

# Detecting Variable Length Anomaly Patterns in Time Series Data

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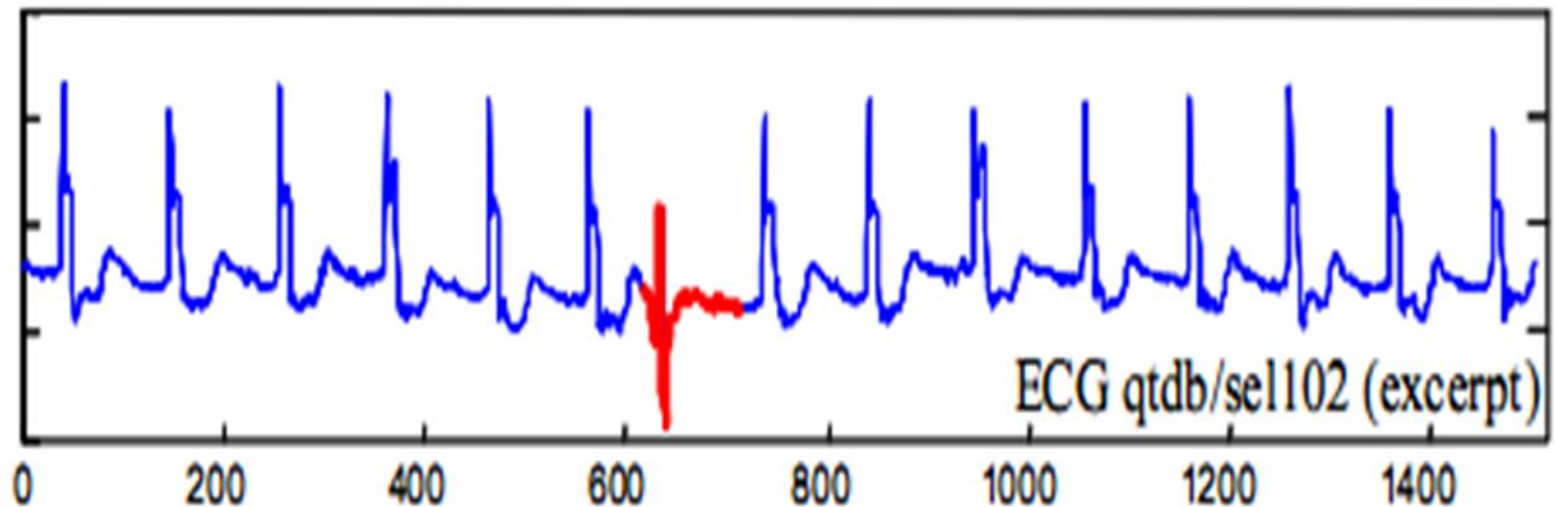
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# Introduction

- A time series is a sequence of data points made over a continuous time interval
- Anomaly Patterns are subsequences that do not conform to a well defined notion of normal behavior
- Application areas that explore such time series anomalies include fault diagnostics, intrusion detection, fraud detection, auditing and data cleansing

# Introduction



unusual subsequence in time series data

# Introduction

- Anomalous subsequences detection is a challenging topic
- The exact lengths of the unusual subsequence are vary and are often unknown.
- Best similarity/distance measures which can be used for different types of time series is not easy to determine.
- Time series in real applications are usually long and as the length increases the computational complexity also increases.

# Some Definitions

- *Time Series*: A time series  $T = t_1, t_2, \dots, t_m$  is an ordered set of  $m$  real values measured at equal intervals.
- *Subsequence*: Given a time series  $T$  of length  $m$ , a subsequence  $C$  is a sampling of length  $n < m$  of contiguous positions from  $T$ , i.e.,  $C = t_p, t_{p+1}, \dots, t_{p+n-1}$ , for  $1 \leq p \leq m-n+1$ . Sometimes,  $C$  is denoted as  $(s_p, e_{p+n-1})$ , where  $s_p = t_p$  and  $e_{p+n-1} = t_{p+n-1}$ .

# Some Definitions

- *Distance function*:  $Dist(C, M)$  is a positive value used to measure the difference between two time series  $C$  and  $M$ , based on some measure method
- *k-distance of a pattern*: Given a positive integer  $k$ , a pattern set  $D$  and a pattern  $P \in D$ , the  $k$ -distance of  $P$ , denoted as  $k-dist(P)$ , is defined as the distance between  $P$  and a pattern  $Q \in D$  such that.
  - i) For at least  $k$  patterns  $Q' \in D$  it holds that  $Dist(D, Q') \leq Dist(D, Q)$ .
  - ii) For at most  $k-1$  patterns  $Q' \in D \setminus \{Q\}$  it holds that  $Dist(D, Q') < Dist(D, Q)$ .

# Some Definitions

- *Non-Self Match*: Given a time series  $T$ , its two subsequences  $P$  of length  $n$  starting at position  $p$  and  $Q$  starting at position  $q$ , we say that  $Q$  is a non-self-match to  $P$ , if  $Dist(P, Q) \geq e$  or  $|p - q| \geq n$ , where  $e$  is a given value of distance threshold.
- *Anomaly factor*: For any pattern set  $D$  and a pattern  $P \in D$ ,  $k-dist(D)$  denotes all  $k-dist$  of patterns, anomaly factor of pattern  $P$  defined as the ratio of  $k-dist(P)$  to  $median(k-dist(D))$ .
- *Anomaly Pattern*: Given any pattern set  $D$ , a pattern  $P \in D$ ,  $P$  is anomaly only if its anomaly factor is larger than  $a$ , where  $a$  is the threshold of anomaly pattern



# Related Works

- Hot Sax (E. Keogh et al. 2005 ).
- WAT (Y. Bu et al. 2007).
- Method based on PAA bit representation and clustering (Li et al. 2013)
- Method based on Segmentation and Anomaly Factors (M.Leng et al. 2008)

# HOT SAX

- Find discord of length  $n$ .
- A discord is a subsequence which have largest nearest neighbor distance.
- Use heuristic to improve brute-force.
- Distance function: Euclid distance.

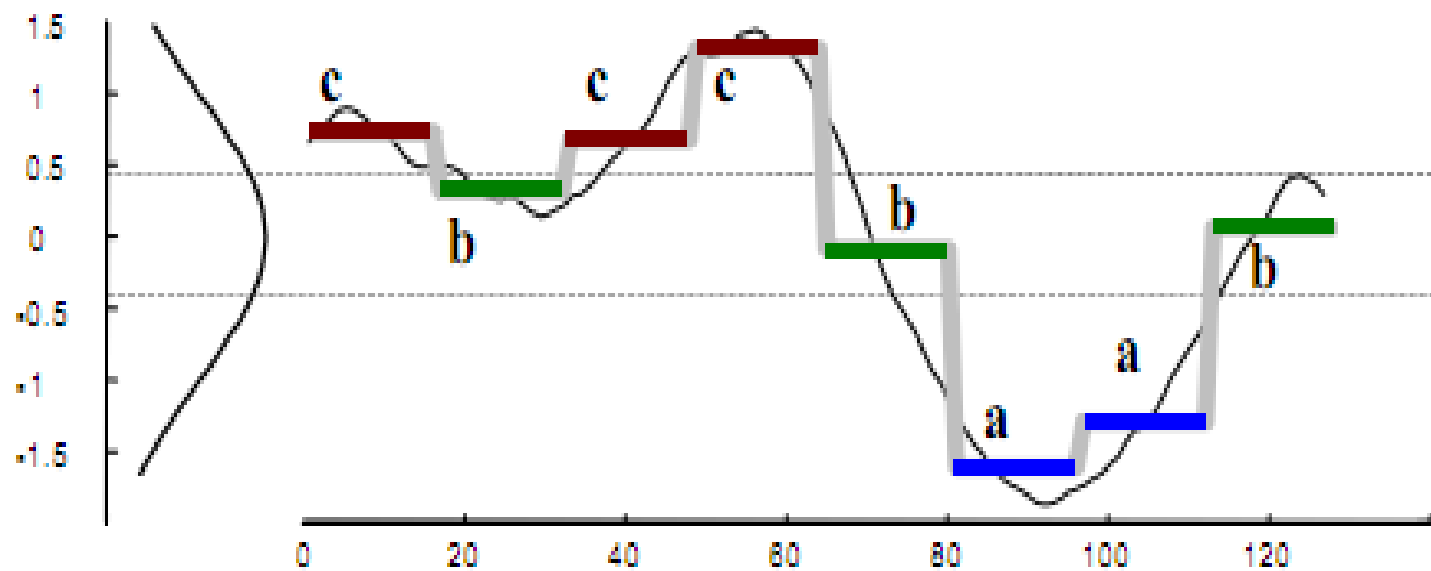
# HOT SAX

```
Function  $[dist, loc] = \text{Heuristic\_Search}(T, n, \text{Outer}, \text{Inner})$   
   $best\_so\_far\_dist = 0$   
   $best\_so\_far\_loc = NaN$   
  For Each  $p$  in  $T$  ordered by heuristic Outer // Begin Outer Loop  
     $nearest\_neighbor\_dist = \text{infinity}$   
    For Each  $q$  in  $T$  ordered by heuristic Inner // Begin Inner Loop  
      IF  $|p - q| \geq n$  // non-self match?  
        IF  $\text{Dist}((t_p, \dots, t_{p+n-1}), (t_q, \dots, t_{q+n-1})) < best\_so\_far\_dist$   
          Break // Break out of Inner Loop  
        End  
        IF  $\text{Dist}((t_p, \dots, t_{p+n-1}), (t_q, \dots, t_{q+n-1})) < nearest\_neighbor\_dist$   
           $nearest\_neighbor\_dist = \text{Dist}((t_p, \dots, t_{p+n-1}), (t_q, \dots, t_{q+n-1}))$   
        End  
      End // End non-self match test  
    End // End Inner Loop  
    IF  $nearest\_neighbor\_dist > best\_so\_far\_dist$   
       $best\_so\_far\_dist = nearest\_neighbor\_dist$   
       $best\_so\_far\_loc = p$   
    End  
  End // End Outer Loop  
  Return  $[best\_so\_far\_dist, best\_so\_far\_loc]$ 
```

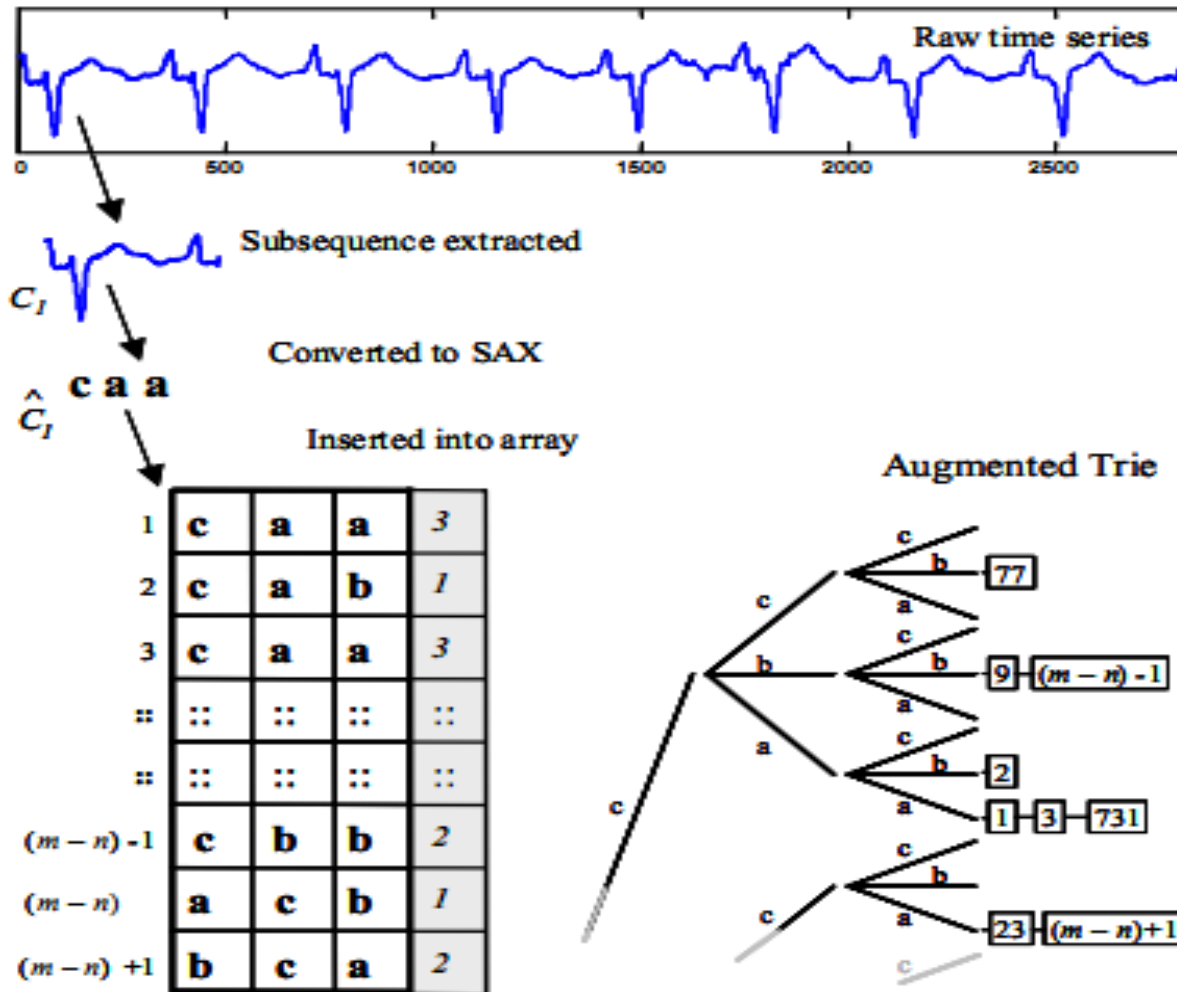
# HOT SAX

- Present subsequences of length  $n$  to SAX words.
  - Use PAA for Dimensionality reduction
- Use an array and an augmented trie to embed all the SAX words

# HOT SAX



# HOT SAX



# HOT SAX

- In outer loop, subsequences presented by words which have smallest appearing frequency in the array are checked first
- In inner loop, subsequences belong to the same left node in the augmented trie with the subsequence chosen at outer loop are checked first.

# WAT

- Same HOT SAX but use Haar wavelet instead of PAA.



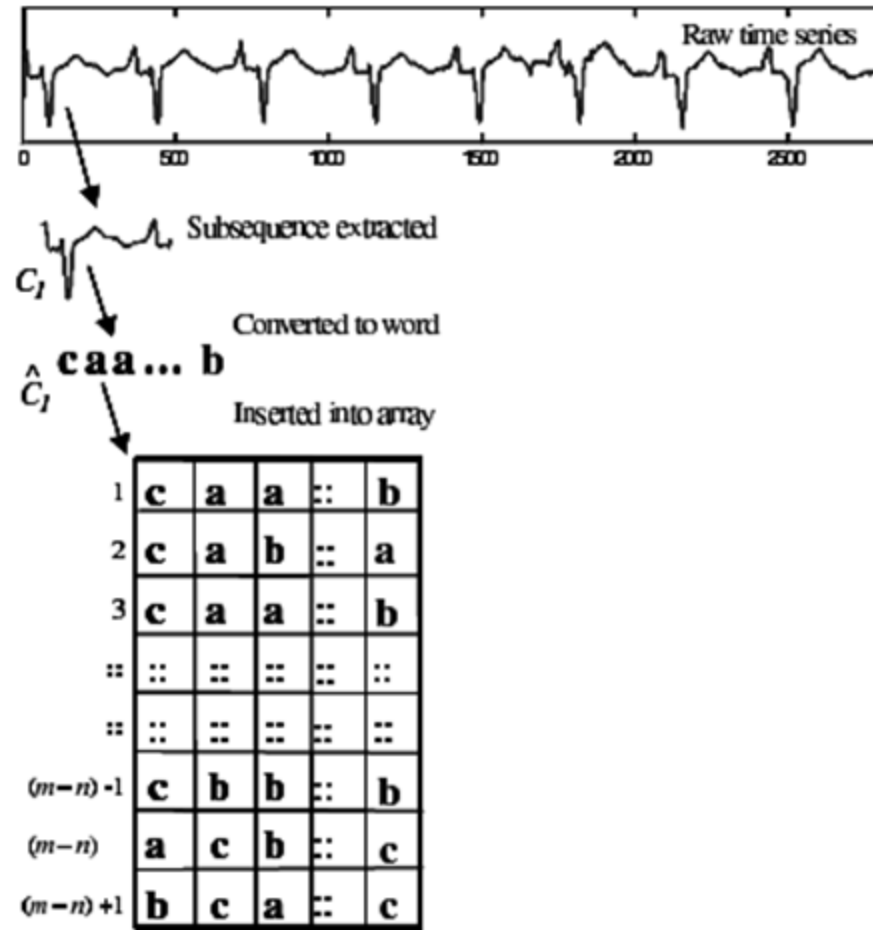
# WAT

- Haar wavelet can be seen as a series of averaging and differencing operations
- Ex:  $f(x) = (9, 7, 3, 5) \rightarrow (6, 2, 1, -1)$

Resolution   Averages   Coefficients

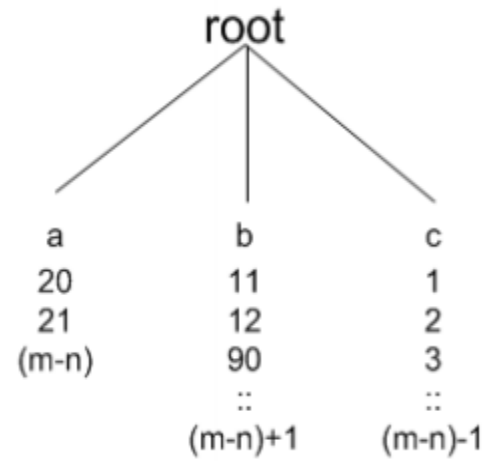
4	(9 7 3 5)	
2	(8 4)	(1 -1)
1	(6)	(2)

# WAT



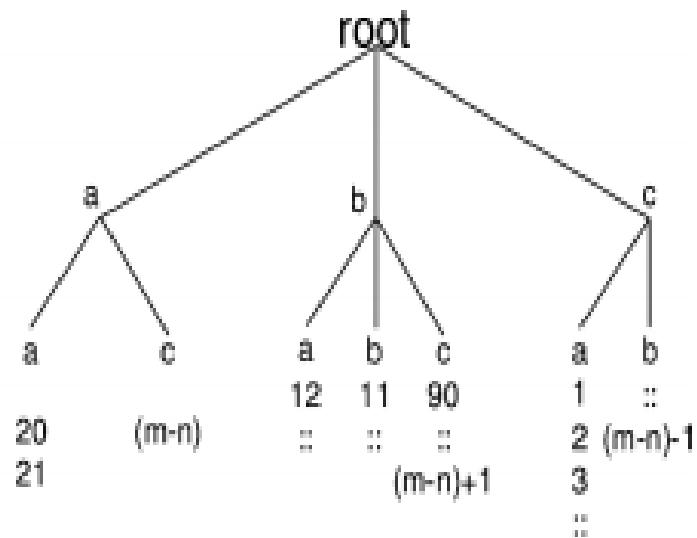
# WAT

1	<b>c</b>	<b>a</b>	<b>a</b>	<b>::</b>	<b>b</b>
2	<b>c</b>	<b>a</b>	<b>b</b>	<b>::</b>	<b>a</b>
3	<b>c</b>	<b>a</b>	<b>a</b>	<b>::</b>	<b>b</b>
::	<b>::</b>	<b>::</b>	<b>::</b>	<b>::</b>	
::	<b>::</b>	<b>::</b>	<b>::</b>	<b>::</b>	
$(m-n)-1$	<b>c</b>	<b>b</b>	<b>b</b>	<b>::</b>	<b>b</b>
$(m-n)$	<b>a</b>	<b>c</b>	<b>b</b>	<b>::</b>	<b>c</b>
$(m-n)+1$	<b>b</b>	<b>c</b>	<b>a</b>	<b>::</b>	<b>c</b>



# WAT

1	<b>c</b>	<b>a</b>	<b>a</b>	::	<b>b</b>
2	<b>c</b>	<b>a</b>	<b>b</b>	::	<b>a</b>
3	<b>c</b>	<b>a</b>	<b>a</b>	::	<b>b</b>
::	::	::	::	::	
::	::	::	::	::	
$(m-n)-1$	<b>c</b>	<b>b</b>	<b>b</b>	::	<b>b</b>
$(m-n)$	<b>a</b>	<b>c</b>	<b>b</b>	::	<b>c</b>
$(m-n)+1$	<b>b</b>	<b>c</b>	<b>a</b>	::	<b>c</b>

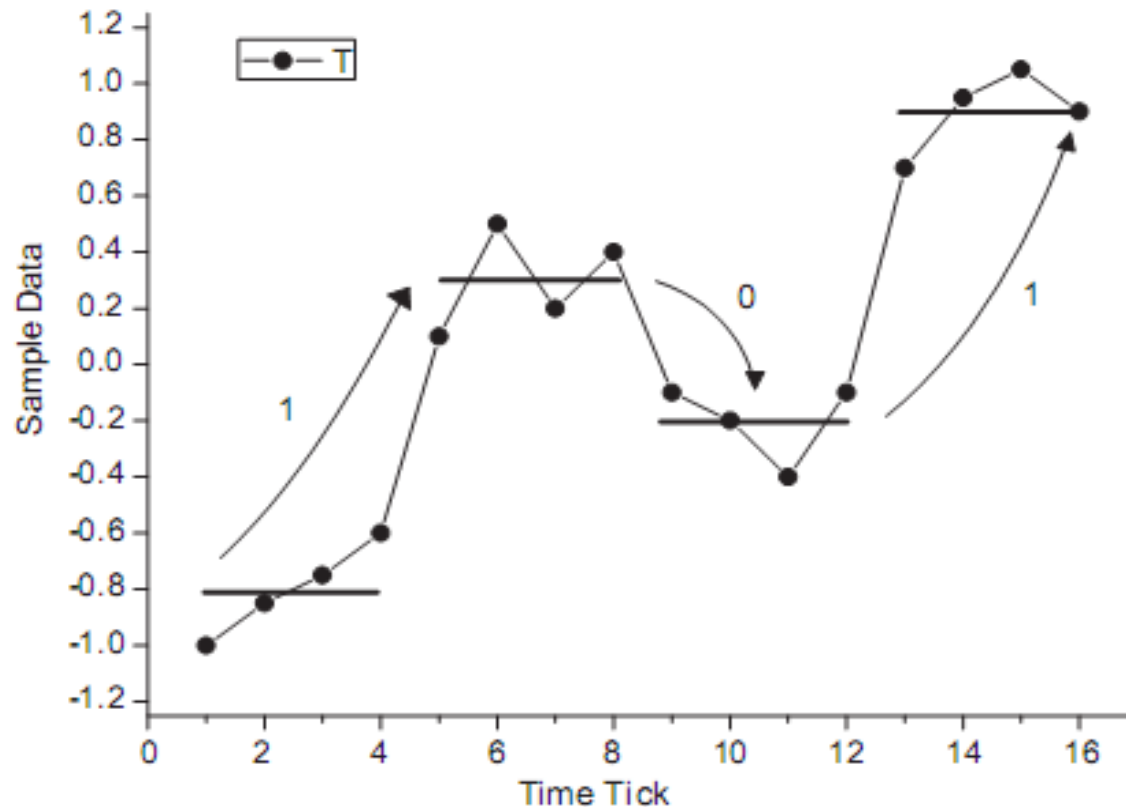


# WAT

- In outer loop, subsequences presented by words which have smallest appearing frequency in the array are checked first
- In inner loop, subsequences belong to the same left node in the tree with the subsequence chosen at outer loop are checked first.

# Method based on PAA bit representation and clustering

- PAA Bit representation of a series



# Method based on PAA bit representation and clustering

- Bit distance between of two PAA bit representation  $s = \{s_1, s_2, \dots, s_{w-1}\}$  and  $t = \{t_1, t_2, \dots, t_{w-1}\}$
- $BitSeries\_Dist(s, t) = \frac{1}{w-1} \sum_1^{w-1} BDist(s_i, t_i)$
- Here,  $BDist(s_i, t_i) = \begin{cases} 1 & \text{if } s_i \neq t_i \\ 0 & \text{otherwise} \end{cases}$

# Method based on PAA bit representation and clustering

- $k$  bit representations are randomly chosen as  $k$  cluster centres
- A bit representation is assigned to the nearest cluster.
- In outer loop, subsequences which have bit representations belong to the smallest cluster (cluster has the fewest elements) is checked first
- In inner loop, subsequences belong to the same cluster with the subsequence chosen at outer loop are checked first.



# Method based on PAA bit representation and clustering

- 3 above algorithms need to know the length of anomaly subsequence in advance.

# Method based on Segmentation and Anomaly Factors

- Proposed by Leng et al. in 2008
- The method includes 2 phase
- Phase 1: use quadratic regression model to segment time series. Regression function is defined as  $f(t) = \beta_0 + \beta_1 t + \beta_2 t^2$ ,
- Phase 2: Calculate anomaly factors for the subsequences and show abnormal patterns.

1. Let  $s_l = 1$  denote the start position of the first segment,  $l$  denote the initial length of each segment,  $m = s_l$  and then update this segment, the updating procedure is as follows. Calculate the values of  $\beta_0$ ,  $\beta_1 t$ , and  $\beta_2$  and the sum of squared errors:

$$SSE = \sum_{i=p}^{p+l-1} (f(i) - t_i)^2$$

2. If  $SSE < \varepsilon_l$  let  $l = l + 1$ , go to step 1, else go to step 3.
3. Let  $(s_l, e_l)$  denote this segment,  $e_l$  denotes the end position of this segment and  $e_l = l - 1$ .
4. Let  $i = 1$ , if  $Dist((s_l, e_l), (s_l + i, e_l + i)) \leq \varepsilon_2$ , increment  $i$  by 1, calculate it, repeat this procedure until  $Dist((s_l, e_l), (s_l + i, e_l + i)) > \varepsilon_2$  or  $i > (e_l - s_l)$ , let  $s_2 = e_l + i$ , the  $s_2$  is the start position of the next segment.
5. Let  $m = s_2$ , calculate  $e_2$  using the steps 1-3, and then obtain the second segment  $(s_2, e_2)$ .
6. Calculate  $s_j$  using the step 4, let  $m = s_j$ , calculate the value of  $e_j$  using the steps 1 – 3, and then obtain the  $j$ -th segment  $(s_j, e_j)$ .
7. Repeat the above procedure until the end position of the segment is  $n$ .

1. Find the maximum value,  $l_{max}$  and minimum value  $l_{min}$  of the lengths of all extracted segments.
2. Calculate the distance matrix  $D = (d_{ij})_{m \times m}$  of segments:

$$d_{ij} = \min_{l_{min} \leq l \leq l_{max}} Dist((s_i, e_i), (s_j, s_j + l)) \quad (1)$$

where  $1 \leq i, j \leq m$  and  $i \neq j$ .

3. Compute k-distance of each segments based on the distance matrix, let  $k-dist(D)$  denote the set of all these distances.
4. Calculate  $median(k-dist(D))$
5. Calculate anomaly factor of each segment, and determine whether each segment is an anomaly pattern or not basing on the anomaly factor threshold  $a$ .
6. Given two anomaly patterns  $(s_i, e_i)$  and  $(s_j, e_j)$ , if they are overlapped, merge them into one pattern.

# Method based on Segmentation and Anomaly Factors

- Don't need to know the length of anomaly patterns in advance
- Use DTW to calculate distance between two subsequences of different length
- Time consuming

# Proposed Algorithms

- Use Homothetic Transformation + Modified Euclidean Distance instead of DTW
- Reduce the number of distance calculating in building distance matrix.
- Proposed a segmentation method based on important extreme points as alternative segmentation method

# New distance function

- A Homothetic Transformation with center  $O$  and ratio  $k$  transforms the point  $M$  to the point  $M'$  such that

$$\overrightarrow{OM'} = k \times \overrightarrow{OM} .$$

# New distance function

- Apply to time series, a homothety transforms a time series  $T$  of length  $n$  ( $T = \{y_1, y_2, \dots, y_n\}$ ) to time series  $T'$  of length  $n'$  by performing the following steps.
- First, compute  $Y\_MAX = MAX(y_1, \dots, y_n)$ ,  $Y\_MIN = MIN(y_1, \dots, y_n)$ .
- Second, set the center  $I$  of homothety with the coordinates  $X\_C = n/2$ ,  $Y\_C = (Y\_MAX + Y\_MIN)/2$ . Next, perform the homothety with the center  $I$  and the ratio  $n'/n$



# New distance function

```
Function Dist(X, Y)  
if (X.length == Y.length) return Euclid(X, Y)  
else  
    MEAN =  $|(\textit{Y.length} + \textit{X.length})| / 2$   
    Y' = Homothety(Y, MEAN)  
    X' = Homothety(X, MEAN)  
    return Euclid(X', Y')
```

# New distance function

- Minimum Euclidean Distance between  $T' = \{T'_1, T'_2, \dots, T'_{N'}\}$  and  $Q' = \{Q'_1, Q'_2, \dots, Q'_{N'}\}$

$$D(T', Q') = \text{Min} \left\{ \sum_{i=1}^{N'} (T'_i - Q'_i - b)^2 \right\}$$

$$b = \frac{1}{N'} \sum_{i=1}^{N'} (Q'_i - T'_i)$$

Time Complexity of new distance function:  
 $O(n)$

# Reduce the number of distance calculating

- Propose parameter  $r$
- Replace formula (1) to

$$d_{ij} = \min_{l_{lower} \leq l \leq l_{upper}} Dist((s_i, e_i), (s_j, s_j + l)) \quad (1')$$

Where

$$l_{upper} = \lfloor l_{avg}(1 + r) \rfloor$$

$$l_{lower} = \lceil l_{avg}(1 - r) \rceil$$

$l_{avg}$  is mean length of the subsequences.

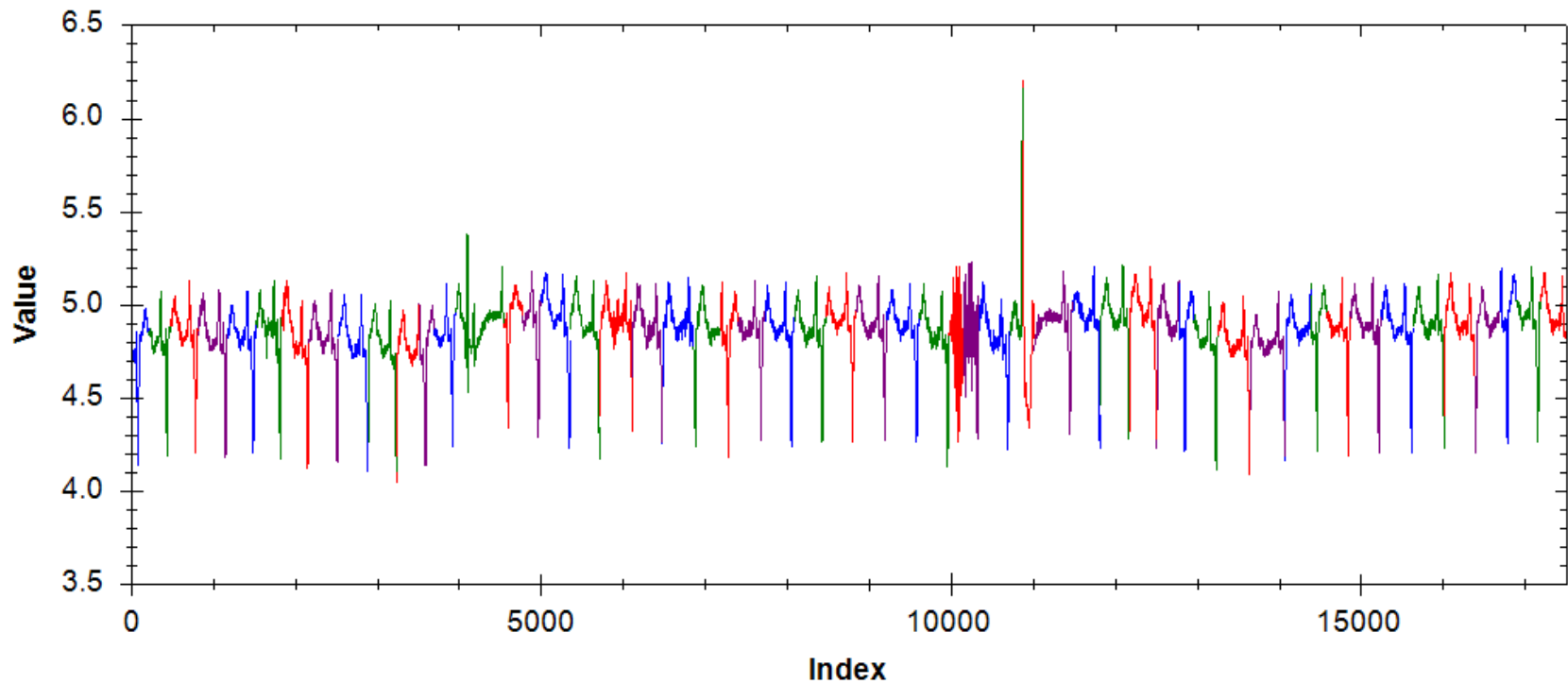
# Alternative segmentation method

- Based on the concepts of Important Extreme Points, proposed by Pratt and Fink in 2002
- Extract all important extreme points of the time series  $T$ . The result of this step is a sequence of extreme points  $EP = (ep_1, ep_2, \dots, ep_l)$
- Compute all the candidate patterns iteratively. A candidate pattern  $CP_i(T)$ ,  $i = 1, 2, \dots, l-2$  is the subsequence of  $T$  that is bounded by extreme points  $ep_i$  and  $ep_{i+2}$ .

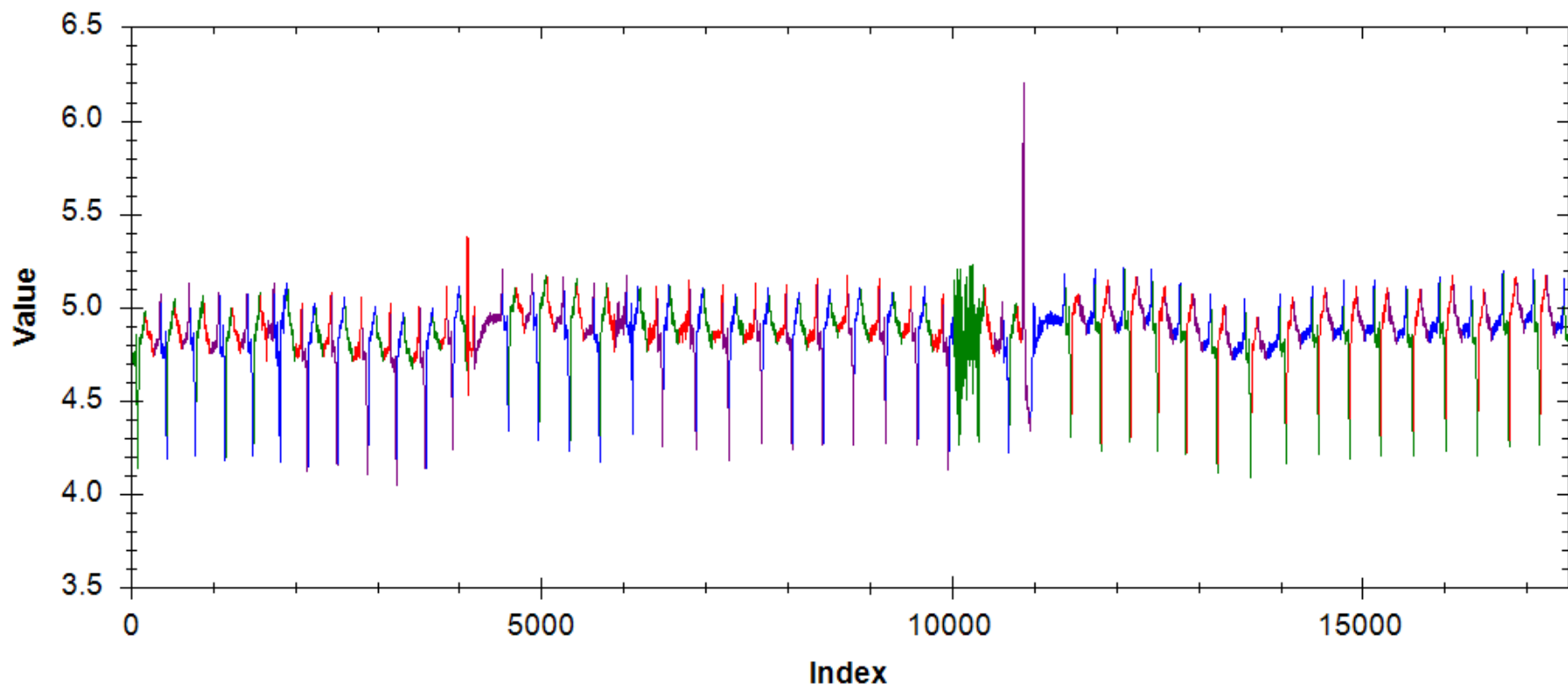
# Alternative segmentation method

- **Parameter:**

- Compression rate:  $R > 1$
- Lower bound for the distance two extracted extreme points: *min\_length*. So the minimum length of a subsequence is  $2min\_length$



Segmentation based on quadratic regression



Segmentation based on important extreme points

# Experimental Evaluation

- VL\_QR| HT: The proposed algorithm with quadratic regression
- VL\_EP| HT: The proposed algorithm with important extreme points.
- Compare results and running time among VL\_QR| HT, VL\_EP| HT and HOT SAX
- Compare running time between VL\_QR| HT with the original algorithm proposed by Leng et al.



# Experimental Evaluation

- Experiments environment : Intel® Core™ 2 Duo 2.0GHz, Ram 3072MB PC
- All algorithms are implemented in Microsoft Visual C# programming language

# Compare VL\_QR| HT, VL\_EP| HT and HOT SAX

- Use 8 time series in diverse domains:  
ECG 108, ECG 308, ERP, Memory, Power Demand In Italy, Dutch Power Demand, Stock20 and TEK16
- To compute location difference of anomaly subsequences use:

$$d = |p - q| / l \times 100$$

where  $p$  is the start position of the anomaly pattern found by proposed algorithm and  $q$  is the start position of the anomaly pattern found by HOT SAX and  $l$  is its length

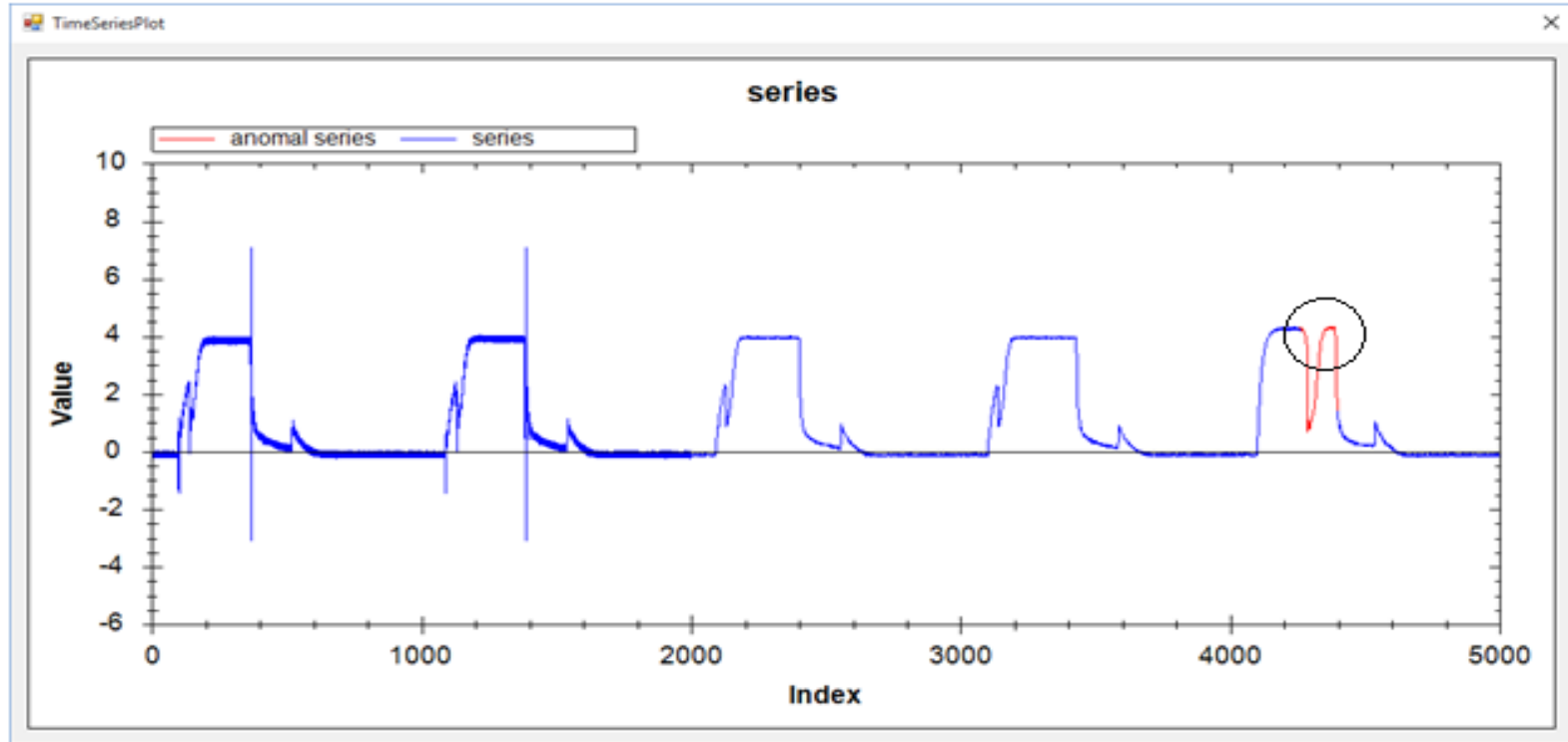
# Parameter values in the three algorithms for each series

Datasets	VL_QR HT	VL_EP HT	HOTSAX
ECG108 (17500 points)	$\varepsilon_1 = 5.0, \varepsilon_2 = 0.3, \alpha = 3.5, r = 0.1$	$R = 1.04, \min\_length = 50, \alpha = 4, r = 0.1$	$n = 600, PAA\_size = 60$
ECG308 (1300 points)	$\varepsilon_1 = 51000, \varepsilon_2 = 50, \alpha = 2.2, r = 0.1$	$R = 1.03, \min\_length = 10, \alpha = 2.5, r = 0.23$	$n = 60, PAA\_size = 6$
ERP (5000 points)	$\varepsilon_1 = 3.0, \varepsilon_2 = 1.0, \alpha = 3.0, r = 0.15$	$R = 1.42, \min\_length = 10, \alpha = 3.5, r = 0.1$	$n = 100, PAA\_size = 10$
Memory (6875 points)	$\varepsilon_1 = 8.0, \varepsilon_2 = 0.1, \alpha = 2.2, r = 0.1$	$R = 1.1, \min\_length = 40, \alpha = 1.6, r = 0.1$	$n = 100, PAA\_size = 20$
Power Demand Italy (7000 points)	$\varepsilon_1 = 100000, \varepsilon_2 = 100, \alpha = 2.5, r = 0.1$	$R = 1.8, \min\_length = 20, \alpha = 3.0, r = 0.1$	$n = 300, PAA\_size = 30$
Dutch Power Demand (9000 points)	$\varepsilon_1 = 31000000, \varepsilon_2 = 500, \alpha = 1.6, r = 0.2$	$R = 1.7, \min\_length = 180, \alpha = 1.33, r = 0.25$	$n = 1200, PAA\_size = 120$
Stock20 (5000 points)	$\varepsilon_1 = 2.0, \varepsilon_2 = 0.1, \alpha = 3.5, r = 0.1$	$R = 1.01, \min\_length = 200, \alpha = 1.5, r = 0.1$	$n = 700, PAA\_size = 70$
TEK1 (5000 points)	$\varepsilon_1 = 14.5, \varepsilon_2 = 0.1, \alpha = 2.0, r = 0.1$	$R = 1.02, \min\_length = 50, \alpha = 1.5, r = 0.1$	$n = 128, PAA\_size = 8$

# Positions of anomaly patterns detected by VL\_QR|HT and VL\_EP|HT and HOT SAX

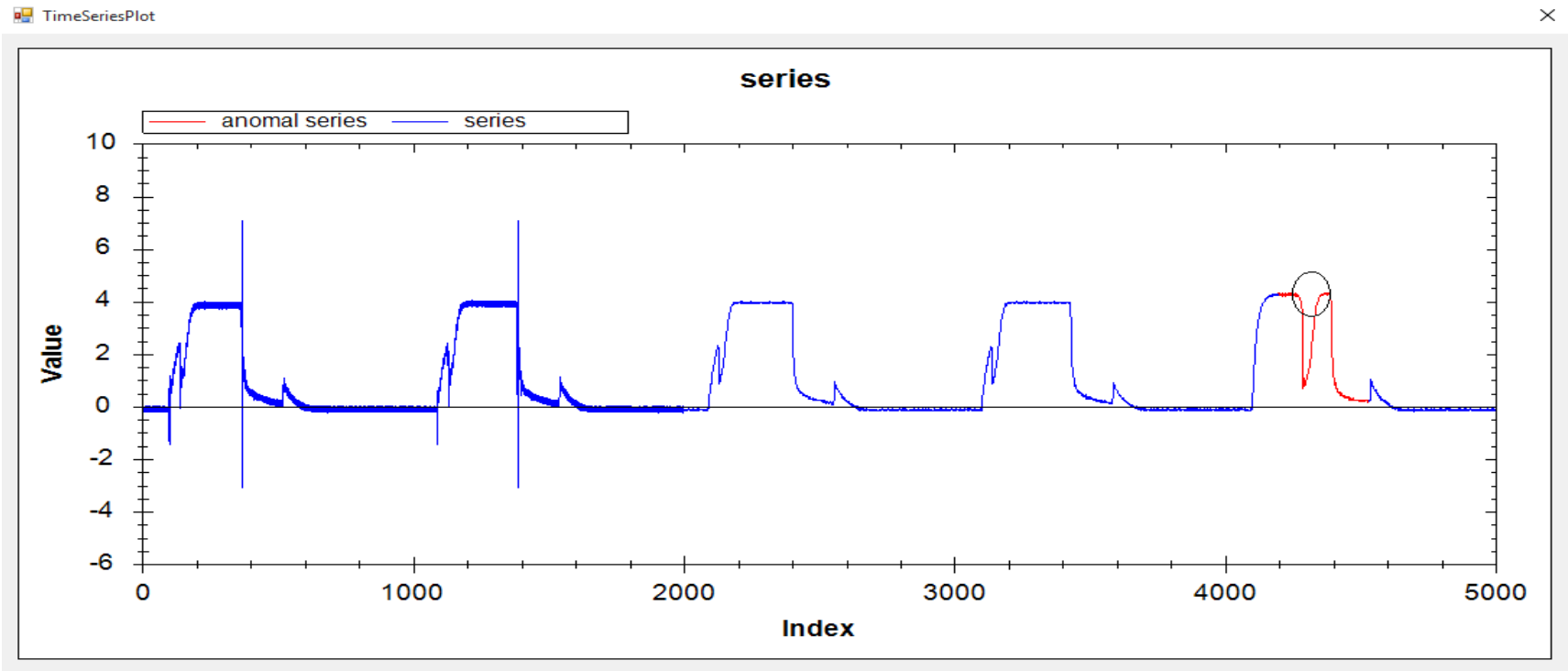
Dataset	VL_QR HT		VL_EP HT		HOT SAX
	Start position	Location difference ( $d$ )	Start position	Location difference ( $d$ )	Start position
ECG108	10875	1.5%	10792	12.3%	10866
ECG308	656	33%	673	5%	676
ERP	2617	32%	2602	17%	2585
Memory	2425	2%	2399	11%	2421
Power Demand Italy	5311	4%	5383	13.33%	5343
Dutch Power Demand	6424	1.3%	6277	11%	6466
Stock20	2961	1.7%	2960	1.6%	2949
TEK16	4262	24%	4197	27%	4231

# Positions of anomaly patterns detected by VL\_QR|HT and VL\_EP|HT and HOT SAX



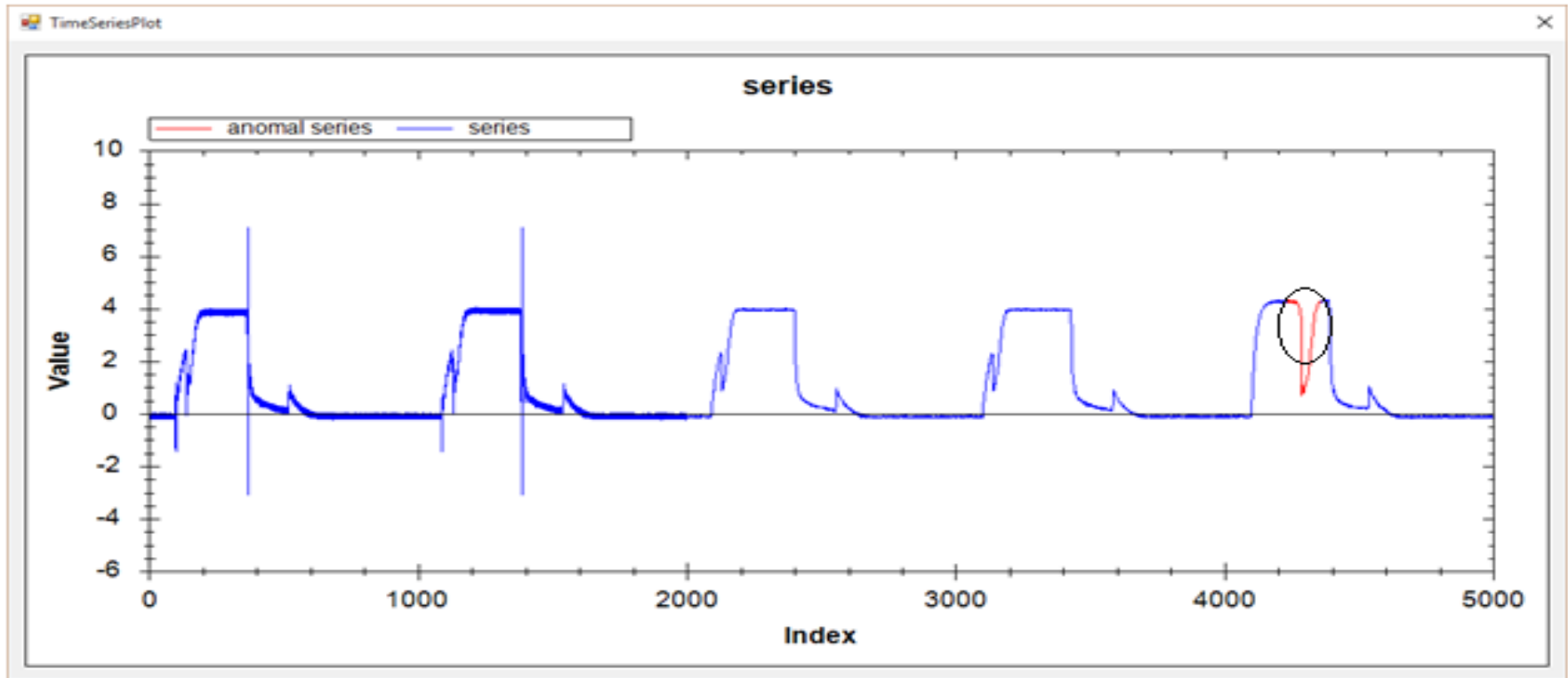
- Result of VL\_QR | HT with series TEK16

# Positions of anomaly patterns detected by VL\_QR|HT and VL\_EP|HT and HOT SAX



- Result of VL\_EP | HT with series TEK16

# Positions of anomaly patterns detected by VL\_QR|HT and VL\_EP|HT and HOT SAX



- Result of HOT SAX with series TEK16

# Running time among VL\_QR|HT, VL\_EP|HT and HOT SAX (s)

Dataset	VL_QR HT	VL_EP HT	HOT SAX
ECG108	10	47	458
ECG308	0.31	2	13
ERP	7	14	60
Memory	4	6	79
Power Demand Italy	9	11	102
Dutch Power Demand	1	9	302
Stock20	1	1	94
TEK16	1	1	52



# Running time between VL\_QR| HT with the original algorithm (s)

Dataset	VL_QR HT	VL_QR DTW
ECG 108 (1000 points)	1	14
ECG 308 (1300 points)	1	11
ERP (1000 points)	1	5
Memory (1000 points)	1	4
Power Demand In Italy (3000 points)	6	200
Dutch Power Demand (3000 points)	1	24
Stock20 (3000 points)	2	194
TEK16 (2000 points)	1	81

# Conclusion

- VL\_QR|HT and VL\_EP|HT bring out a remarkable improvement for the original algorithm in time efficiency without compromising anomaly detection accuracy.
- Experimental results on eight datasets demonstrate these algorithms outperform the VL\_QR|DTW in time efficiency.
- Not easy to estimate the regression error threshold  $\varepsilon_1$  and the non-self match threshold  $\varepsilon_2$  in the segmentation phase of VL\_QR|HT.
- Not easy to find the  $R$  in VL\_EP|HT.
- Future work: apply some more advanced method of time series segmentation.

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Q & A

Thank you!