A Resilient Dynamic Gateway Selection Algorithm Based on Quality Aware Metrics for Smart Grids

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Abstract—**Smart Grids are the evolution of the current electrical power system to meet the challenge of increasing demands for energy by fully integrating the electrical power grid with data communication networks. The challenge faced by this kind of network is to fulfill reliability and resilience requirements in order to meet various types of services and applications. Wireless mesh networks can provide scalability and resilience to this communication network, but there are issues that need to be addressed in order for them to be used in practical smart grids. This paper proposes an algorithm for dynamic selection of gateways in a multihoming smart grid network, improving performance when a gateway's failure occurs. It uses a probabilistic approach for choosing gateways with reliable paths. Our evaluations indicate that the proposed algorithm makes the routing protocol more robust and resilient against gateway failure compared to existing algorithms for dynamic gateway selection.**

Keywords—*wireless mesh networks, gateway selection, smart grid communications.*

I. INTRODUCTION

The current electrical power system has an outdated hierarchical architecture that does not meet the future demands of energy consumption due to various limitations, such as limited generation capacity, one-way flow of energy and control, low and deficient communication and reliability problems [1]. An evolution of the existing electrical power system aims at solving these problems, improving efficiency and reliability, integrating the use of renewable energy produced by consumers, departing from the current one-way flow and deploying a two-way flow for energy and communication [1], [2]. A two-way communication infrastructure is essential for smart grids [3], because it needs to send commands and to receive information from its components and sensors in real time allowing monitoring, maintenance and control.

Smart grids have specific requirements of delay, bandwidth, reliability and time response for each distinct application in their different fields [3]. The Advanced Metering Infrastructure (AMI) is fundamental and is the first step to realize a smart grid [4], [5]. Its requirements should provide robustness and resilience to prevent or recover from problems, providing stability and reliability to the AMI network.

This communication may use available wired or wireless technologies that support the exchange information between components of the AMI [6], [7]. Different types of technologies can be used: cellular technology [8], WiMAX, ZigBee [8], RF Mesh [9], IEEE 802.11-based Wireless Mesh Networks

(WMN) and Power Line Communication (PLC) [10]. PLC is a promising wired technology [6], but it has limitations. In case of failures, such as physical disruption of power lines, it would not be possible to maintain communication between AMI components [11]. Wireless networks offer more benefits than wired networks, such as lower cost, ease deployment and signal availability in large areas [8].

Among all wireless technologies, WMN has advantages compared to single-hop infrastructured network architectures, since it communicates in a multi-hop way that extends the coverage of the network and allows communication with alternative paths in case of failures [9], [12]. WMN, however, must be adapted to the communication requirements of AMI, where hundreds of meters communicate with the Utility's headend through a Data Aggregation Point (DAP). DAPs are the gateways of this network. Typically, an AMI is constituted of networks connecting meters in the same neighborhood to single DAP. Each DAP is connected to the headend through AMI wide area network. The large number of nodes is the main challenge for the WMN [13], since over 100 smart meters may be associated to one single DAP. If all meters send data simultaneously, this can cause congestion in the network.

A way to mitigate this problem is the use of multiple DAPs. The routing protocol must be able to find reliable routes to improve performance and meet the requirements of the AMI network. Given the problems faced by routing protocols in WMN to comply with AMI requirements, we propose an algorithm that dynamically selects DAPs for communication between meters and the headend. In this problem, we assume that each meter can connect, through multiple hops, to a set of DAPs. The main goal of this algorithm, called Dynamic DAP Selection Algorithm (DDSA), is to increase the reliability, robustness and resiliency allowing meters to use multiple DAPs, thus improving performance in the presence of failure.

The organization of the paper is as follows. Section II describes the AMI communication network, its challenges and problems. Section III presents the related work. Section IV proposes DDSA and describes its principle. Section V presents the results obtained in case of failures and Section VI concludes the paper and presents ideas for future work.

II. BACKGROUND

According to Farhangi [1], nearly 8% of all generated energy is lost along the transmission lines and 20% of the total generation capacity is only to support peak demands,

which represent only 5% of the total demand . About 90% of power outages and disturbances are related to the distribution subsystem, thus the success of a smart grid depends on the deployment of a reliable interconnected distribution subsystem.

The AMI aims to improve reliability and changes the paradigm to one where customer demands adjust to the power generation. The AMI is basically composed of smart meters, DAPs and Utility's headend, all interconnected by communication networks. The headend is connected to multiple DAPs, which in turn have connections to multiple smart meters. The meters send measurement data to the headend through a DAP and this traffic is characterized by the exchange of short messages. These messages have a payload that varies from tens to hundreds of bytes and are sent periodically, typically in a 15 minutes interval [7], [14]. Furthermore, the headend can send commands and requests to meters through DAPs.

According to Gungor et al [15], each meter requires a band from 10 to 100 Kb/s and the latency should be less than 2000 ms. Since investments in the power sector are long-lasting, it is desirable that the AMI should also support long-term operations [5]. New demands for information may arise, making the requirements more stringent such as latency that should be less than hundreds of milliseconds in applications that need information in real time [5], [16].

The AMI traffic can be classified into regular or on-demand. It is regular when data is automatically sent by the meters at predetermined intervals [7], [17] and constitutes the majority of data traffic flowing through the AMI [5]. The on-demand traffic is composed of alert messages from meters, command and control sent by the headend to meters and the responses to these commands [17]. In the latter type of traffic, an increase in network congestion can occur due to the request for sending information by the headend to a large number of meters. The DAP is a natural single point of failure, because all traffic between meters and the headend (or vice versa) flows through it. Hence a DAP failure would inhibit the entire network from working.

The residential density determines the amount of meters per area, which according to [17] can be classified into rural, suburban or urban scenario, with densities varying from 10 to 2000 meters per km^2. The external environment conditions in combination with the number of meters will determine the level of interference and attenuation in communication between meters and DAPs.

Due to the peculiarities of the AMI network and due to a large amount of smart meters, there is a possibility of problems such as loops and broken routes [18], causing a degradation in performance that may lead to unmet communication requirements. Thus the routing protocol should handle this variation, providing an acceptable level of service regardless of WMN density. Another problem is the increased amount of collisions near the DAP because all packets are forwarded to it [6].

III. RELATED WORK

The work in [11] proposes the use of WMN where multiple domains of mesh networks are connected by a WiMAX backbone. This architecture provides redundant paths between meters mitigating problems like broken routes due to node failure increasing their resilience. However, since this work considers only one DAP acting as gateway in each WMN domain, if it becomes unavailable there will be no communication between meters and the headend. This is the same problem studied by [9], where the WMN consists of meters, routers and collectors. The meters communicate with routers or directly to collectors, and the latter controls up to 25,000 meters and routers on a single network.

The work in [19] makes use of multiple gateways to increase WMN resilience, because in addition to providing redundant paths, it also provides gateway redundancy. This approach is applied to WMN that serves as the backbone for internet access but uses only the gateway with the best path.

The works in [20] and [21] are designed to meet the requirements of AMI networks and make use of multiple DAPs for communication between meters and headend, modifying the HWMP protocol (Wireless Hybrid mesh Protocol). Although the work in [21] solves some deficiencies of the HWMP protocol, it still suffers other problems such as route instability and loops. According to the authors, this is a characteristic of the distributed backpressure system adopted by them. However, neither of them has evaluated the protocol behavior in an environment with DAP failure, nor allowing transmission rate adaptation, which increases the problem of route instability. They use a base protocol that has scalability problems due to the congestion caused by control messages [22] making it difficult to use in AMI.

Our proposal, DDSA, makes use of multiple DAPs for communication between meters and headend, and differs from [20] and [21] because it is designed to improve performance in environments with DAP failure. DDSA is independent from the routing protocol. Moreover, it can be implemented in a protocol that best suits the implementation of the AMI. In DDSA for each new data packet, meters probabilistically choose a DAP, from a set of available DAPs with good quality paths.

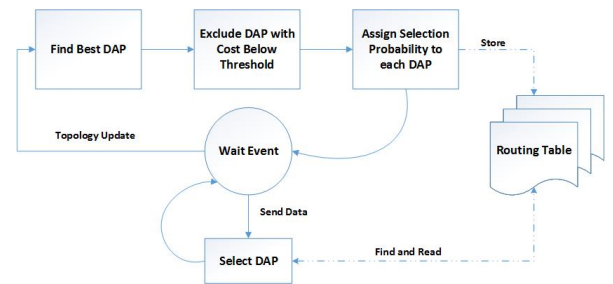## IV. DYNAMIC DAP SELECTION ALGORITHM (DDSA)

Fig. 1. Block Diagram of DDSA

DDSA is based in randomized algorithm and there are two events that start the algorithm: Topology Update and Send Data, as shown in Figure 1. The purpose of Topology Update is maintain routing table entries updated only with a set of best DAPs that have better path cost. Send Data query routing table for each transmission and probabilistically choose a DAP and the better the cost for a given DAP, the higher the probability of selecting it. In the next subsections will be described in more detail the principle of the proposed algorithm.

A. Topology Update

Topology Update is the first event and is started when routing protocol update its topology. After protocol finishes its update, DDSA computes each DAP selection probability and write only the better values in the routing table. To compute this value, DDSA uses the value of the quality metric among smart meter and each DAP in the network. There are three major process that are executed by DDSA in Topology Update: Find Best DAP, Exclude DAP with Cost Below Threshold, and Assign Selection Probability to each DAP.

1) Find Best DAP: Find DAP with best path cost is a necessary task and is needed by next process. While looking for the best DAP a sum of all metric values of the smart meter for all DAPs is computed and will be used to compute DAP selection probability. In algorithm 1 shows how the process work for meter m_i, first the metric value $M_{m_i d_k}$ for each DAP d_k is found and added to Sum. If it is the better metric then is stored in $best\_M$ variable.

> $Sum \leftarrow 0, best\_M \leftarrow 0$
> *//sum of all metric values and select the best DAP*
> **for** *each DAP* d_k **do**
> > $M_{m_i d_k} \leftarrow findMetric(m_i, d_k)$
> > $Sum \leftarrow Sum + M_{m_i d_k}$
> > **if** $best\_M < M_{m_i d_k}$ **then**
> > > $best\_M \leftarrow M_{m_i d_k}$
> >
> > **end**
>
> **end**

Algorithm 1: Find Best DAP.

2) Exclude DAP with Cost Below Threshold: To prevent selection of DAPs with very low metrics, a threshold $\alpha \in [0, 1]$ is employed by the algorithm. The selection probability of the best DAP given by $\frac{best\_M}{Sum}$ (algorithm 2) is multiplied by α, that have values between 0 and 1, resulting in a value γ that is compared with others DAPs probability. If a DAPs probability given by $\frac{M_{(m_i, d_k)}}{Sum}$ is smaller than γ then it is discarded. The threshold α is a parameter of the DDSA algorithm that affects the performance and the behavior of the network. A lower value of α implies in selecting more DAPs wich improving resilience in case of DAP faults, as opposed of higher value that selects only the best set of DAPs wich reduces the resilience.

> *//Compute gamma γ*
> $\gamma \leftarrow \alpha * \frac{best\_M}{Sum}$
> *//Exclude worst DAPs*
> **for** *each DAP* d_k **do**
> > **if** $\frac{M_{(m_i, d_k)}}{Sum} < \gamma$ **then**
> > > $discardDAP(d_k)$
> >
> > **end**
>
> **end**
> *//Recalculate Sum*
> $Sum \leftarrow recalculate\_Sum$

Algorithm 2: Exclude DAP with Cost Below Thresholdl.

After discarding worst DAPs, DDSA recalculate the sum using only the better DAPs metrics. This is necessary to redistribute probabilities of the discarded DAPs among the remaining DAPs and bring fairness to the algorithm.

> *//Compute probability*
> **for** *each DAP* d_k **do**
> > $Prob_{(m_i, d_k)} \leftarrow \frac{M_{(m_i, d_k)}}{Sum}$
> > $updateProbability(m_i, Prob_{(m_i, d_k)})$
>
> **end**

Algorithm 3: Assign Selection Probability to each DAP.

3) Assign Selection Probability to each DAP: For each DAP d_j the probability $P_{(m_i, d_j)}$ is computed by the meter m_i by the expression:

$$P_{(m_i, d_j)} = \frac{M_{(m_i, d_j)}}{\sum\limits_{k=1}^{N} M_{(m_i, d_k)}} , \tag{1}$$

where $M_{(m_i, d_j)}$ is the value of the quality metric of the path (m_i, d_j), which is divided by the sum of all metric values of the meter m_i for all DAPs. The new probability is stored in routing table and can be read by Send Data event. Notice that this expression assumes that the routing metrics assigns higher values for better paths. If the used metric employs a reverse logic, the following expression is used:

$$P_{(m_i, d_j)} = \frac{1/M_{(m_i, d_j)}}{\sum\limits_{k=1}^{N} 1/M_{(m_i, d_k)}} . \tag{2}$$

B. Send Data

> $Prob\_temp \leftarrow 0$
> $Prob\_rand \leftarrow randomUniform(0, 1)$
> *//choosing DAP*
> **for** *each DAP* d_k **do**
> > $prob \leftarrow findProbability(d_k)$
> > $Prob\_temp \leftarrow Prob\_temp + prob$
> > **if** $Prob\_temp >= Prob\_rand$ **then**
> > > $Selected\_DAP \leftarrow d_k$
> > > **break**
> >
> > **end**
>
> **end**
> **retorna** $Selected\_DAP$

Algorithm 4: Send Data algorithm.

The second event is Send Data and is started when a smart meter wants to send information to headend. It starts the process of DAP selection by DDSA that choose probabilistically one in a set of best DAPs that have better path cost. In Algorithm 4 the first task is randomly draw lots a value between 0 or 1 ($Prob\_rand$) that will determine which DAP will be selected. DDSA read DAP probability and sum it to $Prob\_temp$ variable then it is compared to $Prob\_rand$. If $Prob\_temp$ is lesser than $Prob\_rand$ the next DAP probability is read from routing table and added to $Prob\_temp$ until $Prob\_temp$ is greater or equal than $Prob\_rand$ then this DAP is chosen.

V. Performance Evaluation

A. Simulation Environment

The performance of DDSA is evaluated using the ns-2 simulator [23]. To simulate the behavior of an AMI network composed of smart meters and DAPs, ns-2 is set to simulate a suburban external scenario using the shadowing propagation model with the following parameters: path loss exponent = 2.7, standard deviation = 7.4 and reference distance = 1.0, as defined in [17].
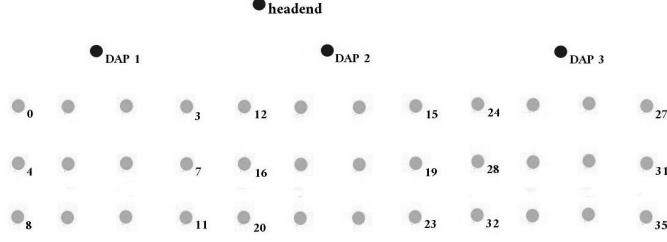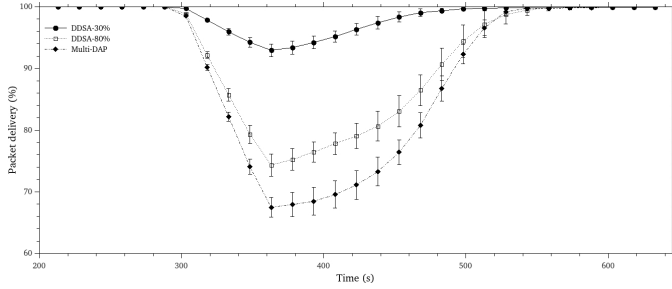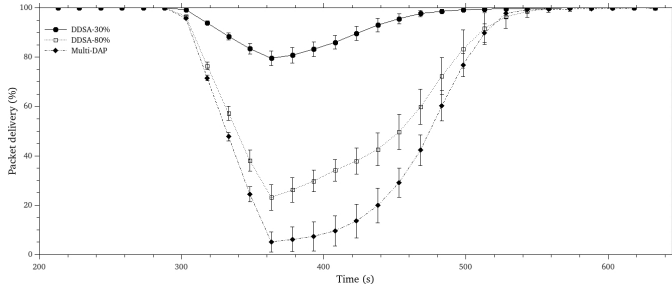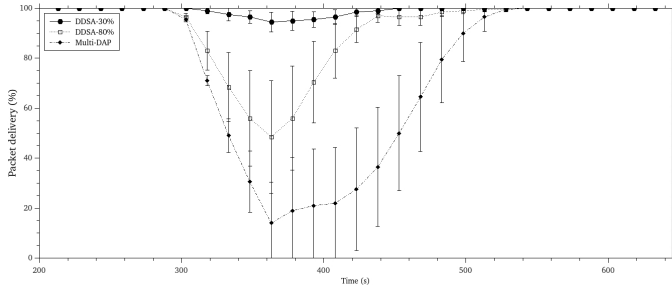

Fig. 2. Scenario used in the simulation.
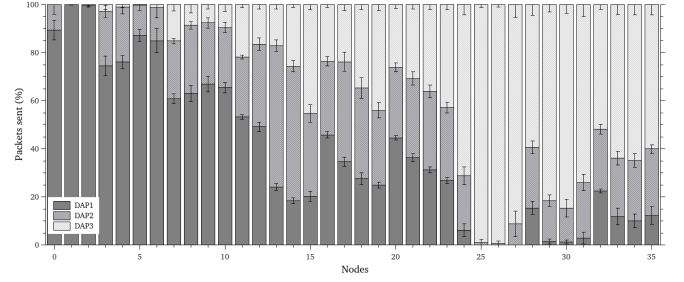

(a) Packet delivery for all nodes
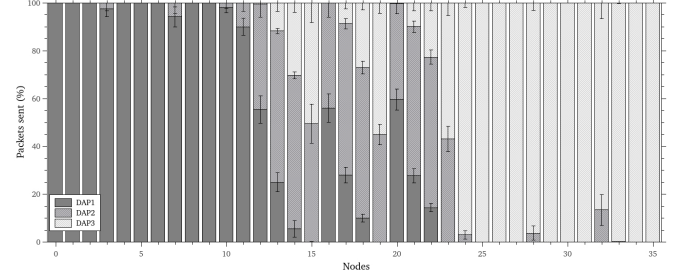

(b) Packet delivery for nodes 12-23

(c) Packet delivery for node 16

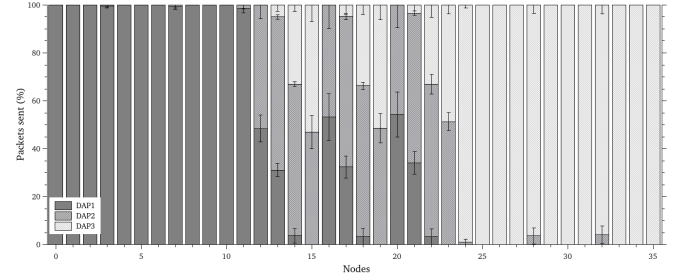Fig. 3. Packet delivery rates.

The used simulation topology is composed of 36 nodes arranged in a grid and 3 DAPs (Fig. 2). To simulate the

(a) DDSA-30%

(b) DDSA-80%

(c) Multi-DAP

Fig. 6. Percentage of DAP usage by each node.

exchange of information between meters and DAPs in a typical application of AMI, a Constant Bit Rate (CBR) UDP traffic is used with fixed packet size of 400B [17] at the rate of 20 packets per minute. Each node sends 10 flows of data in a synchronized way for every three seconds forcing network congestion. The UDP is an unreliable transport protocol, but has some advantages such low latency. To make it more reliable we used an application protocol over UDP that implements its own transport service suitable for AMI traffic [24]. The packet delivery rate was calculated for the application layer. We compute as failure in delivery packet if none of the flows in a round are received by any of the DAPs. The exchange of information starts at time 150 seconds and at 300 seconds DAP 2 fails. A total of 10 simulations were performed with a duration of 650 seconds each and in charts the confidence interval is 95%.

The DDSA was implemented in ns-2 simulator using OLSR [25] as the routing protocol. MARA [26] was employed as the routing metric and rate adaptation. In assessing the results the following solutions were evaluated:

(1)   DDSA with α = 0.3 referred to DDSA-30%;

(2)   DDSA with α = 0.8 referred to DDSA-80%; and
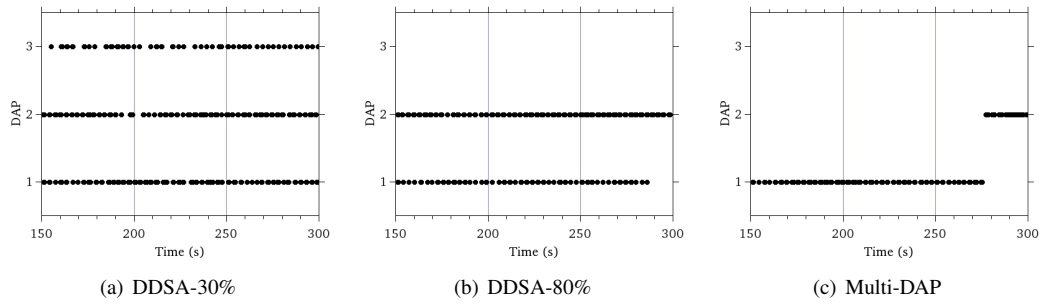
(a) DDSA-30%   (b) DDSA-80%   (c) Multi-DAP

Fig. 4. DAP choice by node 16 before failure.
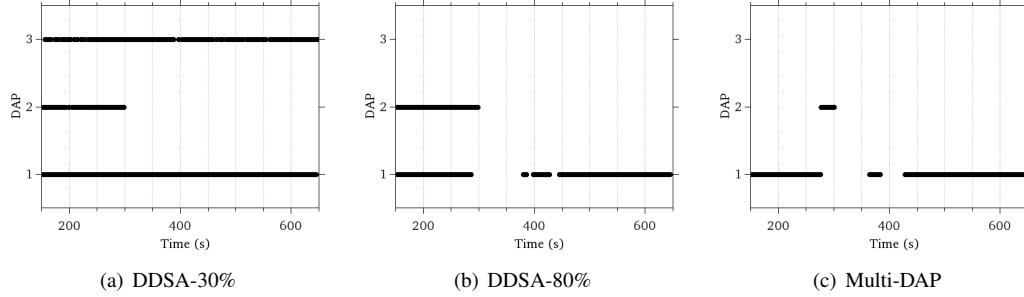

(a) DDSA-30%   (b) DDSA-80%   (c) Multi-DAP

Fig. 5. DAP choice by node 16.

(3)   the mechanism of [19] for dynamic gateway selection that chooses the best DAP according to routing metric at the time of sending the data packet. We refer to this as Multi-DAP.

B. Simulation Results

Figure 3(a) shows the packet delivery rate for all nodes as a function of time, considering the latest 60 seconds. It is noticeable that the performance of DDSA-30% is better than the others proposals after the occurrence of the DAP failure. At time 363 seconds DDSA-30% has about 95% of packet delivery rate, while DDSA-80% has 73% and Multi-DAP has only 69% of packets delivery. This happens because DDSA-30%, mainly in central region of the network (nodes 12 to 23), is less affected by the failure of the DAP 2. This can be noticed in Figure 3(b) that shows, at time of 363 seconds, about 86% of packet delivery rate while DDSA-80% and Multi-DAP have only 20.7% and 11.8%, respectively. This shows that the design choice of distributing packets among DAPs makes it more robust and resilient than others when a failure occurs. Figure 3(c) analyzes the behavior of the packet delivery rate for one node. Node 16 was chosen because it is located geographically between two DAPs. Notice that with DDSA-30%, when DAP 2 fails, the reduction in the rate was not as sharp as for the other proposals, becoming 77% and 85% better than DDSA-80% and Multi-DAP, respectively, at time 363 seconds.

To understand how the behavior of DDSA differs from the behavior of Multi-DAP in terms of DAP selection, in Figure 4 we show the choices of node 16 using the three proposals during the simulations with a single seed for the period prior to the DAP failure. As seen, the choice of α affects the behavior in choosing the DAP. Notice how the lower α (Fig. 4(a)) causes the farthest DAP (DAP 3) to be eventually chosen. The

higher α (Fig. 4(b)) makes choices alternating between DAP 1 and DAP 2 that are closest and have better metrics, excluding DAP 3 for having a quality that is too low. The Multi-DAP (Fig. 4(c)) rarely makes exchanges between DAPs, having only used one different DAP in the last 25 seconds of the graph.

Figure 5 shows the same information of Figure 4, but extends the view to the whole period of simulation. After the occurrence of the DAP failure at time 300 seconds it is verified that the packets from nodes 16 are not received by any DAP for a long time in the simulations with DDSA-80% and Multi-DAP. For DDSA-80%(Fig. 5(b)), the gap to deliver new packets lasted 64 seconds. For Multi-DAP (Fig. 5(c)), the gap lasted 81 seconds. This happens because before the failure they were sending packets only to DAP 2. Note that for DDSA-30% (Fig. 5(a)), there is no noticeable gap because node 16 already balances the load among all DAP.

To check if this behavior is shared by other nodes, we gathered the information shown in Figure 6. We show the DAP choices for all nodes using intensities of grayscale to represent the percentages of DAP choices by nodes. Higher values are represented by darker gray and lower values by lighter gray. Each matrix represents the percentages of choices for one DAP and each element is a representation of the nodes position. Notice that DDSA-30% (Fig. 6(a)) has a much more balanced DAP selection than the other proposals, as shown by the elements with intermediate gray. For nodes in the central region of the network with DDSA-30% the lowest value is 14.2%. For DDSA-80% (Fig. 6(b)) and Multi-DAP (Fig. 6(c)) six nodes (nodes 12, 15, 16, 19, 20, 23) do not choose DAP 1 or DAP 2. Except by nodes 13 and 14, in DDSA-30% all values in the central region are below 40% for DAP 2. For DDSA-80%, all values are above 38% and for Multi-DAP they are above 40%. This demonstrates that DDSA-30% distributes more evenly the packets among DAPs in central region of

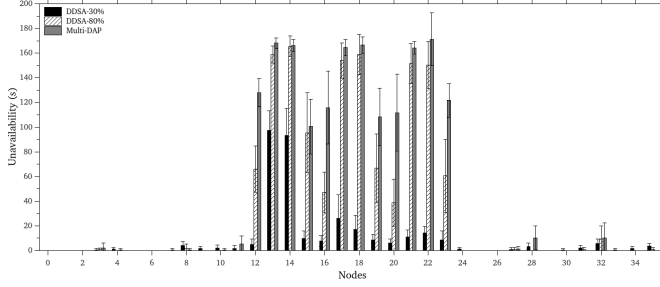network, improving the resilience against DAP failures.

Fig. 7. Unavailability Period.

Figure 7 shows the unavailability period τ for each node with respect to the headend, *i.e.*, the sum of the periods during which each node could not reach the headend through any DAP. It is noticed that the long gap to deliver new packets, as seen for node 16, is repeated for other nodes in the central region of the network. For DDSA-80% and Multi-DAP these nodes suffered long τ to deliver packets to any DAP, while DDSA-30% sustained lower τ. Nodes 13 and 14 are the closest ones to the failed DAP 2, thus suffering more influence of this failure, because its path cost is much better compared to the other DAP. Except for these nodes, the average unavailability period τ_{avg} for the DDSA-30% is 6.2 seconds and the maximum unavailability period τ_{max} is 41.8 seconds. For DDSA-80%, τ_{avg} is 33.7 seconds and τ_{max} is 161.6 seconds, and for Multi-DAP τ_{avg} is 42.6 seconds and τ_{max} is 170.3 seconds.
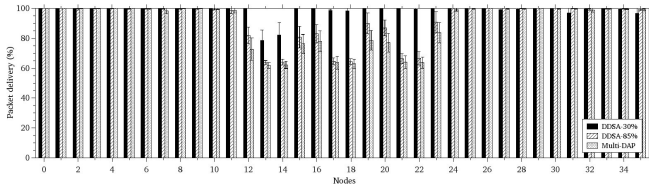

Fig. 8. Packet delivery rate for each node.

Finally, Figure 8 shows the packet delivery rate for all nodes considering all simulation time. It is observed that the central region of the network is the most affected by the failure of the DAP. For these nodes, DDSA-30% is the less affected by the fault with a performance above DDSA-85 % and Multi-DAP, for which most of nodes in the central region showed a low packet delivery rates.

VI. CONCLUSION

This work presented DDSA, a dynamic DAP selection algorithm to increase the reliability and resilience of AMI applications through the use of multiple DAPs in WMN networks. In this kind of network, the DAP has an important role in exchanging information between the meter and the headend because all traffic flows through it. A failure in a DAP inhibits the exchange of information on the AMI network, so alternative routes through other DAPs should be used after failure to sustain communication between meters and headend.

The results obtained showed that the DDSA increased the resilience of routing protocol even suffering with DAP failure

evidenced by lower loss in performance. The results also showed the importance of the choice of the α parameter that influences the routing behavior. Lower values favor resilience, because DDSA distributes packets among more DAPs.

For future work we intend to do a deeper analysis of the impact of the α parameter to observe the behavior of DDSA and find values that bring more resilience or improving performance in routing. We also intend to study the use of dynamic adaptation of the α parameter and to analyze other metrics to improve the performance and resilience in a region with failure.

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