Introduction to Database Systems

Lecture 09: E/R Diagrams (4.1-6) and Constraints (7.1-2)

Database Design

What it is:

 Starting from scratch, design the database schema: relation, attributes, keys, foreign keys, constraints etc.

Why it's hard:

- The database will be in operation for years.
- Updating the schema in production is very hard:
 - schema change modifications are expensive (why?)
 - making the change without introducing any bugs is hard
 - this part is, by far, the most important consideration in practice

Database Design

- Consider issues such as:
 - What entities to model
 - How entities are related
 - What constraints exist in the domain
- Several formalisms exists
 - We discuss E/R diagrams
- Reading: Sec. 4.1-4.6

Database Design Process

Conceptual Model:

product makes company name address

Relational Model:

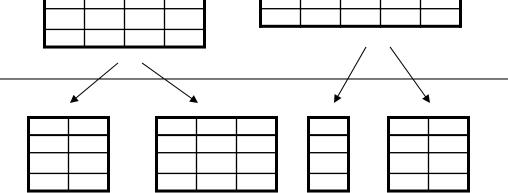
Tables + constraints

And also functional dep.

Normalization:

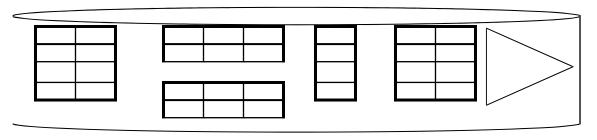
Eliminates anomalies

Conceptual Schema



Physical storage details

Physical Schema



Entity / Relationship Diagrams

- Entity set = a class
 - An entity = an object

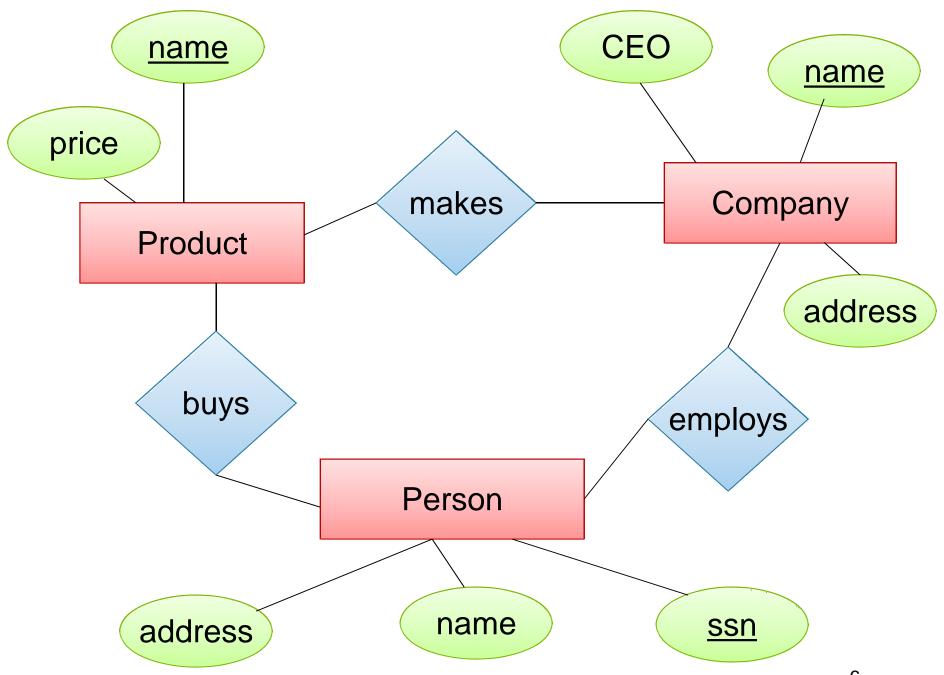
Product

Attribute

city

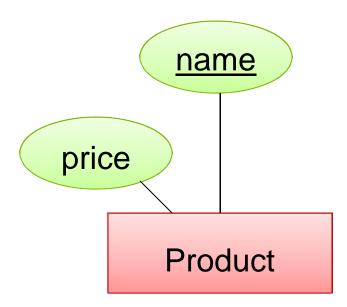
Relationship

makes



Keys in E/R Diagrams

Every entity set must have a key

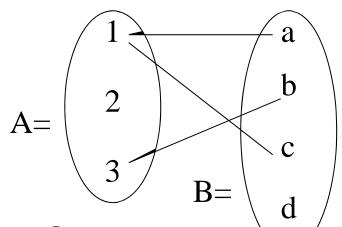


What is a

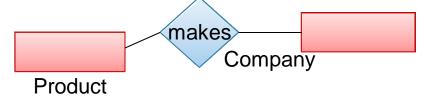
Relation(ship)?

• A mathematical definition:

- - if A, B are sets, then a relation R is a subset of A \times B
- $A = \{1, 2, 3\}, B = \{a, b, c, d\},$ $A \times B = \{(1,a), (1,b), \dots, (3,d)\}$ $R = \{(1,a), (1,c), (3,b)\}$

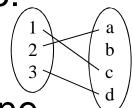


makes is a subset of **Product** × **Company**:

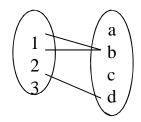


Multiplicity of E/R Relations

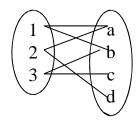
one-one:

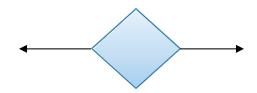


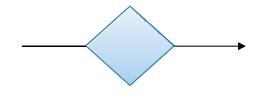
many-one

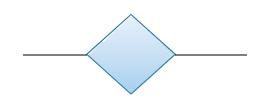


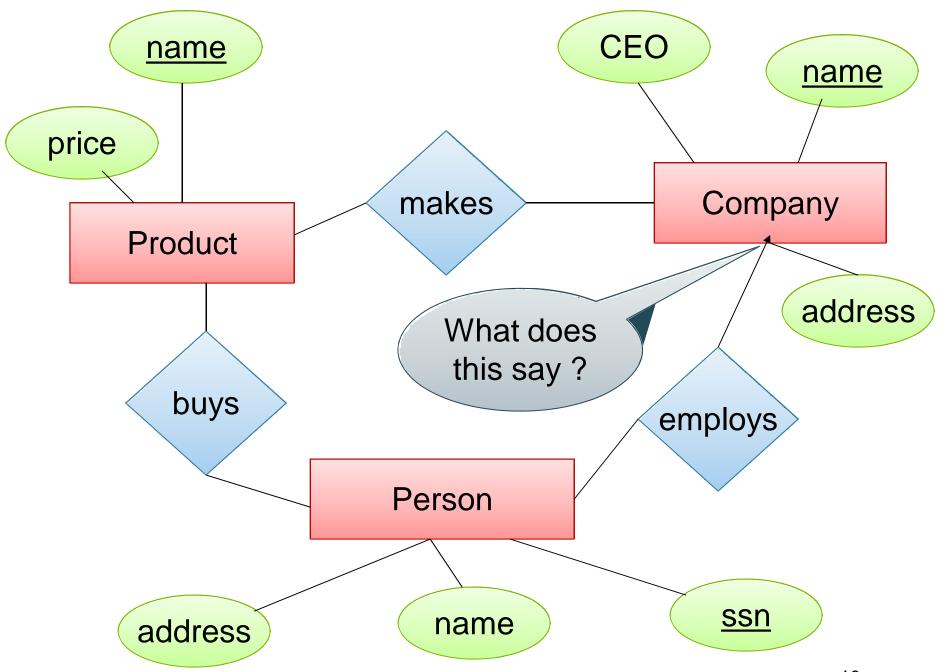
many-many





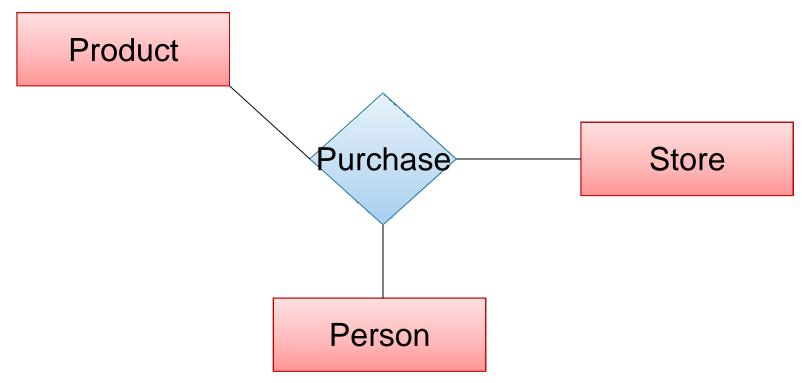






Multi-way Relationships

How do we model a purchase relationship between buyers, products and stores?

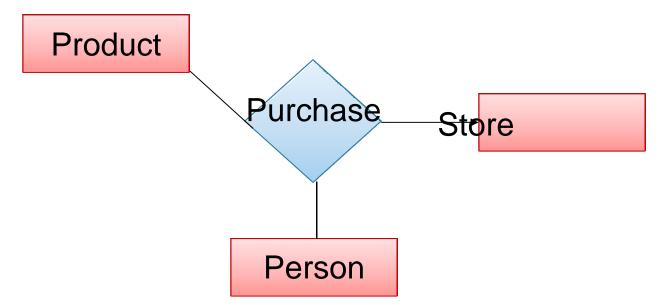


Can still model as a mathematical set (Q. how?)

A. As a set of triples⊆ Person × Product × Store

Arrows in Multiway Relationships

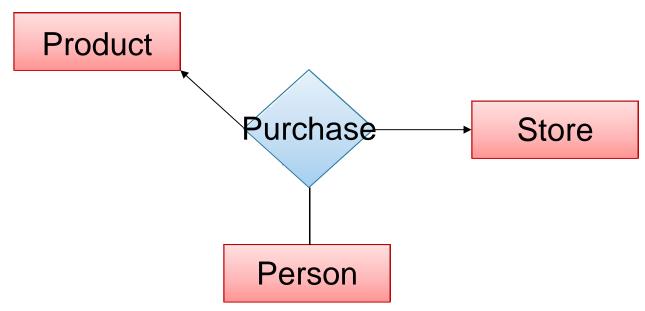
Q: What does the arrow mean?



A: A given person buys a given product from at most one store [Arrow pointing to E means that if we select one entity from each of the other entity sets in the relationship, those entities are related to at most one entity in E]

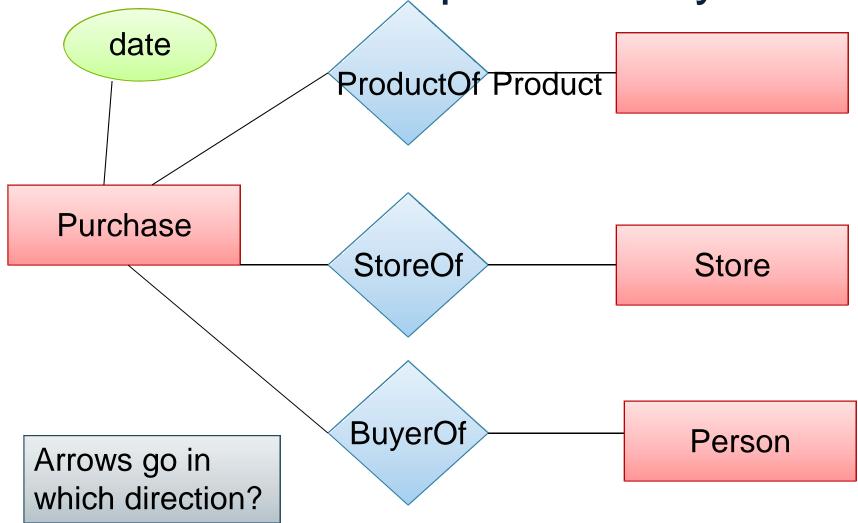
Arrows in Multiway Relationships

Q: What does the arrow mean?

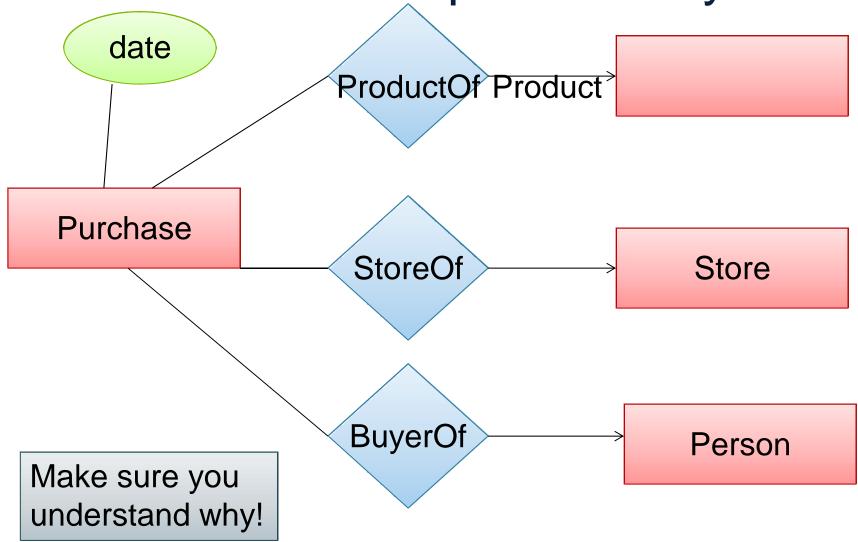


A: A given person buys a given product from at most one store AND every store sells to every person at most one product

Converting Multi-way Relationships to Binary

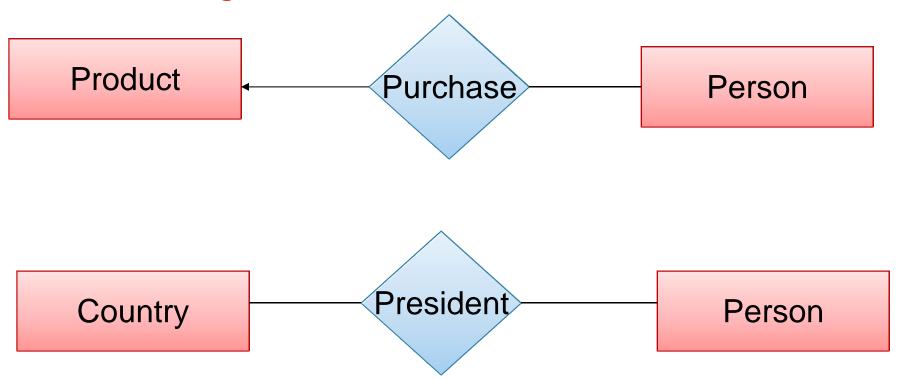


Converting Multi-way Relationships to Binary



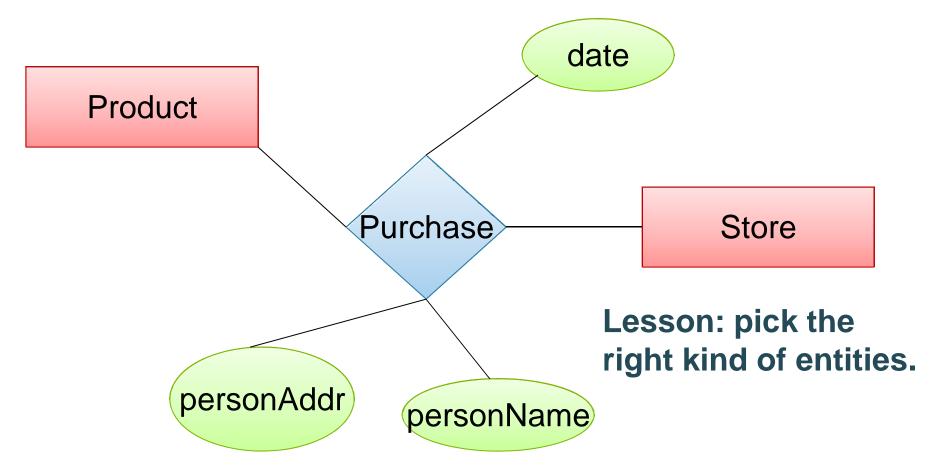
3. Design Principles

What's wrong?

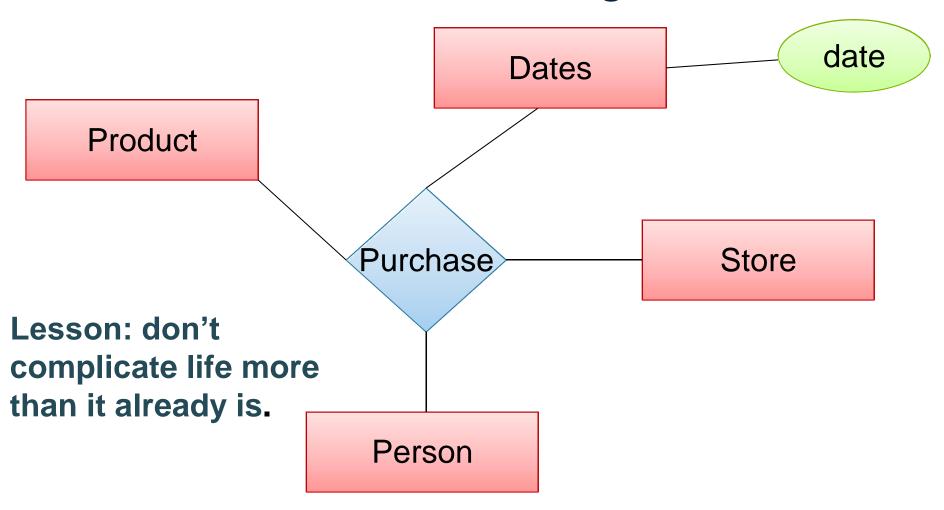


Lesson: be faithful to the specifications of the app!

Design Principles: What's Wrong?



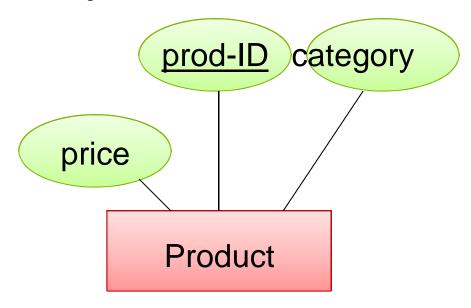
Design Principles: What's Wrong?



From E/R Diagrams to Relational Schema

- Relationship 2 relation

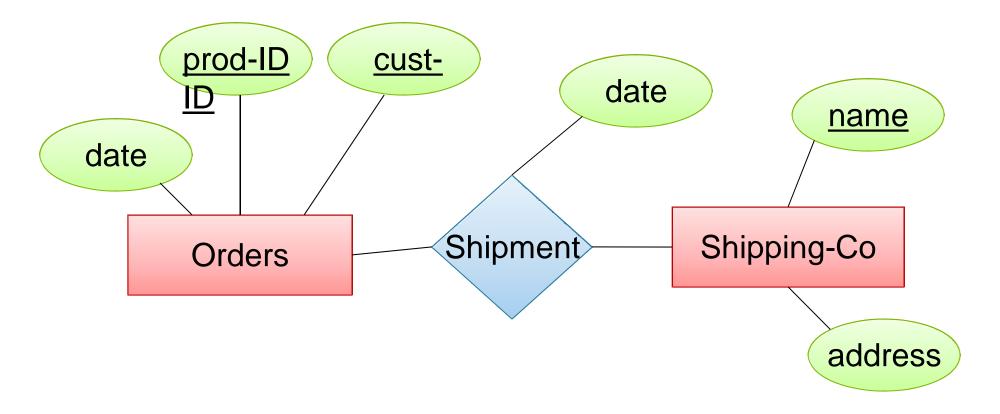
Entity Set to Relation



Product(prod-ID, category, price)

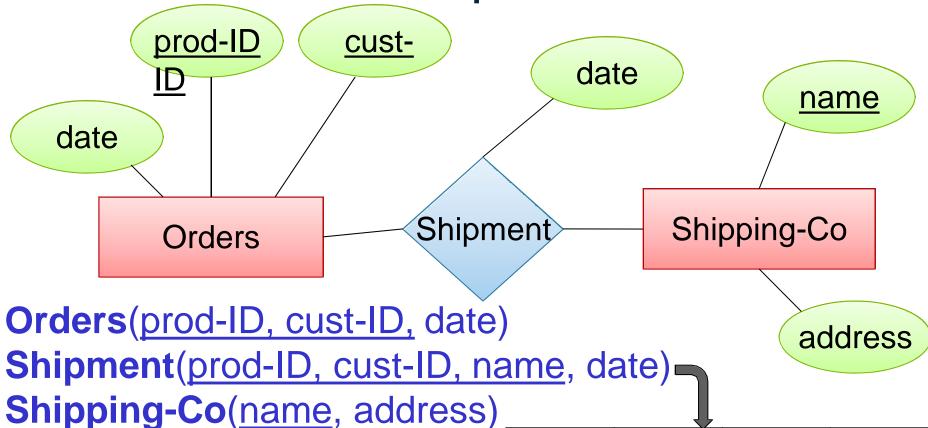
prod-ID	category	price
Gizmo55	Camera	99.99
Pokemn19	Toy	29.99

N-N Relationships to Relations



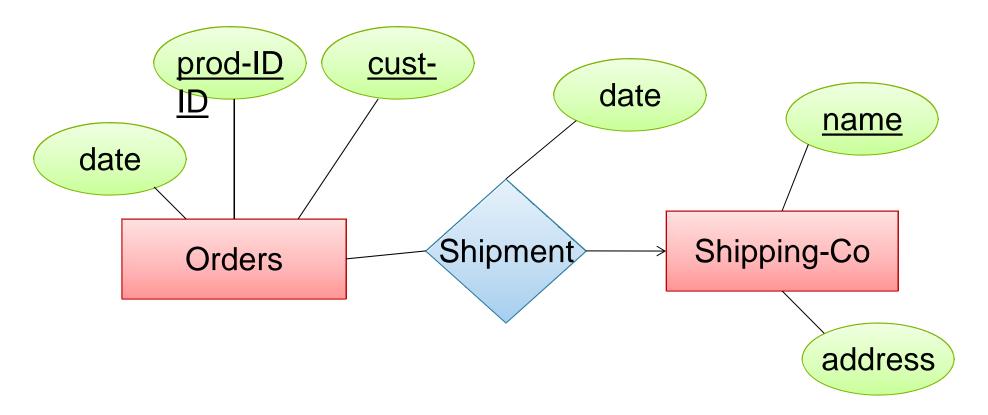
Represent this in relations

N-N Relationships to Relations



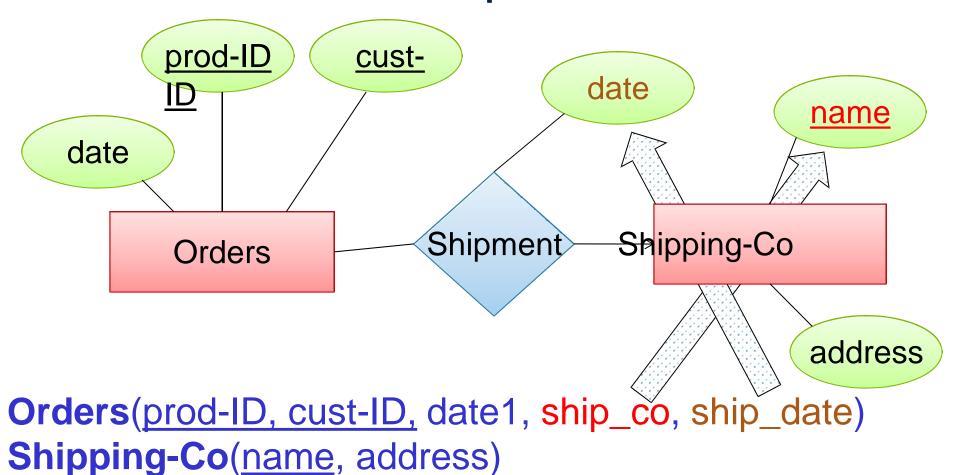
prod-ID	cust-ID	<u>name</u>	date
Gizmo55	Joe12	UPS	4/10/2011
Gizmo55	Joe12	FEDEX	4/9/2011

N-1 Relationships to Relations



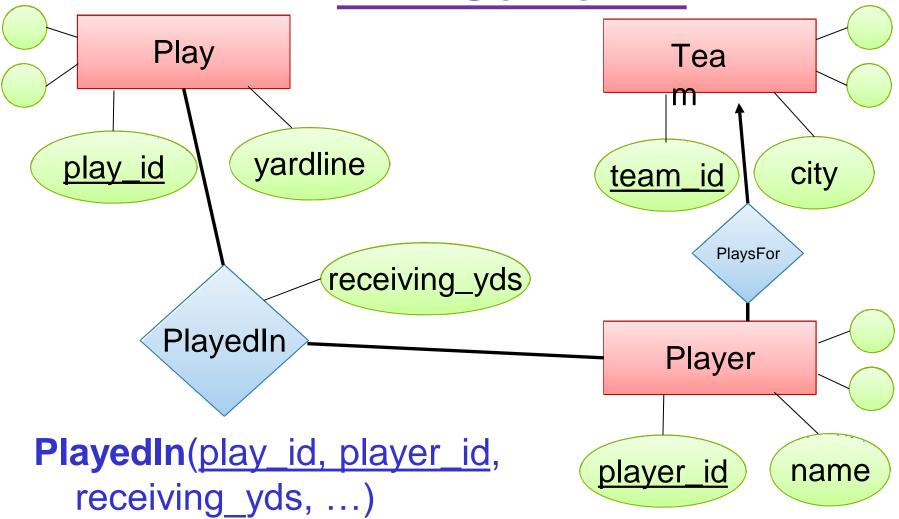
Represent this in relations

N-1 Relationships to Relations



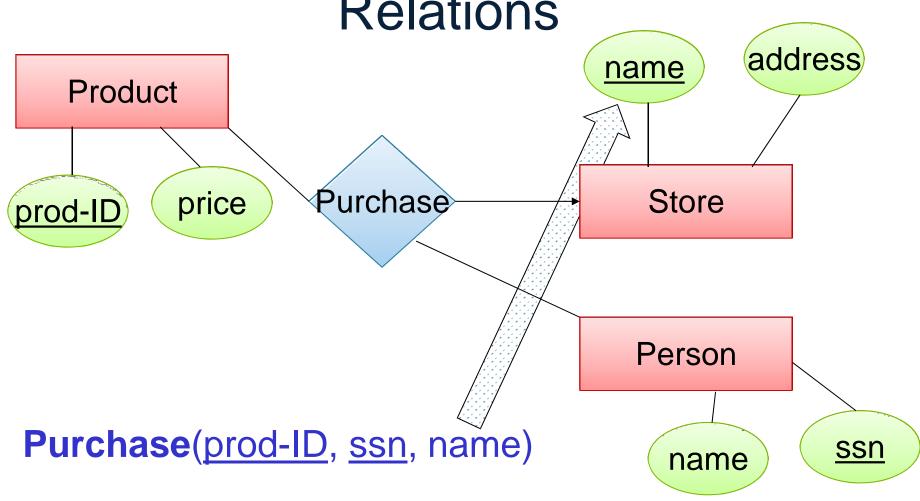
Remember: many-one relationship becomes FK, not relation

Ex: NFL Game DB



(Actually, the key of Play is not play_id. More on this later...)

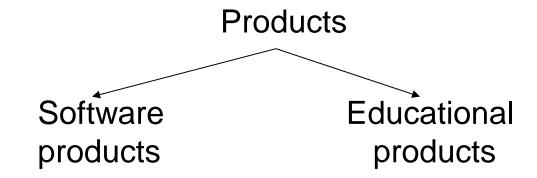
Multi-way Relationships to Relations



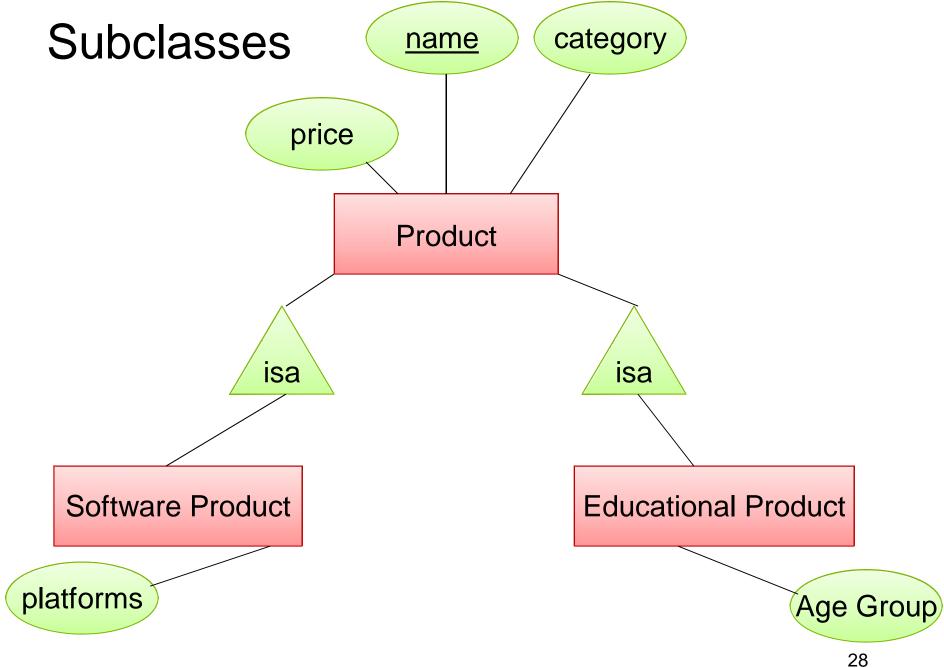
Modeling Subclasses

Some objects in a class may be special

- define a new class
- better: define a subclass



So --- we define subclasses in E/R



Subclasses to Relations (one option)

<u>name</u>

Product

price

qategory

Product

<u>Name</u>	Price	Category
Gizmo	99	gadget
Camera	49	photo
Toy	39	gadget



<u>Name</u>	platforms
Gizmo	unix

Software Product

platforms

isa

Educational Product

isa

Age Group

Ed.Product

<u>Name</u>	Age Group
Gizmo	toddler
Toy	retired

Other ways to convert are possible...

Modeling Union Types with Subclasses

FurniturePiece

Person

Company

Say: each piece of furniture is owned either by a person or by a company

Modeling Union Types with Subclasses

Say: each piece of furniture is owned either by a person or by a company

Solution 1. Acceptable but imperfect (What's wrong?)

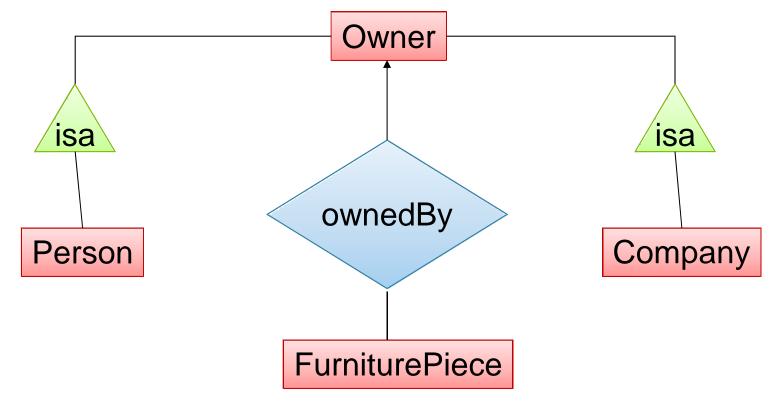
Person FurniturePieceCompany

ownedByPerson

ownedByComp.

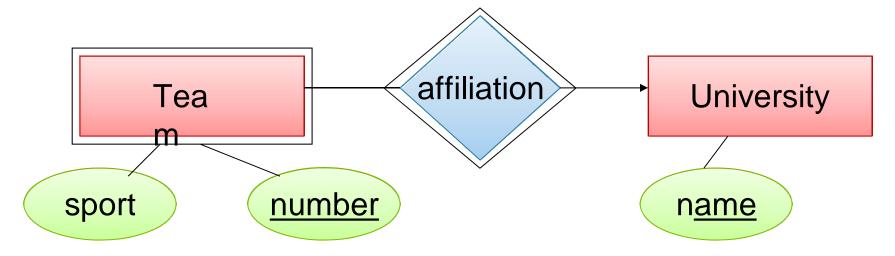
Modeling Union Types with Subclasses

Solution 2: better, more laborious



Weak Entity Sets

Entity sets are weak when their key comes from other classes to which they are related.



Team(sport, <u>number</u>, <u>universityName</u>) University(<u>name</u>)

Integrity Constraints Motivation

An integrity constraint is a condition specified on a database schema that restricts the data that can be stored in an instance of the database.

Most important issue in practice

- ICs help prevent entry of incorrect information
- How? DBMS enforces integrity constraints
 - Allows only legal database instances (i.e., those that satisfy all constraints) to exist
 - Ensures that all necessary checks are always performed and avoids duplicating the verification logic in each application

Constraints in E/R Diagrams

Finding constraints is part of the modeling process. Commonly used constraints:

Keys: social security number uniquely identifies a person.

Single-value constraints: can have only one genetic father

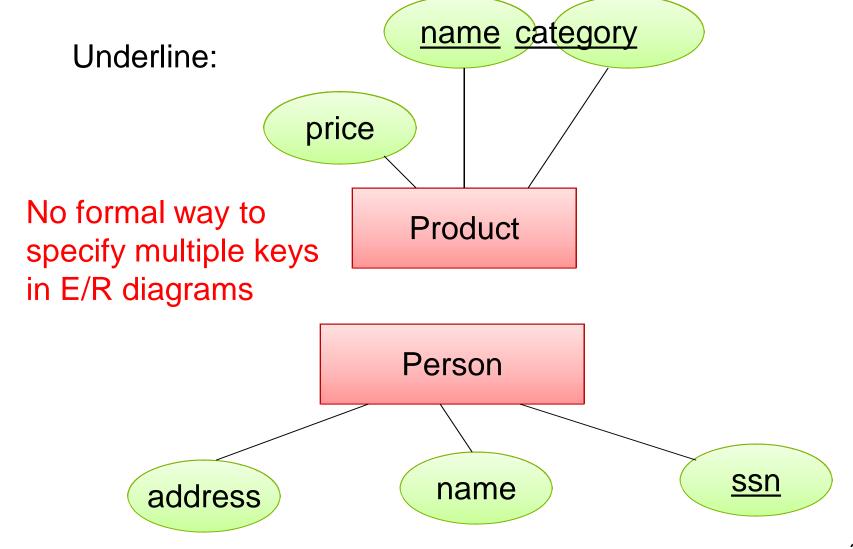
Referential integrity constraints: if you work for a company,

it

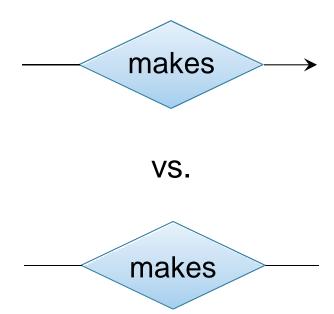
must exist in the database.

Other constraints: peoples' ages are between 0 and 150₃₅ some values should not be NULL

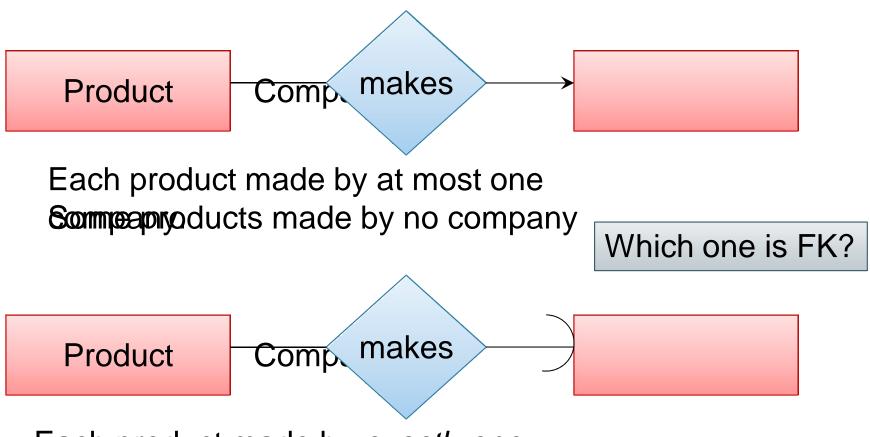
Keys in E/R Diagrams



Single Value Constraints

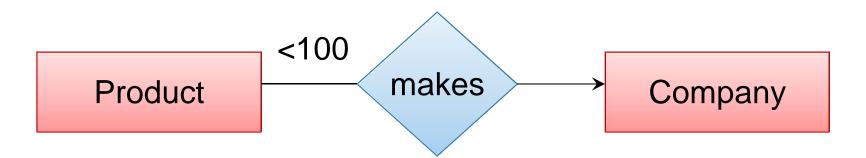


Referential Integrity Constraints



Each product made by *exactly* one company.

Other Constraints



Q: What does this mean?

A: A Company entity cannot be connected

by relationship to more than 99 Product entities

Constraints in SQL

Constraints in SQL:

- Keys, foreign keys
- Attribute-level constraints
- Tuple-level constraints
- Global constraints: assertions

Most complex

simplest

- The more complex the constraint, the harder it is to check and to enforce...
 - (Still, performance is secondary to correctness.)

Key Constraints

Product(<u>name</u>, category)

```
CREATE TABLE Product (
name CHAR(30) PRIMARY KEY,
category VARCHAR(20))
```

OR:

```
CREATE TABLE Product (
name CHAR(30),
category
VARCHAR(20),
PRIMARY KEY (name))
```

Keys with Multiple Attributes

Product(name, category,

price)

```
CREATE TABLE Product (
name CHAR(30),
category VARCHAR(20),
price INT,
PRIMARY KEY (name, category))
```

Name	Category	Price
Gizmo	Gadget	10
Camera	Photo	20
Gizmo	Photo	30
Gizmo	Gadget	40

Other Keys

```
CREATE TABLE Product (
productID CHAR(10),
name CHAR(30),
category VARCHAR(20),
price INT,
PRIMARY KEY (productID),
UNIQUE (name, category))
```

There is at most one PRIMARY KEY; there can be many UNIQUE

Foreign Key Constraints

CREATE TABLE Purchase (
prodName CHAR(30)
REFERENCES Product(name),
date DATETIME)

Referential integrity constraints

prodName is a foreign key to Product(name)

name must be a **key** in Product (i.e., PK or UNIQUE)

May write just Product if name is PK

Foreign Key Constraints

Example with multi-attribute primary key

```
CREATE TABLE Purchase (
    prodName CHAR(30),
    category
    VARCHAR(20), date
    DATETIME,
    FOREIGN KEY (prodName, category)

REFERENCES Product(name, category))
```

(name, category) must be a KEY in Product

What happens when data changes?

Types of updates:

- In Purchase: insert/update
- In Product: delete/update

Product

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

Purchase

ProdName	Store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

What happens when data changes?

SQL has three options for maintaining referential integrity on changes:

- RESTRICT not allowed
- NO ACTION reject bad modifications (default)
- CASCADE after delete/update do delete/update
- SET NULL set foreign-key field to NULL
- SET DEFAULT set FK field to default value
 - need to be declared with column, e.g.,
 CREATE TABLE Product (pid INT DEFAULT 42)

Maintaining Referential Integrity

```
CREATE TABLE Purchase (
    prodName CHAR(30),
    category
    VARCHAR(20), date
    DATETIME,
    FOREIGN KEY (prodName, category)
    REFERENCES Product(name,
    category) ON UPDATE CASCADE
```

Product

ON DELETE SET NULL)
Purchase

Name	Category
Gizmo	gadget
Camera	Photo
OneClick	Photo

ProdName	Category
Gizmo	gadget
Snap	Camera
EasyShoot	Camera

Constraints on attributes:

NOT NULL
CHECK condition

-- obvious meaning...

-- any condition!

Constraints on tuples
 CHECK condition

```
CREATE TABLE Product (
    productID CHAR(10),
    name CHAR(30),
    category VARCHAR(20),
    price INT CHECK (price > 0),
    PRIMARY KEY (productID))
```

```
CREATE TABLE Product (
    productID CHAR(10),
    name CHAR(30),
    category VARCHAR(20)
    CHECK (category in ('toy', 'gadget', 'apparel')),
    price INT CHECK (price > 0),
    PRIMARY KEY (productID))
```

```
CREATE TABLE Product (
    productID
    CHAR(10),
    name CHAR(30) NOT NULL,
    category VARCHAR(20)
    CHECK (category in ('toy', 'gadget', 'apparel')),
    price INT CHECK (price > 0),

PRIMARY KEY (productID))
```

```
CREATE TABLE R (
   A int NOT NULL,
   B int CHECK (B > 50 and B < 100),
   C varchar(20),
   D int,
   CHECK (C >= 'd' or D > 0))
```

What does this constraint do?

CREATE TABLE Purchase (prodName CHAR(30)

CHECK(**Strb@Ca**me IN

Product.name FROM

Product)),

date DATETIME NOT NULL)

What is the difference from Foreign Key?

General Assertions

```
CREATE ASSERTION myAssert
CHECK (NOT EXISTS(
    SELECT Product.name
    FROM Product, Purchase
    WHERE Product.name = Purchase.prodName
    GROUP BY Product.name
    HAVING count(*) > 200) )
```

But most DBMSs do not implement assertions, because it is hard to support them efficiently. Instead, DBMSs provide triggers