

## LINEAR ALGEBRA: ANSWER TO HOMEWORK 9

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### 1. 33.2

During some step  $n$  of Arnoldi iteration,  $h_{n+1,n}$  is encountered.

#### 1.1. a.

**Theorem 1.1.1.**  $AQ_n = Q_{n+1}\hat{H}_n$  can be simplified to  $AQ_n = Q_nH_n$ .

*Proof.* As the last row of  $\hat{H}_n$  is 0, no vector  $Aq_i$  has a component along  $q_{n+1}$ . Therefore the simplification holds.  $\square$

*Remark 1.1.2.* Implications on the structure of the  $A = QHQ^*$ : The submatrix  $H_{n+1:m,1:n}$  is 0.

#### 1.2. b.

**Theorem 1.2.1.**  $K_n$  is an invariant subspace of  $A$ :  $AK_n \subseteq K_n$ .

*Proof.*  $K_n = \langle b, \dots, A_{n-1}b \rangle = \langle q_1, \dots, q_n \rangle$ . So, any  $v \in K_n$  can be written as  $v = Q_nx$ . So,  $Av = AQ_nx = Q_nH_nx$ . So,  $Av$  is in the column space of  $Q_n$ . So,  $Av$  is still in  $K_n$ .

Thus,  $AK_n \subseteq K_n$ .  $\square$

#### 1.3. c.

**Theorem 1.3.1.**  $K_n = \langle b, \dots, A^{n-1}b \rangle$ . Then,  $K_n = K_{n+1} = K_{n+2} \dots$

*Proof.*  $K_{n+1} = \langle K_n, A^n b \rangle$ . But, by the previous theorem, as  $A^{n-1}b \in K_n$ ,  $A^n b = AA^{n-1}b \in K_n$ . Thus,  $K_{n+1} \subseteq K_n$ . But trivially,  $K_n \subseteq K_{n+1}$ . Thus,  $K_n = K_{n+1}$ .

Using the same argument inductively, we see that  $K_n = K_{n+1} = K_{n+2} \dots$   $\square$

#### 1.4. d.

**Theorem 1.4.1.** Each ew of  $H_n$  is an ew of  $A$ .

*Proof.* As we had previously remarked,  $H_{n+1:m,1:n}$  is 0. So, the characteristic polynomial of  $A$ ,  $\det(A - lI)$  can be written as  $\det(H_{1:n,1:n} - lI_n) \det(H_{n+1:m,n+1:m} - lI_{m-n})$ .

So, any value of  $l$  which causes  $\det(H_{1:n,1:n} - lI_n) = 0$  also causes  $\det(A - lI) = 0$ . Thence the result.  $\square$

1.5. e.

**Theorem 1.5.1.** *If  $A$  is non singular, then solution  $x$  to  $Ax=b$  lies in  $K_n$ .*

*Proof.*  $K_n$  is spanned by columns of  $Q_n$ . Let  $Q_{n\perp}$  be a matrix whose columns form an orthonormal basis for the subspace of  $\text{Range}(A)$  orthogonal to  $K_n$ .

$$\begin{aligned} \text{Let: } Q_n y + Q_{n\perp} y' &= x \\ Q_n y + Q_{n\perp} y' &= x \\ Q_n y - x &= -Q_{n\perp} y' \\ A Q_n y - A x &= -A Q_{n\perp} y' \end{aligned}$$

$Q_n y$  is in  $K_n$ ; and  $K_n$  being invariant,  $A Q_n y$  is also in  $K_n$ . Also,  $Ax = b = \|b\| q_1$  is also in  $K_n$ . So,  $A Q_n y - Ax$  is also in  $K_n$ .

[Incomplete].

□

## 2. 36.1

**Theorem 2.0.2.**  *$A$  is real and symmetric.  $r(x) = \frac{x^T A x}{x^T x}$ . Stationary values of  $r(x)$  are the ew of  $A$ . Then, Ritz values at step  $n$  of the Lancos iteration are the stationary values of  $r(x)$  if  $x$  is restricted to  $K_n$ .*

*Proof.* If  $x$  is restricted to  $K_n$ , it can be written as  $Q_n y$ . Then:

$r(x) = r(Q_n y) = \frac{y^T Q_n^T A Q_n y}{y^T y} = \frac{y^T T_n y}{y^T y}$ . Let us denote this by  $r'(y)$ . We note that this is the Rayleigh quotient for the matrix  $T_n$ , and that, whenever  $y$  is an ev this quantity is the corresponding ew.

Following the analysis in the Rayleigh quotients chapter, we see that  $\nabla r'(y) = \frac{2}{y^T y} (T_n y - r'(y) y)$ . This is 0 exactly when  $y$  is an ev of  $T_n$  and  $r'(y)$  is the corresponding ew of  $T_n$ .

So, stationery values of  $r'(y)$  and  $r(x)$  restricted to  $x$  of the form  $x = Q_n y$  are exactly the same:  $r'(y) = 0 \equiv r(Q_n y) = 0 \equiv r(x) = 0$ . □

## 3. 38.5

Minimizing  $f(x) = 2^{-1} x^T A x - x^T b$  using steepest descent:  $p_n = r_n$ .

3.1. a.

**Theorem 3.1.1.**  $\nabla f(x) = -r$ .

*Proof.*

$$\begin{aligned} f(x) &= 2^{-1} x^T A x - x^T b \\ \nabla f(x) &= \nabla 2^{-1} x^T A x - \nabla x^T b \\ &= A x - b \\ &= -r \end{aligned}$$

□

3.2. b.

**Theorem 3.2.1.** *Optimal step*  $a_n = \frac{r_{n-1}^T r_{n-1}}{r_{n-1}^T A r_{n-1}}$ .

*Proof.* We want to find an optimal  $a_n$  which minimize  $f(x_n)$ . Our search direction is  $r_{n-1}$ ; so we want to find  $x_n = x_{n-1} + a_n r_{n-1}$ .

$$\begin{aligned} \nabla f(x_n) &= 0 \\ Ax_n - b &= 0 \\ A(x_{n-1} + a_n r_{n-1}) - b &= 0 \\ b - Ax_{n-1} &= a_n A r_{n-1} \\ r_{n-1} &= a_n A r_{n-1} \\ a_n &= \frac{r_{n-1}^T r_{n-1}}{r_{n-1}^T A r_{n-1}} \end{aligned}$$

□

3.3. c. The full steepest descent iteration:

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 $x_0 = 0, r_0 = b.$ 
foreach  $n = 1, 2 \dots$  do
     $a_n = \frac{r_{n-1}^T r_{n-1}}{r_{n-1}^T A r_{n-1}}$ 
     $x_n = x_{n-1} + a_n r_{n-1}$ 
     $r_n = b - Ax_n$ 
end
    
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4. 38.6

[Incomplete].

5

Following link provides a data structure to store sparse matrices:

[www.cs.utexas.edu/~inderjit/courses/cs383c/sparse\\_matrices.txt](http://www.cs.utexas.edu/~inderjit/courses/cs383c/sparse_matrices.txt)

Write a matlab code using the above specified data structure to compute the matrix-vector product  $y = Ax$  in  $O(nz)$ , where  $nz$  is the number of non-zeros in the sparse matrix  $A$ . Also write a matlab code to compute  $y = A^T x$  in  $O(nz)$ . Note that you are not allowed to store  $A^T$  into a new matrix.

[Incomplete].