

# CS 40 Homework 4

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## 1 Types

- a.  $f : \mathbb{R}^+ \rightarrow \mathbb{R}$
- b.  $f : \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}^+$
- c.  $f : \mathbb{R} \rightarrow [-1, 1]$
- d.  $f : \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{B}$
- e.  $f : \mathbb{Z} \rightarrow \wp(\mathbb{Z}^+)$
- f.  $f : \mathbb{N} \rightarrow \mathbb{Z}^+$

## 2 Relations

Goal: Prove  $\sim$  is an equivalence relation.

Method: Direct proof

- (1) By definition, if  $\sim$  is an equivalence relation, it satisfies the properties of being reflexive, symmetric, and transitive.
- (2) Given:  $a \sim b \leftrightarrow f(a) = f(b)$
- (3) Let  $\exists(a, b, c) \in A(a \neq b \wedge b \neq c \wedge a \neq c)$
- (4) Separating (2) into two implication statements, we know the following: if  $a \sim b$ , then  $f(a) = f(b)$ . Thus, if  $a \sim a$ , then  $f(a) = f(a)$ . This implies that  $\sim$  is reflexive.
- (5) We also know that if  $b \sim a$ , then  $f(b) = f(a)$ . Since the  $=$  operator is symmetric, this means that  $f(a) = f(b)$ . Hence,  $\sim$  is symmetric.
- (6) Finally, if  $a \sim b$  and  $b \sim c$ , then  $f(a) = f(b)$  and  $f(b) = f(c)$ . The  $=$  operator is also transitive; therefore, we can deduce that  $f(a) = f(c)$ . Thus,  $\sim$  is transitive.

By the definition of an equivalence relation, since  $\sim$  is reflexive, symmetric, and transitive, it is therefore an equivalence relation.

### 3 Cancellation

Goal: Prove the below inference rules are valid iff  $f$  is a bijection.

$$\begin{array}{l} f \circ g = f \circ h \quad \vdash \quad g = h \\ g \circ f = h \circ f \quad \vdash \quad g = h \end{array}$$

Method: Proof by cases

We can conduct a proof by cases with 4 cases.

Case (i):  $f$  is both injective and surjective. If  $f$  is both injective and surjective, then  $f^{-1}$  must exist. If we apply  $f^{-1}$  to both sides of the first inference rule, we obtain  $f^{-1}(f(g(x))) = f^{-1}(f(h(x)))$ . This provides the identity  $\iota$ . Therefore,  $f^{-1}$  and  $f$  cancel and we are left with  $g(x) = h(x)$ . The same case applies to the second inference rule; however, the inverse is still applied to  $f$  and we obtain  $\iota$  again.

Case (ii):  $f$  is surjective but not injective. We can define  $f : \mathbb{R} \rightarrow \mathbb{R}^+ \cup \{0\} = x^2$ .  $f$  is therefore surjective since every positive real number and zero can be mapped to by a real number. However, the function isn't injective since not every element in the domain has a distinct element in the codomain. Furthermore, we define  $g(x) = x$  and  $h(x) = |x|$ .  $f(g(x))$  and  $f(h(x))$  will always produce the same output; however,  $g(x) \neq h(x)$ .

Case (iii):  $f$  is injective but not surjective. Since  $f$  isn't surjective, there exists at least one element in the codomain of  $f$  that isn't mapped to by an element in its domain. Thus,  $g$  and  $h$  aren't bounded by that element, so  $g$  and  $h$  can diverge from there.

Case (iv):  $f$  is neither injective nor surjective. Since cases (ii) and (iii) don't hold, the combination of them will not hold either. Therefore, in this case, the two inference rules are invalid.

### 4 Counting

(1) Firstly, we want to prove that  $|\mathbb{Z}| = |\mathbb{N}|$ .

(2) We can prove this by defining a function  $f : \mathbb{Z} \rightarrow \mathbb{N}$ :

$$f(x) = \begin{cases} 2x - 1, & x > 0 \\ -2x, & x \leq 0 \end{cases} \quad (1)$$

Since  $f(x)$  is both injective and surjective, it is bijective. Through the property of bijectivity,  $|\mathbb{Z}| = |\mathbb{N}|$ . Therefore, the set of all integers is countably infinite.

- (3) Next, we can define a set  $S$  as the set of all square numbers. We can also define a function  $g : \mathbb{N} \rightarrow S$  where  $g(x) = x^2$
- (4)  $g(x)$  is a bijective function since there is a unique element in the codomain for each element in the domain and every element in the codomain can be mapped to by some element in the domain.
- (5) Therefore,  $|S| = |\mathbb{N}|$ .
- (6) By transitivity from steps (2) and (5), we can therefore state that  $|\mathbb{Z}| = |S|$ .

Thus, we can conclude that the cardinality of the set of integers is equal to the cardinality of the set of square numbers.

## 5 Cantor's Theorem

- (1) Assume for the sake of contradiction that  $f : S \rightarrow \wp(S)$  is a surjective function.
- (2) Since  $f$  is surjective,  $\exists x \in S (f(x) = D)$ .
- (3) We also know from the property of surjection that  $|S| \geq |\wp(S)|$ .
- (4) Assume that  $x \in D$ . This means that  $x \in f(x)$  which means that  $x \notin D$  from the definition of the construction of  $D$  through its predicate. We've reached a contradiction that  $x$  is an element of  $D$  yet not an element too.
- (5) Let's assume the opposite then that  $x \notin D$ . This means that  $x \notin f(x)$  which means that  $x \in D$ . We've reached the same contradiction where  $x$  is simultaneously in  $D$  and not in  $D$ .
- (6) Therefore, we've shown that there is no  $x \in S$  such that  $f(x) = D \subseteq S$ .  $f$  cannot be surjective then.
- (7) Our initial assumption of  $f$  being surjective was therefore invalid.

Thus, by proof by contradiction,  $|S| < |\wp(S)|$ .