Final Lecture

Eugene Wu

Administrivia

Final next Friday 4-7PM, in regular classroom

Open book, paper notes

Comparable to midterm, slightly longer.

Grading

Normalize all assignments to [0-1]
Weighed sum based on assignment weights
Compute cut-offs, ~ B+ avg
Then add in extra credit

DBMSes in the Wild

Classic Relational

\$\$: Oracle, IBM, Microsoft, Teradata, EMC, etc

Free: MySQL, PostgreSQL

New Relational

In-Memory, Column-store, Streaming

Non-traditional

Search (Google, Bing, Lucene), Scientific, Geographic

NoSQL

Big Data: Hadoop, Spark, etc

Key-value: Mongo, Berkeley DB, Cassandra, etc

DBMS-as-a-Service

Microsoft Azure, Amazon Redshift/RDS, etc...

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Modern Database Systems

Hardware changes affect

Compression is good \rightarrow Column stores

Large scale aggregation queries

"Data Warehouses"

Memory is cheap → In-memory stores

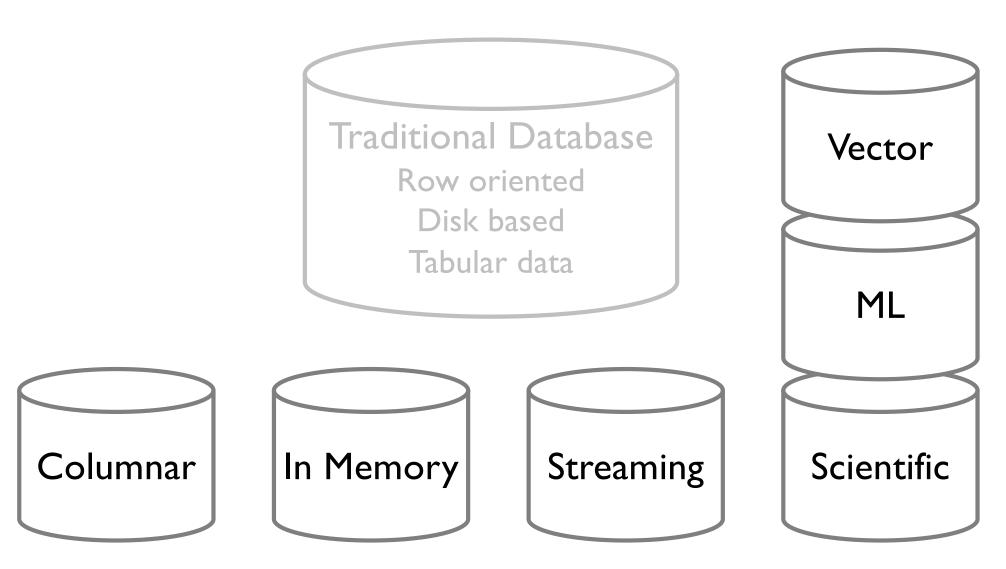
Transactional systems

One Size Fits All

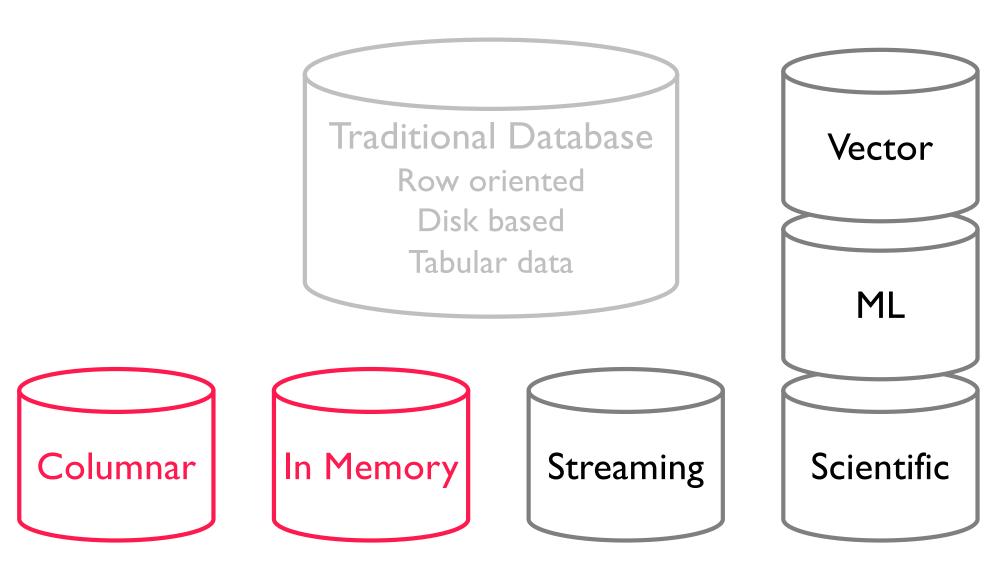
Traditional Database

Row oriented
Disk based
Tabular data

One Size Does Not Fits All



One Size Does Not Fits All



Data Warehouses

Store all historical data for future analysis

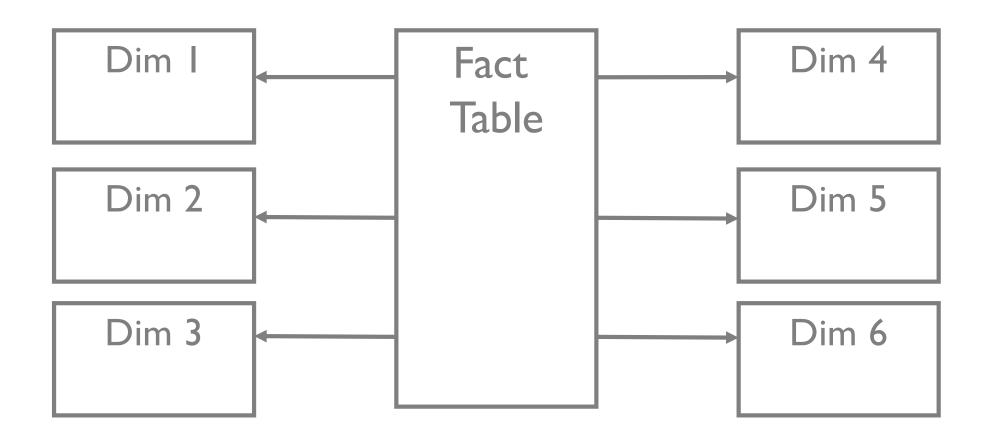
Sales by month over past 20 years

Clicks by youth in texas

Cost by product component

Defacto standard at any company

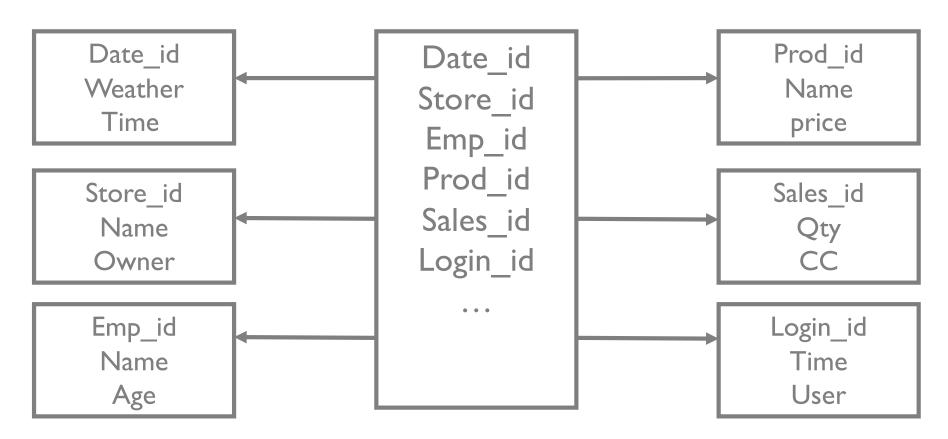
Star Schema



Star Schema

Fact table is "fat"

Queries access ~6 attrs

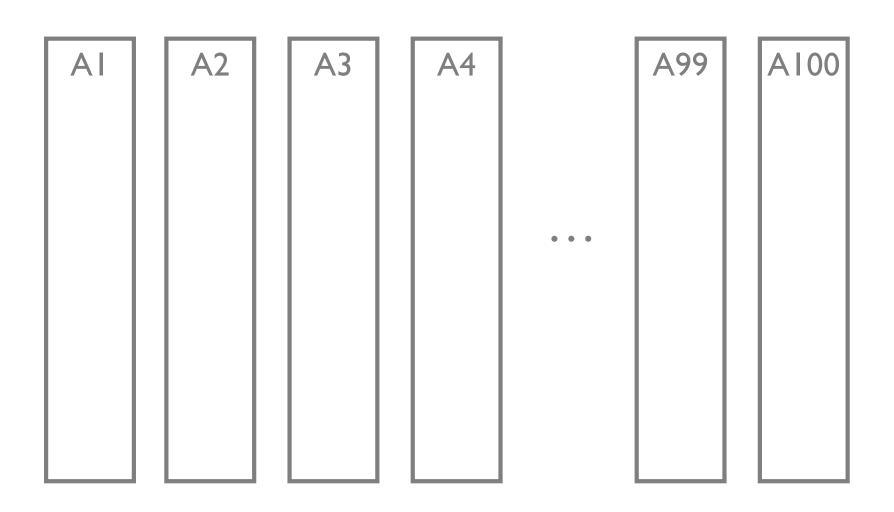


SELECT A2, A3, AVG(A99) GROUP BY A2, A3

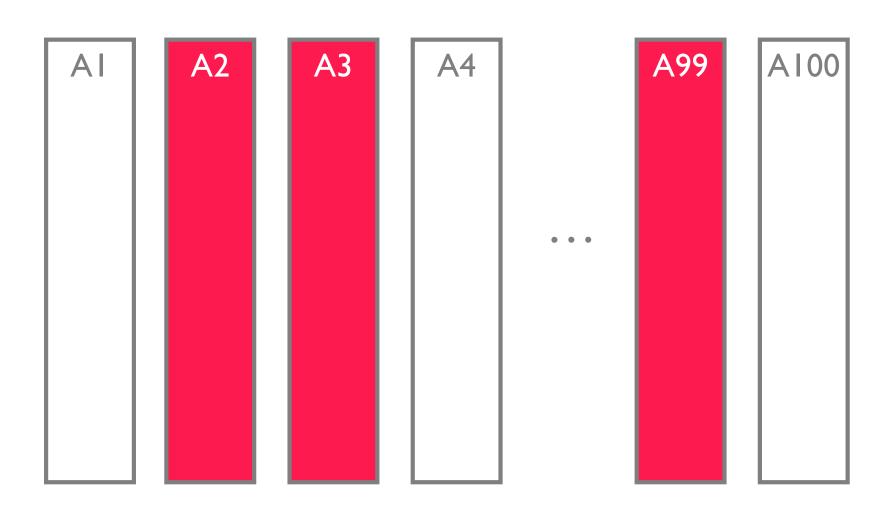
100 attributes

Tuple I
Tuple 2
Tuple 3
Tuple 4
Tuple 5

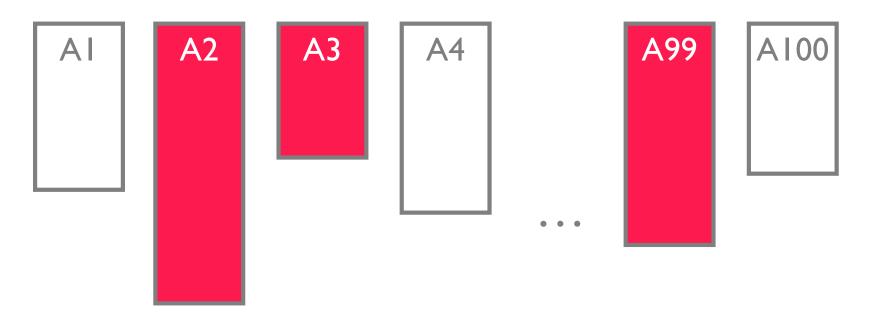
SELECT A2, A3, AVG(A99) GROUP BY A2, A3



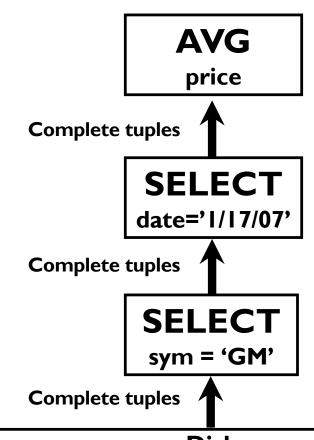
16x less data read. Unfair advantage.



16x less data read. Unfair advantage.Compression better on single columnExecute on compressed data



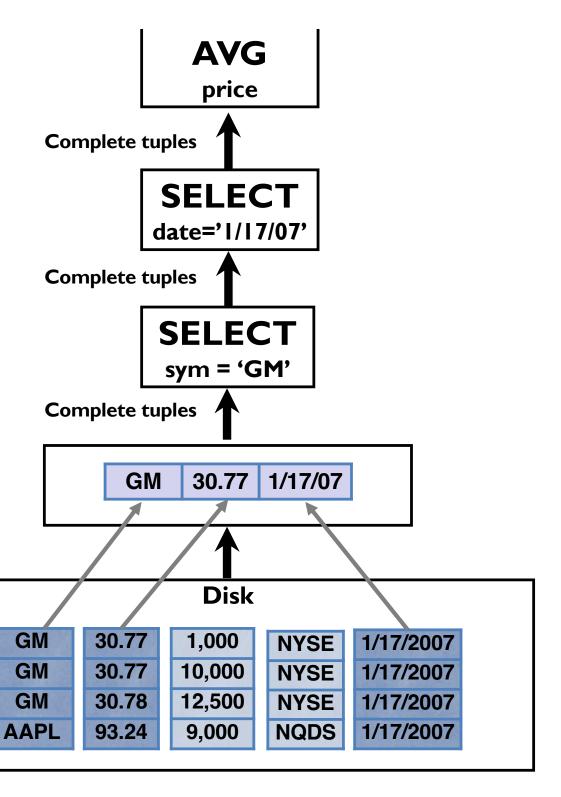
Traditional DBMS



		Disk		
GM	30.77	1,000	NYSE	1/17/2007
GM	30.77	10,000	NYSE	1/17/2007
GM	30.78	12,500	NYSE	1/17/2007
AAPL	93.24	9,000	NQDS	1/17/2007

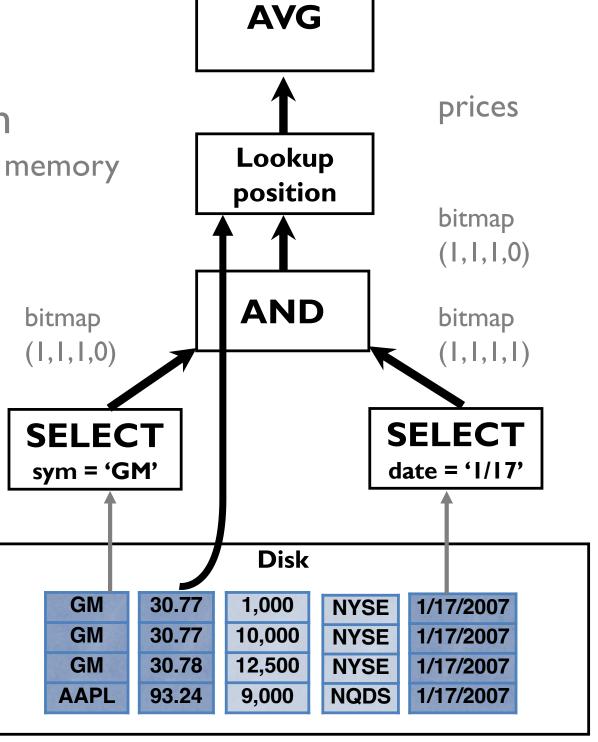
Naïve:

Early Materialization Row oriented execution



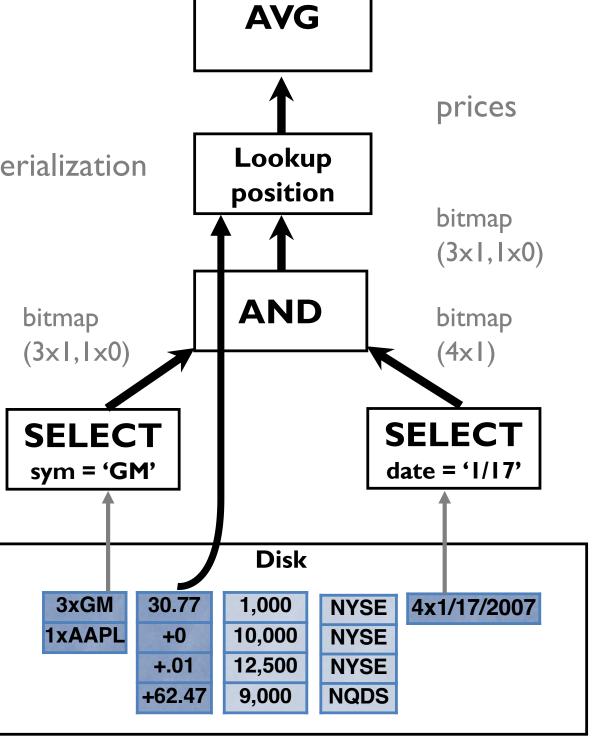
C-Store Late Materialization

Much less data moving in memory



C-Store Compression

Only possible w/ late materialization



Column Stores

Optimized for data warehouses

Store data by attribute/column rather than row

Compression

Compressed query plan execution

50-100x faster than row store

In-Memory DBMSes

Transaction-oriented apps

remove I unit from product move 5 units from org I to org 2 (shopping carts, inventory)

Data stored in memory

Disk only used for check pointing

No concurrency

Active-active replication for fault-tolerance

Traditional Database

Indexes queries go faster

Concurrency queries go faster

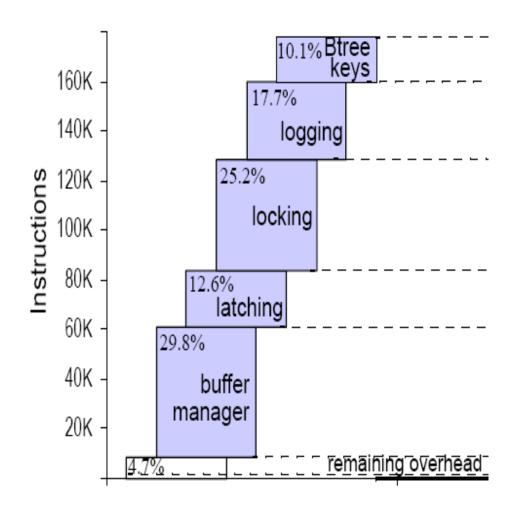
Locking serializability

Logging recovery

Buffer Manager manage pages in memory

Results after removing the components (in # instruction)

Instruction of useful work is only <2% of a memory resident DB



Effect of removing different components for payment (1/6)

commit

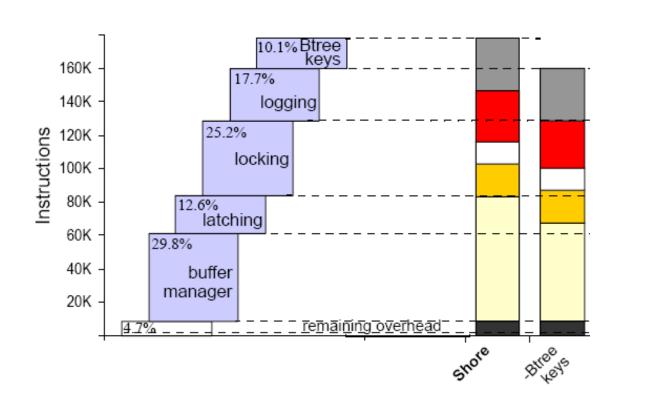
begin

create record

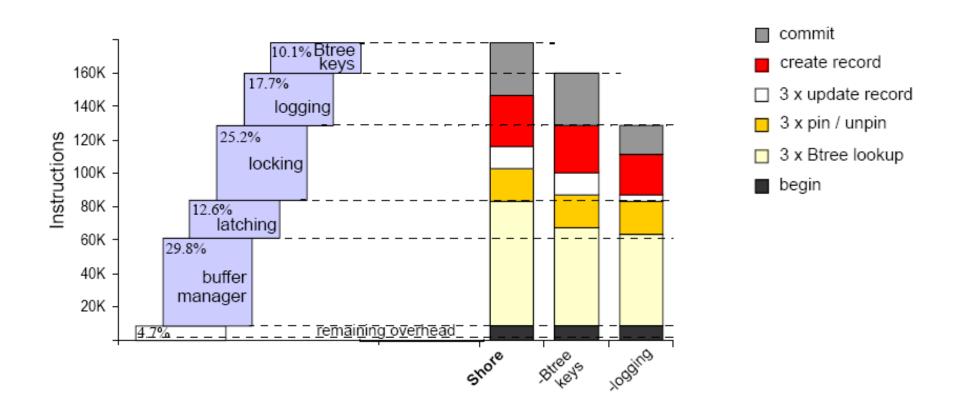
3 x pin / unpin

3 x Btree lookup

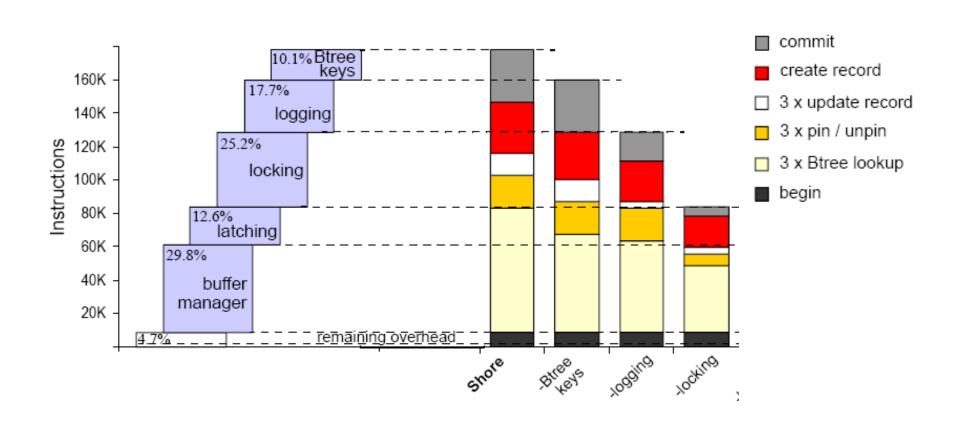
3 x update record



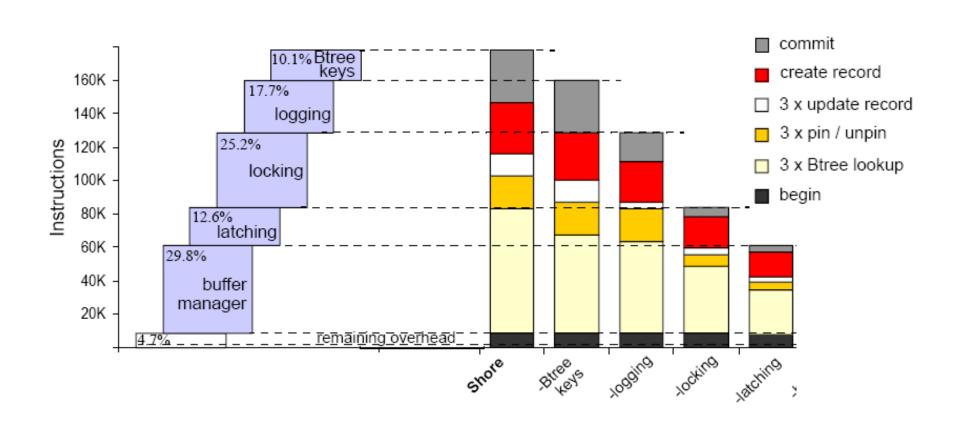
Effect of removing different components for payment (2/6)



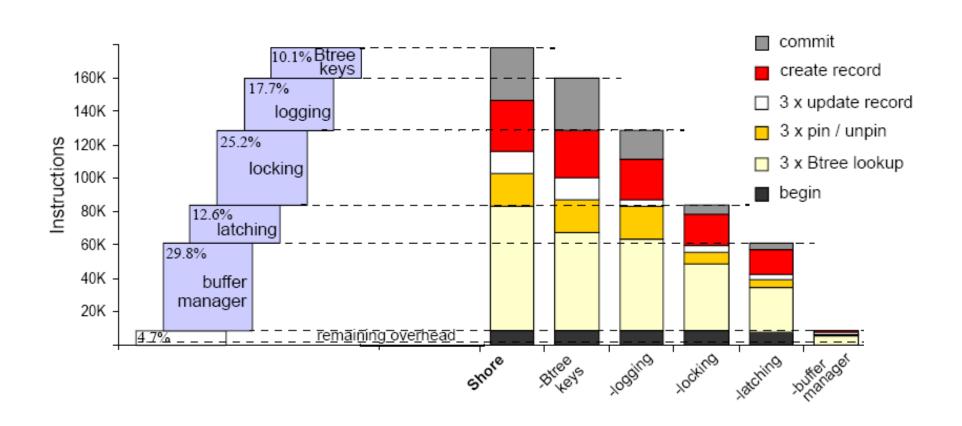
Effect of removing different components for payment (3/6)



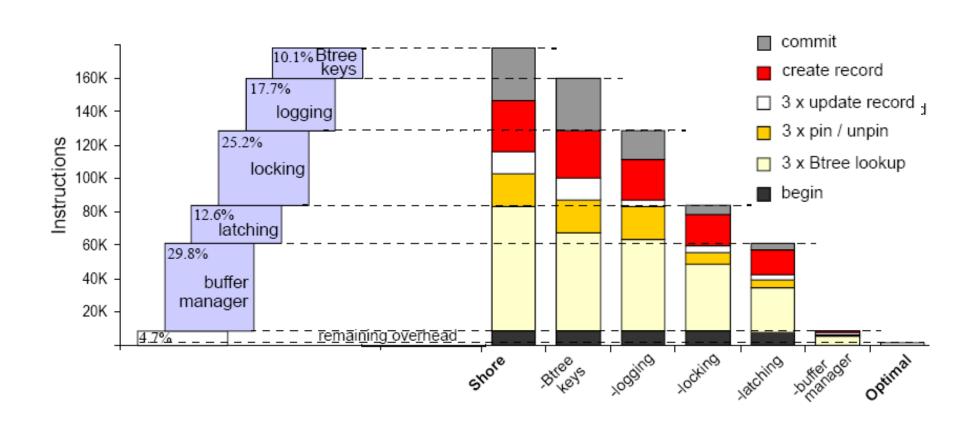
Effect of removing different components for payment (4/6)



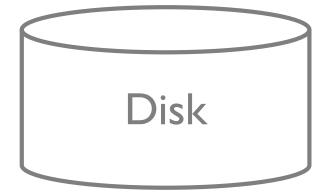
Effect of removing different components for payment (5/6)



Effect of removing different components for payment (6/6)



Barebones Executor



In memory DB

Barebones Executor

What about

Parsing

Concurrency?

Recovery?

In memory DB

Barebones Executor

Stored Procedures

What about

Parsing

Concurrency?

Recovery?

Procedure:

```
p1 = SELECT cost
FROM fact_table
WHERE id = ?
```

Query:

p1(10)

In memory DB

Barebones Executor

What about

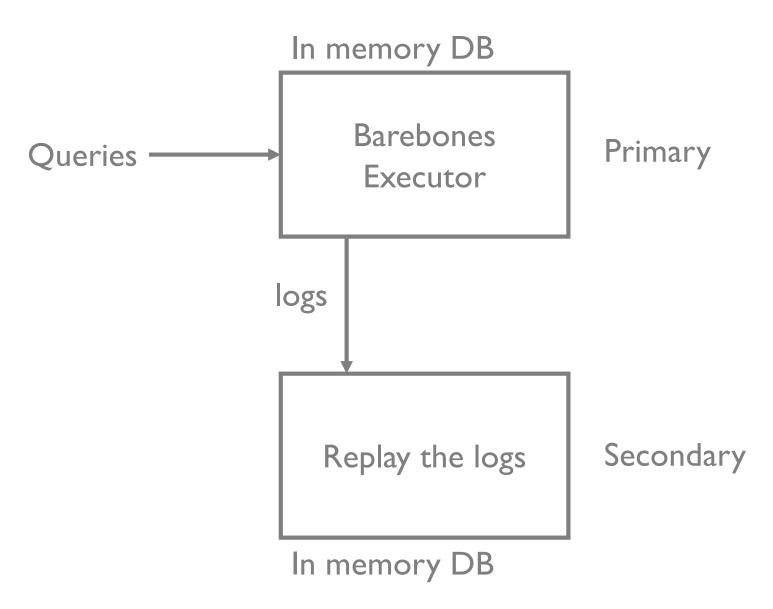
Parsing

Concurrency?

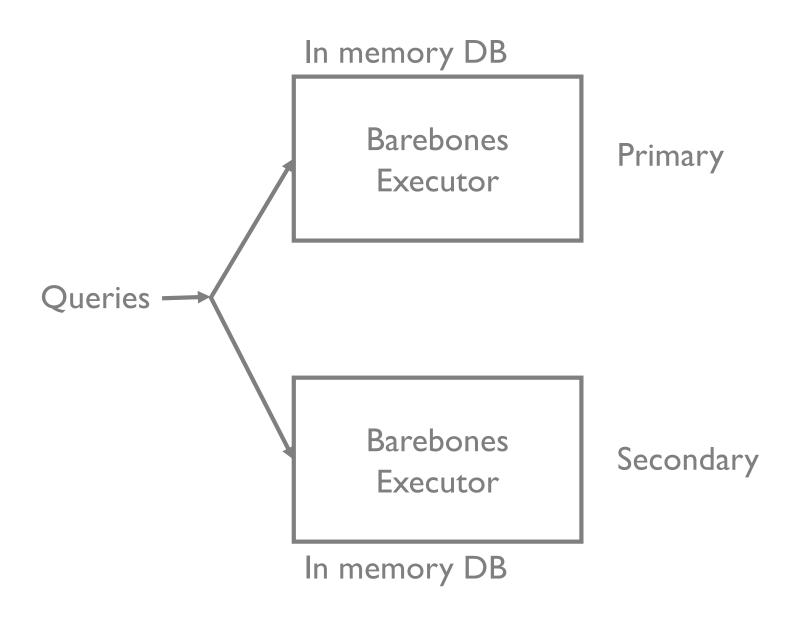
no buffer manager, no concurrency, no locks

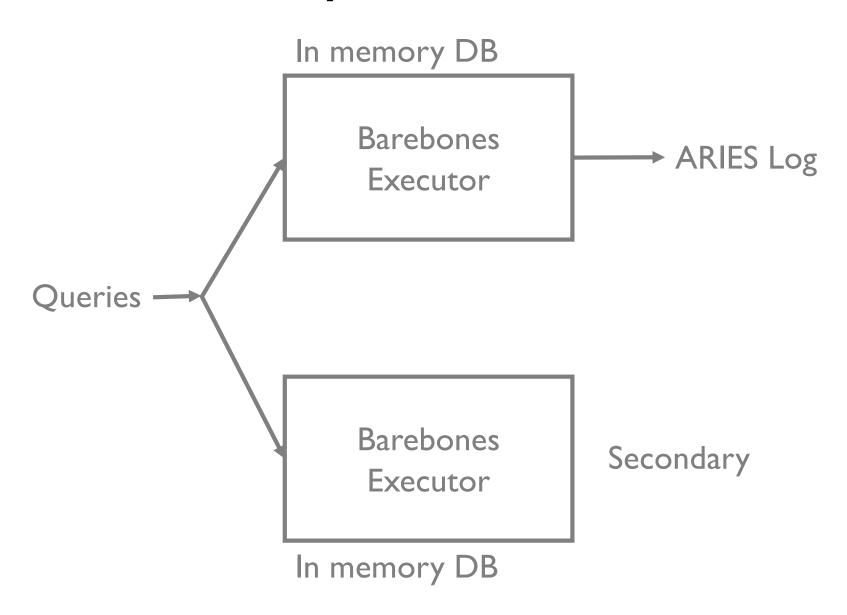
Recovery?

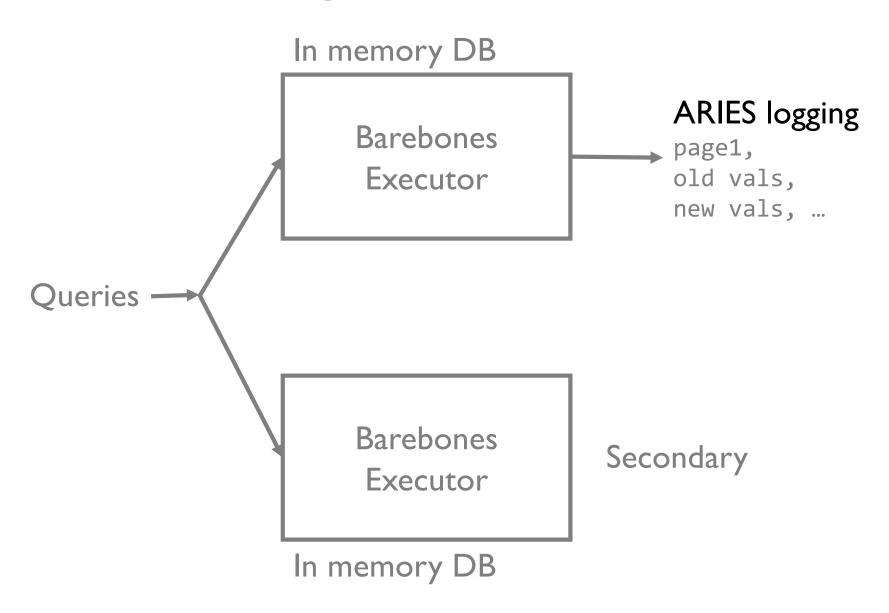
Log Shipping

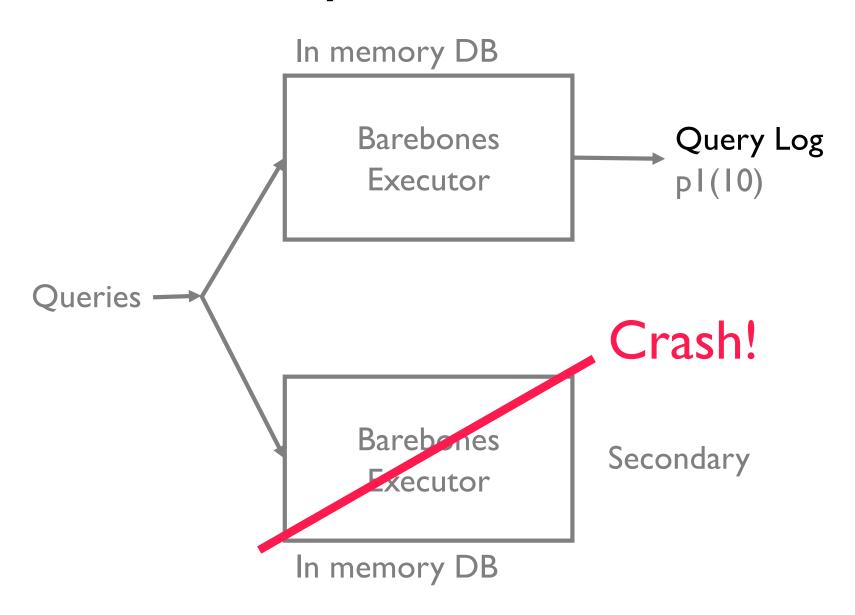


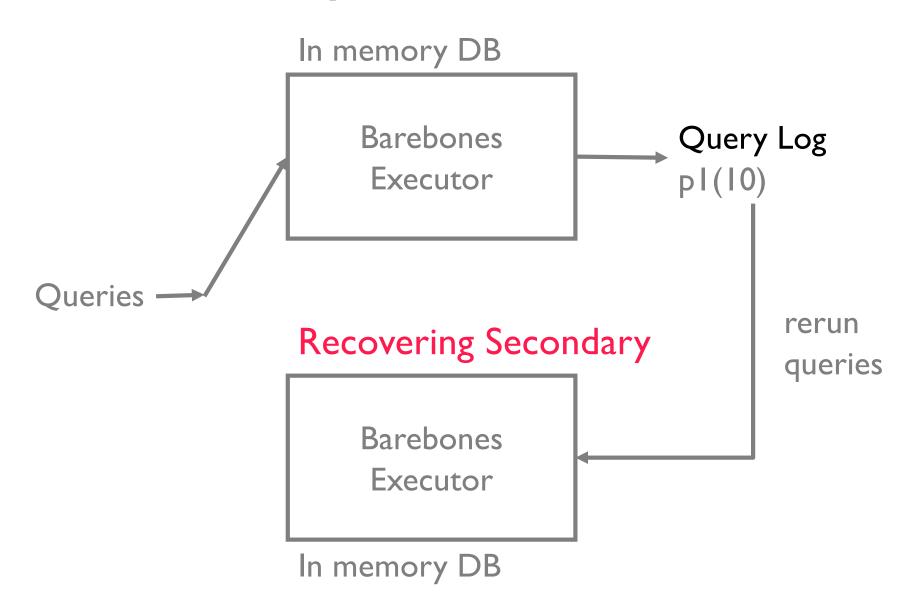
Active-Active











In Practice

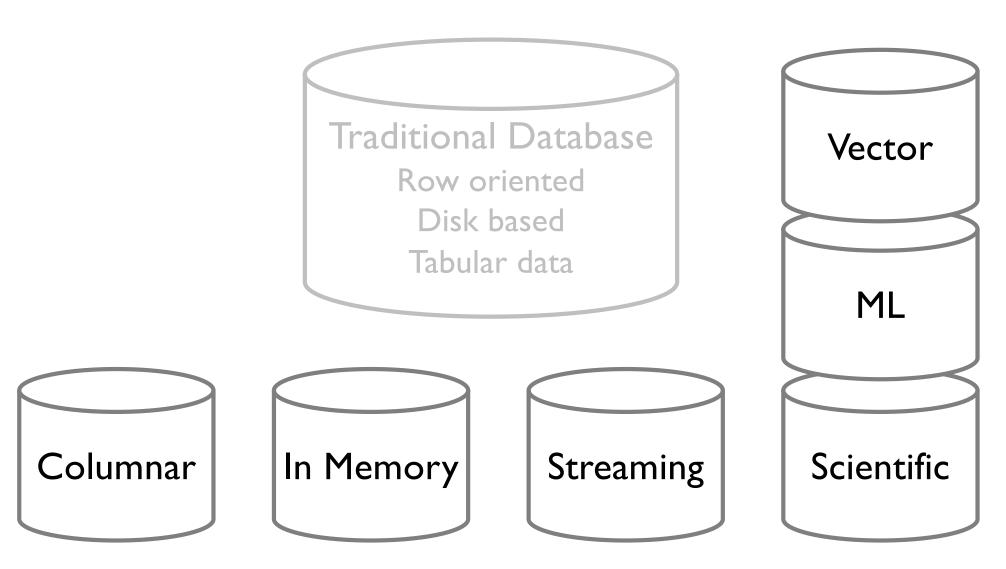
Single machine

- DuckDB: in-memory column store
- SQLite: in-memory row store

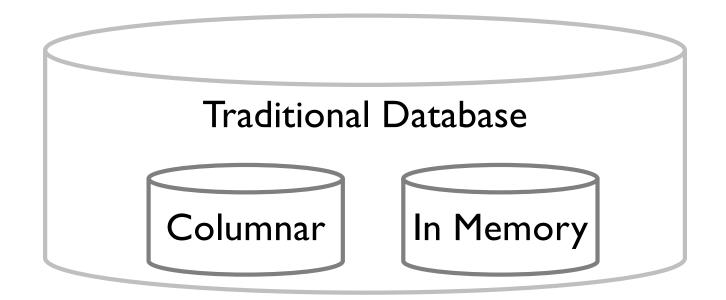
Multi-machine

- Motherduck: cloud-client DuckDB
- Clickhouse, BigQuery, Redshift, ...: columnar
- Cockroach, Spanner, Yugabyte: row

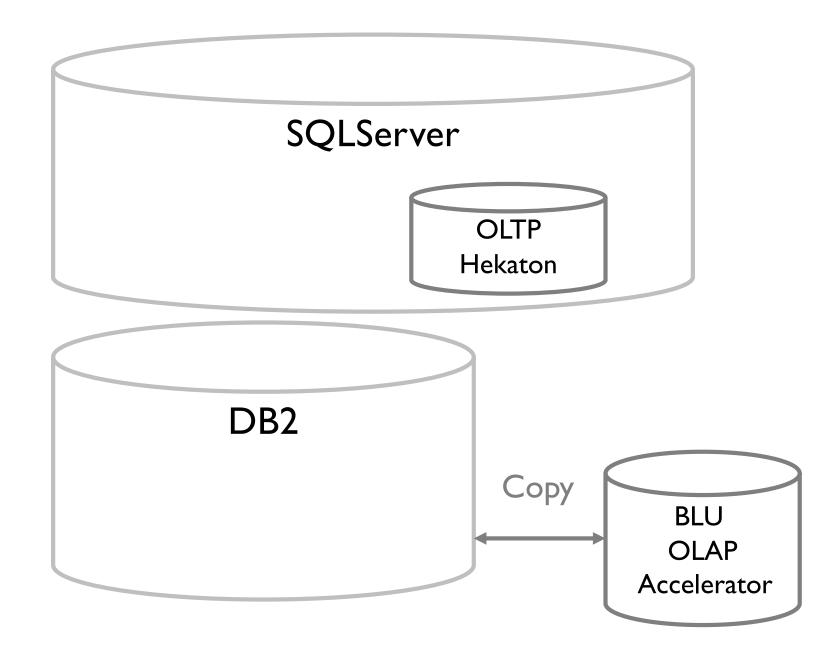
One Size Does Not Fits All



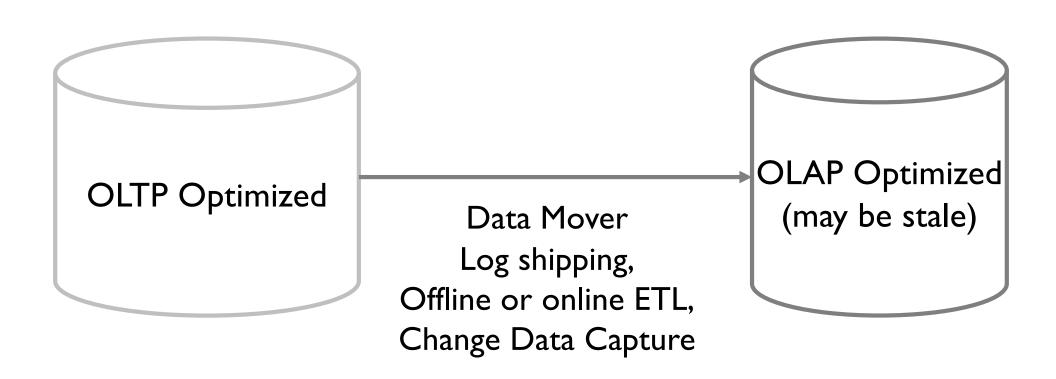
One Size Does Not Fits All



One Size Does Not Fits All



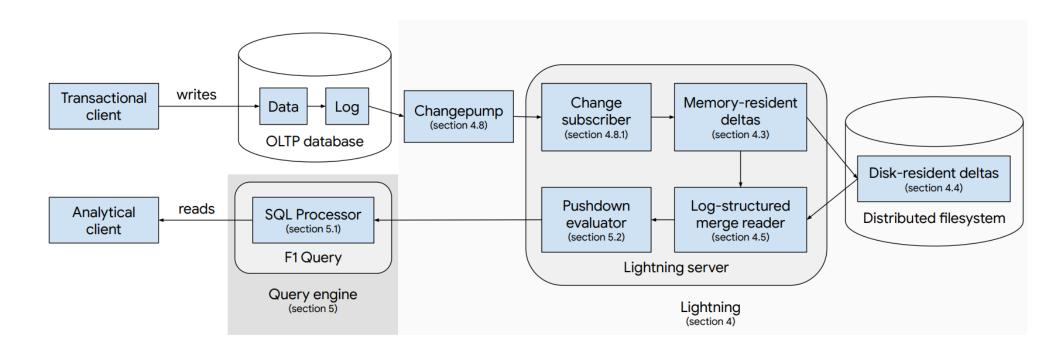
"HTAP"



F1 Lightning: HTAP as a Service

Jiacheng Yang Ian Rae Jun Xu Jeff Shute Zhan Yuan Kelvin Lau Qiang Zeng Xi Zhao Jun Ma Ziyang Chen Yuan Gao Qilin Dong Junxiong Zhou Jeremy Wood Goetz Graefe Jeff Naughton John Cieslewicz Google LLC

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Recent Popular Systems

Vector Stores

- Pinecone, Qdrant, Weaviate, ...
- Nearest neighbor + vector data type

Key-value Stores

B-tree index

OK Let's Step Back

Discussed the how:

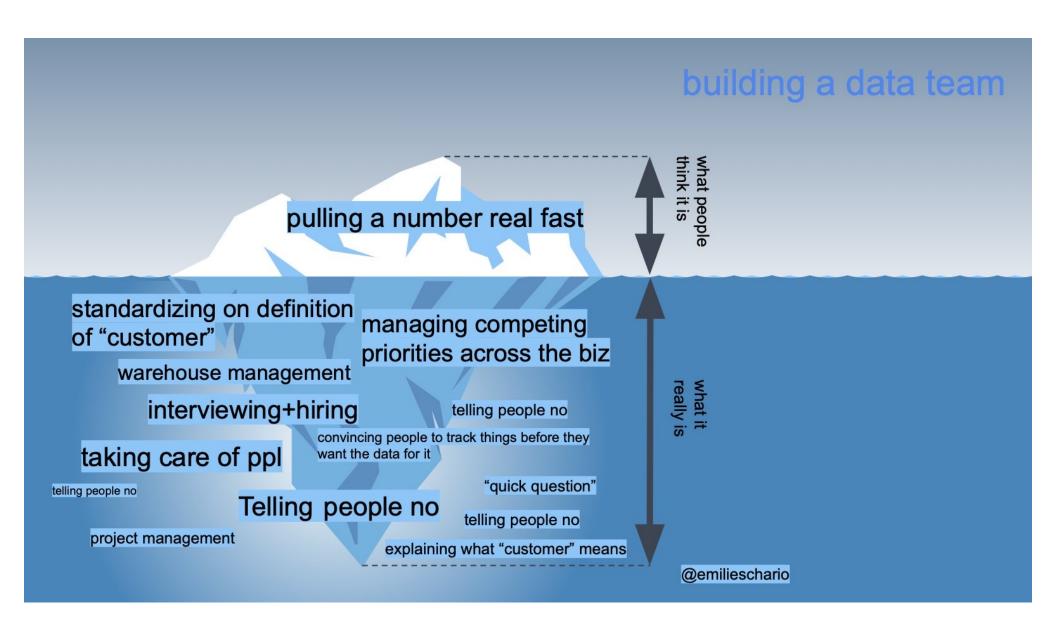
how to model data needed by an application

how to query databases

how databases execute queries

how databases run fast, correctly

To what end?



Production ML

An on-call engineer's biggest nightmare

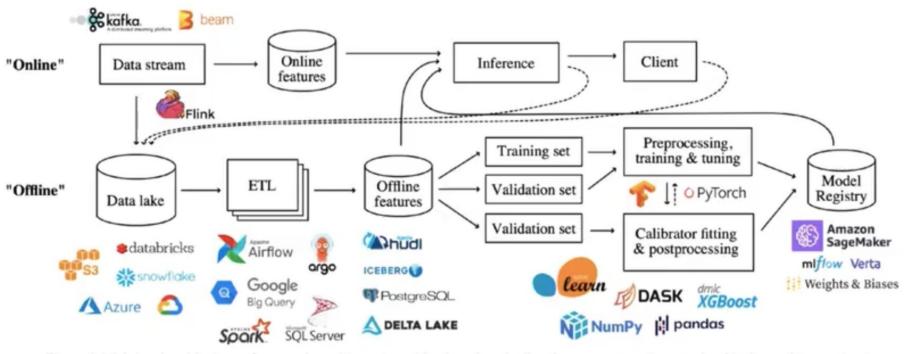
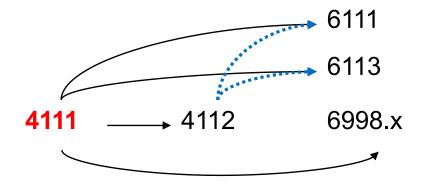


Figure 1: High-level architecture of a generic end-to-end machine learning pipeline. Logos represent a sample of tools used to construct components of the pipeline, illustrating heterogeneity in the tool stack. Shankar et al. 2021

https://www.facebook.com/Engineering/videos/1578607659138164/



4xxx: "here are facts about data management"

4111: basics

4112: gory DBMS internals

6xxx: "how people figured out those facts"

6111: info extraction and web

6113: classic and modern DB research

6998: systems for human data interaction

PhD: "we don't understand anything. Plz help"

Thanks!

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