# 词法分析---NFA转换到DFA

编译原理

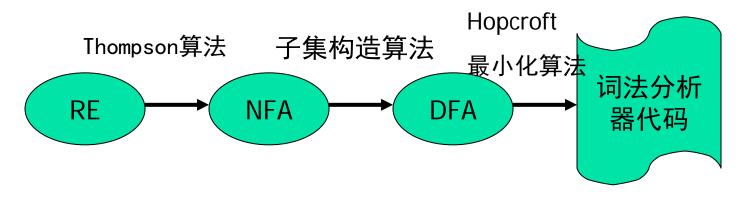
华保健

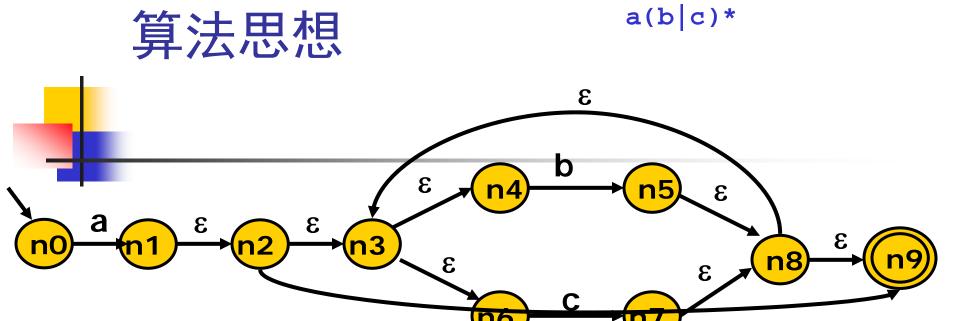
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## 回顾: 自动生成



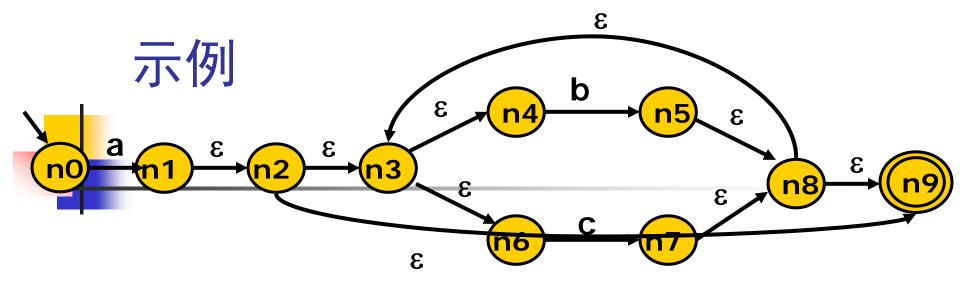
词法分析器





## 子集构造算法

```
(* 子集构造算法:工作表算法 *)
q0 <- eps closure (n0)
Q < - \{q0\}
workList <- q0
while (workList != [])
  remove q from workList
  foreach (character c)
    t <- e-closure (delta (q, c))
    D[q, c] < -t
    if (t\not\in Q)
      add t to Q and workList
```



```
(* 子集构造算法:工作表算法 *)
q0 <- eps_closure (n0)
Q < - \{q0\}
workList <- q0</pre>
while (workList != [])
  remove q from workList
  foreach (character c)
    t <- e-closure (delta (q, c))
    D[q, c] \leftarrow t
    if (t\not\in Q)
      add t to Q and workList
```

### 对算法的讨论

- 不动点算法
  - 算法为什么能够运行终止
- ■时间复杂度
  - 最坏情况O(2N)
  - 但在实际中不常发生
    - 因为并不是每个子集都会出现

## ε-闭包的计算: 深度优先

```
/* \epsilon-closure: 基于深度优先遍历的算法 */
set closure = {};
void eps_closure (x)
  closure += {x}
  foreach (y: x--\varepsilon-> y)
    if (!visited(y))
      eps_closure (y)
                   3
```

## ε-闭包的计算: 宽度优先

```
/* ε-closure: 基于宽度优先的算法 */
set closure = {};
Q = []; // queue
void eps closure (x) =
  O = [x];
  while (Q not empty)
    q <- deQueue (Q)
    closure += q
    foreach (y: q--\varepsilon-> y)
      if (!visited(y))
        enQueue (Q, y)
```