

Termination

ARS ($A \rightarrow$) model of computation

→ a one-step computation

flexible: "one step" big or small

- one operation on the machine
- the body of a whole loop
- ...

$A \xrightarrow{\varphi} N$ measure function

$a \rightarrow b \Rightarrow \varphi(a) > \varphi(b)$

Exles of programs and used measure functions to prove termination.

Let us turn to a theoretical question.

Theory is practice on a higher level.

Practice is application of a method.

Theory is development of a method.

Practice: proving termination of your program
if \exists measure function \Rightarrow terminating singular

Theory: What is a method to prove termination for programs in general.

practical

Question:

If program is terminating
 \Rightarrow F measure function ?

Theoretical question for a programmer
practical question for a developer
of methods to verify termination

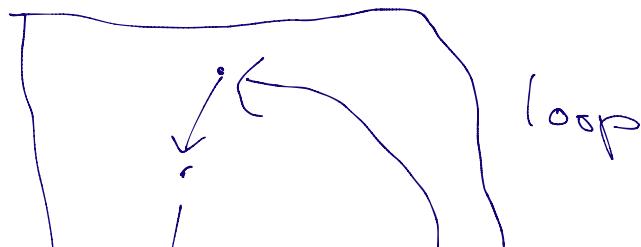
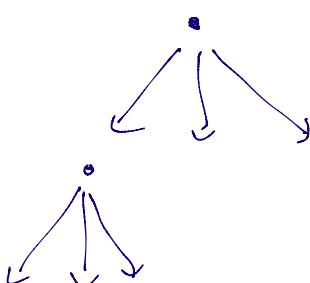
Thm: If terminating & finitely branching

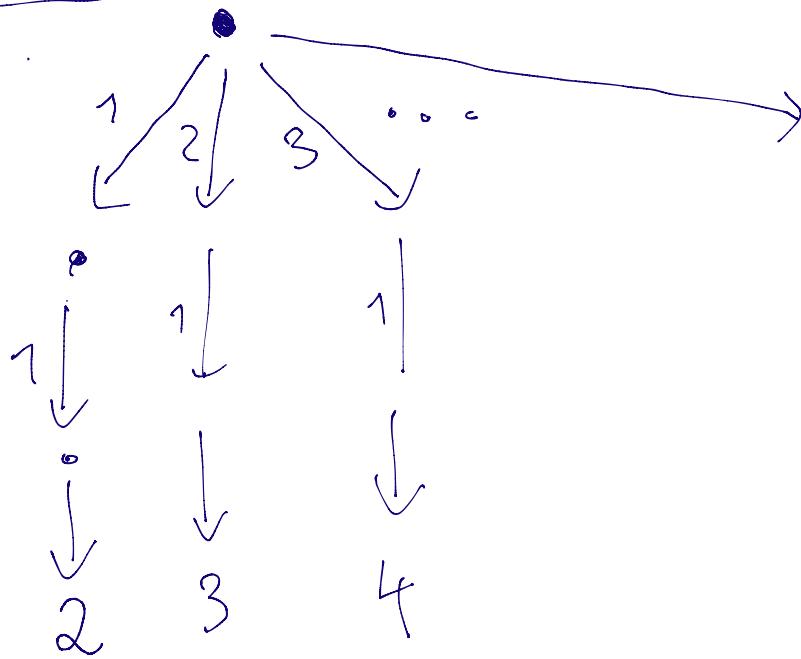
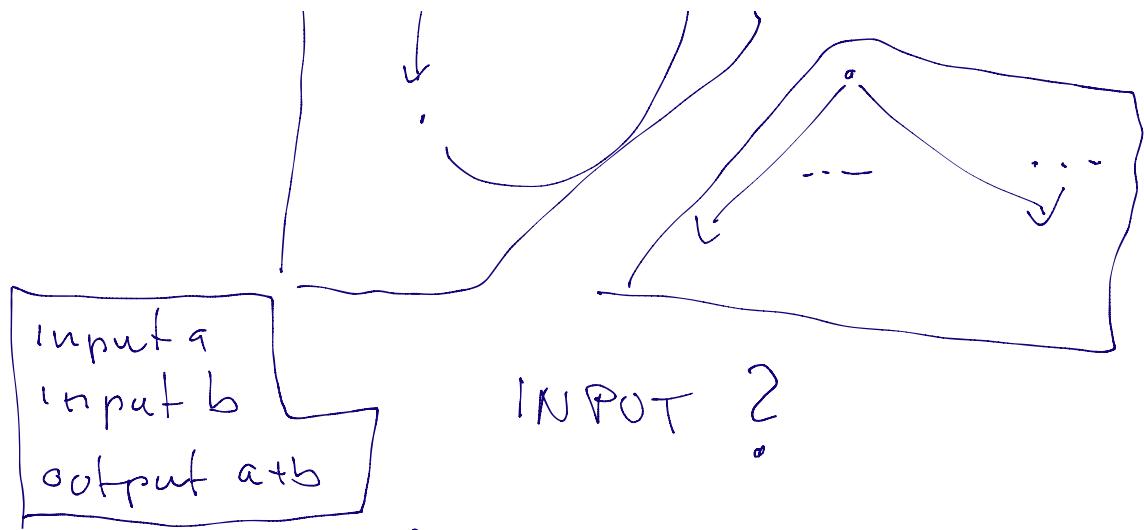
Then F measure function

Q

finitely branching

all $a \in A$ there are
only finitely b
such that $a \rightarrow b$





Claim: Terminat. & fin. br. \Rightarrow J meas. fun.

Assume: (A, \rightarrow) is termination
 (A, \rightarrow) is fin. - br.

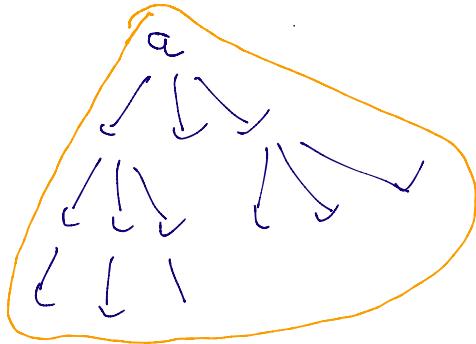
Show: J measure function
 $\varphi : A \rightarrow X$

have an idea is needed

| 15 |

The number
of elements
of S

$$\varphi(a) \stackrel{\text{def}}{=} |\{b \in A \mid a \xrightarrow{*} b\}|$$



Show: φ is a measure function

Choose any a, b $a \rightarrow b \Rightarrow \varphi(a) > \varphi(b)$

Assume $a \rightarrow b$

Show $\varphi(a) > \varphi(b)$

$$(1) |\{c \mid a \xrightarrow{*} c\}| > |\{d \mid b \xrightarrow{*} d\}|$$

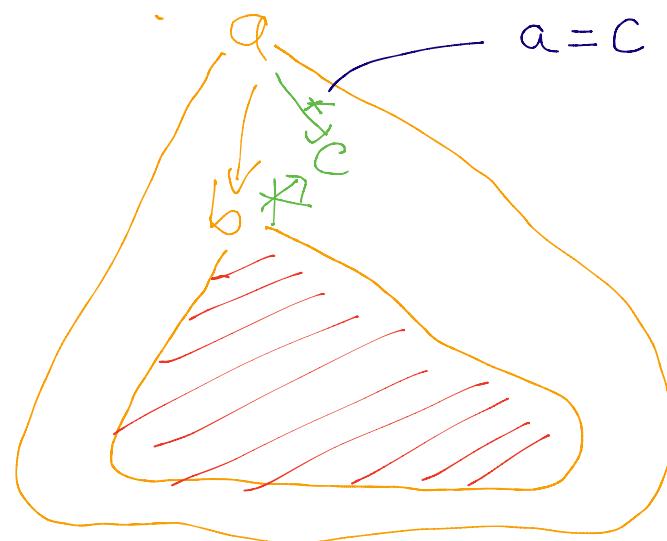
$$(1) \{c \mid a \xrightarrow{*} c\} \supseteq \{d \mid b \xrightarrow{*} d\}$$

$$(2) \exists c \text{ st. } a \xrightarrow{*} c, b \not\xrightarrow{*} c$$

Show(1) Use def $\xrightarrow{*}$

refl & transitive closure

of \rightarrow



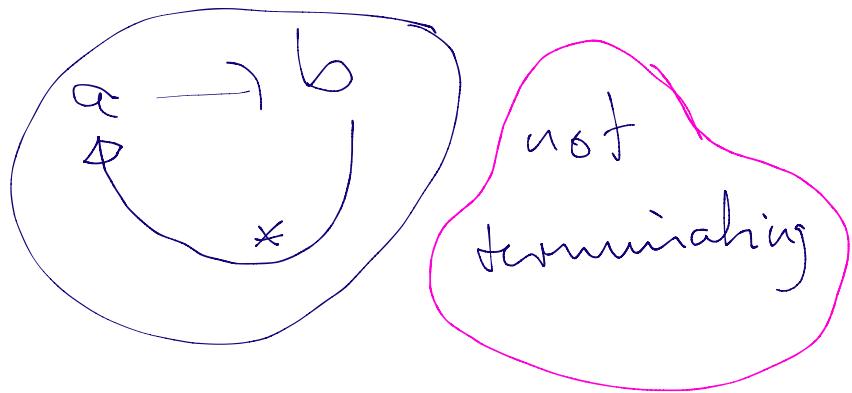
QED for (1)

Show(2) Idea needed

what could be c?

$$c = a$$

remains: $b \xrightarrow{*} a$



Useful observation

$$\rightarrow \subseteq \xrightarrow{*}$$

General remark on proofs.

Any proof of

$$A \Rightarrow B$$

has the structure

Assume A
⑥⁶
⑥⁶
⑧

|
B

Any proof of
not A

has the structure of

| Assume A
|
|
| contradiction