

## PSEUDOCODE

### Breadth-First Search

```
BFS( $G, s$ )
1   for each vertex  $u \in G.V - \{s\}$ 
2        $u.color = \text{WHITE}$ 
3        $u.d = \infty$ 
4        $u.\pi = \text{NIL}$ 
5    $s.color = \text{GRAY}$ 
6    $s.d = 0$ 
7    $s.\pi = \text{NIL}$ 
8    $Q = \emptyset$ 
9   ENQUEUE( $Q, s$ )
10  while  $Q \neq \emptyset$ 
11       $u = \text{DEQUEUE}(Q)$ 
12      for each  $v \in G.Adj[u]$ 
13          if  $v.color == \text{WHITE}$ 
14               $v.color = \text{GRAY}$ 
15               $v.d = u.d + 1$ 
16               $v.\pi = u$ 
17              ENQUEUE( $Q, v$ )
18       $u.color = \text{BLACK}$ 
```

### Depth-First Search

<pre>DFS(<math>G</math>) 1   <b>for</b> each vertex <math>u \in G.V</math> 2       <math>u.color = \text{WHITE}</math> 3       <math>u.\pi = \text{NIL}</math> 4   <math>time = 0</math> 5   <b>for</b> each vertex <math>u \in G.V</math> 6       <b>if</b> <math>u.color == \text{WHITE}</math> 7           DFS-VISIT(<math>G, u</math>)</pre>	<pre>DFS-VISIT(<math>G, u</math>) 1       <math>time = time + 1</math> 2       <math>u.d = time</math> 3       <math>u.color = \text{GRAY}</math> 4       <b>for</b> each <math>v \in G.Adj[u]</math> 5           <b>if</b> <math>v.color == \text{WHITE}</math> 6               <math>v.\pi = u</math> 7               DFS-VISIT(<math>G, v</math>) 8       <math>u.color = \text{BLACK}</math> 9       <math>time = time + 1</math> 10      <math>u.f = time</math></pre>
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### STRONGLY-CONNECTED-COMPONENTS( $G$ )

```
1   call DFS( $G$ ) to compute finishing times  $u.f$  for each vertex  $u$ 
2   compute  $G^T$ 
3   call DFS( $G^T$ ), but in the main loop of DFS, consider the vertices in order of decreasing  $u.f$  (as computed in line 1)
4   output the vertices of each tree in the depth-first forest formed in line 3 as a separate strongly connected component
```

## Heaps

PARENT( $i$ )

1   **return**  $\lfloor i/2 \rfloor$

LEFT( $i$ )

1   **return**  $2i$

RIGHT( $i$ )

1   **return**  $2i + 1$

MAX-HEAPIFY( $A, i$ )

```
1   $l = \text{LEFT}(i)$ 
2   $r = \text{RIGHT}(i)$ 
3  if  $l \leq A.\text{heap-size}$  and  $A[l] > A[i]$ 
4       $\text{largest} = l$ 
5  else  $\text{largest} = i$ 
6  if  $r \leq A.\text{heap-size}$  and  $A[r] > A[\text{largest}]$ 
7       $\text{largest} = r$ 
8  if  $\text{largest} \neq i$ 
9      exchange  $A[i]$  with  $A[\text{largest}]$ 
10     MAX-HEAPIFY( $A, \text{largest}$ )
```

BUILD-MAX-HEAP( $A$ )

```
1   $A.\text{heap-size} = A.\text{length}$ 
2  for  $i = \lfloor A.\text{length} / 2 \rfloor$  downto 1
3      MAX-HEAPIFY( $A, i$ )
```

HEAPSORT( $A$ )

```
1  BUILD-MAX-HEAP( $A$ )
2  for  $i = A.\text{length}$  downto 2
3      exchange  $A[1]$  with  $A[i]$ 
4       $A.\text{heap-size} = A.\text{heap-size} - 1$ 
5      MAX-HEAPIFY( $A, 1$ )
```

HEAP-MAXIMUM( $A$ )

1   **return**  $A[1]$

HEAP-EXTRACT-MAX( $A$ )

```
1  if  $A.\text{heap-size} < 1$ 
2      error "heap underflow"
3   $\text{max} = A[1]$ 
4   $A[1] = A[A.\text{heap-size}]$ 
5   $A.\text{heap-size} = A.\text{heap-size} - 1$ 
6  MAX-HEAPIFY( $A, 1$ )
7  return  $\text{max}$ 
```

HEAP-INCREASE-KEY( $A, i, \text{key}$ )

```
1  if  $\text{key} < A[i]$ 
2      error "new key is smaller than current key"
3   $A[i] = \text{key}$ 
4  while  $i > 1$  and  $A[\text{PARENT}(i)] < A[i]$ 
5      exchange  $A[i]$  with  $A[\text{PARENT}(i)]$ 
6       $i = \text{PARENT}(i)$ 
```

MAX-HEAP-INSERT( $A, \text{key}$ )

```
1   $A.\text{heap-size} = A.\text{heap-size} + 1$ 
2   $A[A.\text{heap-size}] = -\infty$ 
3  HEAP-INCREASE-KEY( $A, A.\text{heap-size}, \text{key}$ )
```

## Disjoint-Set Forests

MAKE-SET( $x$ )

```
1   $x.p = x$ 
2   $x.rank = 0$ 
3
```

FIND-SET( $x$ )

```
1  if  $x \neq x.p$ 
2       $x.p = \text{FIND-SET}(x.p)$ 
3  return  $x.p$ 
```

LINK( $x, y$ )

```
1  if  $x.rank > y.rank$ 
2       $y.p = x$ 
3  else  $x.p = y$ 
4      if  $x.rank == y.rank$ 
5           $y.rank = y.rank + 1$ 
```

UNION( $x, y$ )

```
1  LINK(FIND-SET( $x$ ), FIND-SET( $y$ ))
```

## Minimum Spanning Trees

MST-KRUSKAL( $G, w$ )

```
1   $A = \emptyset$ 
2  for each vertex  $v \in G.V$ 
3      MAKE-SET( $v$ )
4  sort the edges of  $G.E$  into nondecreasing order by weight  $w$ 
5  for each edge  $(u, v) \in G.E$ , taken in nondecreasing order by weight
6  if FIND-SET( $u$ )  $\neq$  FIND-SET( $v$ )
7       $A = A \cup \{(u, v)\}$ 
8      UNION( $u, v$ )
9  return  $A$ 
```

MST-PRIM( $G, w, r$ )

```
1  for each  $u \in G.V$ 
2       $u.key = \infty$ 
3       $u.\pi = \text{NIL}$ 
4   $r.key = 0$ 
5   $Q = G.V$ 
6  while  $Q \neq \emptyset$ 
7       $u = \text{EXTRACT-MIN}(Q)$ 
8      for each  $v \in G.Adj[u]$ 
9          if  $v \in Q$  and  $w(u, v) < v.key$ 
10              $v.\pi = u$ 
11              $v.key = w(u, v)$ 
```

## Shortest Paths in Weighted Graphs

INITIALIZE-SINGLE-SOURCE( $G, s$ )

```
1  for each vertex  $v \in G.V$ 
2       $v.d = \infty$ 
3       $v.\pi = \text{NIL}$ 
4   $s.d = 0$ 
```

RELAX( $u, v, w$ )

```
1  if  $v.d > u.d + w(u, v)$ 
2       $v.d = u.d + w(u, v)$ 
3       $v.\pi = u$ 
```

DIJKSTRA( $G, w, s$ )

```
1  INITIALIZE-SINGLE-SOURCE( $G, s$ )
2   $S = \emptyset$ 
3   $Q = G.V$ 
4  while  $Q \neq \emptyset$ 
5       $u = \text{EXTRACT-MIN}(Q)$ 
6       $S = S \cup \{u\}$ 
7      for each vertex  $v \in G.Adj[u]$ 
8          RELAX( $u, v, w$ )
```

## Hash Tables Using Open Addressing

HASH-INSERT( $T, k$ )

```
1   $i = 0$ 
2  repeat
3       $j = h(k, i)$ 
4      if  $T[j] = \text{NIL}$  or  $T[j] = \text{DELETED}$ 
5           $T[j] = k$ 
6          return  $j$ 
7      else  $i = i + 1$ 
8  until  $i == m$ 
9  error "hash table overflow"
```

HASH-SEARCH( $T, k$ )

```
1   $i = 0$ 
2  repeat
3       $j = h(k, i)$ 
4      if  $T[j] == k$ 
5          return  $j$ 
6       $i = i + 1$ 
7  until  $T[j] = \text{NIL}$  or  $i == m$ 
8  return NIL
```

HASH-DELETE( $T, k$ )

```
1   $i = \text{HASH-SEARCH}(T, k)$ 
2  if  $i \neq \text{NIL}$ 
3       $T[i] = \text{DELETED}$ 
```

## Red-Black Trees

LEFT-ROTATE( $T, x$ )

```
1   $y = x.right$ 
2   $x.right = y.left$ 
3  if  $y.left \neq T.nil$ 
4       $y.left.p = x$ 
5   $y.p = x.p$ 
6  if  $x.p == T.nil$ 
7       $T.root = y$ 
8  elseif  $x == x.p.left$ 
9       $x.p.left = y$ 
10 else  $x.p.right = y$ 
11  $y.left = x$ 
12  $x.p = y$ 
```

RB-TRANSPLANT( $T, u, v$ )

```
1  if  $u.p == T.nil$ 
2       $T.root = v$ 
3  elseif  $u == u.p.left$ 
4       $u.p.left = v$ 
5  else  $u.p.right = v$ 
6   $v.p = u.p$ 
```

RB-INSERT( $T, z$ )

```
1   $y = T.nil$ 
2   $x = T.root$ 
3  while  $x \neq T.nil$ 
4       $y = x$ 
5      if  $z.key < x.key$ 
6           $x = x.left$ 
7      else  $x = x.right$ 
8   $z.p = y$ 
9  if  $y == T.nil$ 
10      $T.root = z$ 
11 elseif  $z.key < y.key$ 
12      $y.left = z$ 
13 else  $y.right = z$ 
14  $z.left = T.nil$ 
15  $z.right = T.nil$ 
16  $z.color = RED$ 
17 RB-INSERT-FIXUP( $T, z$ )
```

RB-INSERT-FIXUP( $T, z$ )

```
1  while  $z.p.color == RED$ 
2      if  $z.p == z.p.p.left$ 
3           $y = z.p.p.right$ 
4          if  $y.color == RED$ 
5               $z.p.color = BLACK$ 
6               $y.color = BLACK$ 
7               $z.p.p.color = RED$ 
8               $z = z.p.p$ 
9          else if  $z == z.p.right$ 
10              $z = z.p$ 
11             LEFT-ROTATE( $T, z$ )
12              $z.p.color = BLACK$ 
13              $z.p.p.color = RED$ 
14             RIGHT-ROTATE( $T, z.p$ )
15         else (same as then clause with “right” and “left” exchanged)
16      $T.root.color = BLACK$ 
```

RB-DELETE( $T, z$ )

```
1   $y = z$ 
2   $y\text{-original-color} = y.\text{color}$ 
3  if  $z.\text{left} == T.\text{nil}$ 
4       $x = z.\text{right}$ 
5      RB-TRANSPLANT( $T, z, z.\text{right}$ )
6  elseif  $z.\text{right} == T.\text{nil}$ 
7       $x = z.\text{left}$ 
8      RB-TRANSPLANT( $T, z, z.\text{left}$ )
9  else  $y = \text{TREE-MINIMUM}(z.\text{right})$ 
10      $y\text{-original-color} = y.\text{color}$ 
11      $x = y.\text{right}$ 
12     if  $y.p == z$ 
13          $x.p = y$ 
14     else RB-TRANSPLANT( $T, y, y.\text{right}$ )
15          $y.\text{right} = z.\text{right}$ 
16          $y.\text{right}.p = y$ 
17     RB-TRANSPLANT( $T, z, y$ )
18      $y.\text{left} = z.\text{left}$ 
19      $y.\text{left}.p = y$ 
20      $y.\text{color} = z.\text{color}$ 
21 if  $y\text{-original-color} == \text{BLACK}$ 
22     RB-DELETE-FIXUP( $T, x$ )
```

RB-DELETE-FIXUP( $T, x$ )

```
1  while  $x \neq T.\text{root}$  and  $x.\text{color} == \text{BLACK}$ 
2      if  $x == x.p.\text{left}$ 
3           $w = x.p.\text{right}$ 
4          if  $w.\text{color} == \text{RED}$ 
5               $w.\text{color} = \text{BLACK}$ 
6               $x.p.\text{color} = \text{RED}$ 
7              LEFT-ROTATE( $T, x.p$ )
8               $w = x.p.\text{right}$ 
9          if  $w.\text{left}.color == \text{BLACK}$  and  $w.\text{right}.color == \text{BLACK}$ 
10              $w.\text{color} = \text{RED}$ 
11              $x = x.p$ 
12          else if  $w.\text{right}.color == \text{BLACK}$ 
13               $w.\text{left}.color = \text{BLACK}$ 
14               $w.\text{color} = \text{RED}$ 
15              RIGHT-ROTATE( $T, w$ )
16               $w = x.p.\text{right}$ 
17               $w.\text{color} = x.p.\text{color}$ 
18               $x.p.\text{color} = \text{BLACK}$ 
19               $w.\text{right}.color = \text{BLACK}$ 
20              LEFT-ROTATE( $T, x.p$ )
21               $x = T.\text{root}$ 
22          else (same as then clause with “right” and “left” exchanged)
23       $x.\text{color} = \text{BLACK}$ 
```