TASK-8; Normalizing database using functional dependencies uft BCNf To normalize a database using functional dependencies (fDs) up to Boyce codd womal form (Benuf) with the help of a toble wormalization toolas by Aim: -> functional Defendency [FD],

A functional defendency LFD) describes a relation slup between

attributes in a relation. If attribute A determines attribute Normalization
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Normalization
18 the Process of organizing dataina database to B, we write A-7B Normal forom condition All Attributes contain atomic 1 Nf cindivisible valug IN 1NF and all NON 19ey 2 Nf attsibutes are fully defends IN 2NF and no transitive defendency exists for every FDX=>Y, x must be a superkey BCNf examples; consider realtion Ristudent ID, course ID, in stouctor, foom) and the fortiving 705 ; 1, Student ID, CourseID -> instandor, Room 2, instructors -> Poor Hell; Int · INF

gtep2:2Nf the key is I student IP, course ID) no hopial delendences => Alxoody inont, steps: 3NF check for toursitive defendencier. Instructor -> 500m (transitive demondren) Decompose . R, (Student ID, (GUSSEID, in Stoudos) . Re l'instautor (Room) stely; BCNF check BCNf Condition! . In By totadent ID, cows e ID) is a 14ey ->OK. . In Re Instructor - Room, and instructor is the ney-ox

FORMANCE (5) Then selation is successfully normalized 40%

FUCE - coold Normal form (BCNF).