

# RDBMS AND SQL RELATIONAL ALGEBRA

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The primary goal of a DBMS is to provide a way to store and retrieve database information that is both *convenient* and *efficient*. - Silberschatz, Korth and Sudarshan.

# Preliminaries

# Quiz

- A **relation** R from a set A to set B is a subset of the cartesian product  $A \times B$ . **True/False?**

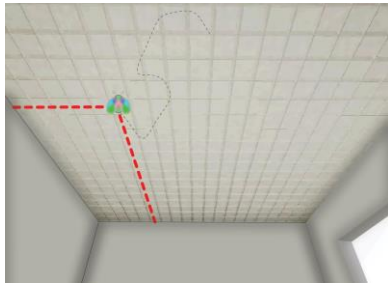
$$\begin{array}{c} \{\text{blue circle}, \text{black circle}, \text{red circle}\} \\ \text{set A} \end{array} \times \begin{array}{c} \{\text{blue triangle}, \text{red triangle}\} \\ \text{set B} \end{array} = \begin{array}{c} \{ (\text{blue circle}, \text{red triangle}), (\text{blue circle}, \text{blue triangle}), \\ (\text{black circle}, \text{red triangle}), (\text{black circle}, \text{blue triangle}), \\ (\text{red circle}, \text{red triangle}), (\text{red circle}, \text{blue triangle}) \} \\ \text{set of all ordered pairs, } A \times B \\ A \times B = \{ (a, b) \mid a \in A \text{ and } b \in B \} \end{array}$$

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Let's say a relation exists between the reds:

$$\text{Relation R} = \{ (\text{red circle}, \text{red triangle}) \}$$

# The Fly on the Ceiling - Descartes



French mathematician René  
Descartes (1596-1650)

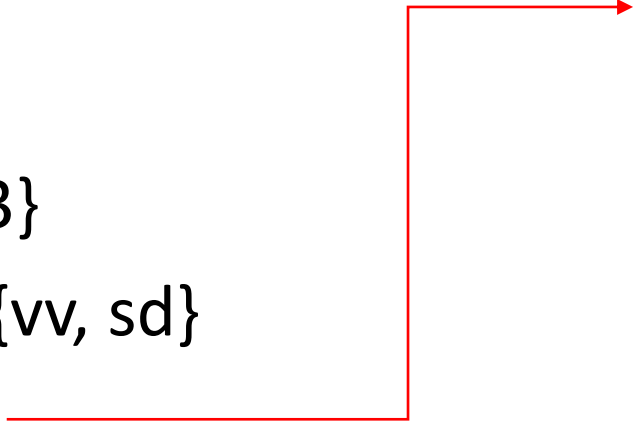
Image Source:

<http://sites.psu.edu/solvingproblems>

<https://wild.maths.org/ren%C3%A9-descartes-and-fly-ceiling>

vectorstock.com

# A Relation

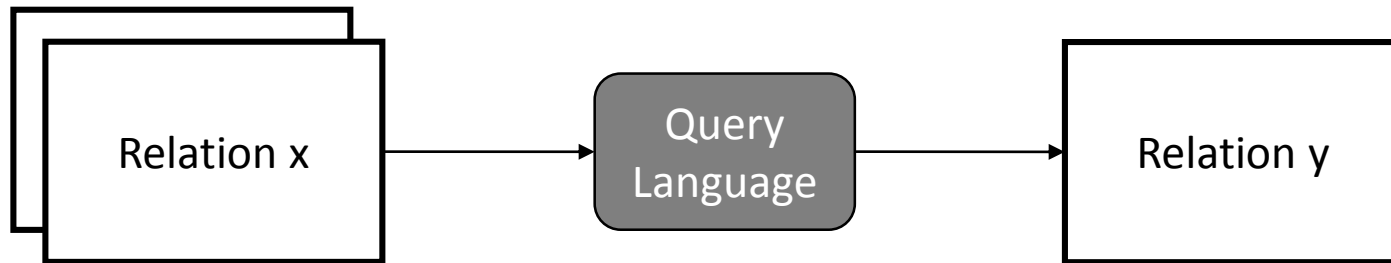
- Let the set,  $id = \{1, 2, 3\}$
- Let the set,  $names = \{vv, sd\}$
- What is **id x names**? 
- We have a **relation** if we assign a sequential id to each name.

id	name
1	sd
1	vv
2	sd
2	vv
3	sd
3	vv

id	name
1	sd
2	vv

# Relational Algebra and Query Languages

# Query Languages



## **Procedural Language**

Relational Algebra

## **Declarative Languages**

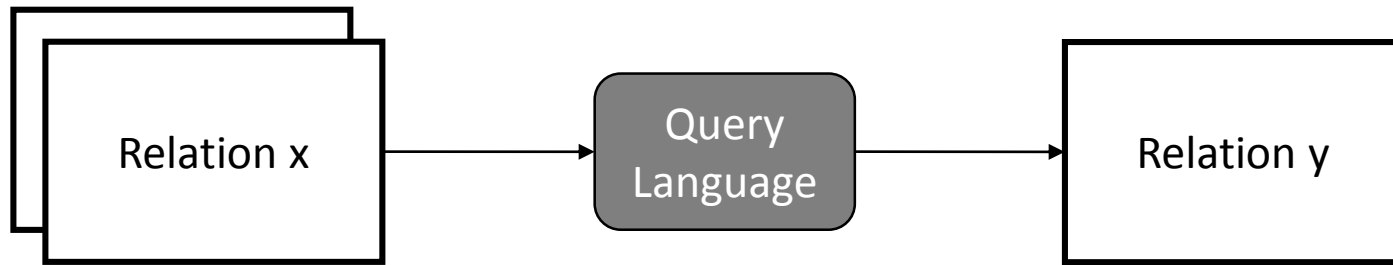
Tuple Relational Calculus

Domain Relational Calculus

## **Popular Language**

SQL

# Relational Algebra



## **Fundamental Operations**

select, project, rename  
set difference, cartesian product, union

## **More Operations**

set intersection, natural join, assignment



# Select Operation

- Notation:  $\sigma_p(r)$  where  $p$  is called the selection predicate

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
76766	Crick	Biology	72000
10101	Srinivasan	Comp. Sci.	65000
58583	Califieri	History	62000
83821	Brandt	Comp. Sci.	92000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
76543	Singh	Finance	80000

(a) The *instructor* table

- Example of selection:

$\sigma_{\text{dept\_name}=\text{"Physics"}}(\text{instructor})$

# Select Operation

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
76766	Crick	Biology	72000
10101	Srinivasan	Comp. Sci.	65000
58583	Califieri	History	62000
83821	Brandt	Comp. Sci.	92000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
76543	Singh	Finance	80000

(a) The *instructor* table

- Example of selection:

$\sigma_{\text{dept\_name}=\text{"Physics"}}(\text{instructor})$

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
22222	Einstein	Physics	95000
33456	Gold	Physics	87000

# Project Operation

- Notation:  $\Pi_{A_1, A_2, \dots, A_k}(r)$  where  $A_i$  are attribute names

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
76766	Crick	Biology	72000
10101	Srinivasan	Comp. Sci.	65000
58583	Califieri	History	62000
83821	Brandt	Comp. Sci.	92000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
76543	Singh	Finance	80000

(a) The *instructor* table

- Example of projection:

$$\Pi_{ID, name, salary}(instructor)$$

- Duplicate rows removed from result, since relations are sets

# Projection

$\Pi_{ID, name, salary}(instructor)$

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
22222	Einstein	Physics	95000
12121	Wu	Finance	90000
32343	El Said	History	60000
45565	Katz	Comp. Sci.	75000
98345	Kim	Elec. Eng.	80000
76766	Crick	Biology	72000
10101	Srinivasan	Comp. Sci.	65000
58583	Califieri	History	62000
83821	Brandt	Comp. Sci.	92000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
76543	Singh	Finance	80000

(a) The *instructor* table



<i>ID</i>	<i>name</i>	<i>salary</i>
10101	Srinivasan	65000
12121	Wu	90000
15151	Mozart	40000
22222	Einstein	95000
32343	El Said	60000
33456	Gold	87000
45565	Katz	75000
58583	Califieri	62000
76543	Singh	80000
76766	Crick	72000
83821	Brandt	92000
98345	Kim	80000

# Union Operation

- Notation:  $r \cup s$ . Defined as:

$$r \cup s = \{t \mid t \in r \text{ or } t \in s\}$$

- For  $r \cup s$  to be valid:
  - $r, s$  must have the *same* **arity** (same number of attributes)
  - The attribute domains must be **compatible** (example: 2<sup>nd</sup> column of  $r$  deals with the same type of values as does the 2<sup>nd</sup> column of  $s$ )
- Example: to find all courses taught in the Fall 2009 semester, or in the Spring 2010 semester, or in both

$\Pi_{course\_id} (\sigma_{semester="Fall" \wedge year=2009} (section)) \cup$

$\Pi_{course\_id} (\sigma_{semester="Spring" \wedge year=2010} (section))$

# Union, Selection and Projection

- $\Pi_{course\_id} (\sigma_{semester="Fall" \wedge year=2009}(section)) \cup$   
 $\Pi_{course\_id} (\sigma_{semester="Spring" \wedge year=2010}(section))$

course_id	sec_id	semester	year	building	room_number	time_slot_id
BIO-101	1	Summer	2009	Painter	514	B
BIO-301	1	Summer	2010	Painter	514	A
CS-101	1	Fall	2009	Packard	101	H
CS-101	1	Spring	2010	Packard	101	F
CS-190	1	Spring	2009	Taylor	3128	E
CS-190	2	Spring	2009	Taylor	3128	A
CS-315	1	Spring	2010	Watson	120	D
CS-319	1	Spring	2010	Watson	100	B
CS-319	2	Spring	2010	Taylor	3128	C
CS-347	1	Fall	2009	Taylor	3128	A
EE-181	1	Spring	2009	Taylor	3128	C
FIN-201	1	Spring	2010	Packard	101	B
HIS-351	1	Spring	2010	Painter	514	C
MU-199	1	Spring	2010	Packard	101	D
PHY-101	1	Fall	2009	Watson	100	A



course_id
CS-101
CS-315
CS-319
CS-347
FIN-201
HIS-351
MU-199
PHY-101

# Set Difference Operation

- Notation:  $r - s$ . Defined as

$$r - s = \{t \mid t \in r \textbf{ and } t \notin s\}$$

- Example: to find all courses taught in the Fall 2009 semester, but not in the Spring 2010 semester

$$\Pi_{course\_id} (\sigma_{semester="Fall" \wedge year=2009}(section)) - \Pi_{course\_id} (\sigma_{semester="Spring" \wedge year=2010}(section))$$

# Union, Selection and Projection

- $\Pi_{course\_id} (\sigma_{semester="Fall" \wedge year=2009}(section)) - \Pi_{course\_id} (\sigma_{semester="Spring" \wedge year=2010}(section))$

course_id	sec_id	semester	year	building	room_number	time_slot_id
BIO-101	1	Summer	2009	Painter	514	B
BIO-301	1	Summer	2010	Painter	514	A
CS-101	1	Fall	2009	Packard	101	H
CS-101	1	Spring	2010	Packard	101	F
CS-190	1	Spring	2009	Taylor	3128	E
CS-190	2	Spring	2009	Taylor	3128	A
CS-315	1	Spring	2010	Watson	120	D
CS-319	1	Spring	2010	Watson	100	B
CS-319	2	Spring	2010	Taylor	3128	C
CS-347	1	Fall	2009	Taylor	3128	A
EE-181	1	Spring	2009	Taylor	3128	C
FIN-201	1	Spring	2010	Packard	101	B
HIS-351	1	Spring	2010	Painter	514	C
MU-199	1	Spring	2010	Packard	101	D
PHY-101	1	Fall	2009	Watson	100	A



course_id
CS-347
PHY-101



# Set-Intersection Operation

- Notation:  $r \cap s$
- Defined as:
- $r \cap s = \{ t \mid t \in r \textbf{ and } t \in s \}$

**Quiz: True/False?**

$$r \cap s = r - (r - s)$$

# Cartesian-Product Operation

- Notation  $r \times s$

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
33456	Gold	Physics	87000
45565	Katz	Comp. Sci.	75000
58583	Califieri	History	62000
76543	Singh	Finance	80000
76766	Crick	Biology	72000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000

**instructor relation**

<i>ID</i>	<i>course_id</i>	<i>sec_id</i>	<i>semester</i>	<i>year</i>
10101	CS-101	1	Fall	2009
10101	CS-315	1	Spring	2010
10101	CS-347	1	Fall	2009
12121	FIN-201	1	Spring	2010
15151	MU-199	1	Spring	2010
22222	PHY-101	1	Fall	2009
32343	HIS-351	1	Spring	2010
45565	CS-101	1	Spring	2010
45565	CS-319	1	Spring	2010
76766	BIO-101	1	Summer	2009
76766	BIO-301	1	Summer	2010
83821	CS-190	1	Spring	2009
83821	CS-190	2	Spring	2009
83821	CS-319	2	Spring	2010
98345	EE-181	1	Spring	2009

**teaches relation**

The *instructor*  $X$  teaches table

[illegible]

# Join Operation

- The Cartesian-Product *instructor X teaches* associates every tuple of instructor with every tuple of teaches.
- Most of the resulting rows have information about instructors who did NOT teach a particular course.
- To get only those tuples of *instructor X teaches* that pertain to instructors and the courses that they taught, we write:

$$\sigma_{instructor.id = teaches.id} (instructor \times teaches)$$

## Join Operation (Cont.)

- $\sigma_{instructor.id = teaches.id}$  (*instructor* x *teaches*)

<i>instructor.ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>	<i>teaches.ID</i>	<i>course_id</i>	<i>sec_id</i>	<i>semester</i>	<i>year</i>
10101	Srinivasan	Comp. Sci.	65000	10101	CS-101	1	Fall	2017
10101	Srinivasan	Comp. Sci.	65000	10101	CS-315	1	Spring	2018
10101	Srinivasan	Comp. Sci.	65000	10101	CS-347	1	Fall	2017
12121	Wu	Finance	90000	12121	FIN-201	1	Spring	2018
15151	Mozart	Music	40000	15151	MU-199	1	Spring	2018
22222	Einstein	Physics	95000	22222	PHY-101	1	Fall	2017
32343	El Said	History	60000	32343	HIS-351	1	Spring	2018
45565	Katz	Comp. Sci.	75000	45565	CS-101	1	Spring	2018
45565	Katz	Comp. Sci.	75000	45565	CS-319	1	Spring	2018
76766	Crick	Biology	72000	76766	BIO-101	1	Summer	2017
76766	Crick	Biology	72000	76766	BIO-301	1	Summer	2018
83821	Brandt	Comp. Sci.	92000	83821	CS-190	1	Spring	2017
83821	Brandt	Comp. Sci.	92000	83821	CS-190	2	Spring	2017
83821	Brandt	Comp. Sci.	92000	83821	CS-319	2	Spring	2018
98345	Kim	Elec. Eng.	80000	98345	EE-181	1	Spring	2017

# Join Operation (Cont.)

- The **join** operation allows us to combine a select operation and a Cartesian-Product operation into a single operation.
- Consider relations  $r (R)$  and  $s (S)$ . Let “theta” be a predicate on attributes in the schema  $R \cup S$ . The join operation  $r \bowtie_{\theta} s$  is defined as follows:

$$r \bowtie_{\theta} s = \sigma_{\theta} (r \times s)$$

- Thus

$$\sigma_{instructor.id = teaches.id} (instructor \times teaches)$$

- Can equivalently be written as

$$instructor \bowtie_{instructor.id = teaches.id} teaches.$$

# Rename Operation

- Allows us to refer to a relation by more than one name.
- Example:

$$\rho_x(E)$$

returns the expression  $E$  under the name  $X$

- If a relational-algebra expression  $E$  has arity  $n$ , then

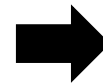
$$\rho_{x(A_1, A_2, \dots, A_n)}(E)$$

returns the result of expression  $E$  under the name  $X$ , and with the attributes renamed to  $A_1, A_2, \dots, A_n$ .

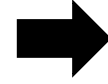
# Find the highest salary in the university

- Steps to reach the solution:

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
33456	Gold	Physics	87000
45565	Katz	Comp. Sci.	75000
58583	Califieri	History	62000
76543	Singh	Finance	80000
76766	Crick	Biology	72000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000



<i>salary</i>
65000
90000
40000
95000
60000
87000
75000
62000
72000
80000
92000



**95000**



# Find the highest salary in the university

- Compute instructor x instructor
- Compute

$$\Pi_{instructor.salary} (\sigma_{instructor.salary < d.salary} (instructor \times \rho_d (instructor)))$$

- Compute

$$\Pi_{salary} (instructor) - \Pi_{instructor.salary} (\sigma_{instructor.salary < d.salary} (instructor \times \rho_d (instructor)))$$

**We use rename operation to distinguish the two salary attributes.**

# Relational Algebra

- A **basic expression** in the relational algebra consists of either one of the following:
  - A relation in the database
  - A constant relation
- Let  $E_1$  and  $E_2$  be **relational-algebra expressions**; the following are all relational-algebra expressions:
  - $E_1 \cup E_2$
  - $E_1 - E_2$
  - $E_1 \times E_2$
  - $\sigma_p(E_1)$ ,  $P$  is a predicate on attributes in  $E_1$
  - $\Pi_s(E_1)$ ,  $S$  is a list consisting of some of the attributes in  $E_1$
  - $\rho_x(E_1)$ ,  $x$  is the new name for the result of  $E_1$

# Fundamental Operations

The select, project, rename, set difference, cartesian product, and union are sufficient to express queries!

# Extended Operations

- Set Intersection

- Find the set of all courses taught in both the Fall 2017 and the Spring 2018 semesters

$$\Pi_{course\_id} (\sigma_{semester="Fall" \wedge year=2017}(section)) \cap \Pi_{course\_id} (\sigma_{semester="Spring" \wedge year=2018}(section))$$

Result

<i>course_id</i>
CS-101

- Assignment

- Find all instructor in the “Physics” and Music department.

$Physics \leftarrow \sigma_{dept\_name="Physics"}(instructor)$

$Music \leftarrow \sigma_{dept\_name="Music"}(instructor)$

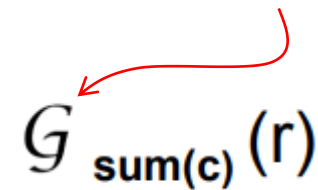
$Physics \cup Music$

# Aggregate Functions

- Improves ease of use
- Most common functions:

Functions
sum
max
min
avg
count

Calligraphic G Notation



The diagram shows the calligraphic G notation  $G_{\text{sum}(c)}(r)$ . A red arrow points from the text "Calligraphic G Notation" to the calligraphic 'G' symbol.

Compute sum of all values of attribute c on relation r.

# Summary

- Very much like normal algebra ( $x - y$ ). We use relations instead of numbers as basic expressions.
- Basis for commercial query languages such as SQL.
- Three major components:
  - Fundamental Operations
  - Extended Operations
  - Aggregate Functions