

Assignment - 1

Name : OP.Pavani

VTU : 30283

Course code : 10211CA202

Course title : Data Base management
systems.



3. Third Normal form (3NF):

Rule:

* Table must be in 2NF

* NO transitive dependency (non-key attribute doesn't depend on other non-key attribute)

Example (Before 3NF):

Student ID	Student Name	Dept ID	Dept Name
2	Ravi	D ₁	Maths.

Here, DeptName depends on DeptID & not directly on StudentID after 3NF

Should Student table

Student ID	Student Name	Dept ID
2	Ravi	D ₁

Department table

Dept ID	Dept Name
D ₁	Maths.

4. Boyce-Codd Normal form (BCNF):

Rule:

* Table must be in 3NF

* for every functional dependency ($X \rightarrow Y$), X should be a superkey

Example:

Course	Instructor	Room
DBMS	Raj	R ₁
DBMS	Ravi	R ₁

Explain about deadlock and its handling in DBMS

A deadlock is a condition in a multiprogramming system where two or more processes are waiting indefinitely for resources held by each other, causing all of them to remain blocked forever.

Eg:

- process P₁ holds Resource R₁ and waits for R₂
 - process P₂ holds Resource R₂ and wait for R₁
- Both process wait forever deadlock occurs

Necessary conditions for deadlock

A deadlock can occur only if all these four conditions hold simultaneously.

1. Mutual Exclusion: only one process can use a resource at a time

2. Hold and Wait: A process holding a resource waiting for others.

3. No preemption: Resources cannot be forcibly taken from a process

4. Circular Wait: A circular chain process exists, each waiting for a resource held by the next process

Deadlock Handling methods:

1. Deadlock Prevention: Design the system in such a way that at least one of the four necessary conditions never occurs.

Eg: Don't allow hold-and-wait

i) Normalization and its various types of Normalization
Normalization is a process of organizing the data in database to avoid data redundancy and improve data integrity.

It divides large tables into smaller related tables and links them using foreign keys.

Objectives:

1. To eliminate data redundancy.
2. To avoid update, insert, and delete anomalies.
3. To ensure data consistency and integrity.
4. To make database structure simple and efficient.

Types of Normalization:

Rule:

- * Each column should contain atomic (indivisible) values
- * No repeating groups or arrays are allowed

Ex: (Before 1NF)

Student ID	Name	Subjects
1	Ravi	Math, Science

(After 1NF):

Student ID	Name	Subject
1	Ravi	Math
1	Ravi	Science

Explanation:

- * Inner query finds Dept ID of CSE \rightarrow 101
- * Outer query selects all Students with Dept ID 101

Result: Ravi, Arun.

Example: ~ Nested query with IN

Find Students who belong to CSE or ECE

```
SELECT Name  
FROM STUDENTS  
WHERE DEPTID INL.
```

```
SELECT DEPT ID  
FROM Department.  
WHERE DeptName IN ('CSE', 'ECE')
```

```
);
```

2. Joins:

- A join is used to combine data from two or more tables based on a related column

Types of joins:

a) INNER JOIN

- * Returns only the matching rows from both tables

```
SELECT STUDENTS.Name,  
Department.DPT_Name  
from STUDENTS
```

RAID storage and its types

RAID stands for Redundant Array of Independent Disks. It is a data storage technology that combines multiple physical hard drives into a single logical unit to improve performance, fault tolerance, and data reliability.

RAID is mainly used in servers and high performance systems where data availability and speed are important.

objectives of RAID

- * To increase storage capacity

- * To improve speed

- * To provide data redundancy

- * To enhance system reliability.

Working principle:

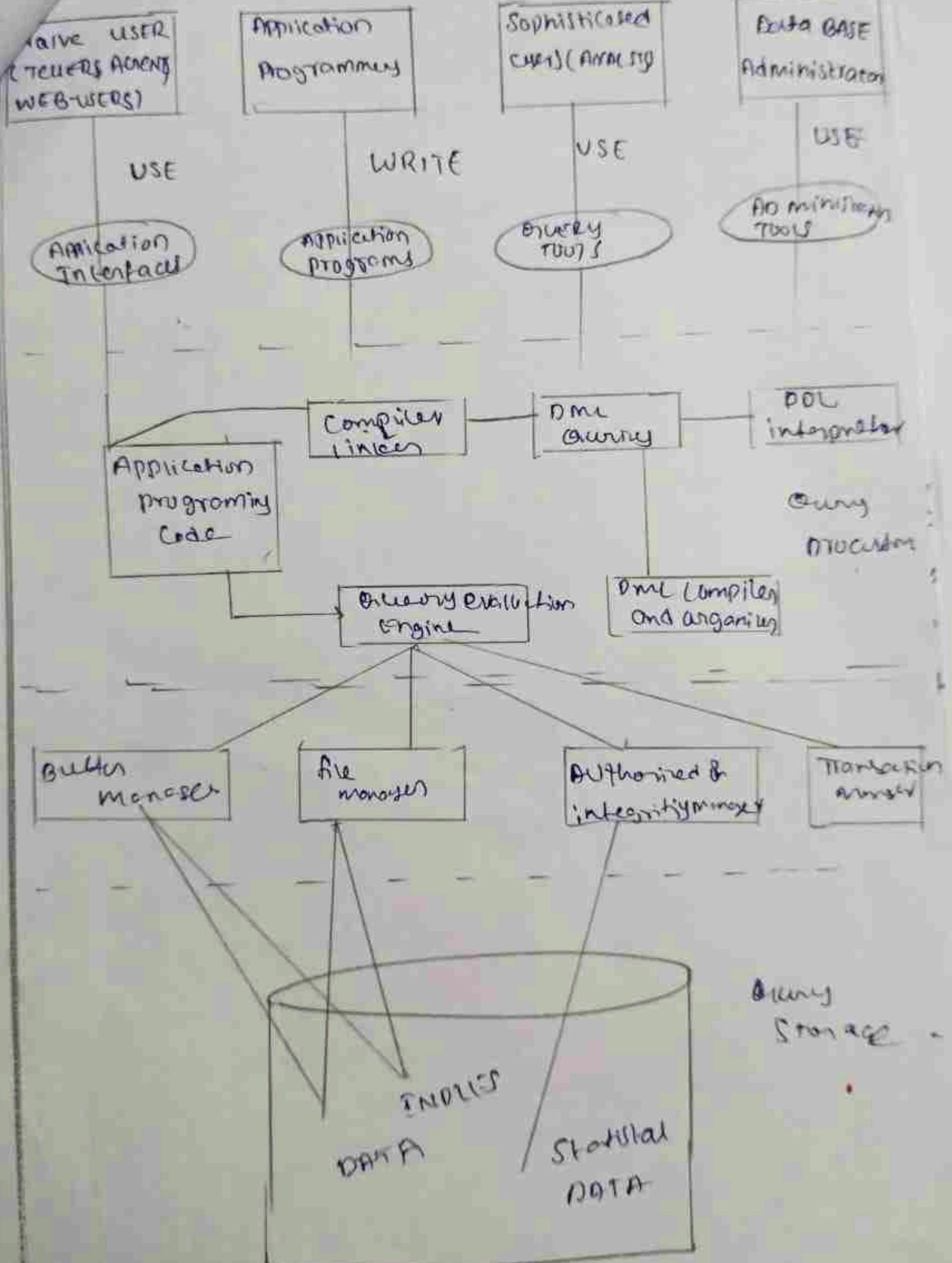
RAID uses a technique called data striping (dividing data into blocks and storing them across multiple disks) and parity (extra information used for recovery if one disk fails).

Types of RAID levels:

RAID 0 (striping): Data is divided into blocks and stored across multiple disks.

RAID 1 (mirroring): Some data is copied on two or more disks.

RAID 2 (bit-level striping): Data is striped at bit level with error correction code (ECC).



- 4. **deadlock avoidance**: maintaining their resources.
means prevent the system objects can an
infect state.
- 5. **deadlock detection and recovery**: detect deadlock to
occur, then taking an organization and the
resources.
- 6. **deadlock tolerance**: The system assumes that
deadlock never occurred in most of because deadlocks
are rare.

Deadlock handling is essential in operating systems
to ensure smooth execution of process and efficient
resource utilization. Different techniques are used
depending on system requirements and complexity.



Explain in detail about Data Base architecture with neat diagram.

Database Architecture describes how a database management system is designed, how its components are arranged, and how users interact with the database.

It defines:

- * Physical Structure
- * Logical Structure
- * External Interaction.

Key Components of Database Architecture

1. Users : End users or applications that request data.
2. DBMS software : manages data storage, retrieval, security, & currency.
3. Database (Storage) - stores actual data in files, tables, and indexes.
4. Application layer: handles business logic and rules.

explain in detail about nested queries & solve with suitable examples

- * A nested query (or subquery) is a query inside another query.
- * The inner query runs first and its result is used by the outer query.
- * They are usually written in the where, having or from clause.

Example :

Student ID	Name	Age	Dept ID	Dept Name	Phone No
1	Ravi	20	101	CSE	9876543210
2	Sita	21	102	ECE	9876765432
3	Arjun	22	103	Mechanical	9876543211
4	Meena	23	101	Civil	9182763212

Example : simple Nested Query -

Find the students who belong to CSE department

SELECT Name

FROM Students

WHERE Dept ID =

SELECT Dept ID

FROM Departments

WHERE Department = 'CSE'

)

RAID 3 (Byte-level striping with parity): Data is striped at byte level and a separate disk stores parity info.

RAID 4 (Block-level striping with parity): Parity information is stored on a dedicated disk.

RAID 5 (Block-level striping with distributed parity): parity is distributed across all disks.

RAID 6 (Double Distributed parity): Similar to RAID 5 but uses two parity blocks for extra protection.

RAID 10 (combination of RAID 1 + RAID 0): Combined mirroring and striping.

RAID storage technology provides a balance between performance and reliability.

The selection of RAID level depends on system needs - whether speed / cost or data safety is the priority.

RAID 1



3) presentation layer (client layer).

- * user interface: web browser, mobile app, desktop app.

- * collects user input and shows results

(b) application layer (Business logic layer)

- * middle tier that processes requests

- * performs authentication, validation (calculation and send)

(c) Database layer (DBMS layer).

- * stores actual data.

- * DBMS like MySQL, Oracle, PostgreSQL manages data, concurrency, backup, and security.

Advantages

- * High security (users never directly access database).

- * Scalability (supports thousands of users).

- * Better performance with caching & load balancing

Assignment

Name: OP Ravoni

vrn: 30283

sub: DBMS

5
Ravoni

on students . Dept ID = DepartmentID. Dept ID;
ID didn't exist , Dept name would be null

c) Right join

- * opposite of LEFT JOIN
- * Returns all rows from Departments,
and matched Students.

full outer join

- * Returns all rows from both tables.
(match or no match.)
 - * In my SQL we simulate it with Union
- ```
SELECT Students.Name,
Departments.Dept Name
FROM Students
LEFT JOIN Departments
ON Students.Dept ID = Departments
ON Students.Dept ID = Departments.Dept ID;
Result:
```

Ravi CSE

Sita EEE

Arun CSE

Meena Mechanical

(left join (outer join))

Returns all rows from the left table (Students)  
and matched rows from departments.

if no match → NULL

```
SELECT Students.Name,
Departments.Dept Name
FROM Students
LEFT JOIN Departments;
```

by fifth normal form (5NF):

with

a table must be in 4NF

it should not have "join dependency".

Ex:-

Relation R = {Subject, Lecture, Semester}

| Subject   | Lecturer | Semester   |
|-----------|----------|------------|
| Computer  | Anshika  | Semester 1 |
| Computer  | John     | Semester 1 |
| Math      | John     | Semester 1 |
| Math      | John     | Semester 1 |
| Chemistry | Praveen  | Semester 1 |

R<sub>1</sub> = {Semester, Subject}

R<sub>2</sub> = {Subject, Lecturer}

R<sub>3</sub> = {Semester, Lecture}

All three relations are now in 5NF

Assignment

Name :- op. pavani

Wtu :- 20283

subj :- OBMS

(3)

(5)

Fourth Normal form (4NF):

Rule:

\* Table must be in 3NF

\* There should be no multi valued dependencies

Ex:

| Student | Hobby   | language |
|---------|---------|----------|
| Ravi    | Cricket | English  |
| Ravi    | MUSIC   | Hindi    |

Here, "Hobby" & "language" are independent multi-valued facts.

After 4NF:

Student Hobby:

| Student | Hobby   |
|---------|---------|
| Ravi    | cricket |
| Ravi    | MUSIC   |

Student language:

==

| Student | language |
|---------|----------|
| Ravi    | English  |
| Ravi    | Hindi    |

Q) Second Normal form(2NF);

Rule:

- \* Table must be in 1NF
- \* No partial dependency - i.e., a non-key attribute should not depend on part of a composite key.

Ex: (Before 2NF):

| Student ID | Course ID      | Student Name | Course Name |
|------------|----------------|--------------|-------------|
| 1          | C <sub>1</sub> | Ravi         | DBMS        |

Here:

- \* Student ID + Course ID = Primary Key.
- \* Student Name depends only on Student ID (Partial dependence)

After(2NF): Student table:

| Student ID | Student Name |
|------------|--------------|
| 1          | Ravi         |

Course table:

| Course ID      | Course Name |
|----------------|-------------|
| C <sub>1</sub> | DBMS        |

Enrollment table:

| Student ID | Course ID      |
|------------|----------------|
| 1          | C <sub>1</sub> |