DESIGN AND ANALYSIS OF ALGORITHMS LAB3

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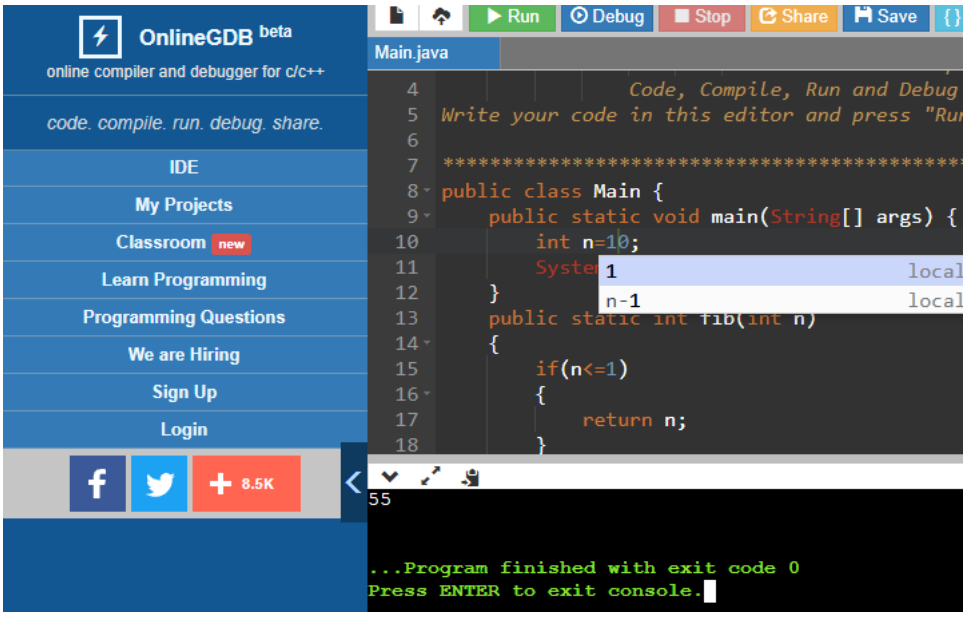
Naïve Fibonacci series:

Code:

```
public class Main {


   public static void main(String[] args) {
      int n=300;
      System.out.println(+ fib(n));
   }
   public static int fib(int n)
   {
      if(n<=1)
      {
         return n;
      }
      else
      {
         return fib(n-1)+fib(n-2);
      }
   }


}
```

Output:

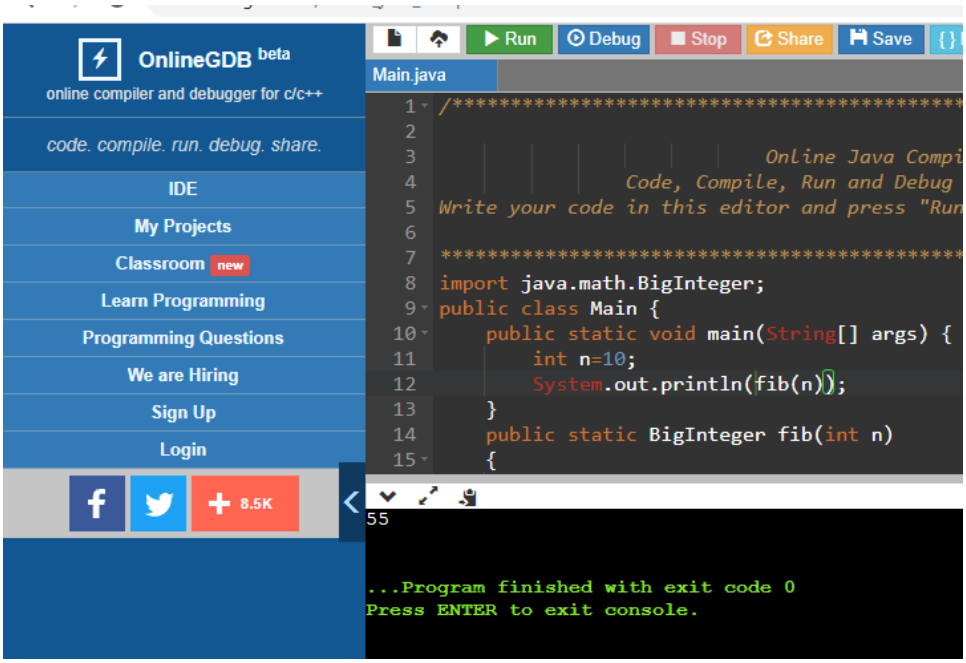


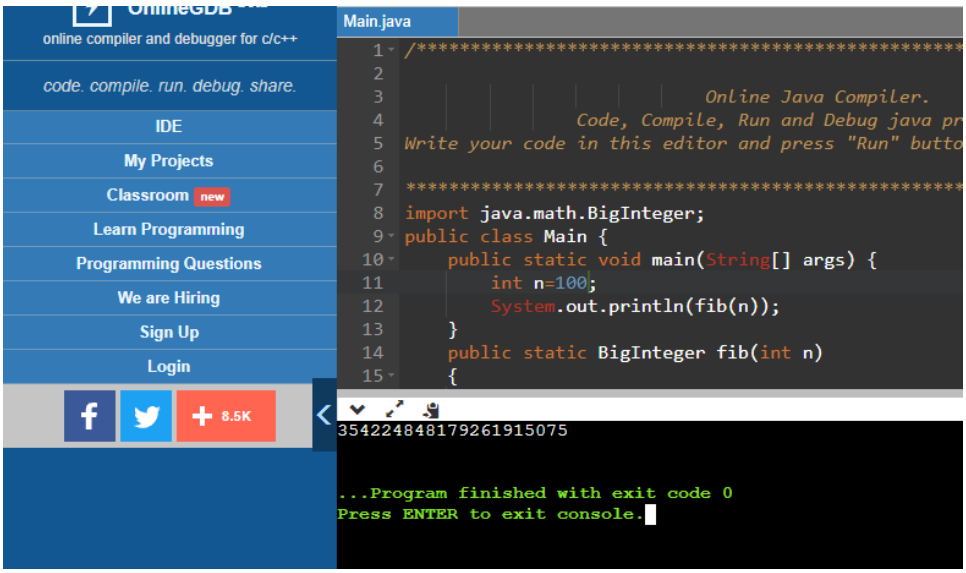
Effective Fibonacci series:

```
import java.math.BigInteger;

public class Main {

    public static void main(String[] args) {

        int n=100;

        System.out.println(fib(n));

    }
```

```
    public static BigInteger fib(int n)
    {
        BigInteger a=BigInteger.valueOf(0);
        BigInteger b=BigInteger.valueOf(1);
        BigInteger c=BigInteger.valueOf(1);
        for(int i=2;i<=n;i++)
        {
            c=a.add(b);
            a=b;
            b=c;
        }
       return (b);
    }

}
```

Output:





Asymptotic Analysis:

Naive algorithm:-

```
If n<=1
  return n
else
  return FibRecurs(n-1) + FibRecurs(n-2)
```

$$T(n) = \begin{cases} 2 & \text{if } n \le 1 \\ T(n-1) + T(n-2) + 3 & \text{else} \end{cases}$$

$$T(n) = T(n-1) + T(n-2)$$

$$x^n = x^{n-1} + x^{n-2}$$

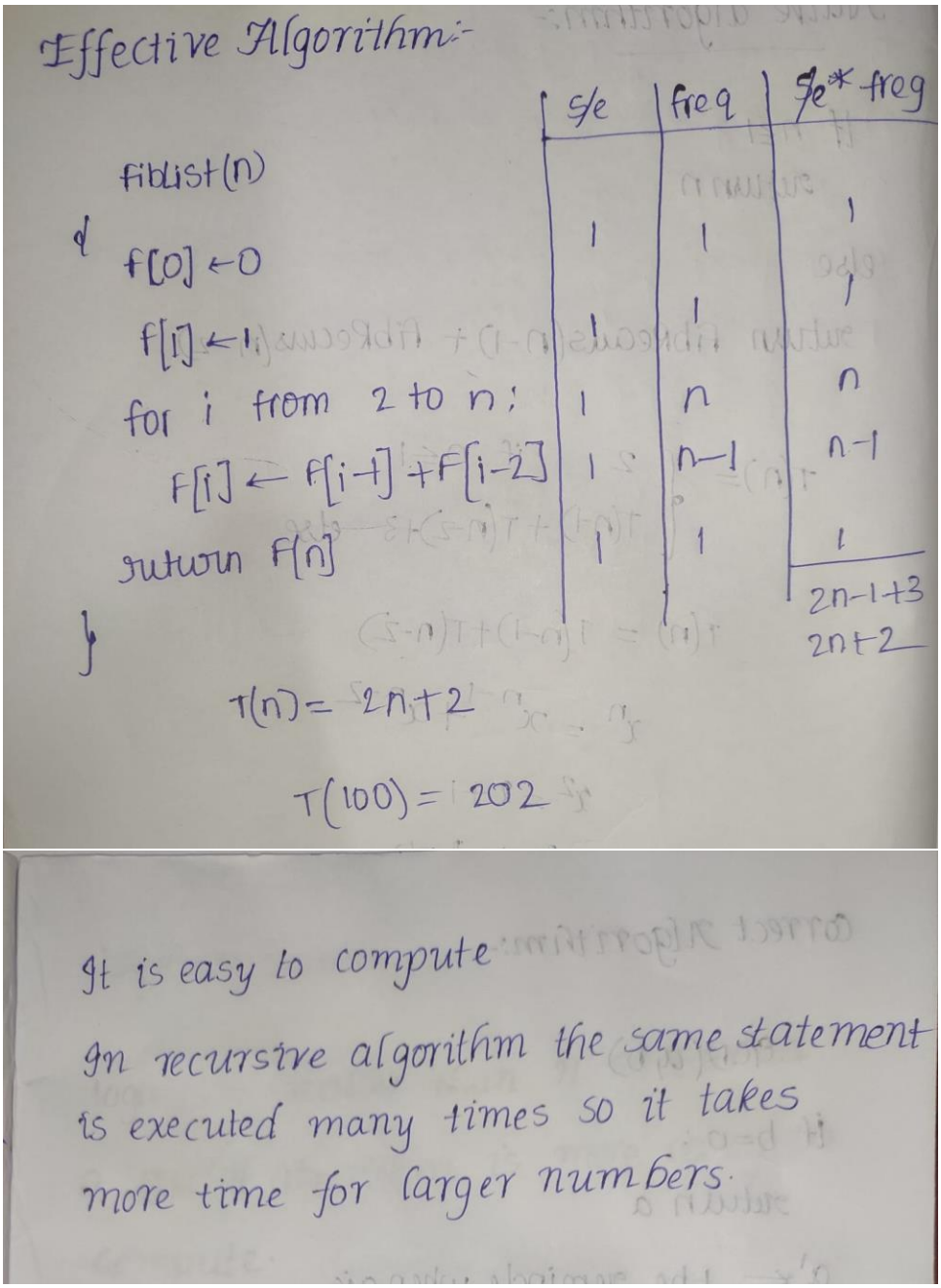
$$x^2 = x + 1$$

$$x^2 - x - 1 = 0$$

$$x = \left(\frac{1+\text{root}(5)}{2}\right)^n \text{ or } \frac{(1-\text{root}(5))^n}{2^n}$$

Time complexity = $\frac{1+(\text{root}(5))^n}{2^n}$

$$T(100) \simeq 1.77 \cdot 10^{21}$$

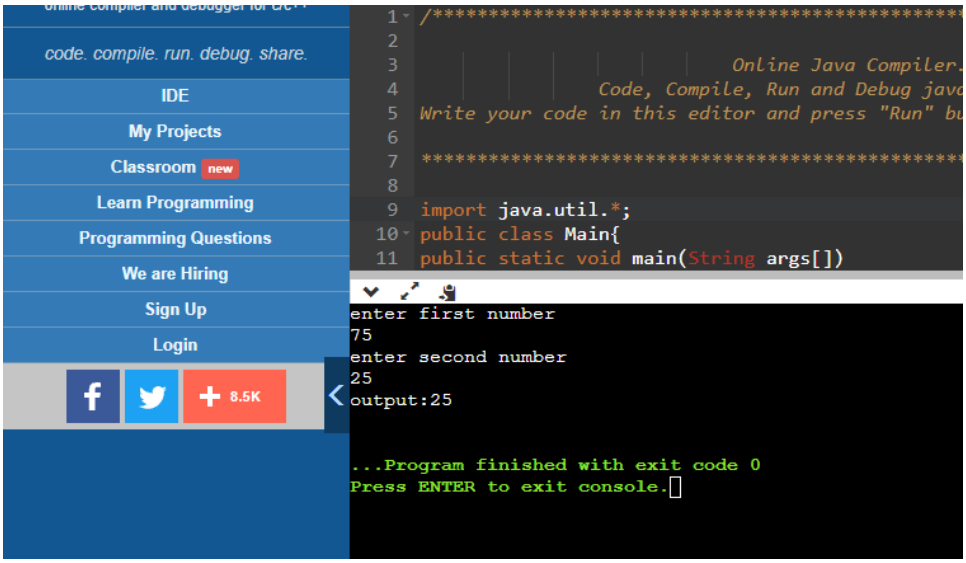


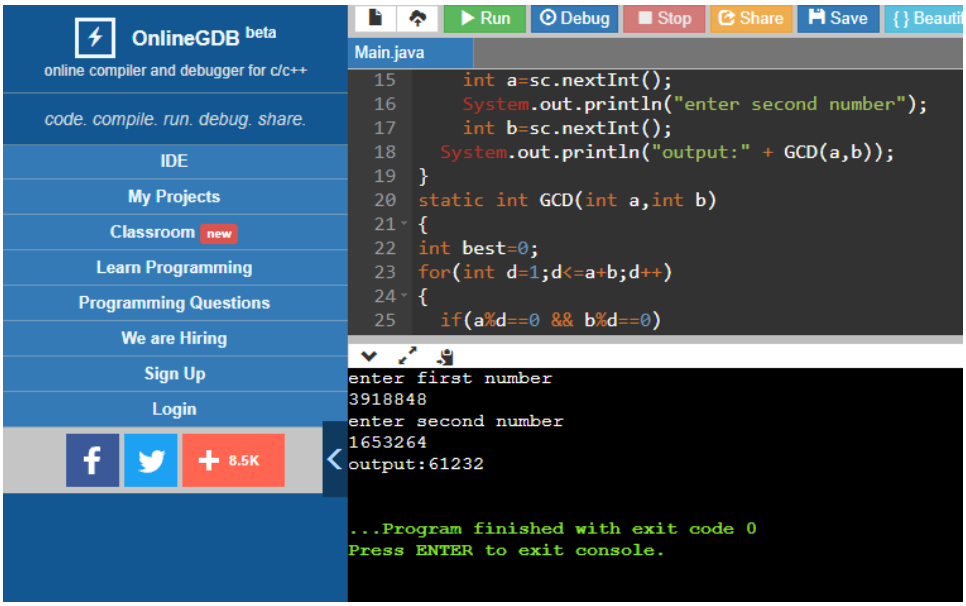
Naïve GCD:

```
import java.util.*;
public class Main{
public static void main(String args[])
{
```

```
Scanner sc=new Scanner(System.in);
System.out.println("enter first number");
int a=sc.nextInt();
System.out.println("enter second number");
int b=sc.nextInt();
  System.out.println("output:" + GCD(a,b));
}
static int GCD(int a,int b)
{
int best=0;
for(int d=1;d<=a+b;d++)
{
  if(a%d==0 && b%d==0)
  {
  best=d;
  }
}return best;
}
}
```

Output:



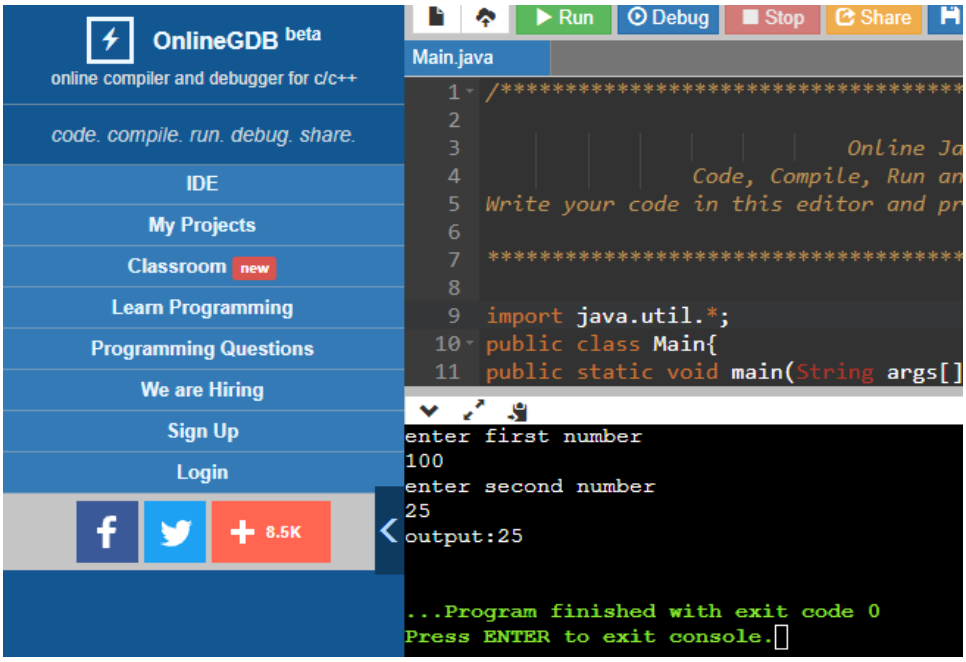


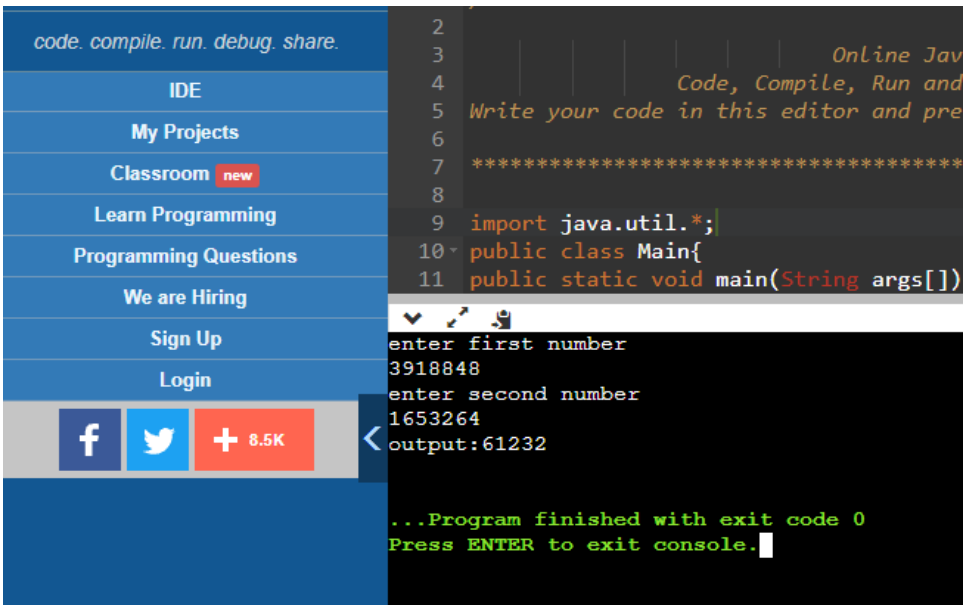
Euclidian GCD:

```
import java.util.*;

public class Main{

public static void main(String args[])

{

Scanner sc=new Scanner(System.in);
```

```
System.out.println("enter first number");
int a=sc.nextInt();
System.out.println("enter second number");
int b=sc.nextInt();
  System.out.println("output:" + euclidGCD(a,b));
}
static int euclidGCD(int a,int b)
{
if(b==0)
{
return a;
}
int a1;
a1=a%b;
return euclidGCD(b,a1);
}
}
```

Output:





Asymptotic Analysis:

Naive GCD:-

```
GCD (a,b)
  best ← 0
  for d from 1 to a+b:
    if d/a and d/b:
      best ← d
  return best
```

	s/e	freq	s/e * freq
	1	1	1
	1	a+b+1	a+b+1
	1	1	1
	1	1	1
	⋮	⋮	⋮
	1	1	1

$$T(n) = a+b+1+1+1$$
$$\leq a+b+4$$
$$O(n)$$

correct Algorithm:-

```
GCD(a, b)
if b=0:
  return a
a' ← the reminde when a is
        divided by b
return GCD(b, a')
```

$T(n) = T(n/2) + 2$

$= T(n/4) + 2 + 2$

$\vdots$

$T(n) = T(n/2^k) + 2 + 2 + \cdots 2$

$= T(1) + 2 \cdot k$

$= 2 \log_2 n + 1$

$O(\log_2 n)$

$\log n$ is smaller than n so correct

euclid algorithm is more easy to compute.