SLP & SQP Methods

LI-PING ZHANG

Department of Mathematical Sciences
Tsinghua University, Beijing 100084

Office: New Science Building #A302, Tel: 62798531

E-mail: lipingzhang@tsinghua.edu.cn



- Successive Linear Programming Method
- Frank-Wolfe Method
- Successive Quadratic Programming Method

消華大学

基本思想: 将非线性规划线性化, 通过解线性规 划问题来求解原问题的近似解。

(NP)
$$\begin{cases} \min f(x) \\ s.t. \quad g_j(x) \ge 0 \quad j = 1, \dots, m \\ h_j(x) = 0 \quad j = 1, \dots, l \end{cases}$$

其中 $x \in \mathbb{R}^n$, f, g_i , h_i 均存在一阶连续偏导数。

基本思想:

将(NP)中的目标函数f(x)和约束函数 $g_i(x)$, $h_j(x)$ 线性化,并对变量的取值范围加以限制,从而得到线性近似规划,用单纯形方法求解此线性规划问题,把其最优解作为(NP)的解的近似。

设 $x^{(k)}$ 是原问题的可行解,将 $f(x), g_i(x),$

 $h_i(x)$ 在 $x^{(k)}$ 点 Taylor展开。

$$\begin{cases} \min f(x^{(k)}) + \nabla f(x^{(k)})^T (x - x^{(k)}) \\ s.t. \quad g_i(x^{(k)}) + \nabla g_i(x^{(k)})^T (x - x^{(k)}) \ge 0, i = 1, \dots m \\ h_j(x^{(k)}) + \nabla h_j(x^{(k)})^T (x - x^{(k)}) = 0, j = 1, \dots l \end{cases}$$



$$\begin{cases} \min \ f(x^{(k)}) + \nabla f(x^{(k)})^T (x - x^{(k)}) \\ s.t. \quad g_i(x^{(k)}) + \nabla g_i(x^{(k)})^T (x - x^{(k)}) \geq 0, i = 1, \cdots m \\ h_j(x^{(k)}) + \nabla h_j(x^{(k)})^T (x - x^{(k)}) = 0, j = 1, \cdots l \\ |x_j - x_j^{(k)}| \leq \delta_j^{(k)}, \quad j = 1, 2, \cdots, n \\ \text{为了保证有最优解,限制在紧集上} \end{cases}$$

道道等大学

步骤

- 1. 给定初始可行解 $x^{(1)}$, $\delta_{j}^{(1)}$, $j = 1, 2, \dots, n$,缩小误差 $\beta \in (0,1)$,允许误差 ε_{1} , ε_{2} ,置k := 1。
- 2. 求解线性规划问题:

$$\begin{cases} \min f(x^{(k)}) + \nabla f(x^{(k)})^T (x - x^{(k)}) \\ s.t. & g_i(x^{(k)}) + \nabla g_i(x^{(k)})^T (x - x^{(k)}) \ge 0, i = 1, \dots m \\ & h_j(x^{(k)}) + \nabla h_j(x^{(k)})^T (x - x^{(k)}) = 0, j = 1, \dots l \\ & |x_j - x_j^{(k)}| \le \delta_j^{(k)}, \quad j = 1, 2, \dots, n \end{cases}$$

得最优解求.

清華大学

3. 若 \overline{x} 是(*NP*)的可行解,则令 $x^{(k+1)} = \overline{x}$,转4; 否则,置 $\delta_{j}^{(k)} := \beta \delta_{j}^{(k)}$, $j = 1, 2, \dots, n$,返回2。

4. 若 $|f(x^{(k+1)}) - f(x^{(k)})| < \varepsilon_1$,且 $||x^{(k+1)} - x^{(k)}|| < \varepsilon_2$ 或 $|\delta_j^{(k)}| < \varepsilon_2$, $j = 1, 2, \dots, n$,则点 $x^{(k+1)}$ 为近似最优解;否则,令 $\delta_j^{(k+1)} := \delta_j^{(k)}$, $j = 1, 2, \dots, n$,置k := k+1,返回2。



Frank-Wolfe方 法

$$\begin{cases} \min f(x) \\ s.t. & Ax \ge b \\ Ex = e \end{cases}$$
 线性约束

$$A_{m \times n}$$
, $r(A) = m$, $E_{l \times n}$, $f(x)$ 连续可微
令 $S = \{x \mid Ax \ge b, Ex = e, x \in R^n\}$

基本思想: 在每次迭代中,将目标函数f(x)线性化,通过解线性规划求得下降可行方向,进而沿此方向在可行域内做一维搜索。

清華大学

任取 $x^{(k)} \in S$,在 $x^{(k)}$ 处以f(x)的一阶Taylor展开式作为f(x)的线性逼近函数:

$$f_L(x) = f(x^{(k)}) + \nabla f(x^{(k)})^T (x - x^{(k)})$$

= $f(x^{(k)}) - \nabla f(x^{(k)})^T x^{(k)} + \nabla f(x^{(k)})^T x$

求解线性规划问题

$$\min_{S.t.} f_L(x) \\ s.t. \quad x \in S$$
 (2)
$$\begin{cases} \min \nabla f(x^{(k)})^T x \\ s.t. \quad x \in S \end{cases}$$

设问题(2)存在有限最优解y(k).

宣传華大学

求解(2)有下列两种情形之一:

- 1. \overline{A} $\nabla f(x^{(k)})^T (y^{(k)} x^{(k)}) = 0$,则停止迭代, $x^{(k)}$ 是 (1)的KKT点。则由yk最优知xk为最优解
- 2. 若∇ $f(x^{(k)})^T(v^{(k)}-x^{(k)}) \neq 0$,则必有

$$\nabla f(x^{(k)})^T (y^{(k)} - x^{(k)}) < 0$$

:: S是凸集, $:: \forall \lambda \in (0,1)$, 有

$$\nabla f(x^{(k)})^{T}(y^{(k)} - x^{(k)}) < 0$$

$$\Rightarrow y^{(k)} - x^{(k)} \Rightarrow x^{(k)} \Rightarrow$$

$$\lambda y^{(k)} + (1 - \lambda)x^{(k)} = x^{(k)} + \lambda(y^{(k)} - x^{(k)}) \in S$$

 $\Rightarrow v^{(k)} - x^{(k)} \to x^{(k)}$ 处的可行方向。

 $\Rightarrow v^{(k)} - x^{(k)} \to x^{(k)}$ 处的下降可行方向。

首華大学

定理: 设y^(k)是(2)的最优解,且满足

$$\nabla f(x^{(k)})^T (y^{(k)} - x^{(k)}) = 0$$
,则 $x^{(k)}$ 是(1)的 KKT 点。

证明: 由
$$\nabla f(x^{(k)})^T (y^{(k)} - x^{(k)}) = 0$$
,得
$$\nabla f(x^{(k)})^T y^{(k)} = \nabla f(x^{(k)})^T x^{(k)}$$

- $:: y^{(k)}$ 是(2)的最优解,且 $x^{(k)}$ ∈ S
- $\therefore x^{(k)}$ 是(2)的KKT点 $\Rightarrow \exists w \ge 0 (w \in E^m), v \in E^n$ 使得

$$\begin{cases} \nabla f(x^{(k)}) - A^T w - E^T v = 0 \\ w^T (Ax^{(k)} - b) = 0 \\ Ex^{(k)} = e \end{cases}$$

$$\begin{cases} \min \nabla f(x^{(k)})^T x \\ s.t. & x \in S \end{cases}$$

这也是(1)的KKT条件: $x^{(k)}$ 是(1)的KKT点。

(本大学

步骤

- 1.给定初始可行点 $x^{(1)}$, 允许误差 $\varepsilon > 0$, 置 k = 1。
- 2. 求解下面的线性规划问题得到最优解y(k)。

$$\begin{cases} \min \nabla f(x^{(k)})^T x \\ s.t. \quad x \in S \end{cases}$$

- 3. 若 $\nabla f(x^{(k)})^T (y^{(k)} x^{(k)}) \le \varepsilon$,则停止,得 $x^{(k)}$; 否则转 4。
- 4. 从 $x^{(k)}$ 出发,沿方向 $y^{(k)} x^{(k)}$ 进行一维搜索:

$$\begin{cases} \min \ f(x^{(k)} + \lambda(y^{(k)} - x^{(k)})) \\ s.t. \quad 0 \le \lambda \le 1 \end{cases}$$

得最优解 λ_k .

$$5.x^{(k+1)} = x^{(k)} + \lambda_k(y^{(k)} - x^{(k)})$$
, 置 $k := k+1$, 返回 2。

例:
$$\begin{cases} \min f(x) = 4x_1^2 + (x_2 - 2)^2 \\ s.t. -2 \le x_1 \le 2 \\ -1 \le x_2 \le 1 \end{cases}$$

解: 取初始点 $x^{(1)} = (-2, -1)^T$

第一次迭代
$$\nabla f(x^{(1)}) = (-16, -6)^T$$

$$\Re \begin{cases}
\min \nabla f(x^{(1)})^T x = -16x_1 - 6x_2 \\
s.t. -2 \le x_1 \le 2 \\
-1 \le x_2 \le 1
\end{cases}$$

得
$$y^{(1)} = (2,1)^T$$
.

清華大学

得
$$\lambda = 0.56$$

$$\therefore x^{(2)} = x^{(1)} + \lambda_1 (y^{(1)} - x^{(1)}) = (0.24, 0.12)^T$$

第二次迭 代

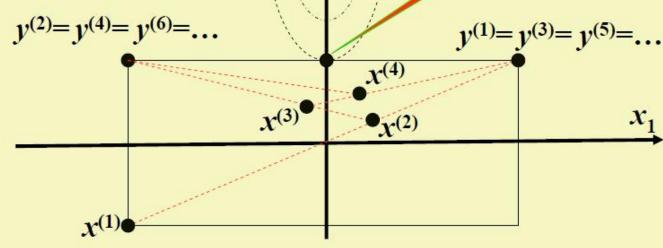
$$\Re \begin{cases}
\min \nabla f(x^{(2)})^T x = 1.92x_1 - 3.76x_2 \\
s.t. -2 \le x_1 \le 2 \\
-1 \le x_2 \le 1
\end{cases}$$

得
$$y^{(2)} = (-2,1)^T$$
.

首者大学

从 $x^{(2)}$ 出发,沿 $y^{(2)}-x^{(2)}$ 作一维搜索,得 $x^{(3)}$,

这样继续下去,得



Main idea of SQP

Successive quadratic programming (SQP) method was first proposed by Wilson (1963) and developed by Han (1976) and Powell (1977). Hence, it is also called the Wilson-Han-Powell method. This method is based on directly solving the Lagrange first-order necessary conditions of constrained optimization problems. The main idea is to solve the KKT system by using quasi-Newton method at each iteration and to reformulate the KKT system into a quadratic programming subproblem.

Lagrange-Newton method

Consider the following equality-constrained nonlinear programming problem

$$\min \qquad f(x) \\
s.t. \qquad h_i(x) = 0, \quad i \in \mathcal{E}.$$
(1)

The KKT system of (1) is written as

$$\nabla f(x) - \sum_{i \in \mathcal{E}} \lambda_i^T \nabla h_i(x) = 0$$

$$-h_i(x) = 0, \ i \in \mathcal{E}$$
(2)

and its Lagrange function is

$$L(x,\lambda) = f(x) - \lambda^T h(x).$$

We use Newton method to solve (2) and the current iterate is (x^k, λ^k) . Let $h_i(x), i \in \mathcal{E}$ group into a vector h(x) and let $\nabla h_i(x), i \in \mathcal{E}$ group into a

matrix H(x). Then the Newton-step procedure is

$$\nabla_x^2 L(x^k, \lambda^k)(x - x^k) - H(x^k)(\lambda - \lambda^k) = -(\nabla f(x^k) - H(x^k)\lambda^k)$$
$$H(x^k)(x - x^k) = -h(x^k),$$

which can be rewritten as

$$\frac{\nabla_x^2 L(x^k, \lambda^k) d - H(x^k) \lambda = -\nabla f(x^k)}{H(x^k) d = -h(x^k).}$$
(3)

Solving the above system of linear equations, we obtain (d^k, λ^{k+1}) and then derive a new iterate $(x^{k+1} = x^k + d^k, \lambda^{k+1})$. This is called Lagrange-Newton method.

SQP method

直接求(3)不方便因此将其转化为别的问题的KKT条件

Obviously, (3) is equivalent to the following equality-constrained quadratic programming problem

min
$$f(x^k) + \nabla f(x^k)^T d + \frac{1}{2} d^T \nabla_x^2 L(x^k, \lambda^k) d$$
s.t.
$$h_i(x^k) + \nabla h_i(x^k)^T d = 0, \quad i \in \mathcal{E}$$

Thus, we can extend the Lagrange-Newton method to solve the mixed constrained nonlinear programming:

min
$$f(x)$$

$$s.t. h_i(x) = 0, i \in \mathcal{E}$$

$$h_i(x) \ge 0, i \in \mathcal{I}.$$

$$(4)$$

Here, the direction d^k is obtained by solving the quadratic program

min
$$f(x^k) + \nabla f(x^k)^T d + \frac{1}{2} d^T \nabla_x^2 L(x^k, \lambda^k) d$$
s.t.
$$h_i(x^k) + \nabla h_i(x^k)^T d = 0, \quad i \in \mathcal{E}$$

$$h_i(x^k) + \nabla h_i(x^k)^T d \ge 0, \quad i \in \mathcal{I},$$

where

$$L(x,\lambda) = f(x) - \sum_{i \in \mathcal{E} \cup \mathcal{I}} \lambda_i h_i(x).$$

If $d^k = 0$, then x^k is a KKT point of (4).

There are two disadvantages: to compute the Hessian matrix of the Lagrange function $L(x,\lambda)$; to guarantee the positive definiteness of the Hessian matrix. Hence, we use the approximation of $\nabla_x^2 L(x^k,\lambda^k)$ based on the quasi-Newton method. Thus, we describe the SQP method as follows.

- **0.** Given a starting point $x^0 \in \mathcal{R}^n$ and an initial positive definite matrix B_0 . Set k=0.
- **1.** Compute d^k by solving the quadratic program

$$\min \qquad f(x^k) + \nabla f(x^k)^T d + \frac{1}{2} d^T B_k d$$

$$Q(x^k, B_k) \quad s.t. \qquad h_i(x^k) + \nabla h_i(x^k)^T d = 0, \quad i \in \mathcal{E}$$

$$h_i(x^k) + \nabla h_i(x^k)^T d \ge 0, \quad i \in \mathcal{I}.$$

函数的目标函数(此时相当于

无约束)来确定

- **2.** If $d^k=0$, stop; x^k is a KKT point of (4). Otherwise, let $x^{k+1}=x^k+\alpha_k d^k$ where α_k is determined by a certain line search.
- 3. Update B_k to form B_{k+1} such that B_{k+1} is positive definite. Set k=k+1 and go to Step 1. 这里的精确搜索可以由加上罚

可以利用这个投影矩阵,考虑有约束的规划的搜索

Update B_{k}

Let

$$L(x,\lambda) = f(x) - \sum_{i \in \mathcal{E} \cup \mathcal{I}} \lambda_i h_i(x),$$

where $\lambda_i,\ i\in\mathcal{E}\cup\mathcal{I}$ are optimal multipliers of Q(x,B). Generally, we use the BFGS formulation to update B_k

$$B_{k+1} = B_k + \frac{\gamma_k \gamma_k^T}{\gamma_k^T s_k} - \frac{B_k s_k s_k^T B_k}{s_k^T B_k s_k},$$

where

$$s_k = x^{k+1} - x^k, \quad \gamma_k = \nabla_x L(x^{k+1}, \lambda^{k+1}) - \nabla_x L(x^k, \lambda^k).$$

理论上需要 グバック 但是数値上需要大于某个定値

To guarantee the positive definiteness, Powell (1978) introduced the following formulation:

$$B_{k+1} = B_k + \frac{\eta_k \eta_k^T}{\eta_k^T s_k} - \frac{B_k s_k s_k^T B_k}{s_k^T B_k s_k},$$

where $\eta_k = \theta_k \gamma_k + (1 - \theta_k) B_k s_k$, and

$$\theta_k = \left\{ \begin{array}{ll} 1, & \text{if } s_k^T \gamma_k \geq 0.2 s_k^T B_k s_k \\ \frac{0.8 s_k^T B_k s_k}{s_k^T B_k s_k - s_k^T \gamma_k}, & \text{otherwise} \end{array} \right.$$

Clearly, $s_k^T \eta_k \ge 0.2 s_k^T B_k s_k$. When B_k is positive definite, so is B_{k+1} from Theorem 6 in Lecture Note #9 (Page 67).

Remarks

- SQP method is globally convergent and has the rate of superlinear convergence under some suitable assumptions.
- Maratos observed that the unit step-size is not reached even at a superlinearly convergent step. Thus, the rate of convergence of the SQP method is only linear. This phenomenon is called "Maratos effect".