

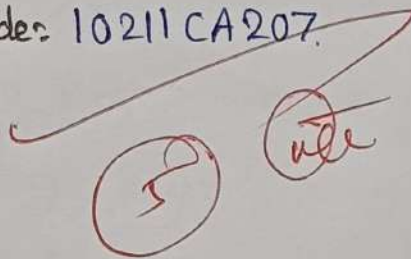
# — Data Base Management System —

## Assignment

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Course Code: 10211CA207.



1. Explain in detail about data base system architecture with a neat diagram.

The database system Architecture can be divided into three main parts: Users, Query Processor and Storage manager, with the disk storage at the bottom.

#### 1. Users:

Different types of user interact with the database:

- Naive-users (browsers, agents, web users) → use application interfaces.
- Application Programmers → write application programs.
- Sophisticated users (analysts) → use query tools.
- Database Administrators (DBAs) → use the administration tools.

#### 2. Query Processor.

This is responsible for interpreting and executing the queries.

- Application program object code:  
Generated by compiler and linker
- DML Queries: Data Manipulation language.  
(E.g: SELECT, INSERT, UPDATE, DELETE)
- DDL Interpreter: Interprets schema definition (table, constraints).

- DML Compiler and organizer: Checks syntax and optimizes queries.

- Query Evaluation Engine: Executes optimized queries.

### 3. Storage Manager:

The Storage manager controls how data is stored and Retrieved

- Buffer Manager: Minimizes disk I/O by storing frequently accessed data in memory

- File manager: Manages allocation of space and file structures.

- Authorization and Integrity: Ensures data security and Integrity constraints.

- Transaction Manager: Ensures consistency concurrency control and recovery.

### 4. Disk storage:

This is the physical layer where actual data resides

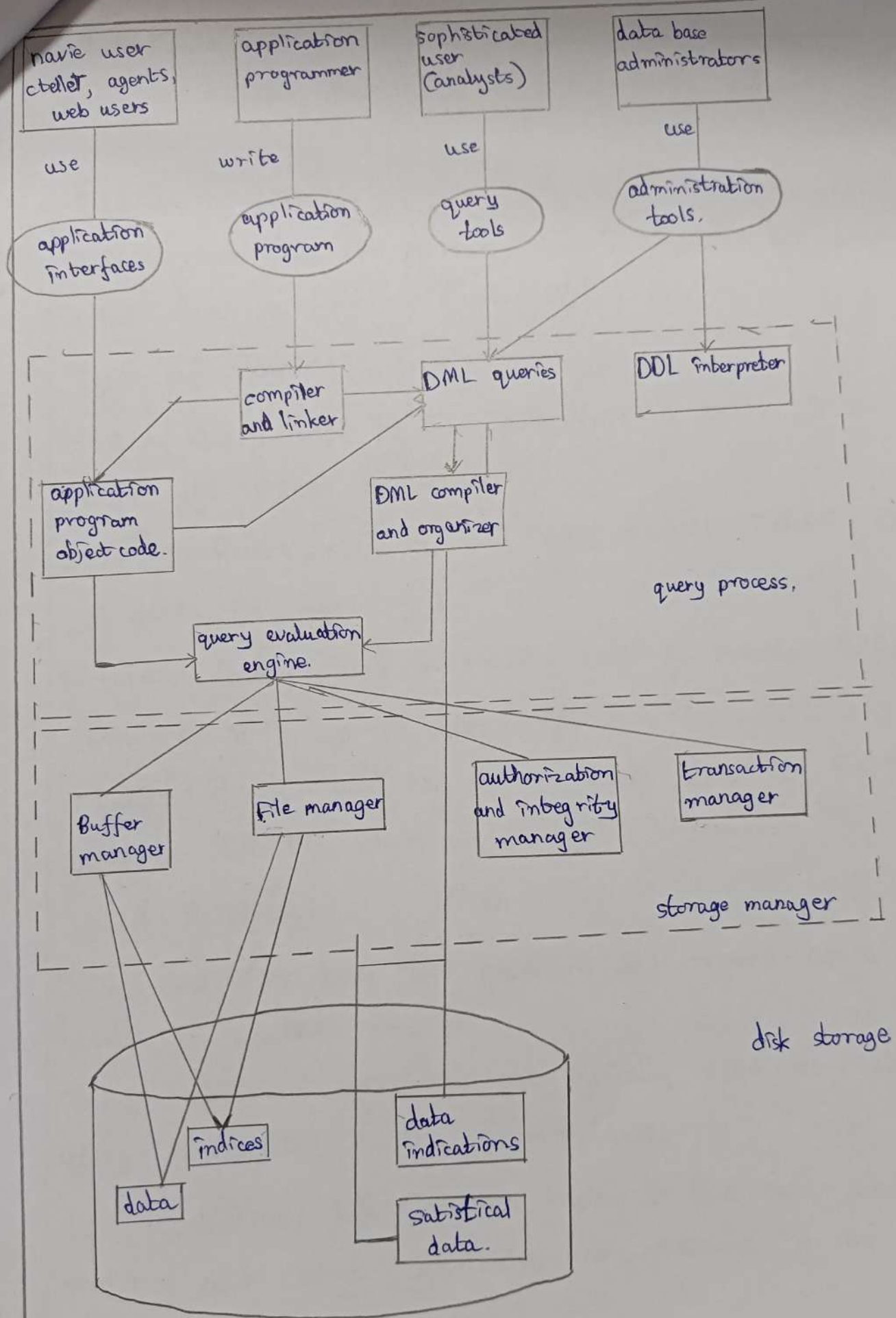
- Data: Tables and records.

- Indices: Used for fast searching

- Data Dictionary: Stores meta data (information about tables, schemes, users).

- Statistical Data: Used for query optimization.





2. Explain various Queries and Joins with suitable example.

### Queries in DBMS:

Queries are used to retrieve, insert, update and delete data from a database using SQL.

### Types of Queries:

1. Select Query - used to fetch data from a table `SELECT name, age FROM students`.
2. Insert Query - used to add new records `INSERT INTO students (id, name, age)`.
3. Update Query - used to modify existing records `UPDATE student SET age = 21 WHERE id`.
4. Delete Query - used to remove records `DELETE FROM students WHERE id = 1`;

### Joined Relations:

Join operations take two relations and return as a result to another relation.

These additional operations are typically used as sub query expressions in the from clause.

Join conditions - defines which tuples in the two relations match, and what attributes are present in the result of the join



Join type - defines how tuples in each relation that do not match any tuple in the other relation are treated.

Join types
inner join
left outer join
right outer join
full outer join

Join conditions.
natural
on < predicate >
using (A <sub>1</sub> , A <sub>2</sub> , ..., A <sub>n</sub> )

Database Example - Join

Instructor

ID	Name	dept_name
10101	Srinivasan	compt sci
12121	Venkate	Finance
15151	Mozart	Music

Teaches

ID	Course_Id
10101	CS_101
12121	Fin_201
15151	Bio_101

#### • Inner Join:

Returns only the matching rows from both tables based on a given conditions.

Example: Display Instructors who are teaching atleast one course

```
SELECT instructor.ID, name,
Course_id.
```

FROM instructor

INNER JOIN teaches

ON instructor.ID = teaches.ID,

ID	Name	Course id
10101	Srinivasan	CS-101
12121	Venkat	Fin-201

• Left Outer Join:

Returns all rows from the left table and the matching rows from the right table. If no match, NULLs are shown for right table columns.

Example: List all instructors including those who are not teaching any courses.

SELECT instructor.ID, name, course\_id

FROM instructor

LEFT JOIN teaches.

ON instructor.ID = teaches.ID;

ID	Name	Course_id
10101	Srinivasan	CS-101
12121	Venkat	Fin-201
15151	Mozart	NULL

• Right outer Join:

Returns all rows from the right table and the matching rows from the left table. If no match NULLs are shown for left table columns.



Example: Show all courses and their instructor including courses with no assigned instructor  
`SELECT instructor_ID, name, course_id.`

`FROM instructor`

`RIGHT JOIN teaches`

`ON instructor.ID = teaches.ID;`

• FULL Outer Join:

Returns all rows from both tables, matching rows where available and filling NULLs where there is no match.

Example: Display all instructor and all courses even if there is no matching entry in either table.

Left join:

`SELECT instructor_ID, name, course_id`

`FROM instructor`

`LEFT JOIN teaches ON instructor.ID = teaches.ID.`

Union:

`SELECT instructor_ID, name, course_id.`

`FROM instructor.`

`RIGHT JOIN teaches ON instructor.ID = teaches.ID.`

ID	Name	Course_Id.
10101	Srinivasan	CS_101
12121	Venkab	Fin-201
15131	Mozart	NULL
76766	NULL	Bio-101



## EQUAL JOIN:

A type of INNER JOIN that uses an equality (=) operator to match rows.

Example: Find Instructor - Course associations using equality condition.

```
SELECT Instructor_ID, name, Course_Id  
FROM Instructor, teaches  
WHERE Instructor_ID = teaches_ID;
```

ID	Name	Course-Id.
10101	Srinivasan	CS-101
12121	Venkat	Fin-201

## Cross Join:

Returns the cartesian product of two tables every row from the first table joined with every row from the second table.

Example: Generate all possible combinations of Instructor and courses.

```
SELECT name, course_id,  
FROM Instructor,  
CROSS JOIN teaches.
```

Name	Course_id
Srinivasan	CS_101
Srinivasan	Fin_201
Srinivasan	Bro_101
Venkat	CS_101
Venkat	Fin_201
Venkat	Bro_101
Mozart	CS_101
Mozart	Fin_201
Mozart	Bro_101



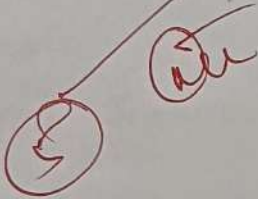
# ← Data Base Management System —

Assignment.

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3) Explain about RAID storage and its types? &

RAID (Redundant Array of Independent ~~disks~~) is a data storage technology that combines multiple physical hard drives into a single logical unit.

The main goals are,

- Improved performance.
- Increased storage capacity
- Data redundancy (fault tolerance).

RAID is commonly used in servers, data centres, and storage systems to ensure data reliability and faster access.

2. Key features.

Feature	Description.
Redundancy	Data is duplicated or spread to protect against drive failure.
Performance	Multiple disks (can read) write data simultaneously.
Fault Tolerance.	System continues to function even if one drive fails.



Striping

Data is copied identically on two or more disks.

Mirroring

Data is copied identically on two or more disks.

Parity

Error checking information is stored to recover data in disk failure.

Types of RAID.

RAID Level: RAID 0 (striping)

Technique used: Data divided into blocks and spread across disks.

Description: Fast performance (no redundancy).

Advantages: High speed

Disadvantages: No fault tolerance. If one disk fails, all data lost.

RAID Level: RAID 1 (mirroring)

Technique used: Same data (copied on two disks)

Description: Provides redundancy.

Advantages: High reliability, simple recovery.

Disadvantages: Storage cost doubles.

RAID level: Bit level stripping with error correction case (ECC).

Technique used: Rarely used

Description: Rarely used

Advantages: Error correction possible.

Disadvantages: Expensive, complex.

RAID Level: RAID 3

Technique used: Byte level stripping with dedicated parity disk.

Description: All drives work together.

Advantages: Good for large sequential data.

Disadvantages: Parity disk may be a bottleneck.

RAID level: RAID 4

Technique used: Block level stripping with dedicated parity disk.

Description: Parity used for recovery

Advantages: Fast read performance.

Disadvantages: Single parity disk bottle neck.



### RAID Level: RAID 5.

Technique used: Block level striping with distributed parity

Description: Most common RAID Level.

Advantages: Good balance b/w performance and redundancy

Disadvantages: Complex rebuild if a disk fails.

### RAID Level: RAID 6.

Technique used: Like RAID 5 but with two parity blocks.

Description: Higher fault tolerance.

Advantages: Can tolerate two disk failure

Disadvantages: Slower write performance.

### RAID Level: RAID 10

Technique used: Combination of RAID 1 & RAID 0

Description:

Advantages:

Disadvantages:

## 2. Deadlock and its handling?

A deadlock occurs in a database when two or more transactions are waiting for each other to release locks on resources, preventing further progress.

Example:

\* Transaction T1 locks resource A and waits for Resource B,  
\* Transaction T2 locks resource B and waits for Resource A.

→ Both wait forever → deadlock.

Conditions for Deadlock. (Coffman Conditions).

1. Mutual Exclusion:  
Only one transaction can use a resource at a time.
2. Hold and wait:  
A transaction holds one resource and waits for another.
3. No preemption:  
Resources cannot be forcibly taken away.
4. Circular wait:  
A circular chain of waiting transactions exists.



## Deadlock Handling Techniques

Method

Description

### 1. Dead Lock Prevention.

Avoids deadlocks by denying at least one Coffman condition.

Example: Assign resource ordering or use time stamp based methods [wait-die, wound-wait].

### 2. Deadlock Avoidance,

The system checks before the granting a lock to ensure no deadlock will occur.

Example: Banker's algorithm checks safe states.

### 3. Deadlock detection,

The system allows deadlocks to occur but periodically checks for them using a wait-for graph (WFG).

### 4. Deadlock Recovery

Once detected, resolve by rolling back one or more transaction victim selects



1. Normalization and its various types of normalization.  
 Normalization and its various types;  
 Normalization is a process in database management system  
 DBMS used to organize data in a database to  
 reduce redundancy (duplicate) data and improve data  
 integrity.

1) First Normal Form (1NF).

→ Each cell must contain atomic (indivisible) values.

→ No repeating groups or arrays allowed.

Ex:

St-ID	St-Name	St-Phone	St-Group
123	Rakesh	9843298760	Physics
125	Amrita	91354628198 8763549321	stat
126	Archana	7362543210 869453210	Chemistry

New table:

St-ID	St-Name	St-Phone	St-Group
123	Rakesh	984073512	Physics
125	Amrita	9135462818	stat
125	Amrita	876354932	stat

## Second Normal Form (2NF).

- In the 2NF, first table must be an 1NF.
- In the 2NF, all non key attributes are fully functionally dependent on the primary key.
- Every non key attributes are full FFD on key attribute.

If  $P \rightarrow A$  holds, then those should not be any proper subset of  $Q$  of  $P$ .

$Q \rightarrow A$

St-ID	st-Name	Prof-ID	Prof-Name	Grade
101	ABC	2	Sameer	4
102	XYZ	3	Mrangan	6
103	PQR	1	Sushmitha	5.

st-ID is Primary key.

and no multi valued are there so satisfies 1NF.

St-ID	st-Name
101	ABC
102	XYZ
103	PQR

Prof-ID	Prof-Name
1	Sushmitha
2	Sameer
3	Mrangan

Grades:

St-ID	Prof-ID	Grade
101	2	4
102	3	6
103	1	5

### 3. Third Normal Form (3NF)

→ A relation will be in 3NF.

→ It not contains any transitive dependency.

→ The non-key attributes should not have inter dependencies among them, and the non-key attributes should fully functionally depend on key attribute.

Then it is called 3NF principle.

→ By using 3NF to achieve data integrity and

Data duplication.

Transitive dependency.

If  $A \rightarrow B$  &  $B$  functionally dependent on  $A$ )

$B \rightarrow C$

$A \rightarrow C$  ( $C$  is indirectly dependent on  $A$ )

If it is called T.D.



Ex:

St-ID	St-Name	Dept. Name	Dept-Head
101	Alice	CSE	Dr. Lalitha
102	Bob	ECE	Dr. Kumar
103	Carol	CSE	Dr. Rao.

Student table.

St-ID	St-Name	Dept-Name.
101	Alice	CSE
102	Bob	ECE
103	Carol	CSE

Dept-table.

Dept-Name	Dept-Head.
CSE	Dr. Lalitha
ECE	Dr. Kumar
CSE	Dr. Rao.

4. BCNF (Boyce-codd Normal Form).

→ It is advanced version of 3NF.

→ It is in 3NF.

→ For every FD  $A \rightarrow B$   $A \rightarrow$  Super key

→ A should be super key of a table.

Ex:

Student	Course	Teacher
Ramesh	Physics	Kishore
Kumar	Chemistry	Ramu
Sruthi	Maths	Saigeev
Vinay	Physics	Kishore.

Keys  $\rightarrow$  {student, course}  $\rightarrow$  ?

To eliminate redundant data we move to 2 tables

Student	Course	Course	Teacher
Ramesh	Physics	Physics	Kishore.
Kumar	Chemistry	Chemistry	Ramu.
Sruthi	Maths	Maths	Saigeev
Vinay	Physics	Physics	

Fourth Normal Form (4NF)

$\rightarrow$  First relation in BCNF (and) 3.NF

$\rightarrow$  It may not contain more than one multivalued attribute

For a dependency.

$A \rightarrow B$ .

$\rightarrow$  If for a single value of A multiple values of B exists, then the relation will be multivalued dependency.

St-ID	Course	Hobbies.
111	Maths	Dancing
111	Computer	Singing
222	Chemistry	Dancing
444	Sanskrit	Crickat
555	Physics	Hockey.
St-ID		
	Course	St-ID Hobbies.
111	Maths	111 Dancing
111	Computer	111 Singing
222	Chemistry	222 Dancing
444	Sanskrit	444 Crickat
555	Physics	555 Hockey.