

PCG for the RKHS mode- k CP subproblem with missing entries

We solve for $W \in \mathbb{R}^{n \times r}$ in

$$[(Z \otimes K)^\top SS^\top (Z \otimes K) + \lambda(I_r \otimes K)] \text{vec}(W) = (I_r \otimes K) \text{vec}(B),$$

with $K \in \mathbb{R}^{n \times n}$ a symmetric kernel matrix, $Z \in \mathbb{R}^{M \times r}$ the Khatri–Rao product of the other CP factors, and $S \in \mathbb{R}^{N \times q}$ selecting the q observed entries of the mode- k unfolding ($N = nM$). We assume $n, r < q \ll N$.

1. Why (P)CG applies

Let $P \equiv SS^\top$ (a diagonal mask). Define

$$A = (Z \otimes K)^\top P(Z \otimes K) + \lambda(I_r \otimes K), \quad b = (I_r \otimes K) \text{vec}(B) = \text{vec}(KB).$$

A is symmetric. If $K \succ 0$ and $\lambda > 0$, then $A \succ 0$, so CG applies; preconditioning reduces iterations. Direct dense solve costs $O((nr)^3) = O(n^3r^3)$ and requires forming A , while PCG only needs (i) matvecs $x \mapsto Ax$ and (ii) applying a cheap approximation $M^{-1} \approx A^{-1}$.

2. Matvec in $O(n^2r + qr)$ without forming A

Represent $x = \text{vec}(X)$ with $X \in \mathbb{R}^{n \times r}$. Use

$$(Z \otimes K) \text{vec}(X) = \text{vec}(KXZ^\top).$$

Store the observed index list in unfolding coordinates as pairs (i_t, j_t) for $t = 1, \dots, q$ with $i_t \in [n]$, $j_t \in [M]$. Then $S^\top \text{vec}(U) = (U_{i_t, j_t})_{t=1}^q$ (gather) and S scatters a length- q vector back into an $n \times M$ sparse matrix.

Given X :

1. Compute $G \leftarrow KX$ ($O(n^2r)$).
2. For each observed entry t : compute (or fetch) the row $z_t \equiv Z_{j_t,:}$ and the scalar $u_t \leftarrow G_{i_t,:} z_t^\top$ ($O(r)$).
3. Accumulate $H \in \mathbb{R}^{n \times r}$ via $H_{i_t,:} += u_t z_t$ for $t = 1, \dots, q$ ($O(qr)$).
4. Return $Y \leftarrow KH + \lambda G$ and output $\text{vec}(Y)$ ($O(n^2r)$).

This equals Ax because the first term is $(Z \otimes K)^\top \text{vec}(\tilde{U})$ with \tilde{U} the masked matrix whose nonzeros are u_t at (i_t, j_t) , and $(Z \otimes K)^\top \text{vec}(\tilde{U}) = \text{vec}(K\tilde{U}Z) = \text{vec}(KH)$. Equivalently, define $\mathcal{L}(X) \equiv S^\top \text{vec}(KXZ^\top)$ so that $A = \mathcal{L}^\top \mathcal{L} + \lambda(I \otimes K)$; its adjoint is $\mathcal{L}^\top u = \text{vec}(KUZ)$ where U is the sparse matrix with $\text{vec}(U) = Su$, which is implemented by the same scatter/accumulate loop.

Avoiding explicit Z . M can be huge; instead compute each needed z_t on the fly as a Hadamard product of the other factor rows using the observed multi-index. Optionally cache all z_t once in $O(qr)$ memory.

3. RHS in $O(n^2r + qr)$

Compute $B = TZ$ by a sparse MTTKRP over the q observed entries: for each observed value t_t at (i_t, j_t) , do $B_{i_t,:} += t_t z_t$ ($O(qr)$), then compute KB ($O(n^2r)$).

4. Preconditioner exploiting Kronecker structure

A standard preconditioner replaces the mask by a scaled identity ($P \approx \alpha I$ with $\alpha = q/N$), motivated by $\mathbb{E}[(Z \otimes K)^\top P(Z \otimes K)] = \alpha(Z^\top Z \otimes K^2)$ under uniform sampling, yielding

$$A_0 = \alpha[(Z \otimes K)^\top (Z \otimes K)] + \lambda(I \otimes K) = \alpha(Z^\top Z) \otimes (K^2) + \lambda(I \otimes K).$$

Let $G \equiv Z^\top Z$ (computable without forming Z via Hadamard products of the factor Gram matrices). With eigendecompositions $K = U\Lambda U^\top$ and $G = V\Sigma V^\top$, A_0 diagonalizes in the Kronecker basis $(V \otimes U)$ and applying A_0^{-1} to $\text{vec}(X)$ reduces to:

$$\hat{X} \leftarrow U^\top X V, \quad \hat{X}_{b,a} \leftarrow \hat{X}_{b,a}/(\alpha\sigma_a\lambda_b^2 + \lambda\lambda_b), \quad X \leftarrow U\hat{X}V^\top.$$

Cost per application: $O(n^2r + nr^2)$ after one-time setup $O(n^3 + r^3)$. In matrix form, $A_0 \text{vec}(X) = \text{vec}(K(\alpha K X G + \lambda X))$, so applying A_0^{-1} amounts to solving the Sylvester-type equation $\alpha K X G + \lambda X = K^{-1}R$.

5. Complexity

Let m be the PCG iteration count (typically $m \ll nr$ with a good preconditioner). Per iteration:

$$\text{matvec } Ax : O(n^2r + qr), \quad \text{preconditioner } M^{-1} : O(n^2r + nr^2), \quad \text{BLAS-1} : O(nr).$$

Total solve: $O(m(n^2r + qr + nr^2))$ time and $O(nr + q)$ memory (plus optional $O(qr)$ cache), with no $O(N)$ computation.