

## Basic definitions: Strings of non-empty sets

Work with strings  $s = \alpha_1 \cdots \alpha_n$  of sets  $\alpha_i$ , boxed. E.g.,  $\square$  is  $\emptyset = \{\}$  *qua* symbol.

$$(\square\text{-drop}) \quad d(s) := s \text{ with each } \alpha_i = \square \text{ deleted}$$

For any set  $A$ ,

$$\begin{aligned} (A\text{-strips}) \quad \mathcal{L}_A &:= \{d(s) \mid s \in (2^A)^*\} = (2^A - \{\square\})^* \\ (A\text{-reduct}) \quad \rho_A(s) &:= (\alpha_1 \cap A) \cdots (\alpha_n \cap A) \\ (A\text{-projection}) \quad d_A(s) &:= d(\rho_A(s)) \\ (A\text{-preimage}) \quad \llbracket(A, s)\rrbracket_{A'} &:= \{s' \in \mathcal{L}_{A'} \mid d_A(s') = s\} \end{aligned}$$

### Superposition as intersection with granularity

$$\begin{aligned} s \&_{A,A'} s' &:= \llbracket(A, s)\rrbracket_{A \cup A'} \cap \llbracket(A', s')\rrbracket_{A \cup A'} \\ &= \{s'' \in \mathcal{L}_{A \cup A'} \mid d_A(s'') = s \text{ and } d_{A'}(s'') = s'\} \end{aligned}$$

which lifts to languages (i.e. string sets)  $L, L'$

$$\begin{aligned} L \&_{A,A'} L' &:= \bigcup_{s \in L} \bigcup_{s' \in L'} s \&_{A,A'} s' \\ \llbracket(A, L)\rrbracket_{A'} &:= \{s' \in \mathcal{L}_{A'} \mid d_A(s') \in L\} = \bigcup_{s \in L} \llbracket(A, s)\rrbracket_{A'} \end{aligned}$$

so that

$$\begin{aligned} L \&_{A,A'} L' &= \llbracket(A, L)\rrbracket_{A \cup A'} \cap \llbracket(A', L')\rrbracket_{A \cup A'} \\ s \&_{A,A'} s' &= \{s\} \&_{A,A'} \{s'\} \end{aligned}$$

and the special case  $A = A'$  leads to intersection

$$\begin{aligned} L \&_{A,A} L' &= \llbracket(A, L)\rrbracket_A \cap \llbracket(A, L')\rrbracket_A \\ &= \llbracket(A, L \cap L')\rrbracket_A \\ &= L \cap L' \text{ provided } L \cup L' \subseteq \mathcal{L}_A \\ s \&_{A,A} s' &= \begin{cases} \{s\} & \text{for } s = s' \in \mathcal{L}_A \\ \emptyset & \text{otherwise.} \end{cases} \end{aligned}$$

Different granularities  $A \neq A'$  are useful in Allen interval relations, where a string  $s$  has default granularity

$$(\text{vocabulary}) \quad \text{voc}(\alpha_1 \cdots \alpha_n) := \bigcup_{i=1}^n \alpha_i$$

and often  $\text{voc}(s) \neq \text{voc}(s')$  for  $s \neq s'$

$$\begin{aligned} \boxed{l(i)} \boxed{r(i)} \&_{\{l(i), r(i)\}, \{l(j), r(j)\}} \boxed{l(j)} \boxed{r(j)} &= 13 \text{ strings, one per Allen} \\ &\text{relation; e.g. } \text{meet}(i, j) \\ \boxed{l(i)} \boxed{r(i), l(j)} \boxed{r(j)} & \end{aligned}$$

**Computing superposition** Generate  $\&_{A,A'} \subseteq \mathcal{L}_A \times \mathcal{L}_{A'} \times \mathcal{L}_{A \cup A'}$  so that

$$L \&_{A,A'} L' = \{s'' \mid (\exists s \in L)(\exists s' \in L') \&_{A,A'}(s, s', s'')\}$$

Below:  $\alpha \subseteq A$ ,  $\alpha' \subseteq A'$  (making the side conditions vacuous if  $A \cap A' = \emptyset^1$ )

$$\overline{\&_{A,A'}(\epsilon, \epsilon, \epsilon)} \quad (s0)$$

$$\frac{\&_{A,A'}(s, s', s'')}{\&_{A,A'}(\alpha s, \alpha' s', (\alpha \cup \alpha') s'')} \alpha \cap A' \subseteq \alpha' \quad \alpha' \cap A \subseteq \alpha \quad (s1)$$

$$\frac{\&_{A,A'}(s, s', s'')}{\&_{A,A'}(\alpha s, s', \alpha s'')} \alpha \cap A' = \emptyset \quad (d1)$$

$$\frac{\&_{A,A'}(s, s', s'')}{\&_{A,A'}(s, \alpha' s', \alpha' s'')} \alpha' \cap A = \emptyset \quad (d2)$$

**Python implementations** superp.py and for Allen relations, allenD.py

```
>>> from allenD import *
>>> test()
Allen relations from [{ '10' }, { 'r0' }] superposed with [{ '11' }, { 'r1' }]
b [{ '10' }, { 'r0' }, { '11' }, { 'r1' }]
m [{ '10' }, { '11' }, { 'r0' }, { 'r1' }]
o [{ '10' }, { '11' }, { 'r0' }, { 'r1' }]
fi [{ '10' }, { '11' }, { 'r1' }, { 'r0' }]
di [{ '10' }, { '11' }, { 'r1' }, { 'r0' }]
s [{ '10' }, { '11' }, { 'r0' }, { 'r1' }]
eq [{ '10' }, { '11' }, { 'r1' }, { 'r0' }]
si [{ '10' }, { '11' }, { 'r1' }, { 'r0' }]
d [{ '11' }, { '10' }, { 'r0' }, { 'r1' }]
f [{ '11' }, { '10' }, { 'r1' }, { 'r0' }]
oi [{ '11' }, { '10' }, { 'r1' }, { 'r0' }]
mi [{ '11' }, { '10' }, { 'r1' }, { 'r0' }]
bi [{ '11' }, { 'r1' }, { '10' }, { 'r0' }]
Allen transitivity table tt(r1,r2)
e.g. tt("b",d") = ['b', 'm', 'o', 's', 'd'] by superposing
[{'10'}, {'r0'}, {'11'}, {'r1'}] (depicting 0 b 1) and
[{'12'}, {'11'}, {'r1'}, {'r2'}] (depicting 1 d 2)
0 b 2 from [{ '10' }, { 'r0' }, { '12' }, { '11' }, { 'r1' }, { 'r2' }]
0 m 2 from [{ '10' }, { '12' }, { 'r0' }, { '11' }, { 'r1' }, { 'r2' }]
0 o 2 from [{ '10' }, { '12' }, { 'r0' }, { '11' }, { 'r1' }, { 'r2' }]
0 s 2 from [{ '12' }, { '10' }, { 'r0' }, { '11' }, { 'r1' }, { 'r2' }]
0 d 2 from [{ '12' }, { '10' }, { 'r0' }, { '11' }, { 'r1' }, { 'r2' }]
>>>
```

<sup>1</sup>In which case  $s \&_{A,A'} s'$  is  $\bigcup_{n \geq 0} \{d(s_1 \sqcup s_2) \mid s_1 \in \mathcal{L}_{A,n}(s) \text{ and } s_2 \in \mathcal{L}_{A',n}(s')\}$  where  $\mathcal{L}_{B,n}(x)$  is  $\{y \in (2^B)^* \mid \text{length}(y) = n \text{ and } d(y) = x\}$  and  $\sqcup$  is componentwise union

$$\alpha_1 \cdots \alpha_n \sqcup \alpha'_1 \cdots \alpha'_n := (\alpha_1 \cup \alpha'_1) \cdots (\alpha_n \cup \alpha'_n).$$