





# LECTURE 5: EXCEPTION & INTERRUPT

## **Learning Goals**





- Describe detail about the Cortex-M interrupt handling (to be more specific: NVIC controller).
- Describe about the exception entry/return sequence and how to optimize the interrupt latency.





- 1. Overview of interrupt management
- 2. Vector Table
- 3. Details of NVIC & SCB registers
- 4. Exception Sequence
- 5. Exception handling optimization
- 6. Summary





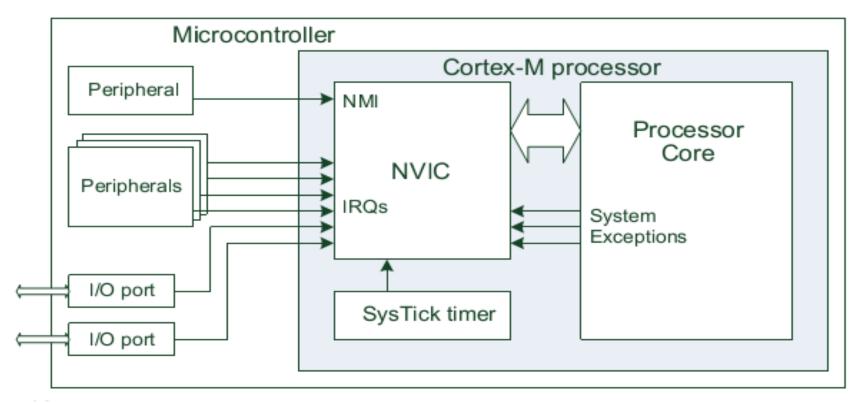
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### Overview of interrupt management

Exception sources



Various sources of exceptions in a typical microcontroller





#### Overview of interrupt management

#### System Exceptions and Interrupt

Excp. Number	Exception Type	Priority
1	Reset	-3 (highest)
2	NMI	-2
3	Hard Fault	-1
4	MemManage Fault	Programmable
5	Bus Fault	Programmable
6	Usage Fault	Programmable
7-10	Reserved	NA
11	SVC	Programmable
12	Debug Monitor	Programmable

Excp. Number	Exception Type	Priority
13	Reserved	NA
14	PendSV	Programmable
15	SysTick	Programmable
16	Interrupt #0	Programmable
17	Interrupt #1	Programmable
•••		
47	Interrupt #31	Programmable
***		
255	Interrupt #239	Programmable





### Overview of interrupt management

#### **Exception Properties**

- Exception number: Identification for the exception
- Vector address: Exception entry point in memory
- Priority level: Determines the order in which multiple pending exceptions are handled





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#### **Vector Table**

#### Vector Table:

- After powerup, vector table is located at 0x0000000
- Vector table contains:
  - Initial value for Main Stack Pointer
  - Handler vector addresses

Address	Exception Number	Value (Word Size)
0x00000000	-	MSP initial value
0x0000004	1	Reset vector
0x0000008	2	NMI handler starting address
0x000000C	3	Hard fault handler starting address
		Other handler starting address



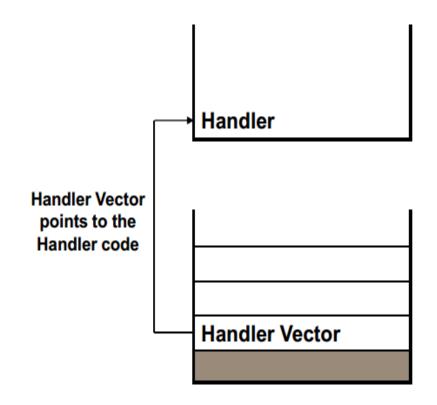


#### **Vector Table**

#### **Vector Table Usage:**

In the case of an exception, the core:

- Reads the vector handler address for the exception from the vector table
- Branches to the handler



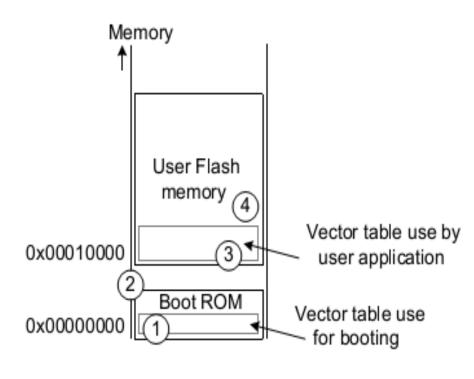




#### **Vector Table & Vector Table Relocation**

#### **Vector Table Relocation:**

- Vector table can be relocated to change interrupt handlers at run-time
- Some cases using vector table relocation feature:
  - Boot loader
  - Applications load to RAM
  - Dynamic changing of Vector







#### **Vector Table**

#### **Vector Table Implementation (Startup.s & linker file):**

```
: Vector Table Mapped to Address 0 at Reset
                           RESET, DATA, READONLY
                  AREA
                          Vectors
                  EXPORT
                          ___Vectors End
                  EXPORT
                            Vectors Size
                          __initial_sp : Top of Stack
 Vectors
                  \mathbf{D}\mathbf{G}\mathbf{D}
                  DCD
                           Reset Handler : Reset Handler
                           NMI Handler : NMI Handler
                  DCD
                  \mathbf{D}\mathbf{G}\mathbf{D}
                          HardFault Handler
                                                : Hard Fault Handler
                  \mathbf{D}\mathbf{G}\mathbf{D}
                           0 : Reserved.
                  DCD
                           0 : Reserved
                           0 : Reserved
                  DCD
                  DCD.
                              : Reserved
  *** Scatter-Loading Description File generated by uVision ***
LR IROM1 0x00000000 0x00020000 { : load region size region
  ER IROM1 0 \times 000000000 \times 000020000 { ; load address = execution address
  *.o (RESET. +First)

*(InRoot$$Sections)

   ANY (+R0)
```

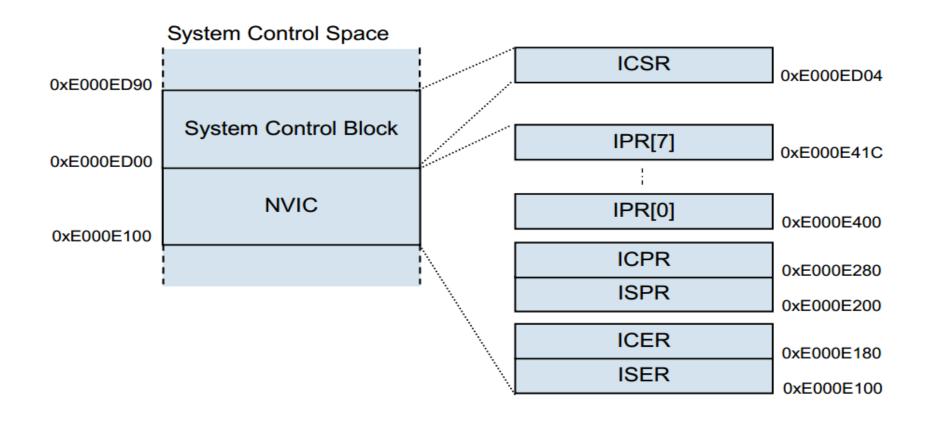




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**SCB & NVIC Register Set for Cortex M0** 





#### **NVIC** registers

Interrupt Priority Level Register (IPR0 – IPR7):

Address	Name	Туре	Reset Value	Descriptions
0xE000E400	PRIORITY0	R/W	0x00000000	Priority level for interrupt 0 to 3 [31:30] Interrupt priority 3 [23:22] Interrupt priority 2 [15:14] Interrupt priority 1 [7:6] Interrupt priority 0

Address	Name	Туре	Reset Value	Descriptions
0xE000E404	PRIORITY1	R/W	0x00000000	Priority level for interrupt 4 to 7
0xE000E408	PRIORITY2	R/W	0x00000000	Priority level for interrupt 8 to 11
0xE000E40C	PRIORITY3	R/W	0x00000000	Priority level for interrupt 12 to 15
0xE000E410	PRIORITY4	R/W	0x00000000	Priority level for interrupt 16 to 19
0xE000E414	PRIORITY5	R/W	0x00000000	Priority level for interrupt 20 to 23
0xE000E418	PRIORITY6	R/W	0x00000000	Priority level for interrupt 24 to 27
0xE000E41C	PRIORITY7	R/W	0x00000000	Priority level for interrupt 28 to 31





#### **NVIC** registers

#### • Interrupt Clear Pending Register

Address	Name	Туре	Reset Value	Descriptions
0xE000E280	CLRPEND	R/W	0x00000000	Clear pending for interrupt 0 to 31; write 1 to clear bit to 0, write 0 has no effect Bit[0] for Interrupt #0 (exception #16)
				Bit[31] for Interrupt #31 (exception #47) Read value indicates the current pending status

#### Interrupt Set Pending Register

Address	Name	Туре	Reset Value	Descriptions
0xE000E200	SETPEND	R/W	0x00000000	Set pending for Interrupts 0 to 31; write 1 to set bit to 1, write 0 has no effect.  Bit[0] for Interrupt #0 (exception #16)  Bit[1] for Interrupt #1 (exception #17)
				Bit[31] for Interrupt #31 (exception #47) Read value indicates the current pending status





#### **NVIC** registers

Interrupt Set Enable Register

Address	Name	Туре	Reset Value	Descriptions
0xE000E100	SETENA	R/W	0x0000000	Set enable for Interrupts 0 to 31; write 1 to set bit to 1, write 0 has no effect Bit[0] for Interrupt #0 (exception #16) Bit[1] for Interrupt #1 (exception #17) Bit[31] for Interrupt #31 (exception #47) Read value indicates the current enable status

#### • Interrupt Clear Enable Register

Address	Name	Туре	Reset Value	Descriptions
0xE000E180	CLRENA	R/W	0x00000000	Clear enable for Interrupts 0 to 31;write 1 to clear bit to 0, write 0 has no effect Bit[0] for Interrupt #0 (exception #16)
1				***





#### **SCB** registers

- Interrupt control and state register
  - Set and clear the pending status of system exceptions including SysTick, PendSV, and NMI.
  - Determine the currently executing exception/interrupt number by reading VECTACTIVE
- Vector table offset Register

Bits	Name	Туре	Reset Value	Description
31:30	Reserved	-	-	Not implemented, tied to 0
29	TBLBASE	RW	0	Table base in Code (0) or RAM (1)
28:7	TBLOFF	R/W	0	Table offset value from Code region or RAM region





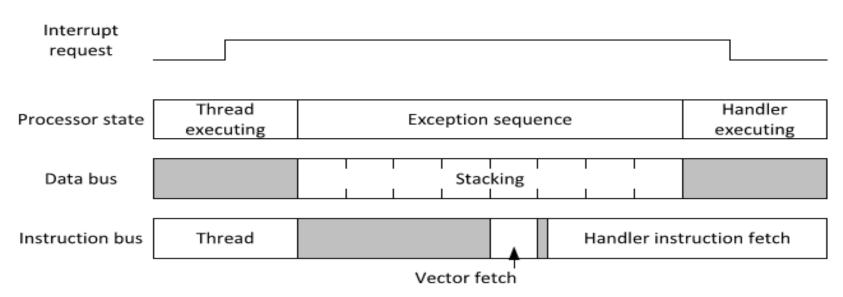
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#### An exception entrance sequence contains:

- 1. Stacking of a number of registers
- Fetching the exception vector
- 3. Fetching the instructions for the exception handler to be executed.
- 4. Update NVIC and core registers







#### Exception handler execution:

1. The processor is in Handler mode

2. Nested Interrupt: If a higher-priority exception arrives, the currently executing handler will be suspended and pre-empted





#### Exception return:

Instructions that can be Used for Triggering Exception Return				
Return Instruction	Description			
BX <reg></reg>	If the EXC_RETURN value is still in the LR when the exception handler ends, we can use the "BX LR" instruction to perform the exception return.			
POP {PC}, or POP {, PC}	Very often the value of LR is pushed to the stack after entering the exception handler. We can use the POP instruction, either a single register POP or a POP operation with multiple registers including the PC, to put the EXC_RETURN value to the Program Counter (PC). This will cause the processor to perform the exception return.			
Load (LDR) or Load multiple (LDM)	It is possible to produce an exception return using the LDR or LDM instructions with PC as the destination register.			





#### Exception return:

- Unstacking: Restore Registers in the stack memory
- Update NVIC & core Registers
- Fetch the instructions of the previously interrupted program.





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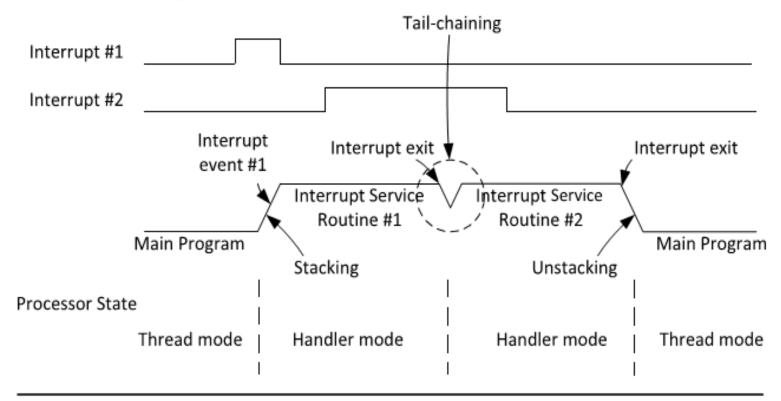


- Tail chaining
- Late arrival
- Pop preemption
- Lazy stacking (Only support for Cortex M4 with FPT)





#### Tail chaining:

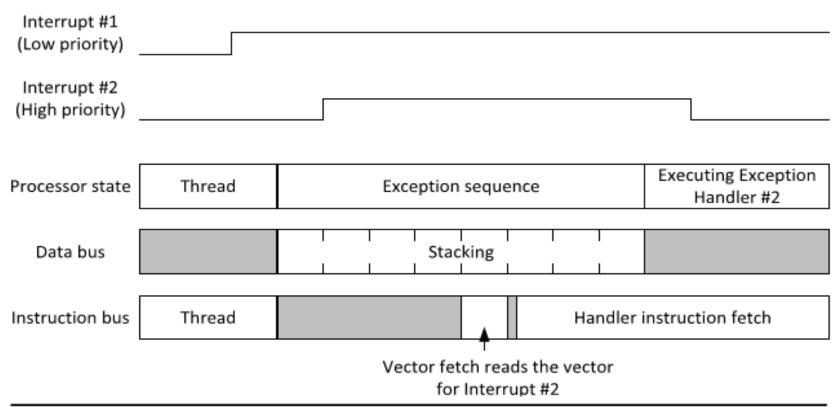


Tail chaining of exception





#### Late arrival:

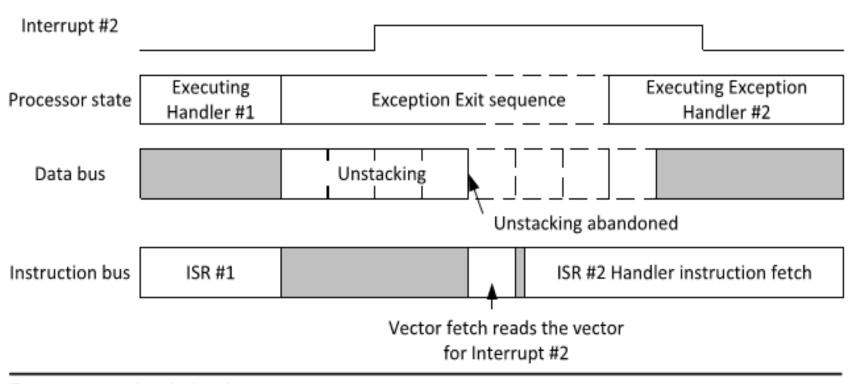


Late arrival exception behavior





#### Pop preemption:



Pop pre-emption behavior





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## Summary





- Cortex-M provides a Nested Vectored Interrupt Controller (NVIC) for interrupt handling.
- Exception entry sequence: Staking, Update registers, Fetch vector number & instruction of exception handler.
- Exception return sequence: Unstacking, Fetch & execute from the restored return address
- Exception handling optimization: Tail chaining, Late arrival, pop preemption.

### **Question & Answer**





# Thanks for attention!

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