

# LAMBDA CALCULUS

VIVATSATHORN T.

## Lambda Calculus

Syntax :  $\lambda(\text{variable}).(\text{body})$  for abstraction

eg.  $\lambda x. x+10 \leftrightarrow f(x) = x+10$

$$\lambda x. \lambda y. x+y \leftrightarrow f(x,y) = x+y$$

Application :  $(\lambda x. x+10) 20 = 20+10 = 30$

$$\begin{aligned} (\lambda x. \lambda y. x+y) 10 ((\lambda x. x+10) 20) &\rightarrow (\lambda y. 10+y) ((\lambda x. x+10) 20) \\ &\rightarrow 10 + ((\lambda x. x+10) 20) \\ &\rightarrow 10 + (20+10) = 40 \end{aligned}$$

Evaluate 'innermost'.

Renaming :  $\lambda x. [z/x]x \rightarrow \lambda z. z$   
( $\alpha$ -Equiv.)

Substitute free occurrences

$$\lambda x. x =_{\alpha} \lambda y. y \quad : \quad \lambda x. x \rightarrow \lambda y. [y/x]x$$

$\beta$ -Equiv.  
(Body Substitution)  $(\lambda x. M) N \rightarrow [N/x]M$

Recursion  $f(n) = \text{if } n=0 \text{ then } 1 \text{ else } n \cdot f(n-1) \rightarrow \text{factorial}$

$$\text{let } f = \lambda n. \text{if } n=0 \text{ then } 1 \text{ else } n \cdot f(n-1)$$

$$\text{Let } G = \lambda f. \lambda n. \text{if } n=0 \text{ then } 1 \text{ else } n \cdot f(n-1)$$

$$\therefore Gf = \lambda n. \text{if } n=0 \text{ then } 1 \text{ else } n \cdot f(n-1) = f$$

$$\therefore f = G(f)$$

fixed point notation  $(f \text{ is fixed point of } G)$

e.g.  $Y = \lambda f. (\lambda x. f(x x)) (\lambda x. f(x x))$

$$Yf =_{\beta} (\lambda x. f(x x)) (\lambda x. f(x x)) =_{\alpha} (\lambda y. f(y y)) (\lambda x. f(x x))$$

$$=_{\beta} f(\underbrace{(\lambda x. f(x x)) (\lambda x. f(x x))}_{Yf})$$

$$\therefore Yf = f(Yf)$$

fixed point of  $f$

on  $G$  (factorial),  $YG$  is fixed point on  $G$

$$\begin{aligned} f(2) &= YG(2) \\ &= G(YG)2 \\ &= (\lambda f. \lambda n. \text{if } n=0 \text{ then } 1 \text{ else } n \cdot f(n-1)) (YG) 2 \\ &= \lambda n. \text{if } n=0 \text{ then } 1 \text{ else } n(YG(n-1)) 2 \\ &= 2(YG(2-1)) \\ &= 2 \cdot YG(1) \\ &= 2 \cdot G(YG(1)) \\ &= 2 \{ \lambda n. \text{if } n=0 \text{ then } 1 \text{ else } n \cdot YG(n-1) \} \\ &= 2 \cdot 1 \cdot YG(1-1) \\ &= 2 \cdot 1 \cdot YG(0) \dots = 2 \cdot 1 \cdot 1 \quad \checkmark \end{aligned}$$

$\beta$ -reduction : e.g.  $\lambda y. y+1 =_{\beta} (\lambda y. y+1)x = x+1$

\* Function evaluation is called "Confluence" (the property that all reduction sequences lead to the same "normal form")