

— Data Base Management System —

Assignment

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1. Explain in detail about data base system architecture with a neat diagram.

The database system Architecture can be divided into three main parts: Users, Query Processor and Storage manager, with the disk storage at the bottom.

1. Users:

Different types of user interact with the database:

- Naive users (tellers, agents, web users) → use application interfaces.
- Application Programmers → write application programs.
- Sophisticated users (analysts) → use query tools.
- Database Administrators (DBAs) → use the administration tools.

2. Query Processor:

This is responsible for interpreting and executing the queries.

- Application program object code:

Generated by compiler and linker

- DML Queries: Data Manipulation language.

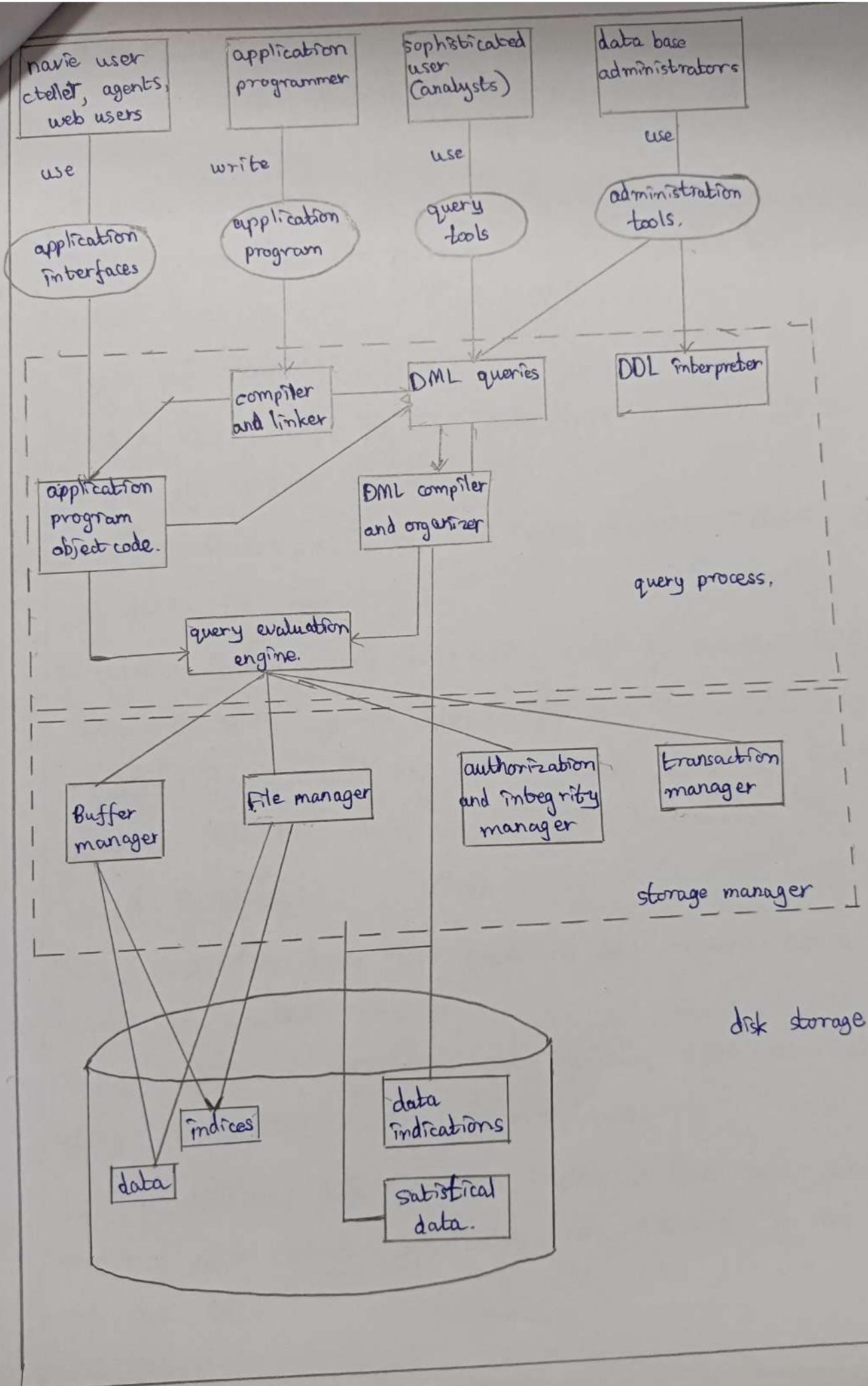
(E.g: SELECT, INSERT, UPDATE, DELETE)

- DDL Interpreter: Interprets schema definition (table, constraints).

- DML Compiler and optimizer: Checks syntax and optimizes queries.
 - Query Evaluation Engine: Executes optimized queries.
3. Storage Manager:
- The storage manager controls how data is stored and retrieved
- Buffer Manager: Minizes disk I/O by storing frequently accessed data in memory
 - File manager: Manages allocation of space and file structures.
 - Authorization and Integrity: Ensures data security and integrity constraints.
 - Transaction Manager: Ensures consistency, concurrency control and recovery.

4. Disk storage:

- This is the physical layer where actual data resides
- Data: Tables and records.
 - Indices: Used for fast searching
 - Data Dictionary: Stores meta data (information about tables, schemes, users).
 - Statistical Data: Used for query optimization.



2. Explain various Queries and Joins with suitable example:-

Queries in DBMS:

Queries are used to retrieve, insert, update and delete data from a database using SQL.

Types of Queries:

1. Select Query - used to fetch data from a table
SELECT name, age FROM students;

2. Insert Query - used to add new records
INSERT INTO students (id, name, age).

3. Update Query - used to modify existing records
UPDATE student SET age=21 WHERE id.

4. Delete Query - used to remove records
DELETE FROM students WHERE id=1;

Joined Relations:

Join operations take two relations and return as a result to another relation.

These additional operations are typically used as sub query expressions in the from clause.

Join conditions - defines which tuples in the two relations match, and what attributes are present in the result of the join

Join type - defines how tuples in each relation that do not match any tuple in the other relation are treated.

| Join types |
|------------------|
| inner join |
| left outer join |
| right outer join |
| full outer join |

| Join conditions. |
|--------------------------------|
| natural |
| on < predicate > |
| using (A_1, A_2, \dots, A_n) |

Database Example - Join

Instructor

| ID | Name | dept_name |
|-------|------------|-----------|
| 10101 | Srinivasan | compt sci |
| 12121 | Venkat | Finance |
| 15151 | Mozart | Music |

Teaches,

| ID | Course_Id |
|-------|-----------|
| 10101 | CS-101 |
| 12121 | Fm-201 |
| 15151 | Bio-101 |

• Inner Join:

Returns only the matching rows from both tables based on a given conditions.

Example: Display instructors who are teaching atleast one course

```
SELECT Instructor.ID, name,
       Course_id.
```

FROM instructor
 INNER JOIN teaches
 ON Instructor.ID = teaches.ID,

| ID | Name | Course ID |
|-------|------------|-----------|
| 10101 | Srinivasan | CS-101 |
| 12121 | Venkat | Fin-201 |

Left Outer Join:

Returns all rows from the left table and the matching rows from the right table. If no match, NULLs are shown for right table columns.

Example: List all instructors including those who are not teaching any courses.

SELECT instructor.ID, name, course_id
 FROM instructor
 LEFT JOIN teaches,

ON Instructor.ID = teaches.ID;

| ID | Name | Course_id |
|-------|------------|-----------|
| 10101 | Srinivasan | CS-101 |
| 12121 | Venkat | Fin-201 |
| 15151 | Mozart | NULL |

Right outer Join:

Returns all rows from the right table and the matching rows from the left table. If no match NULLs are shown for left table columns.

Example: Show all courses and their instructor including courses with no assigned instructor
SELECT instructor.ID, name, course_id.

FROM instructor

RIGHT JOIN teaches

ON instructor.ID = teaches.ID;

FULL Outer Join:

Returns all rows from both tables, matching rows where available and filling NULLs where there is no match.

Example: Display all instructor and all courses even if there is no matching entry in either table.

Left join:

SELECT instructor.ID, name, course_id

FROM instructor

LEFT JOIN teaches ON instructor.ID = teaches.ID.

Union:

SELECT instructor.ID, name, course_id.

FROM instructor

RIGHT JOIN teaches ON instructor.ID = teaches.ID.

| ID | Name | Course-ID. |
|-------|------------|------------|
| 10101 | Srinivasan | CS-101 |
| 12121 | Venkat | Fin-201 |
| 15151 | Mozart | NULL |
| 76766 | NULL | Bio-101 |

EQUI JOIN:

A type of INNER JOIN that uses an equality ($=$) operator to match rows.

Example: Find Instructor-Course associations using equality condition.

```
SELECT Instructor.ID, name, Course_id  
FROM Instructor, Teaches  
WHERE Instructor.ID = teaches.ID;
```

| ID | Name | Course_id |
|-------|------------|-----------|
| 10101 | Srinivasan | CS-101 |
| 12121 | Venkat | FIN-201 |

Cross Join:

Returns the cartesian product of two tables every rows from the first table joined with every row from the second table.

Example: Generate all possible combinations of Instructor and courses.

```
SELECT name, course_id,  
FROM Instructor,  
CROSS JOIN teaches.
```

| Name | Course_id |
|------------|-----------|
| Srinivasan | CS-101 |
| Srinivasan | Fin-201 |
| Srinivasan | BIO-101 |
| Venkat | CS-101 |
| Venkat | Fin-201 |
| Venkat | BIO-101 |
| Mozart | CS-101 |
| Mozart | Fin-201 |
| Mozart | BIO-101 |

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3) Explain about RAID storage and its types?

RAID (Redundant Array of Independent Disks) is a data storage technology that combines multiple physical hard drives into a single logical unit.

The main goals are:

- Improved performance.
- Increased storage capacity
- Data redundancy (fault tolerance).

RAID is commonly used in servers, data centres, and storage systems to ensure data reliability and faster access.

2. Key features.

| Feature | Description. |
|------------------|--|
| Redundancy | Data is duplicated or spread to protect against drive failure. |
| Performance | Multiple disks (can read) write data simultaneously. |
| Fault Tolerance. | System continues to function even if one drive fails. |

Striping

• Data is copied identically on two or more disks.

Mirroring

Data is copied identically on two or more details.

Parity

Error checking information is stored to recover data in disk failure.

Types of RAID.

RAID Level: RAID 0 (striping)

Technique used: Data divided into blocks and spread across disks.

Description: Fast performance (no redundancy).

Advantages: High speed

Disadvantages: No fault tolerance - If one disk fails, all data lost.

RAID Level: RAID 1 (mirroring)

Technique used: Same data copied on two disks.

Description: Provides redundancy.

Advantages: High reliability, simple recovery.

Disadvantages: Storage cost doubles.

RAID level: Bit level stripping with error correction code (ECC).

Technique used: Rarely used

Description: Rarely used

Advantages: Error correction possible.

Disadvantages: Expensive, complex.

RAID Level: RAID 3

Technique used: Byte level stripping with dedicated parity disk.

Description: All drives work together.

Advantages: Good for large sequential data.

Disadvantages: Parity disk may be a bottleneck.

RAID level: RAID 4

Technique used: Block level stripping with dedicated parity disk.

Description: Parity used for recovery

Advantages: Fast read performance.

Disadvantages: Single parity disk bottleneck.

RAID Level: RAID 5.

Technique used: Block level striping with distributed parity

Description: Most common RAID Level.

Advantages: Good balance b/w performance and redundancy

Disadvantages: Complex rebuild if a disk fails.

RAID Level: RAID 6.

Technique used: Like RAID 5 but with two parity blocks.

Description: Higher fault tolerance.

Advantages: Can tolerate two disk failure

Disadvantages: Slower write performance.

RAID Level: RAID 1

Technique used: Combination of RAID 1 & RAID 0

Description:

Advantages:

Disadvantages:

2. Deadlock and its handling?

A deadlock occurs in a database when two or more transactions are waiting for each other to release locks on resources, preventing further progress

Example:

- * Transaction T_1 locks resource A and waits for Resource B.
- * Transaction T_2 locks resource B and waits for Resource A.

\Rightarrow Both wait forever \rightarrow deadlock.

Conditions for Deadlock. (Coffman Conditions).

1. Mutual Exclusion:

Only one transaction can use a resource at a time.

2. Hold and wait:

A transaction holds one resource and waits for another.

3. Non preemption:

Resources cannot be forcefully taken away.

4. Circular wait:

A circular chain of waiting transactions exists.

Deadlock Handling Techniques

Method

Description

1. Dead lock Prevention.

Avoids deadlocks by denying at least one coffman condition.

Example: Assign resource ordering or use time stamp based methods (wait-die, wound-wait).

2. Dead lock Avoidance.

The system checks before the wanting a lock to ensure no deadlock will occur.

Example: Banker's algorithm checks safe states.

3. Deadlock detection.

The system allows deadlocks to occurs but periodically checks for them using a wait-for graph (WFG).

4. Deadlock Recovery

Once detected, resolve by rolling back one or more transaction victim selects

1. Normalization and its various types of normalization.

Normalization and its various types:

Normalization is a process in database management system. DBMS used to organize data in a database to reduce redundancy (duplicate) data and improve data integrity.

1) First Normal Form (1NF).

→ Each cell must contain atomic (indivisible) values.
→ No repeating groups or arrays allowed.

Ex:

| St-ID | St-Name | St-Phone | St-Group |
|-------|---------|---------------------------|-----------|
| 123 | Rakesh | 9843298760 | Physics |
| 125 | Amritा | 91354628198 8763549321 | stat |
| 126 | Archana | 7362543210 869453210 | Chemistry |

New table:

| St-ID | St-Name | St-Phone | St-Group |
|-------|---------|--------------------------|----------|
| 123 | Rakesh | 984073512 | Physics |
| 125 | Amritा | 91354628198 876354932 | stat |
| 125 | Amritा | | stat |

Second Normal Form (2NF).

- In the 2NF, first table must be an 1NF.
- In the 2NF, all non key attributes are fully functionally dependent on the primary key.
- Every non key attributes are full FHD on key attribute.

If $P \rightarrow A$ holds, then those should not be any proper subset - of Q of P .

$$Q \rightarrow A$$

| St-ID | St-Name | Prof-ID | Prof-Name | Grade |
|-------|---------|---------|-----------|-------|
| 101 | ABC | 2 | Sameer | 4 |
| 102 | XYZ | 3 | Niranjan | 6 |
| 103 | PQR | 1 | Sushmitha | 5. |

St-ID → Primary key.

And no multi valued are there so satisfies 1NF.

| St-ID | St-Name |
|-------|---------|
| 101 | ABC |
| 102 | XYZ |
| 103 | PQR |

| Prof-ID | Prof-Name |
|---------|-----------|
| 1 | Sushmitha |
| 2 | Sameer |
| 3 | Niranjan |

Grades:

| St. ID | Prof-ID | Grade. |
|--------|---------|--------|
| 101 | 2 | 4 |
| 102 | 3 | 6 |
| 103 | 1 | 5 |

3. Third Normal Form (3NF)

- A relation will be an 2NF.
- It not contains any transitive dependency.
- The non-key attributes should not have inter dependences among them, and the non-key attributes should fully functionally depend on key attribute.
- Then it is called 3 NF principle.
- By using 3NF to achieve data integrity and Data duplication.

Transitive dependency.

If $A \rightarrow B$ (B functionally dependent on A)

$$B \rightarrow C$$

$A \rightarrow C$ (C is indirectly dependent on A)

It is called T.D.

Ex:

| St-ID | St-Name | Dept-Name | Dept-Head. |
|-------|---------|-----------|-------------|
| 101 | Alice | CSE | Dr. Lalitha |
| 102 | Bob | ECE | Dr. Kumar |
| 103 | Carol | CSE | Dr. Rao. |

Student table.

| St-ID | St-Name | Dept-Name |
|-------|---------|-----------|
| 101 | Alice | CSE |
| 102 | Bob | ECE |
| 103 | Carol | CSE |

Dept-table.

| Dept.-Name | Dept.-Head. |
|------------|-------------|
| CSE | Dr. Lalitha |
| ECE | Dr. Kumar |
| CSE | Dr. Rao. |

4. BCNF (Boyce-Codd Normal Form).

→ It is advanced version of 3NF.

→ It is in 3NF.

→ For every FD $A \rightarrow B$ $A \rightarrow$ super key

→ A should be super key of a table.

Ex:

Student Course Teacher.

Ramesh Physics Kishore

Kumar Chemistry Ramu

Surathi Maths Sarjeet

Vinay Physics Kishore.

Keys $\rightarrow \{$ student, course $\}$ \rightarrow .

To eliminate redundant data we move to 2 tables

Student Course Course Teacher.

Ramesh Physics Physics Kishore.

Kumar Chemistry Chemistry Ramu.

Surathi Maths Maths Sarjeet

Vinay Physics.

Fourth Normal Form (4NF)

relation in BCNF (or) 3.NF

\rightarrow First relation in BCNF (or) 3.NF

\rightarrow It may not contain more than one multivalued

attribute

For a dependency.

$A \rightarrow B$.

\rightarrow If for a single value of A multiple values of B exists, then the relation will be multi-valued dependency.

St-ID

Course

Hobbies.

111

111

111

222

444

555

St-ID

Course

St-ID

Hobbies.

111

111

111

222

444

555

111

111

111

222

444

555

Dancing

Singing

Dancing

Cricket

Hockey,

Dancing

Singing

Dancing

Cricket

Hockey,