Maximizing Client Throughput

Table of Contents

- 3. Introduction
- 4. Benchmarking System
 - Key Metrics Tracked
 - How the Benchmarking System Works
- 5. Enhancing the Rate Limiter
 - Original Implementation
 - Identified Issues
 - Proposed Solution
 - Observation: Performance Improvement
 - Potential Issue: 429 Errors Due to Latency
 - Improved Version: Adaptive Buffering
- 6. Improving the Queue System
 - o Issue
 - Solution: Queue Manager with Dead Letter Queue (DLQ)
 - Lifecycle with Queue Manager
 - Monitoring the Queue
- 7. Addressing Queue Bloating
- 8. Exploring Multithreading
 - Rationale
 - Changes to the Current Code
- 9. Comparison Between Asynchronous and Threading Client
 - Baseline Comparison
 - Observation
 - Summary
- 10. Overview and Modifications Summary
 - Modifying the Original Client
- 11. Final Thoughts

Introduction

This document provides a review and optimization of a client application aimed at maximizing throughput while adhering to server-imposed rate limits. The client interacts with a server using multiple API keys and is designed to send as many requests as possible without exceeding rate limits or causing request failures. It includes the issues identified in the existing implementation, the design choices made to address them, and the impact of these changes on performance.

Link to the github repository here

Benchmarking System

Key Metrics Tracked

1. Total Successful Requests

• Why: Indicates system stability and efficiency under load.

2. Total Failed Requests

• Why: Identifies network or timeout issues to ensure retries and prevent bottlenecks.

3. Average Latency (ms)

• Why: Ensures fast response times and improves user experience.

4. Throughput (TPS)

• Why: Measures the system's capacity to handle high traffic without degradation.

How the Benchmarking System Works

• Recording Successful/Failed Requests:

Logs successes with latencies and tracks failures to monitor stability.

• Calculating Average Latency:

Computes the mean latency of all successful requests to assess performance.

Calculating Throughput:

Measures successful requests per second from the start of the benchmark.

• Metrics Printing:

Regularly prints key metrics for real-time feedback during execution.

Enhancing the Rate Limiter

Original Implementation

```
class RateLimiter:
    def __init__(self, per_second_rate, min_duration_ms_between_requests):
        self.__per_second_rate = per_second_rate
        self.__min_duration_ms_between_requests =
min_duration_ms_between_requests
        self.__last_request_time = 0
        self.__request_times = [0] * per_second_rate
        self.__curr_idx = 0
    @contextlib.asynccontextmanager
    async def acquire(self, timeout_ms=0):
        enter_ms = timestamp_ms()
        while True:
            now = timestamp_ms()
            if now - enter_ms > timeout_ms > 0:
                raise RateLimiterTimeout()
            # Fixed Interval Check
```

Identified Issues

The current rate limiter has two conditional statements that introduce brief sleeps to control the rate at which requests are sent:

1. **Fixed Interval Check:** Ensures that the time interval between consecutive requests does not fall below a predefined minimum (min_duration_ms_between_requests).

```
if now - self.__last_request_time <=
self.__min_duration_ms_between_requests:
    await asyncio.sleep(0.001)
    continue</pre>
```

2. **Circular Buffer Check:** Ensures that no more than a specified number of requests are sent within any 1-second window.

```
if now - self.__request_times[self.__curr_idx] <= 1000:
    await ascynio.sleep(0.001)
    continue</pre>
```

Redundancy in Checks:

• Both checks aim to control request rates.

Proposed Solution

We propose removing the **Fixed Interval Check**, allowing the **Circular Buffer Check** to regulate the request rate effectively. This change reduces unnecessary context switching and improves performance. The **Circular Buffer Check** handles both bursty and constant-rate traffic efficiently.

Importance of Handling Both Traffic Types:

• **Bursty Traffic:** Useful in scenarios like high-frequency trading, where multiple actions need to be performed quickly within a short time frame. For example, executing a series of buy/sell actions to capitalize on rapid price movements in a volatile market.

• Constant-Rate Traffic: Ideal for retrieving data at regular intervals to maintain accuracy and consistency. This applies to oracle services or pricing feeds, where information must be regularly updated to reflect real-time market conditions.

Observation: Performance Improvement

Benchmark Results:

After removing the Fixed Interval Check, throughput increased from ~74 TPS to ~85 TPS.

Reasoning:

- The **Fixed Interval Check** introduces **context switching**. Each time this check fails, the coroutine **yields control** back to the event loop, resulting in **overhead** because:
 - 1. The **state** of the coroutine must be saved.
 - 2. The **event loop switches** to another coroutine.
 - 3. Later, the original coroutine's state is restored to continue execution.
- Removing this check reduces context switching and scheduling delays, improving the efficiency of coroutine execution.
- The Circular Buffer Check alone ensures requests are properly spaced across a 1-second window, allowing the client to handle both burst and constant-rate traffic without unnecessary pauses. This results in smoother, more efficient request handling. Additionally, the server side rate limiter also uses circular buffer to track the requests, adhering to the same rate limiting rules will be optimal.

Potential Issue: 429 Errors Due to Latency

While removing the **Fixed Interval Check** improves throughput, it may still lead to **429 errors**. This is due to **timestamp discrepancies** between the **client** and **server** rate limiters, caused by **incoming latency differences** on the server-side.

How the Issue Occurs:

- The client-side rate limiter records the difference between the current request and the 20th previous request as 1000 ms (1 second), passing the client-side check.
- However, due to latency variations, the server might calculate the time difference as less than
 1000 ms, causing the request to fail the servers rate limit and result in a 429 Error.

Example Walkthrough:

1. 1st Request:

- Sent at **0 ms** from the client.
- Server timestamp: Received at 40 ms (due to 40 ms latency).
- Recorded on the server's circular buffer: 40 ms.

2. 20th Request:

- Sent at 1000 ms on the client (passes the client's rate limit).
- Server timestamp: Received at 1036 ms (with 36 ms latency).
- Recorded on the server's circular buffer: 1036 ms.

3. Circular Buffer Check on the Server:

Difference between the 1st and 20th request:
 1036 ms - 40 ms = 996 ms, which is less than 1000 ms.

Since the **server's calculation** shows the difference as **996 ms** instead of **1000 ms**, the **server rejects the request with a 429 Error**. To prevent **429 errors**, it's important to account for **latency variability** between requests. Adding a **latency buffer** (e.g., an extra 50-100 ms delay) on the client side can help ensure that **timestamp discrepancies** don't cause requests to fail the **server's rate limiter**.

Improved Version: Adaptive Buffering

```
class RateLimiter:
   def __init__(self, per_second_rate, min_duration_ms_between_requests):
        self.__per_second_rate = per_second_rate
        self.__min_duration_ms_between_requests =
min_duration_ms_between_requests
        self.__request_times = [0] * per_second_rate
        self.__curr_idx = 0
        self.__latency_window = deque(maxlen=100) # record of the last
100 latencies
        self. buffer = 40 # initial buffer (ms)
        self.__min_buffer = 30 # min buffer (ms)
        self.__max_buffer = 50 # max buffer (ms)
   def update_buffer(self):
        # calculate a moving average of the recent latencies
        if len(self.__latency_window) > 0:
            avg_latency = sum(self.__latency_window) /
len(self.__latency_window)
            # adjust buffer based on average latency
            self.__buffer = min(self.__max_buffer, max(self.__min_buffer,
int(avg_latency * 1.1)))
   def record_latency(self, latency):
        self.__latency_window.append(latency)
        self.update_buffer()
```

```
@contextlib.asynccontextmanager
    async def acquire(self, timeout ms=0):
        enter ms = timestamp ms()
        buffer = self.__buffer
        initial buffer = self. min duration ms between requests *
self.__per_second_rate
        while True:
            now = timestamp ms()
            if now - enter_ms > timeout_ms > 0:
                raise RateLimiterTimeout()
            if now - self.__request_times[self.__curr_idx] <=</pre>
initial buffer + buffer:
                # sleep the exact remaining time to the next second
                sleep_time = (initial_buffer + buffer - (now -
self. request times[self. curr idx])) / 1000
                await asyncio.sleep(sleep time)
                continue
            break
        self.__request_times[self.__curr_idx] = now
        self.__curr_idx = (self.__curr_idx + 1) % self.__per_second_rate
        yield self
```

Explanation of Adaptive Buffering

This enhanced version introduces **adaptive buffering**, which fine-tunes the buffer size based on real-time **latency trends**. This ensures the system stays compliant with the server's rate limits while minimizing unnecessary delays and maximizing throughput.

Key Enhancements

1. Latency Window for Adaptive Buffering:

- A deque stores the last 100 latencies to track trends in request delays.
- The **buffer size** is updated dynamically using a **moving average** of recorded latencies.
- The formula used is:

```
buffer = min(self.__max_buffer, max(self.__min_buffer,
int(avg_latency * 1.1)))
```

• This ensures the buffer size adjusts proportionally to observed latency while staying within a safe range (30 ms to 50 ms based on server's MAX_LATENCY_MS).

2. Dynamic Sleep Calculation:

 If the time between requests violates the rate limit, the exact remaining time needed to comply is calculated:

```
sleep_time = (initial_buffer + buffer - (now -
self.__request_times[self.__curr_idx])) / 1000
```

• This approach avoids **excessive sleeping** and ensures the request rate is optimized.

Why Adaptive Buffering Matters

Even though the **maximum server-side latency** is known (e.g., MAX_LATENCY_MS = 50 ms), real-world latency often varies. **Hardcoding a fixed buffer** is suboptimal, as it can either:

- Undershoot latency: Leading to premature requests and 429 errors.
- Overshoot latency: Reducing throughput by waiting longer than necessary.

With adaptive buffering, the system continuously learns from recent trends and dynamically adjusts the buffer to balance performance and compliance.

[!NOTE] In some cases of extreme network fluctuation, adpative buffering may still result in 429 Error. Hence a further improvement we can consider is to implement exponential backoffs mechanism to further manage retries.

Conclusion

- 1. Remove Fix Interval Checks and only keeping the Circular Buffer Checks
- 2. Account for the server's latency and dynamically adjusts it for additional buffer when sending requests, to prevent rate limiting errors (429).

Improving Queue System

Current Implementation

```
async def exchange_facing_worker(url: str, api_key: str, queue: Queue,
logger: logging.Logger):
    rate_limiter = RateLimiter(PER_SEC_RATE, DURATION_MS_BETWEEN_REQUESTS)
    async with aiohttp.ClientSession() as session:
        while True:
            request: Request = await queue.get()
                  remaining_ttl = REQUEST_TTL_MS - (timestamp_ms() -
request.create_time)

    if remaining_ttl <= 0:
                  logger.warning(f"Ignoring request {request.req_id} due to
expired TTL.")

        continue
        try:
                  async with rate_limiter.acquire(timeout_ms=remaining_ttl):</pre>
```

```
async with async_timeout.timeout(1.0):
                        nonce = timestamp ms()
                        data = {'api_key': api_key, 'nonce': nonce,
'req_id': request.req_id}
                        async with session.request('GET', url, data=data)
as resp:
                            json = await resp.json()
                            if json['status'] == 'OK':
                                logger.info(f"API response: status
{resp.status}, resp {json}")
                            else:
                                logger.warning(f"API response: status
{resp.status}, resp {json}")
            except RateLimiterTimeout:
                logger.warning(f"Ignoring request {request.reg id} due to
TTL in rate limiter.")
```

Issue

The current implementation allows **more requests to be generated** than the client can process. This leads to **expired TTLs** for unprocessed requests in the queue. When the **remaining_ttl** drops below 0 or when there is a **request timeout**, the request is **dropped**. While this prevents the queue from becoming clogged, it also results in **wasted resources** and **lost requests** that could have been retried.

Solution: Queue Manager with Dead Letter Queue (DLQ)

To improve request management, we introduce a **Queue Manager** that utilizes:

- 1. Main Queue: Processes requests under normal operation.
- 2. **Dead Letter Queue (DLQ):** Stores failed or timed-out requests for **retry** or further processing. This helps ensure that no valid request is wasted, even if it initially fails or exceeds its TTL.

This strategy improves **resilience** by providing better queue state management. Requests are **re-prioritized** from the DLQ, minimizing dropped requests and ensuring all requests receive multiple attempts before being discarded. In the event that the request hits the max retry limit we specified, we will store the req_id in self_graveyard for manual processing, debugging purposes and to prevent sending redundant request, as we can assume that these requests are invalid.

How Do Retry Requests Get Prioritized Over New Requests?

• Cooldown Window for New Requests:

When new requests are generated, a **short cooldown period** is introduced between them. This window allows the **Queue Manager** to **re-slot retry requests** into the queue before more new requests are created.

• Not a Strict Priority Queue:

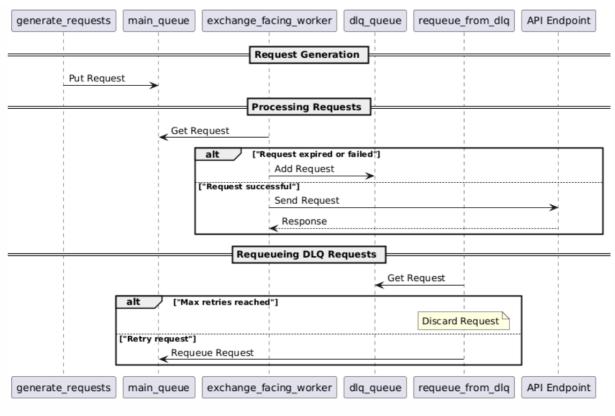
While this is not a strict priority queue (where retry requests are inserted at the front), the Queue

Manager ensures that retry requests are slotted into the queue as soon as possible, taking advantage of gaps in new request generation.

This design ensures that **retry requests** are **handled promptly** without requiring a complex priority queue, optimizing for both **simplicity** and **timeliness**.

Lifecycle with Queue Manager

This diagram illustrates the interaction between the queue (in Queue Manager), DLQ (in Queue Manager) and workers.



Monitoring the Queue

Using the below metrics we are able to get the current state of the DLQ and Main Queue to check if there is an existing bottleneck in processing the requests inside the queue, allowing us to check the risk of the requests inside the main queue timing out due to expired TTL.

- Queue Monitoring: Track main queue and DLQ sizes, processing rates, and retry statistics.
- Queue Sizes: Monitor and log the sizes of the main queue and DLQ at regular intervals.
- **Graveyard Metrics**: Monitor the number of requests that exceed the maximum retry limit and are moved to the graveyard.

```
--- Accumulated Benchmark Metrics ---
2024-10-20 16:40:49,644 - stats - INFO - Elapsed Time: 5.00 seconds
2024-10-20 16:40:49,644 - stats - INFO - Total Successful Requests: 414
```

```
2024-10-20 16:40:49,644 - stats - INFO - Total Failed Requests: 0
2024-10-20 16:40:49,644 - stats - INFO - Total Throughput: 82.78 req/sec
2024-10-20 16:40:49,644 - stats - INFO - Average Latency: 207.68 ms
2024-10-20 16:40:49,644 - stats - INFO - Queue Sizes - Main: 26, DLQ: 0,
Gravevard: 0
2024-10-20 16:40:49,644 - stats - INFO - Average Queue Sizes - Main:
13.00, DLQ: 0.00
2024-10-20 16:40:54,645 - stats - INFO -
--- Accumulated Benchmark Metrics ---
2024-10-20 16:40:54,645 - stats - INFO - Elapsed Time: 10.00 seconds
2024-10-20 16:40:54,645 - stats - INFO - Total Successful Requests: 834
2024-10-20 16:40:54,645 - stats - INFO - Total Failed Requests: 0
2024-10-20 16:40:54,645 - stats - INFO - Total Throughput: 83.38 req/sec
2024-10-20 16:40:54,645 - stats - INFO - Average Latency: 322.23 ms
2024-10-20 16:40:54,645 - stats - INFO - Queue Sizes - Main: 40, DLQ: 0,
Graveyard: 0
2024-10-20 16:40:54,645 - stats - INFO - Average Queue Sizes - Main:
22.00, DLQ: 0.00
2024-10-20 16:40:59,646 - stats - INFO -
--- Accumulated Benchmark Metrics ---
2024-10-20 16:40:59,646 - stats - INFO - Elapsed Time: 15.00 seconds
2024-10-20 16:40:59,646 - stats - INFO - Total Successful Requests: 1262
2024-10-20 16:40:59,646 - stats - INFO - Total Failed Requests: 0
2024-10-20 16:40:59,646 - stats - INFO - Total Throughput: 84.12 req/sec
2024-10-20 16:40:59,646 - stats - INFO - Average Latency: 430.42 ms
2024-10-20 16:40:59,646 - stats - INFO - Queue Sizes - Main: 97, DLQ: 0,
Graveyard: 0
2024-10-20 16:40:59,646 - stats - INFO - Average Queue Sizes - Main:
40.75, DLQ: 0.00
2024-10-20 16:41:04,647 - stats - INFO -
```

Addressing the Root Cause: Bloating of the Main Queue

- Currently, the rate of generating requests exceeds the rate of dequeuing requests, causing
 queue bloating and leading to TTL expirations. We are limited to only 5 API keys, each associated
 with a worker consuming requests from the queue. This restriction results in a processing rate that is
 too slow to handle the incoming traffic effectively, causing requests to expire before being
 processed.
- The situation is further exacerbated by latencies on both the client and server sides due to rate
 limiting constraints, which slow down the consumption of requests from the queue. While the Dead
 Letter Queue (DLQ) allows failed requests to be retried, it does not address the root issue: the
 imbalance between the rate of request generation and processing.

To mitigate this problem, we can consider several workarounds:

1. Implementing Backpressure on generate_requests (): By controlling the rate at which requests are generated based on the current state of the queue, we can prevent the system from being overwhelmed. This approach assumes that we are able to modify the generate_requests () function.

```
# Apply backpressure if the queue is full
if queue.qsize() >= max_queue_size:
    print(f"Queue is full ({queue.qsize()}), pausing request
generation.")
    await asyncio.sleep(0.5) # Wait before checking again
    continue
```

2. **Implementing Multithreading:** By implementing multithreading, we are able to better manage the rate of consumption of the request, resolving the issue of the expired requests TTL.

Exploring Multithreading

Rationale

To manage request processing rate and reduce TTL expirations, we introduced multithreading. Multiple workers accessing the same queue increase the rate at which requests are dequeued, preventing the main queue from overloading.

[!Note] See folder for implementation details.

Changes to the Current Code

1. Implement Multithreading Processes:

- Request Generator: Run generate_requests() in a separate thread (assuming modifications are allowed).
- Metrics Printing: Move metrics logging to a dedicated thread.
- Exchange Facing Workers: Run multiple threads for workers handling API requests.

2. Implement Locking on Shared Resources:

- Queue Manager: Use locks when accessing shared data structures like the graveyard to ensure thread safety.
- **Thread-Safe Rate Limiter:** Modify the rate limiter to be thread-safe if multiple threads share API keys.

[!Note] We use timestamps for nonces, as the GIL (Global Interpreter Lock) ensures only one thread executes Python code at a time, minimizing the chance of generating the same timestamp. However, under high request volumes, there's still a risk of **nonce collisions** if multiple requests fall within the same millisecond. To avoid future issues, consider combining timestamps with a **thread-local counter** or **UUID** for guaranteed uniqueness.

Comparison Between Asynchronous and Multithreading Client

Below is a comparison of our current implementations using asynchronous programming and multithreading.

Baseline of Comparison

Asynchronous:

- 1. 5 Coroutines (one for each API key).
- 2. 1 Coroutine to generate requests.
- 3. 1 Coroutine for the Queue Manager to requeue from the DLQ.
- 4. 2 Coroutines for monitoring and benchmarking.

Multithreading:

- 1. 5 Threads (one for each API key).
- 2. 1 Thread to generate requests.
- 3. 1 Thread for the Queue Manager to requeue from the DLQ.
- 4. 2 Threads for monitoring and benchmarking.

Observation

Asynchronous Client (~84 TPS):

- Achieves higher throughput due to efficient handling of concurrent I/O-bound operations.
- Issue: Requests accumulate in the queue, leading to TTL expirations.

```
ASYNCHRONOUS CLIENT STATS
Elapsed Time: 25.01 seconds
Total Successful Requests: 2102
Total Failed Requests: 13
Total Throughput: 84.06 req/sec
Average Latency: 769.45 ms
Queue Sizes - Main: 136, DLQ: 0, Graveyard: 13
Average Queue Sizes - Main: 81.67, DLQ: 0.00
```

Multithreading Client (~77 TPS):

- Slightly lower throughput due to thread overhead.
- Advantage: Multiple threads dequeue requests simultaneously, keeping the queue size small and reducing TTL expirations.

```
MULTITHREADING CLIENT STATS
Elapsed Time: 25.00 seconds
Total Successful Requests: 1939
Total Failed Requests: 0
Total Throughput: 77.48 req/sec
Average Latency: 75.12 ms
Queue Sizes - Main: 1, DLQ: 0, Graveyard: 0
Average Queue Sizes - Main: 1.00, DLQ: 0.00
```

CPU Utilization

[!NOTE] See here for utilization graph for both clients

• Asynchronous client is slightly less CPU intensive than Multithreading client.

Possible Explanations

CPU Utilization and GIL Impact:

The Global Interpreter Lock (GIL) in Python limits the performance of multi-threaded programs by
allowing only one thread to execute Python bytecode at a time, even on multi-core CPUs. This makes
the Multithreading client more CPU-intensive since threads contend for the GIL. In contrast,
Asynchronous client bypasses the GIL to some extent by focusing on I/O-bound tasks with a single
event loop, resulting in slightly lower CPU usage.

Throughput (TPS) and GIL Constraints:

The GIL can restrict the Multithreading client;s throughput, contributing to its lower 77.48 TPS, as only one thread can execute Python code at a time. However, because threading is more effective at handling blocking I/O, it maintains stable performance. Asynchronous, despite achieving higher 84 TPS, leverages non-blocking I/O and cooperative multitasking, allowing it to manage more requests without being restricted by the GIL.

Graveyard and Queue Size Management:

While Asynchronous excels in throughput, the accumulation of 136 requests in its queue indicates
that the event loop struggles under heavy load, leading to 13 TTL expirations. On the other hand,
the Asynchronous client avoids queue buildup, keeping it at 1 request, with no expirations, as
multiple threads can process requests independently. However, the effectiveness of threading is
limited by the GIL, which can reduce its scalability for CPU-bound workloads.

Overview

Modifying the Original Client for Improved Performance

To enhance throughput and optimize request management, we made the following modifications:

1. Removed Redundant Waits:

- Eliminated the Fixed Interval Check in the rate limiter to reduce unnecessary pauses and context switching.
- Relied on the Circular Buffer Check for rate limiting, improving throughput.

2. Introduced a Queue Manager:

- Added a Queue Manager to handle the Main Queue and DLQ.
- o Implemented Graveyard Tracking for requests exceeding the maximum retry count.
- Prioritized DLQ requests over new ones to ensure time-sensitive retries are processed promptly.

Using the multi-threaded client can offer greater reliability by ensuring guaranteed request delivery at the expense of throughput (~77 TPS). Its preemptive multitasking model ensures that requests are processed without delay, even when individual threads are blocked on I/O. This makes threading the better option when consistent, on-time delivery is critical, as it avoids TTL expirations through concurrent dequeuing and independent thread execution.

However, threading incurs additional overhead from context switching and synchronization mechanisms, leading to higher memory usage. This can limit performance under heavy loads, especially in high-frequency I/O-bound tasks where Asynchronous excels. Asynchronous's non-blocking, single-threaded event loop minimizes these costs and achieves higher throughput (~84 TPS) by avoiding thread management overhead.

With **adaptive rate control** to regulate request generation, the **Asynchronous client** can effectively prevent **TTL expirations**, making it the **superior long-term solution**. Its ability to **scale efficiently** under high concurrency, while consuming **less memory**, ensures that it will outperform threading as workloads grow.

In summary, threading currently provides guaranteed reliability, but with proper request regulation, Asynchronous becomes the optimal solution due to its speed, efficiency, and scalability.

Link to github repo here