

## Assignment - 2

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VTU : 30283

Course code : 10211CA202

Course title : Data Base management  
Systems.



### 3. Third Normal form (3NF):

Rule:

\* Table must be in 2NF

\* No transitive dependency (non-key attribute should not depend on another non-key attribute)

Example (Achieve 3NF):

Student ID	Student Name	Dept ID	Dept Name
1	Ravi	D <sub>1</sub>	Maths

Here, Dept Name depends on Dept ID, not directly on Student ID.

Should Student table.

Student ID	Student Name	Dept ID
1	Ravi	D <sub>1</sub>

Department table

Dept ID	Dept Name
D <sub>1</sub>	Maths

### 4. Boyce-Codd Normal form (BCNF):

Rule:

\* Table must be in 3NF

\* For every functional dependency  $(X \rightarrow Y)$ , X should be a superkey

Example:

Course	Instructor	Room
DBMS	Raj	R <sub>1</sub>
DBMS	Ravi	R <sub>1</sub>



Explain about deadlock and its handling methods

A deadlock is a condition in a multiprogramming system where two or more processes are waiting indefinitely for resources held by each other, causing all of them to remain blocked forever.

Ex:

- Process  $P_1$  holds Resources  $R_1$  and waits for  $R_2$

- Process  $P_2$  holds Resources  $R_2$  and waits for  $R_1$

→ Both processes wait forever deadlock occurs.

Necessary conditions for deadlock

A deadlock can occur only if all these four conditions hold simultaneously.

1. Mutual Exclusion: only one process can use a resource at a time.

2. Hold and wait: A process holding a resource waiting for others.

3. No preemption: Resources cannot be forcibly taken from a process.

4. Circular wait: A circular chain of processes exists, each waiting for a resource held by the next process.

Deadlock Handling methods:

1. Deadlock Prevention: Design the system in such a way that at least one of the four necessary conditions never occurs.

Ex: Don't allow hold-and-wait.



n) Normalization and its various types of Normalization  
Normalization is a process of organizing the data in database to avoid data redundancy and improve data integrity.

It divides large tables into smaller related tables and links them using foreign keys.

Objectives:

1. To eliminate data redundancy.
2. To avoid update, insert, and delete anomalies.
3. To ensure data consistency and integrity.
4. To make database structure simple and efficient.

Types of Normalization:

Rule:

- \* Each column should contain atomic (indivisible) values
- \* No repeating groups or arrays are allowed

Ex: (Before 1NF)

Student ID	Name	Subjects
1	Ravi	Math, Science

(After 1NF):

Student ID	Name	Subject
1	Ravi	Math
1	Ravi	Science



## Explanation:

- \* Inner query finds Dept ID of CSE  $\rightarrow 101$
- \* Outer query selects all students with Dept ID 101

Result: Row, Arun.

Example: ~ Nested query with IN

Find students who belong to CSE or ECE

```
SELECT Name  
FROM STUDENTS  
WHERE DEPTID IN (
```

```
SELECT DEPTID
```

```
FROM Department  
WHERE DeptName IN ('CSE', 'ECE')
```

```
);
```

## 2. Join:

- A join is used to combine data from two or more tables based on a related column

### Types of Join:

a) INNER join

- \* Returns only the matching rows from

both tables

```
SELECT Students.Name,  
Department.DeptName  
FROM Students
```



## RAID storage and its types

RAID stands for Redundant Array of Independent Disks.

It is a data storage technology that combines multiple physical hard drives into a single logical unit to improve performance, fault tolerance, and data reliability.

RAID is mainly used in servers and high performance systems where data availability and speed are important.

### Objectives of RAID

- \* To increase storage capacity

- \* To improve speed

- \* To provide data redundancy

- \* To enhance system reliability.

### Working principle:

RAID uses a technique called data striping (dividing data into blocks and storing them across multiple disks) and parity (extra information used for recovery if one disk fails).

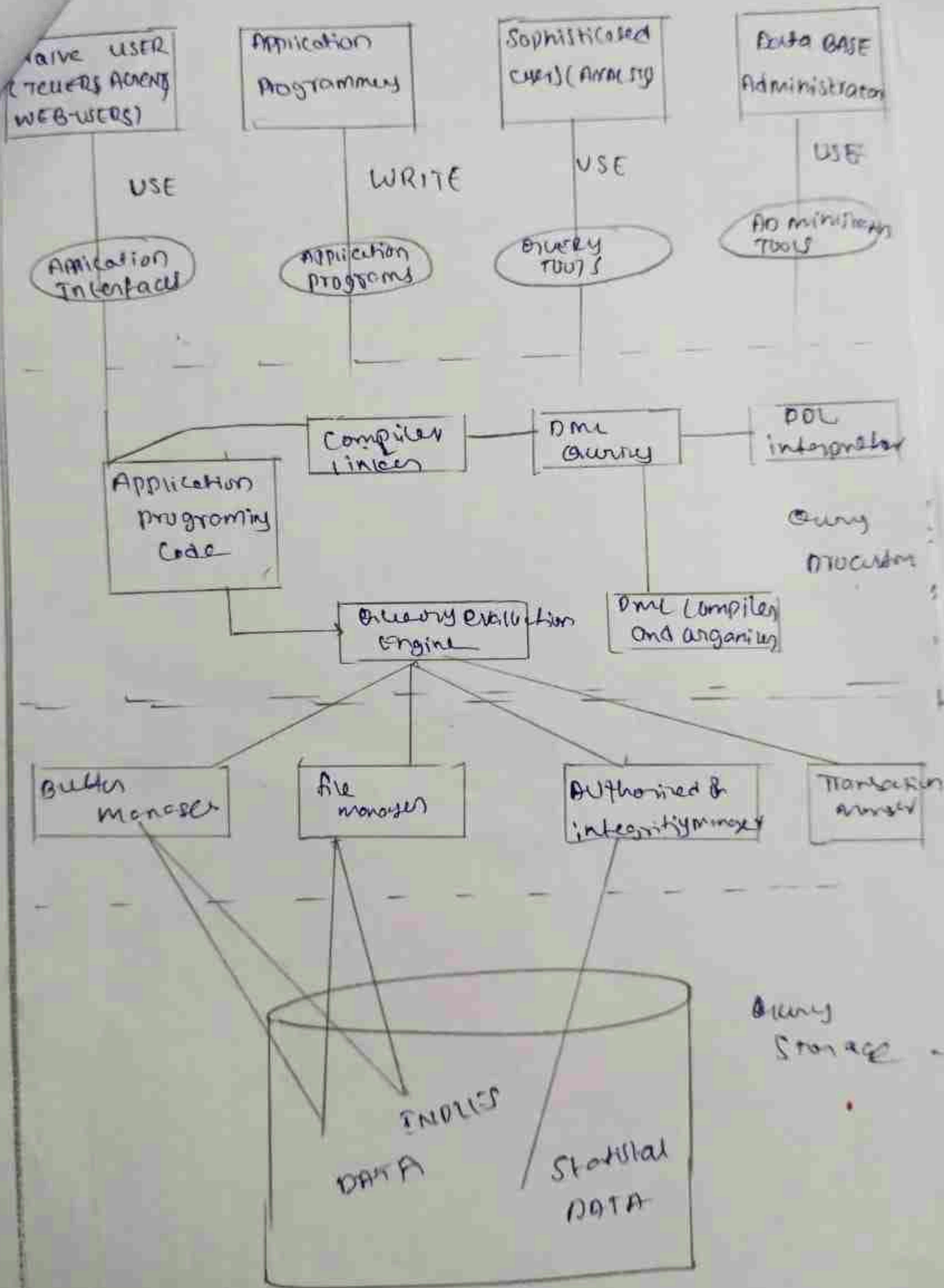
### Types of RAID levels:-

RAID 0 (Striping): Data is divided into blocks and stored across multiple disks.

RAID 1 (Mirroring): Some data is copied on two or more disks.

RAID 2 (Bit-level striping): Data is striped at bit level with error correction code (ECC).







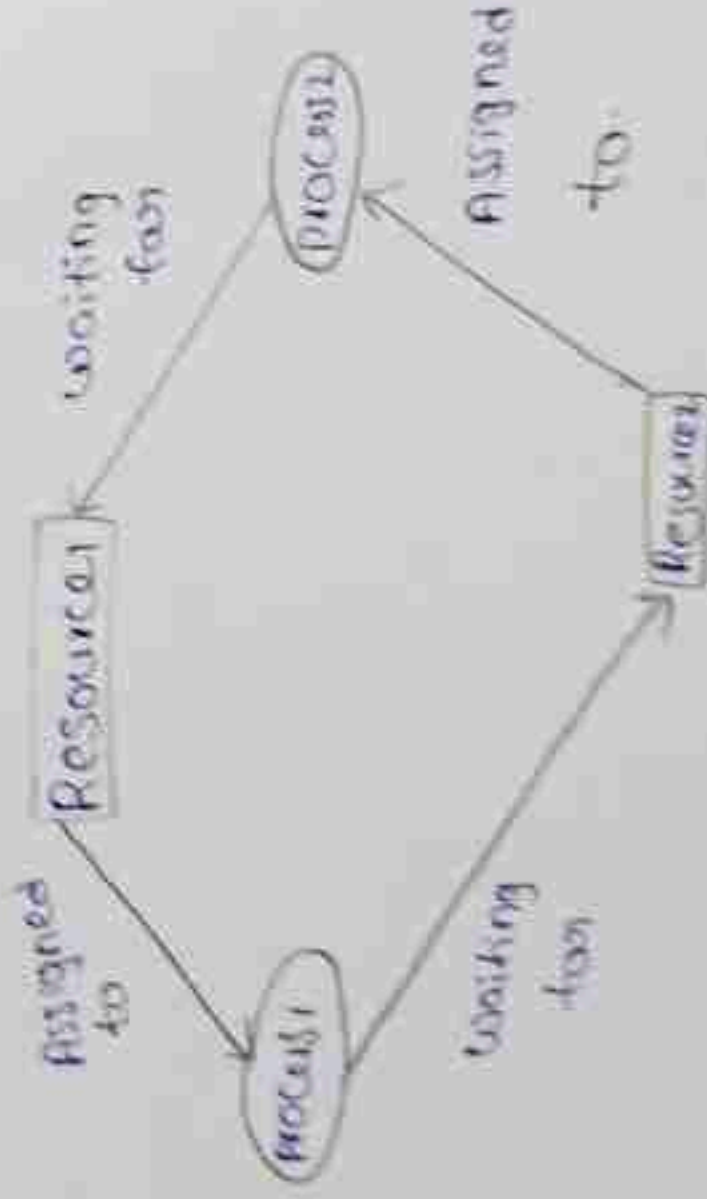
4. Deadlock Avoidance: Dynamically check resource requests to ensure the system never enters an unsafe state.

Ex: Banker's algorithm.

3. Deadlock Detection and Recovery: Allow deadlock to occur detect it using an algorithm and then recover.

4. Deadlock Ignorance: The system assumes that deadlock never occurs in most OS because deadlocks are rare.

Deadlock handling is essential in operating systems to ensure smooth execution of process and efficient resource utilization. Different techniques are used depending on system requirements and complexity.





Explain in detail about Data Base architecture with neat diagram.

Database Architecture describes how a database management system is designed, how its components are arranged, and how users interact with the database.

It defines:

- \* physical Structure
- \* logical structure
- \* External Interaction.

Key Components of Database Architecture

1. Users : End users or applications that request data.
2. DBMS software : manages data storage, retrieval, security, & concurrency.
3. Database (storage) - stores actual data in files, tables, and indexes.
4. Application layer: handles business logic and rules.



Explain in detail about nested queries & solve  
with suitable examples

\* A nested query (or subquery) is a query inside  
another query.

\* The inner query runs first and its result is used  
by the outer query.

\* They are usually written in the where, having  
or from clause.

Example :

Student ID	Name	Age	Dept ID	Dept Name	Phone No
1	Ravi	20	101	CSE	9876543210
2	Sita	21	102	ECE	9876765432
3	Arjun	22	103	Mechanical	9876543210
4	Meena	23	101	Civil	9876543210

Example 1 : Simple Nested Query -

Find the students who belong to CSE department

```
SELECT Name  
FROM students
```

```
WHERE Dept ID = 1
```

```
SELECT Dept ID  
FROM Departments  
WHERE Department Name = 'CSE'  
)
```



RAID 3 (Byte-level Striping with parity): Data is striped at byte level and a separate disk stores parity bit.

RAID 4 (Block-level Striping with parity): Parity information is stored on a dedicated disk.

RAID 5 (Block-level Striping with distributed parity):

Parity is distributed across all disks.

RAID 6 (Double Distributed Parity): Similar to RAID 5 but uses two parity blocks for extra protection.

RAID 10 (combination of RAID 1 + RAID 0): Combined mirroring and striping.

RAID storage technology provides a balance between performance and reliability.

The selection of RAID level depends on system needs, whether speed/cost and data safety is the priority.

RAID 1





(a) presentation layer (Client layer).

- \* User interface: web browser, mobile app, desktop app.

- \* Collects user input and Shows results

(b) Application Layer (Business logic layer)

- \* middle tier that processes requests

- \* performs authentication, validation, calculation and sends

(c) Database Layer (DBMS layer).

- \* Stores actual data.

- \* DBMS like MySQL, Oracle, PostgreSQL manages data, concurrency, backup, and security.

Advantages

- \* High security (Users never directly access database).

- \* Scalability (Supports thousands of users).

- \* Better Performance with Caching & load balancing



# Assignment

Name:- OP Pavanif

VTU :- 30283

Sub :- DBMS

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VTU



on students. Dept ID = Department. Dept ID;  
ID didn't exist, Dept Name would be null

(c) Right join

\* Opposite of LEFT join

\* Returns all rows from departments,  
and matched students.



## Full Outer Join

\* Returns all rows from both tables.  
(match or no match.)

\* In my SQL we simulate it with Union  
SELECT Students. Name,  
Departments. Dept Name

FROM Students

LEFT JOIN Departments

ON Students. Dept ID = Departments

ON Students. Dept ID = Departments Dept ID;

Result's

Ravi CSE

Sita ECE

Rajan CSE

Meena Mechanical

Left join (outer join).

Returns all rows from the left table (Students)  
and matched rows from departments.

if no match  $\rightarrow$  NULL

SELECT Students. Name,

Departments. Dept Name

FROM Students.

Left Join Departments.



5) Fifth Normal form (5NF):

Rule

\* Table must be in 4NF

\* Should not have join dependency.

Ex:-

Relation P = {Subject, Lecturer, Semester}

Subject	Lecturer	Semester
Computer	Anshika	Semester 1
Computer	John	Semester 1
Math	John	Semester 1
Math	John	Semester 1
Chemistry	Proven	Semester 1

$P_1 = \{Semester, Subject\}$

$P_2 = \{Subject, Lecturer\}$

$P_3 = \{Semester, Lecturer\}$

All three relations are now in 5NF



# Assignment

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Sub:-

BBMS







Fourth Normal form (4NF):

ex:

\*Table must be in BCNF

\*There should be no multi valued dependencies

Ex:

Student	Hobby	Language
Ravi	Cricket	English
Ravi	music	Hindi

Here, "Hobby" & "Language" are independent + multi-valued facts.

After 4NF:

Student Hobby:

Student	Hobby
Ravi	cricket
Ravi	music

Student language:

=

Student	Language
Ravi	English
Ravi	Hindi

Q) Second Normal form (2NF):

Rule:

- \* Table must be in 1NF
- \* No partial dependency - i.e., a non-key attribute should not depend on part of a composite key

Ex: (Before 2NF):

Student ID	Course ID	Student Name	Course Name
1	C <sub>1</sub>	Ravi	DBMS

Here:

- \* Student ID + Course ID = Primary Key.
- \* Student Name depends only on Student ID (Partial dependence)

After (2NF): Student table:

Student ID	Student Name
1	Ravi

Course table:

Course ID	Course Name
C <sub>1</sub>	DBMS

Enrollment table:

Student ID	Course ID
1	C <sub>1</sub>