and an input/output architecture with suitable connections is used instead. Each input unit has an encoding vector in which input weights are tuned to fire for a specific key and drive a connected output vector that is optimized to estimate the value associated with the related key.

underlying the implementation of such associative memories are available in [14].

As for the previous structure, this obviously generalizes to weighted symbols $\hat{\mathbf{s}}_i$ and $\hat{\mathbf{a}}_n$ approximate input $\mathbf{s}_\bullet \simeq \hat{\mathbf{s}}_i$, allowing $\hat{\mathbf{u}}_s$ to retrieve the object \mathbf{o}_i weighted by either the modality weighting or the input approximation, indistinctly.

There are several solutions $\hat{\mathbf{u}}_s$ to define such binding, unbinding, and commutator operators. A proposed solution is developed in Appendix A after the work in [51], which was completed by [13]. This design choice is guided by the fact that we need to avoid spurious inferences: ~~With~~ a commutative operator (such as the convolution operator), $\mathbf{B}_{\mathbf{o}_i} \mathbf{s}_i$ would equal $\mathbf{B}_{\mathbf{s}_i} \mathbf{o}_i$, which could generate non-sense deductions (e.g., ~~in-for~~ a driver-vehicle map, ~~it-this~~ would mean that if Ming-Yue drives a bicycle, then the bicycle drives Ming-Yue unless some additional mechanism is considered to avoid such nonsense). The proposed ~~Vector-Derived-Transformation-Binding~~ (VTB) algebra avoids such caveats (see [16] for a recent comparison of different VSAs)³⁵.

This associative memory mechanism has an interesting biological interpretation: It implements an associative memory in the biological sense, with the association ~~storing-stored~~ in a distributed way, and activating the associative memory with an input \mathbf{s}_\bullet allows $\hat{\mathbf{u}}_s$ to retrieve the associated symbol. This is what happens in several biological mechanisms, as reviewed, for instance, in [15].

In particular, a structure of the form:

$$\mathbf{m} \stackrel{\text{def}}{=} \sum_n B_{\mathbf{s}_i} \mathbf{s}_i$$

~~mapping-that maps~~ an object onto itself allows the retrieval of an exact symbol from an approximate input, solving the caveats induced by using only a superposition mechanism, ~~that were~~ presented previously. This is exactly what is expected in an associative encoder (e.g., a Hopfield network); if a symbol is close to an existing symbol, the associative memory will output a weighted version of the symbol.

At the computer programming level, this corresponds to a “map” container (e.g., a Map in JavaScript or a dictionary() in Python), again with only insertion and retrieval methods, ~~and~~ without intrinsic iterators.

~~A-To take this a~~ step further, we can propose a complementary functionality, defining an additional symbol “something” τ whose numerical grounding is fixed to any new random vector σ_τ , ~~that is~~ never used elsewhere. This allows us to enhance the information to be obtained, as follows: Each time ~~a piece of information~~ $\mathbf{s}_i \rightarrow \mathbf{o}_i$ ~~information~~ is added, we also add $\mathbf{s}_i \rightarrow \sigma$ and $\sigma \rightarrow \mathbf{o}_i$, i.e., we make explicit the fact that \mathbf{s}_i and \mathbf{o}_i are defined in this table, which can be retrieved in one step, without ~~requiring-the need~~ to enumerate the different elements. In such a case,

$$\mathbf{m}_{\mathbf{s}_j} \stackrel{\text{def}}{=} \sum_{i, \mathbf{s}_j = \mathbf{s}_i} \mathbf{o}_i = P_{\sigma^\perp} B_{\mathbf{s}_i} \mathbf{m},$$

³⁵ An alternative to VTB algebra is called the-MBAT algebra; ~~requiring it requires~~ matrix inversion instead of transposition, ~~and~~ thus ~~it is~~ less efficient.

~~writing where~~ $P_{\sigma^\perp} \stackrel{\text{def}}{=} \mathbf{I} - \sigma \sigma^T$ is the projection onto the orthogonal of σ , i.e., we must eliminate the symbol “something” from the expected values.

B.3 Indexed and chained list

B.3.1 *Indexes cConstruction of indexes*

In order to define an indexed list, we need indexes, i.e., a mechanism that generates ordinal values. Our main purpose here is to make explicit that what has been developed using convolution operators [53] still holds with VTB. We fix the symbol grounding of the “zero” symbol ν_0 , which is never used elsewhere, and define the following recursively:

$$\nu_{n+1} \stackrel{\text{def}}{=} B_{\nu_0} \nu_n,$$

i.e., the $(n+1)$ -th ordinal value is obtained by binding the n -th, and we easily obtain, from a few algebra operations:

$$B_{\nu_p} \nu_q = B_{\nu_q} \nu_p = B_{\nu_{p+q}} \nu_0, \quad B_{\nu_p} \nu_q^\sim = B_{\nu_q^\sim} \nu_p \simeq B_{\nu_{p-q}} \nu_0.$$

~~In~~ particular, $\nu_{n-1} \simeq B_{\nu_0^\sim} \nu_n$, so that the definition holds for $n \in \mathcal{Z}$.

Here, we only consider the minimal material needed to build an indexed list; ~~while~~ numerical information in the brain is a much more complex subject [54] beyond the scope of this work.

B.3.2 *Indexed list*

We can now define an indexed list or array, often called a vector, since the previous mechanism allows us to generate a “counter” that ~~could can~~ be incremented or decremented using the binding or unbinding operator.

To this end, an associative map indexed by these ordinals can be managed as a list, ~~which~~ whose values can be enumerated. Such a representation is also present at several cognitive levels when considering temporal sequences, actions, or any enumeration. This is also the tool that allows us to enumerate all elements of a symbol set \mathcal{S} , which was defined previously ~~defined~~, or the subjects of an associative map.

To make this mechanism explicit, let us consider a list $\mathbf{l} \stackrel{\text{def}}{=} \sum_i B_{\nu_i} \mathbf{l}_i$, and a variable index \mathbf{k} : A construct of the form:

```
for  $\mathbf{k} \leftarrow \nu_0$ ; while  $\|B_{\mathbf{k}} \mathbf{l}\| > 0$ ; next  $\mathbf{k} \leftarrow B_{\nu_0} \mathbf{k}$  do
   $\mathbf{l}_i \leftarrow B_{\mathbf{k}^\sim} \mathbf{l}$ 
  ..../..
end for
```

allows us to enumerate³⁶ all elements, this being indeed only an algorithmic ersatz to illustrate the mechanism beyond the biologically plausible implementation of sequential memory organization.

³⁶In fact, considering $\mathbf{l} \stackrel{\text{def}}{=} \sum_i B_{\nu_i} \mathbf{l}_i + B_{\nu_{-1}} \lambda$, where λ is the list length, which is updated when ~~adding or deleting~~ an element is added or deleted, would improve the algorithmic ersatz implementation, which is not the issue here.

At the biological plausibility level, following [15], we may consider that the brain memory can have three kinds of memory: associative, sequential, ~~or~~ and hierarchical (~~named~~ called schematic by the author of [15]) memory. ~~a~~ All three ~~being~~ memory types are present and required for cognitive processes. The VSA approach provides both associative and sequential memory. Let us consider the third ~~one type of~~ memory, which has not, to the best of our knowledge, been addressed with regard to VSAs.

At the computer programming level, this corresponds to an extensible “array” (e.g., a `std::vector` in C++ or `java.util.AbstractList` in Java), with basic edition and retrieval methods available.

B.3.3 Chained list

We can also define a chained list using an associative memory of the form:

$$\begin{aligned} \text{first} &\rightarrow \text{second} \\ \text{second} &\rightarrow \text{third} \\ &\dots \\ \text{last} &\rightarrow \text{END} + \text{first} \end{aligned},$$

where every value of the list acts as a key to the value of its successor in the list, thus enumerating the values. END is a predefined specific symbol allowing that makes it possible to know when the list ends; that we can superpose to a pointer to the first value in case we need to iterate through the entire list again.

We also could ~~also~~ have considered multiple binding³⁷, as proposed in [13].

Appendix C RDFS entailment rules implementation

As a side- product of the present development, we would like to illustrate the computational capability of the proposed framework; by briefly showing that our biologically plausible mechanism is at least able to perform RDFS³⁸ specification inferences. This is done, of course, without any assumption about the fact that the brain explicitly performs such computations³⁹.

³⁷In such a case, a list of the form $l = [v_1, v_2, \dots]$ is encoded without associative memory as:

$$l = B_{\text{value}} v_1 + B_{\text{next}} (B_{\text{value}} v_2 + B_{\text{next}} (\dots + B_{\text{next}} (\text{list-end}))),$$

allowing us to obtain by unbinding; the list's head value; and its tail value, and allowing us to detect its end. This corresponds, for instance, to the `rdf:first`, `rdf:rest`, and `rdf:nil` symbols of the RDF representation. However, as discussed in Appendix A, chaining unbinding operations is not numerically very robust due to the additional residual noise.

³⁸According This is true according to the <https://www.w3.org/TR/rdf-schema> specification.

³⁹Considering knowledge representation and reasoning, the capabilities of Semantic Web semantic web modeling languages, such as RDFS (Resource Description Framework Schema) and OWL (Web Ontology Language) (OWL) (see, e.g., [55] for a recent didactic reference), make them is a powerful way of solving to solve modeling problems and manipulate high-level data representations. It appears also They also appear to be rather accessible to an educated person, as pointed out in [55], and corresponds for a certain in part to the usual common-sense formalization, for instance, in terms of the notion of class (e.g., Garfield is a cat, that which is an animal, that which is a living organism). The brain is capable of such reasoning, as discussed in the paper

RDFS modeling in a nutshell

In order to represent symbolic information, the RDFS knowledge representation is based on the RDF data model, which represents knowledge as triples, as made explicit, in Fig. 3. More precisely, the *universe of discourse* is made of *resources*, referenced by some universal International Resource Identifier (IRI), i.e., a fixed lexical token. To structure this universe of discourse, we consider:

- (i) *individuals* that refer to real-world concrete or abstract objects;
- (ii) *literals* to characterize individuals using data attributes, i.e., numerical values, character strings, or any structured information, such as dates;
- (iii) *concepts* and *roles* (namely *classes* and *properties*) that allow us to structure the knowledge about individuals.

In fact, in our context, which is outside the web semantic application field, we have to point out that the RDF/RDFS framework, has to be considered with the following variants:

- We conflate *name* with both the IRI and blank node since on the one hand, the blank node can be eliminated⁴⁰, and on the other hand, because we only process the information locally at this stage, thus avoiding considering all issues regarding distributed information between different sources;
- We do not consider (i) semantic web specific literals (e.g., `rdf:XMLLiteral`), or (ii) utility and annotation or other human-targeted properties (e.g., `rdfs:seeAlso`) at this stage;
- We will introduce both containers, i.e., ordered or unordered sequences, and collections, i.e., chained lists, later in these specifications, but in a somehow somewhat different form, adapted to the numerical representation and obvious easy to map on to RDF representations;
- We do not consider all XSD data types, but we will introduce a precise notion of numerical values and will detail describe how to represent structured data in our framework.

Given the capability of stating its ability to state facts, the RDFS framework allows us to structure concepts, using the following construct:

- The notion of class inheritance (e.g., if Tom is a cat, and cats are animals, Tom is an animal), allows us to define a hierarchical taxonomy of classes, structure the objects in categories, and infer all that is possible from this taxonomy.
- The notion of property inheritance (e.g., if Tom is the brother of Jerry, Tom is in the same family as Jerry), the property of being in the same family is a super-property of being the brothers, which allows structuring us to structure properties, and also infer new properties by inheritance.
- The notion of domain and range (e.g., if Tom is the brother of Jerry, it this also means that Tom is a boy), that allows classes for subjects and/or objects of a category.

introduction, through although the data representation and processing mechanism is obviously different from what is performed in semantic web modelers and reasoners, while the brain is mainly doing generally performs induction and abduction, more than deduction, as reviewed here.

⁴⁰ Using It can be eliminated using a standard process related to skolemization.

The language also allows us to define additional information, such as human-readable elements, but we consider it as a demonstrative subset to consider the main notions reviewed here. **Author query: In the previous sentence, it is not clear what the word “it” is referring to.**

Generality of numeric entailment rule implementation

The RDFS entailment, i.e., all that can be logically deduced from the input information, defines which elements are well-formed and which entailment relations allow ~~deducing us to deduce~~ all derived information. In the case of the RDFS framework, the entailment rules are ~~it is~~ given in Table C1. It appears that each rule can be implemented with the ~~proposed~~ mechanism ~~of proposed~~ in section 4, ~~for instance~~:

- The *rdfs9* class inheritance entailment rule,÷

$$(\$s \text{ rdf:type } \$c_1.) \wedge (\$c_1 \text{ rdfs:subClassOf } \$c_2.) \Rightarrow (\$s \text{ rdf:type } \$c_2.),$$

~~that~~ states that if any subject $\$s$ belongs to the class $\$c_1$ and this class $\$c_1$ is a subclass of $\$c_2$, then $\$s$ also belongs to the class $\$c_2$, as ~~taken as in the~~ major example in subsection 4.3.

- The *rdfs2* rule ~~allowing~~ makes it possible to infer the subject domain class:

$$(\$s_1 \ \$p_1 \ \$o_1.) \wedge (\$p_1 \text{ rdfs:domain } \$o_2.) \Rightarrow (\$s_1 \text{ rdf:type } \$o_2.)$$

yields ~~to~~:

$$(\$s_1 \ \$p_1 \ \$o_1.) \wedge (\$s_2 \ \$p_2 \ \$o_2.) \Rightarrow (\$s_1 \ \tau \text{ rdf:type } \$o_2.),$$

with:

$$\tau \stackrel{\text{def}}{=} (\$p_1 \cdot \$s_2) \ \& \ (\$p_2 \cdot \text{rdfs:domain}).$$

- The *rdfs3* rule ~~allowing~~ makes it possible to infer the object range class:

$$(\$s_1 \ \$p_1 \ \$o_1.) \text{ and } (\$p_1 \text{ rdfs:range } \$o_2.) \Rightarrow (\$o_1 \text{ rdf:type } \$o_2.)$$

yields ~~to~~:

$$(\$s_1 \ \$p_1 \ \$o_1.) \wedge (\$s_2 \ \$p_2 \ \$o_2.) \Rightarrow (\$o_1 \ \tau \text{ rdf:type } \$o_2.),$$

with:

$$\tau \stackrel{\text{def}}{=} (\$p_1 \cdot \$s_2) \ \& \ (\$p_2 \cdot \text{rdfs:range}).$$

- The *rdfs7* rule ~~allows~~ makes it possible to infer sub-property inheritance:

$$(\$s_1 \ \$p_1 \ \$o_1.) \text{ and } (\$p_1 \text{ rdfs:subPropertyOf } \$o_2.) \Rightarrow (\$s_1 \ \$o_2 \ \$o_1.)$$

yields ~~to~~:

$$(\$s_1 \ \$p_1 \ \$o_1.) \wedge (\$s_2 \ \$p_2 \ \$o_2.) \Rightarrow (\$s_1 \ \tau \ \$o_2 \ \$o_1.),$$

with:

$$\tau \stackrel{\text{def}}{=} (\$p_1 \cdot \$s_2) \ \& \ (\$p_2 \cdot \text{rdfs:subPropertyOf}).$$

This easily generalizes to all pertinent rules in ~~†~~Table C1. We also notice that the left-hand side contains only one or two triple patterns, as stated in the main paper. In details, the rules can be described as follows:

- The *se1*, *se2*, *lg*, and *gl* rules are technical rules related to the management of blank nodes and literal allocation, which is not considered in our context. Similarly, the *rdf2* rule applies to XML literals, which are not considered here.
- The *rdf1*, *rdfs1*, *rdf4a*, *rdf4b*, *rdfs8*, *rdfs12*, and *rdfs13* rules are unary rules that state basic ~~meta-property~~ meta-properties of the language; ~~obvious~~ they are easy to implement.
- The *rdfs10* and *rdfs11* rules implement the reflexivity and transitivity of classes (thus, they are very similar to *rdfs9*); ~~rules~~, they are again ~~obvious and similar~~ easy to implement.
- The *rdfs6* and *rdfs5* rules implements the reflexivity and transitivity of properties, while *rdfs7* is equivalent to *rdfs9* but for property inheritance, and thus again they are very similar.

Rule set	Rule N ame	If E contains:	then add:
Simple entailment rules	<i>se1</i>	uuu aaa xxx .	uuu aaa .:nnn ³ where .:nnn identifies a blank node allocated to xxx by rule <i>se1</i> or <i>se2</i> .
	<i>se2</i>	uuu aaa xxx .	.:nnn aaa xxx ³ where .:nnn identifies a blank node allocated to uu by rule <i>se1</i> or <i>se2</i> .
	lg (literal generalization)	uuu aaa lll .	uuu aaa .:nnn ³ where .:nnn identifies a blank node allocated to the literal lll by this rule.
Special case of rule <i>se1</i> for literals (RDFS)	<i>gl</i> (literal instantiation)	uuu aaa .:nnn ³ where .:nnn identifies a blank node allocated to the literal lll by rule <i>lg</i> .	uuu aaa lll .
	<i>nf1</i>	uuu aaa yyy .	aaa rdf:type rdf:Property . .:nnn rdf:type rdf:XMLLiteral ³
	<i>nf2</i>	uuu aaa lll ., where lll is a well-typed XML literal .	where .:nnn identifies a blank node allocated to lll by rule <i>lg</i> . .:nnn rdf:type rdfs:Literal ³
RDF entailment rules	<i>nf3</i>	uuu aaa lll ., where lll is a plain literal (with or without a language tag).	where .:nnn identifies a blank node allocated to lll by rule <i>lg</i> . .:nnn rdf:type rdfs:Literal ³
	<i>nf4</i>	aaa rdfs:domain xxx .	where .:nnn identifies a blank node allocated to lll by rule <i>lg</i> .
	<i>nf5</i>	uuu aaa yyy .	uuu rdf:type xxx .
	<i>nf6</i>	aaa rdfs:range xxx .	vvv rdf:type xxx .
	<i>nf7</i>	uuu aaa vvv .	uuu rdf:type rdfs:Resource .
	<i>nf8</i>	uuu aaa xxx .	vvv rdf:type rdfs:Resource .
	<i>nf9</i>	uuu aaa vvv .	uuu rdfs:subPropertyOf xxx .
	<i>nf10</i>	vvv rdfs:subPropertyOf xxx .	uuu rdfs:subPropertyOf uu .
	<i>nf11</i>	aaa rdfs:subPropertyOf bbb .	uuu bbb yyy .
	<i>nf12</i>	uuu aaa yyy .	uuu rdfs:subClassOf rdfs:Resource .
	<i>nf13</i>	uuu rdf:type rdfs:Class .	vvv rdf:type xxx .
	<i>nf14</i>	uuu rdfs:subClassOf xxx .	uuu rdfs:subClassOf uu .
	<i>nf15</i>	vvv rdf:type uu .	uuu rdfs:subClassOf xxx .
	<i>nf16</i>	uuu rdf:type rdfs:Class .	uuu rdfs:subPropertyOf rdfs:member .
	<i>nf17</i>	vvv rdfs:subClassOf vvv .	uuu rdfs:subClassOf rdfs:Literal .

Table C1 The RDF/RDFS entailment rules, reproduced from <https://www.w3.org/TR/rdf11-mt>.

Appendix D Implementation at the macroscopic scale

The ~~Vector Symbolic Architecture (VSA)~~, when implemented using the ~~Neural Engineering Framework (NEF)~~, allows a microscopic simulation of the neuronal processes, at the spiking neuronal network level, of the memorization and processing operations, for which we ~~have~~ developed an abstract symbolic description previously. At a higher scale, when the VSA is implemented as described in Appendix A using linear algebra and permutation operations, we are at a mesoscopic scale, allowing us to perform the same operations, ~~but~~ without explicitizing the neuronal state value and evolution. This is one major advantage of this class of approaches.

A To take this a step further, at a higher macroscopic scale, we could directly consider the previous operations, predicting the results of the different algebraic operations without explicitly working at the vector component level. Let's us described how this approach can be designed and implemented, using what could be called an “algorithmic ersatz.”

Symbol indexing

In the VSA, each symbol of the vocabulary is associated with a d -dimensional random vector. At the macroscopic scale, we only need to register each vector \mathbf{x}_k ~~by using~~ an integer number k , incremented for each new symbol. Weighted symbols of index i ~~of and~~ vector number k_i also have ~~also~~ a “belief” value $\tau_i \in [-1, 1]$, as discussed in subsection 2.2, that is equal to 1 by default. They are ~~are~~ estimated up to a certain normal centered additive noise $\nu(\sigma)$ of standard deviation σ_i , which is equal to 0 by default, when no approximate operation has been applied ~~on to~~ the symbol. Two symbols may thus have the same vector number but ~~with a~~ different belief levels, or ~~a~~ different noise levels. At the input/output level, the human-readable string (s_i) representation of the symbol is ~~associated- utilized~~, but it is not ~~further~~ considered further here. The associative table of symbols is ~~a- thus a~~ simple associative array data structure of (k_i, τ_i, σ_i) , without ~~explicitizing- the explicitization of~~ the vector value $\tau_i \mathbf{x}_{k_i} + \nu(\sigma_i)$.

Symbol noise derivation

At the mesoscopic scale, calculations are made up to the floating-point machine precision, which is not taken into account here. ~~Operations- The operations~~ rely on the fact that we consider random vectors in a high-dimensional space, and thus they are approximately orthogonal up to, up to the ~~1st- first~~ order, a normal centered additive noise. **Author query: The meaning of the phrase “approximately orthogonal up to, up to the first order, a normal centered additive noise” is not clear. Do you mean something like “approximately orthogonal, up to the first order, for normal centered additive noise”? The main operations are the dot- product used**

to calculate the similarity, as detailed in subsection 2.1, and the approximation of the matrix inverse by using its transpose; for unbinding, as detailed in Appendix A.

We must thus consider a level of noise, for each symbol, and update this level of noise after each calculation; this noise, up to the first order, can still be represented by a centered normal distribution. This cannot be neglected, because we also introduce a level of belief value that can be small, and thus is not negligible with respect to the noise level. We write denote by $\sigma_{\bullet} \stackrel{\text{def}}{=} O(1/d)$ the order of magnitude added by an approximate operation, as discussed previously in this paper.

On the one hand, considering the similarity operation between two symbols, we obtain⁴¹ for the dot-product:

$$(\tau_i \mathbf{x}_{k_i} + \nu(\sigma_i)) \cdot (\tau_j \mathbf{x}_{k_j} + \nu(\sigma_j)) = \tau_i \tau_j \delta_{k_i=k_j} + \nu(\sigma_{ij}), \sigma_{ij} < \sigma_i + \sigma_j + \sigma_{\bullet},$$

up to the first order, and considering that the noise is independent of the vector values, up to the 1st-first order. Here, $\delta_{\mathcal{P}} = 1$ if \mathcal{P} is true and it is false otherwise. We can thus perform this operation without explicitly computing the dot-product.

Here, we propose a conservative choice, by proposing an upper-bound for the noise, while the exact value, up to the 1st-first order, of σ_{ij} can also easily be used. This design choice is also conservative with respect to a mesoscopic implementation, because it increases the noise at each operation, whereas at the mesoscopic level, each numerical random vector is drawn once; thus, depending on the combination of operations, the noise may not increase. We

⁴¹The derivation writes-is written as follows:

$$\begin{aligned} ((\tau_i \mathbf{x}_{k_i} + \nu(\sigma_i)) \cdot (\tau_j \mathbf{x}_{k_j} + \nu(\sigma_j))) = \\ \tau_i \tau_j \mathbf{x}_{k_i} \cdot \mathbf{x}_{k_j} + \tau_i \mathbf{x}_{k_i} \cdot \nu(\sigma_j) + \tau_j \mathbf{x}_{k_j} \cdot \nu(\sigma_i) + \nu(\sigma_i) \cdot \nu(\sigma_j). \end{aligned}$$

If $k_i \neq k_j$, then $\mathbf{x}_{k_i} \cdot \mathbf{x}_{k_j} = \nu(\sigma_{\bullet})$ since these random vectors are approximately orthogonal up to a-normal noise of-with a standard deviation with an order of magnitude of σ_{\bullet} [52], whereas if $k_i = k_j$, then $\mathbf{x}_{k_i} \cdot \mathbf{x}_{k_j} = 1$ since these are unary vectors.

Then, $\mathbf{x}_{k_i} \cdot \nu(\sigma_j)$ is the dot product, and it is a random variable of mean $\mathbb{E}[\mathbf{x}_{k_i} \cdot \nu(\sigma_j)] = 0$, since vectors are assumed to be independent of other sources of noise up to the 1st-first order, and of-variance $\mathbb{E}[\mathbf{x}_{k_i}^T \nu(\sigma_j) \nu(\sigma_j)^T \mathbf{x}_{k_i}] = \sigma_j^2$ since the covariance $\nu(\sigma_j) \cdot \nu(\sigma_j)^T = \sigma_j^2 \mathbf{I}$ because the noise is isotropic; meanwhile, $\mathbf{x}_{k_i}^T \mathbf{x}_{k_i} = 1$ since it is a unary vector. Note that here, \mathbf{x}_{k_i} is not a random variable; but it stands for the mean of the random vector drawn.

The derivation for $\mathbf{x}_{k_j} \cdot \nu(\sigma_i)$ is identical.

Assuming that $\nu(\sigma_i)$ and $\nu(\sigma_j)$ are independent and of zero mean, as hypothesized, its related variance is known as-to be $\sigma_i^2 \sigma_j^2$, noticing-that Author query: In the previous sentence, it is not clear what "its" is referring to. The product of these two normal distributions is not a normal distribution; but the instead, it is a linear combination of chi-square distributions, in the general case. However, here, as being-it is a second-order term, with respect to the expected small values of σ_i and σ_j , it is negligible. Collecting these results, we obtain that up to the 1st-first order,÷

$$(\tau_i \mathbf{x}_{k_i} + \nu(\sigma_i)) \cdot (\tau_j \mathbf{x}_{k_j} + \nu(\sigma_j)) = \tau_i \tau_j \delta_{k_i=k_j} + \nu \left(\underbrace{|\tau_j| \sigma_i + |\tau_i| \sigma_j + |\tau_i \tau_j| \sigma_{\bullet}}_{\stackrel{\text{dec}}{=} \sigma_{ij}} \right),$$

yielding the expected result.

consider here that noise must be added at each step, and we wonder if ~~it~~ this is not more realistic than a ~~freeze-frozen~~ noise value, although it would have been possible (but quite computationally heavy) to consider freezing noise values and ~~each-caching~~ them in some table.

On the other hand, considering a symbol of index j ~~binded-bound~~ by a symbol of index i and ~~unbinded-unbound~~ by a symbol of index i' so that $k = k_i = k_{i'}$ in order to ~~be-ensure~~ a valid binding/unbinding operation, we obtain, up to the ~~1st-first~~ order⁴²:

$$\mathbf{B}_{(\tau_i \mathbf{x}_i + \nu(\sigma_i))}(\tau_j \mathbf{x}_j + \nu(\sigma_j)) = \tau_i \tau_j \bar{\mathbf{x}}_{ij} + \nu(\sigma'_{ij}), \sigma'_{ij} \leq (1 + \sigma_i + \sigma_j) \sigma_{\bullet}^{\frac{1}{4}},$$

where $\bar{\mathbf{x}}_{ij} \stackrel{\text{dec}}{=} \mathbb{E}[\mathbf{B}_{\mathbf{x}_i} \mathbf{x}_j]$ is a new vector orthogonal to \mathbf{x}_i and \mathbf{x}_j , while the noise related to the binding operation is integrated into σ'_{ij} .

An ~~As an~~ unbinding operation ~~being-is~~ simply a binding operation with a dual vector, the noise calculation is the same.

Compounded symbols

Given atomic symbols that are randomly drawn, using the enumerating operations given in Appendix A, we have to compute compounded symbols composed through bundling, and binding (or unbinding if the symbol has been permuted), while in order to compute entailment rules, we need to be able to compute the similarity between any of these compounded symbols.

At the macroscopic level, since we use only linear algebra, it is straightforward to define an “oracle” that can calculate the result of all operations, as follows:

- A symbol corresponding to the *bundling* of other symbols is fully defined by the symbol set, and an operation over a bundling (binding or similarity) results from applying the operation ~~on-to~~ each element and considering either the bundling (in the case of binding) or numerical sum (in the case of similarity) of the result.
- A symbol corresponding to the *binding* of one symbol onto another is fully defined by the ~~symbols-pair~~ of symbols, and yields either a reduction, if it is the corresponding unbinding operation, or a binding combination. ~~Details~~

⁴²As made explicit in Appendix A, the binding of two independent vectors \mathbf{y} and \mathbf{x} is a random vector and we can write:

$$\begin{aligned} & \mathbf{B}_{(\tau_i \mathbf{x}_i + \nu(\sigma_i))}(\tau_j \mathbf{x}_j + \nu(\sigma_j)) &= \\ & \tau_i \tau_j \mathbf{B}_{\mathbf{x}_i} \mathbf{x}_j + \tau_i \mathbf{B}_{\mathbf{x}_i} \nu(\sigma_j) + \tau_j \mathbf{B}_{\nu(\sigma_i)} \mathbf{x}_j + \mathbf{B}_{\nu(\sigma_i)} \nu(\sigma_j) &= \\ & \tau_i \tau_j \mathbb{E}[\mathbf{B}_{\mathbf{x}_i} \mathbf{x}_j] + \tau_i \tau_j \nu(1/d^{1/4}) + \tau_i \sigma_j \nu(1/d^{1/4}) + \tau_j \sigma_i \nu(1/d^{1/4}) + \sigma_j \sigma_i \nu(1/d^{1/4}) &\simeq \\ & \tau_i \tau_j \bar{\mathbf{x}}_{ij} + (\tau_i \tau_j + \tau_j \sigma_i + \sigma_j \sigma_i) \nu(1/d^{1/4}) &\simeq \\ & \tau_i \tau_j \bar{\mathbf{x}}_{ij} + \nu \left(\underbrace{(|\tau_i \tau_j| + |\tau_i| \sigma_j + |\tau_j| \sigma_i) \sigma_{\bullet}^{1/4}}_{\sigma'_{ij}} \right), \end{aligned}$$

up to the ~~1st-first~~ order, since $\nu(\sigma)$ is a random vector of magnitude σ , while:

$$\mathbb{E}[\mathbf{B}_{\mathbf{x}_i} \nu(\sigma_j)] = \mathbb{E}[\mathbf{B}_{\nu(\sigma_i)} \mathbf{x}_j] = \mathbb{E}[\mathbf{B}_{\nu(\sigma_i)} \nu(\sigma_j)] = 0,$$

because these random vectors correspond to centered random vectors.

The details of the implementation ~~are involve~~ simple applications of the algebraic rules, and we refer the reader to the documented source itself for further details⁴³.

The commutator operator $\mathbf{B}_{\leftrightarrow}$ ~~or and~~ the composition operator \odot_{\neg} are not considered here because they are not used in ~~the this~~ application, but similar considerations would easily allow us to implement the same approach at a macroscopic scale.

Author query: Looking at the example references in the PDF of the template in the folder “sn-article-template”, none of the example journal article references have “accessed” dates. You may want to remove this information from your references. Additionally, if a “doi” website is provided, you probably do not need to provide any additional URLs in a given reference (such as those preceded by the word “eprint”. For some of the book references, a location of publication is missing (which is why “???” appears in the PDF in these references. A location of publication should be added to the bib file as follows: address = {New York}. At least one book reference (Raczaszek-Leonardi) is also missing a publisher.

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