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13. Generalized distances and mirror descent

- Bregman distance
- properties
- Bregman proximal mapping
- mirror descent

Motivation: proximal gradient method

proximal gradient step for minimizing f(x) = g(x) + h(x) (page 4.4):

$$x_{k+1} = \operatorname{prox}_{t_k h}(x_k - t_k \nabla g(x_k))$$

$$= \operatorname{argmin}_{u} \left(h(u) + g(x_k) + \nabla g(x_k)^T (u - x_k) + \frac{1}{2t_k} ||u - x_k||_2^2 \right)$$

Interpretation: quadratic term represents

- a penalty that forces x_{k+1} to be close to x_k , where linearization of g is accurate
- ullet an approximation of the error term in the linearization of g at x_k

Generalized proximal gradient method

replace $\frac{1}{2}||u-x||_2^2$ with a generalized distance d(u,x):

$$x_{k+1} = \underset{u}{\operatorname{argmin}} \left(h(u) + g(x_k) + \nabla g(x_k)^T (u - x_k) + \frac{1}{t_k} d(u, x_k) \right)$$

Potential benefits

1. "pre-conditioning": use a more accurate model of g(u) around x, ideally

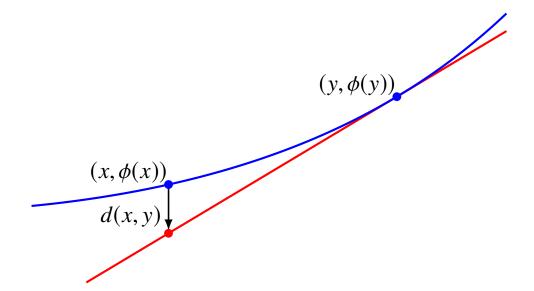
$$\frac{1}{t_k}d(u,x_k) \approx g(u) - g(x_k) - \nabla g(x_k)^T (u - x_k)$$

2. make the generalized proximal mapping (minimizer *u*) easier to compute goal of 1 is to reduce number of iterations; goal of 2 is to reduce cost per iteration

Bregman distance

$$d(x, y) = \phi(x) - \phi(y) - \nabla \phi(y)^{T} (x - y)$$

- ϕ is convex and continuously differentiable on $int(dom \phi)$
- ullet domain of ϕ may include its boundary or a subset of its boundary
- we define the domain of d as $dom d = dom \phi \times int(dom \phi)$
- ullet ϕ is called the *kernel function* or *distance-generating function*



other properties of ϕ will be required but mentioned explicitly (e.g., strict convexity)

Immediate properties

$$d(x, y) = \phi(x) - \phi(y) - \nabla \phi(y)^{T} (x - y)$$

- d(x, y) is convex in x for fixed y
- $d(x, y) \ge 0$, with equality if x = y
- if ϕ is strictly convex, then d(x, y) = 0 only if x = y
- $d(x, y) \neq d(y, x)$ in general

to emphasize lack of symmetry, d is also called a $\it directed \ \it distance$ or $\it divergence$

Squared Euclidean distance (with dom $\phi = \mathbf{R}^n$)

$$\phi(x) = \frac{1}{2}x^T x,$$
 $\nabla \phi(x) = x,$ $d(x, y) = \frac{1}{2}||x - y||_2^2$

General quadratic kernel (with dom $\phi = \mathbf{R}^n$)

$$\phi(x) = \frac{1}{2}x^T A x, \qquad \nabla \phi(x) = A x, \qquad d(x, y) = \frac{1}{2}(x - y)^T A (x - y)$$

- *A* is symmetric positive definite
- in some applications, A is positive semidefinite, but not positive definite

Relative entropy (with dom $\phi = \mathbf{R}_{+}^{n}$)

$$\phi(x) = \sum_{i=1}^{n} x_i \log x_i, \qquad \nabla \phi(x) = \begin{bmatrix} \log x_1 + 1 \\ \vdots \\ \log x_n + 1 \end{bmatrix}$$

$$d(x,y) = \sum_{i=1}^{n} \left(x_i \log \frac{x_i}{y_i} - x_i + y_i \right)$$

Logistic loss divergence (with dom $\phi = [0, 1]^n$)

$$\phi(x) = \sum_{i=1}^{n} (x_i \log x_i + (1 - x_i) \log(1 - x_i)), \qquad \nabla \phi(x) = \begin{bmatrix} \log(x_1/(1 - x_1)) \\ \vdots \\ \log(x_n/(1 - x_n)) \end{bmatrix}$$

$$d(x,y) = \sum_{i=1}^{n} \left(x_i \log \frac{x_i}{y_i} + (1 - x_i) \log \frac{1 - x_i}{1 - y_i} \right)$$

Hellinger divergence (with dom $\phi = [-1, 1]^n$)

$$\phi(x) = -\sum_{i=1}^{n} \sqrt{1 - x_i^2}, \qquad \nabla \phi(x) = \begin{bmatrix} x_1/\sqrt{1 - x_1^2} \\ \vdots \\ x_n/\sqrt{1 - x_n^2} \end{bmatrix}$$

$$d(x,y) = \sum_{i=1}^{n} \left(-\sqrt{1 - x_i^2} + \frac{1 - x_i y_i}{\sqrt{1 - y_i^2}} \right)$$

Logarithmic barrier (with dom $\phi = \mathbf{R}_{++}^n$)

$$\phi(x) = -\sum_{i=1}^{n} \log x_i, \qquad \nabla \phi(x) = \begin{bmatrix} -1/x_1 \\ \vdots \\ -1/x_n \end{bmatrix}, \qquad d(x,y) = \sum_{i=1}^{n} \left(\frac{x_i}{y_i} - \log \frac{x_i}{y_i} - 1 \right)$$

d(x, y) is sometimes called *Itakura–Saito* divergence

Inverse barrier (with dom $\phi = \mathbf{R}_{++}^n$)

$$\phi(x) = \sum_{i=1}^{n} \frac{1}{x_i}, \qquad \nabla \phi(x) = \begin{bmatrix} -1/x_1^2 \\ \vdots \\ -1/x_n^2 \end{bmatrix}, \qquad d(x,y) = \sum_{i=1}^{n} \frac{1}{y_i} \left(\sqrt{\frac{x_i}{y_i}} - \sqrt{\frac{y_i}{x_i}} \right)^2$$

Bregman distances for symmetric matrices

$$d(X,Y) = \phi(X) - \phi(Y) - \operatorname{tr}(\nabla \phi(Y)(X - Y))$$

- kernel ϕ is a convex function on S^n , differentiable on int $(\text{dom }\phi)$
- domain of d is dom $d = \text{dom } \phi \times \text{int } (\text{dom } \phi)$

Relative entropy (with dom $\phi = \mathbf{S}_{++}^n$)

$$\phi(X) = -\log \det X, \qquad \nabla \phi(X) = -X^{-1}$$

$$d(X,Y) = tr(XY^{-1}) - \log \det(XY^{-1}) - n$$

- d(X,Y) is relative entropy between normal distributions N(0,X) and N(0,Y)
- also known as Kullback-Leibler divergence

Bregman distances for symmetric matrices

Matrix entropy (with dom $\phi = \mathbf{S}_{++}^n$):

$$\phi(X) = \operatorname{tr}(X \log X), \qquad \nabla \phi(X) = I + \log X$$

$$d(X,Y) = \operatorname{tr}(X \log X - X \log Y - X + Y)$$

• matrix logarithm $\log X$ is defined as

$$\log X = \sum_{i=1}^{n} (\log \lambda_i) q_i q_i^T$$

if X has eigendecomposition $X = \sum_{i} \lambda_{i} q_{i} q_{i}^{T}$

• d(X,Y) is also known as *quantum relative entropy*

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Three-point identity

for all $x \in \text{dom } \phi$ and $y, z \in \text{int}(\text{dom } \phi)$,

$$d(x,z) = d(x,y) + d(y,z) + (\nabla \phi(y) - \nabla \phi(z))^T (x - y)$$

- easily verified by substituting the definition of d
- if *d* is not symmetric, order of the arguments of *d* in the identity matters
- generalizes the familiar identity for squared Euclidean distance:

$$\frac{1}{2}||x - z||_2^2 = \frac{1}{2}||x - y||_2^2 + \frac{1}{2}||y - z||_2^2 + (y - z)^T(x - y)$$

Strongly convex kernel

we will sometimes assume that ϕ is strongly convex (page 1.19):

$$\phi(x) \ge \phi(y) + \nabla \phi(y)^T (x - y) + \frac{\mu}{2} ||x - y||^2$$

- $\mu > 0$ is strong convexity constant of ϕ for the norm $\|\cdot\|$
- for twice differentiable ϕ , this is equivalent to

$$v^T \nabla^2 \phi(x) v \ge \mu \|v\|^2$$
 for all $x \in \operatorname{int}(\operatorname{dom} \phi)$ and v

(see page 1.18)

• strong convexity of ϕ implies that

$$d(x,y) = \phi(x) - \phi(y) - \nabla \phi(y)^{T} (x - y)$$

$$\geq \frac{\mu}{2} ||x - y||^{2}$$

Regularization with Bregman distance

for given $y \in \operatorname{int}(\operatorname{dom} \phi)$ and convex f, consider

minimize
$$f(x) + d(x, y)$$

- equivalently, minimize $f(x) + \phi(x) \nabla \phi(y)^T x$
- feasible set is $dom f \cap dom \phi$

Optimality condition: $\hat{x} \in \text{dom } f \cap \text{int}(\text{dom } \phi)$ is optimal if and only if

$$f(x) + d(x, y) \ge f(\hat{x}) + d(\hat{x}, y) + d(x, \hat{x}) \quad \text{for all } x \in \text{dom } f \cap \text{dom } \phi$$
 (1)

Equivalent optimality condition: $\hat{x} \in \text{dom } f \cap \text{int}(\text{dom } \phi)$ is optimal if and only if

$$\nabla \phi(y) - \nabla \phi(\hat{x}) \in \partial f(\hat{x}) \tag{2}$$

Proof: we derive optimality conditions for the problem

minimize
$$g(x) + \phi(x)$$
 (3)

with g convex, and apply the results to $g(x) = f(x) - \nabla \phi(y)^T x$

• optimality condition: $\hat{x} \in \text{dom } g \cap \text{int } (\text{dom } \phi)$ is optimal for (3) if and only if

$$g(x) \ge g(\hat{x}) - \nabla \phi(\hat{x})^T (x - \hat{x})$$
 for all $x \in \text{dom } g \cap \text{dom } \phi$ (4)

combined with the 3-point identity this gives the optimality condition (1)

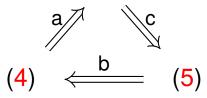
• equivalent optimality condition: $\hat{x} \in \text{dom } g \cap \text{int } (\text{dom } \phi)$ is optimal if and only if

$$-\nabla\phi(\hat{x})\in\partial g(\hat{x})\tag{5}$$

applied to $g(x) = f(x) - \nabla \phi(y)^T x$ this gives the optimality condition (2)

Proof:

optimality of \hat{x}



• implication a follows from convexity of ϕ : if (4) holds, then for all feasible x,

$$g(x) + \phi(x) \ge g(\hat{x}) + \phi(x) - \nabla \phi(\hat{x})^T (x - \hat{x}) \ge g(\hat{x}) + \phi(\hat{x})$$

• implication b: by definition of subgradient, (5) can be written as

$$g(x) \ge g(\hat{x}) - \nabla \phi(\hat{x})^T (x - \hat{x})$$
 for all $x \in \text{dom } g$

• we prove c by contradiction: suppose that for some $x \in \text{dom } g$

$$g(x) < g(\hat{x}) - \nabla \phi(\hat{x})^T (x - \hat{x})$$

define $v = x - \hat{x}$; for small positive t, by convexity of g and Taylor's theorem,

$$g(\hat{x} + tv) + \phi(\hat{x} + tv) \leq g(\hat{x}) + t(g(x) - g(\hat{x})) + \phi(\hat{x} + tv)$$

$$= g(\hat{x}) + \phi(\hat{x}) + t(g(x) - g(\hat{x}) + \nabla\phi(\hat{x})^{T}v) + O(t^{2})$$

$$< g(\hat{x}) + \phi(\hat{x})$$

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Bregman proximal mapping

for convex f and Bregman kernel ϕ , define

$$\operatorname{prox}_{f}^{d}(y, a) = \operatorname{argmin}_{x} \left(f(x) + a^{T}x + d(x, y) \right)$$
$$= \operatorname{argmin}_{x} \left(f(x) + (a - \nabla \phi(y))^{T}x + \phi(x) \right)$$

- first argument y must be in int $(\text{dom }\phi)$
- second argument a can take any value
- we'll use this only if for every y and a, a unique minimizer $x \in \operatorname{int}(\operatorname{dom} \phi)$ exists

Example: quadratic kernel

$$\phi(x) = \frac{1}{2} ||x||_2^2, \qquad d(x, y) = \frac{1}{2} ||x - y||_2^2$$

Bregman proximal mapping can be expressed in terms of standard $prox_f$:

$$\operatorname{prox}_{f}^{d}(y, a) = \operatorname{argmin}_{x} \left(f(x) + a^{T}x + d(x, y) \right)$$
$$= \operatorname{argmin}_{x} \left(f(x) + a^{T}x + \frac{1}{2} ||x - y||_{2}^{2} \right)$$
$$= \operatorname{prox}_{f}(y - a)$$

closedness of f ensures existence and uniqueness (see page 6.2)

Example: relative entropy

$$\phi(x) = \sum_{i=1}^{n} x_i \log x_i, \qquad d(x, y) = \sum_{i=1}^{n} (x_i \log(x_i/y_i) - x_i + y_i)$$

- we take $f = \delta_C$, the indicator of probability simplex $C = \{x \geq 0 \mid \mathbf{1}^T x = 1\}$
- Bregman proximal mapping is

$$\operatorname{prox}_{f}^{d}(y, a) = \underset{\mathbf{1}^{T} x = 1}{\operatorname{argmin}} (a^{T} x + \sum_{i=1}^{n} x_{i} \log(x_{i}/y_{i}))$$

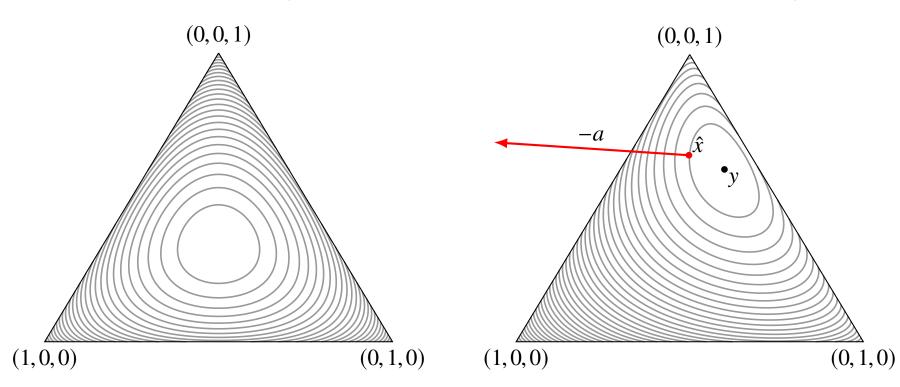
$$= \frac{1}{\sum_{i=1}^{n} y_{i} e^{-a_{i}}} \begin{bmatrix} y_{1} e^{-a_{1}} \\ \vdots \\ y_{n} e^{-a_{n}} \end{bmatrix}$$

• for every y > 0 and a, minimizer in the definition exists, is unique, and positive

Example: relative entropy



Contour lines of d(x, y)



right-hand figure shows

$$\hat{x} = \text{prox}_f^d(y, a) = \operatorname{argmin}(a^T x + d(x, y))$$

for
$$y = (0.1, 0.3, 0.6)$$
 and $a = (-0.540, 0.585, -0.045)$

Optimality condition

apply the optimality conditions for Bregman-regularized problem (page 13.14) to

$$\operatorname{prox}_{f}^{d}(y, a) = \operatorname{argmin}_{x} \left(f(x) + a^{T} x + d(x, y) \right)$$

suppose $\hat{x} \in \text{dom } f \cap \text{int } (\text{dom } \phi)$

• first condition: $\hat{x} = \text{prox}_f^d(y, a)$ if and only if

$$f(x) + a^T x + d(x, y) \ge f(\hat{x}) + a^T \hat{x} + d(\hat{x}, y) + d(x, \hat{x})$$

for all $x \in \text{dom } f \cap \text{dom } \phi$

• second condition: $\hat{x} = \text{prox}_f^d(y, a)$ if and only if

$$\nabla \phi(y) - \nabla \phi(\hat{x}) - a \in \partial f(\hat{x})$$

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Mirror descent

minimize
$$f(x)$$

subject to $x \in C$

- f is a convex function, C is a convex subset of dom f
- ullet we assume f is subdifferentiable on C

Algorithm: choose $x_0 \in C \cap \operatorname{int}(\operatorname{dom} \phi)$ and repeat

$$x_{k+1} = \underset{x \in C}{\operatorname{argmin}} \left(t_k g_k^T x + d(x, x_k) \right), \quad k = 0, 1, \dots$$

 g_k is any subgradient of f at x_k

update can be written as $x_{k+1} = \text{prox}_{\delta_C}^d(x_k, t_k g_k)$ where δ_C is indicator of C

Mirror descent with quadratic kernel

$$x_{k+1} = \underset{x \in C}{\operatorname{argmin}} \left(t_k g_k^T x + d(x, x_k) \right)$$

for $d(x, y) = \frac{1}{2}||x - y||_2^2$, this is the projected subgradient method:

$$x_{k+1} = \underset{x \in C}{\operatorname{argmin}} \left(t g_k^T x + \frac{1}{2} ||x - x_k||_2^2 \right)$$

$$= \underset{x \in C}{\operatorname{argmin}} \frac{1}{2} ||x - x_k + t_k g_k||_2^2$$

$$= P_C(x_k - t_k g_k)$$

where P_C is Euclidean projection on C

Assumptions

- problem on page 13.22 has optimal value f^* , optimal solution $x^* \in C \cap \text{dom } \phi$
- ullet f is Lipschitz continuous on C with respect to some norm $\|\cdot\|$

$$|f(x) - f(y)| \le G||x - y||$$
 for all $x, y \in C$

this is equivalent to $||g||_* \le G$ for all $x \in C$ and $g \in \partial f(x)$ (proof extends proof for Euclidean norm on page 3.4)

• ϕ is 1-strongly convex on C, with respect to the same norm $\|\cdot\|$:

$$d(x, y) \ge \frac{1}{2} ||x - y||^2$$
 for all $x \in \text{dom } \phi \cap C$ and $y \in \text{int}(\text{dom } \phi) \cap C$

Analysis

• apply optimality condition on page 13.21 with $x = x^*$, $y = x_i$, $\hat{x} = x_{i+1}$:

$$d(x^{\star}, x_{i+1}) \leq d(x^{\star}, x_{i}) - d(x_{i+1}, x_{i}) + t_{i}g_{i}^{T}(x_{i} - x_{i+1}) + t_{i}g_{i}^{T}(x^{\star} - x_{i})$$

$$\leq d(x^{\star}, x_{i}) - d(x_{i+1}, x_{i}) + ||t_{i}g_{i}||_{*}||x_{i+1} - x_{i}|| + t_{i}g_{i}^{T}(x^{\star} - x_{i})$$

$$\leq d(x^{\star}, x_{i}) - d(x_{i+1}, x_{i}) + \frac{1}{2}||x_{i+1} - x_{i}||^{2} + \frac{1}{2}||t_{i}g_{i}||_{*}^{2} + t_{i}g_{i}^{T}(x^{\star} - x_{i})$$

last step is arithmetic-geometric mean inequality

apply strong convexity of kernel and definition of subgradient:

$$d(x^*, x_{i+1}) \le d(x^*, x_i) + \frac{1}{2} ||t_i g_i||_*^2 + t_i (f^* - f(x_i))$$

• define $f_{\text{best},k} = \min_{i=0,\dots,k} f(x_i)$ and combine inequalities for $i=0,\dots,k$:

$$(\sum_{i=0}^{k} t_i)(f_{\text{best},k} - f^*) \leq d(x^*, x_0) - d(x^*, x_{k+1}) + \frac{1}{2} \sum_{i=0}^{k} ||t_i g_i||_*^2$$

$$\leq d(x^*, x_0) + \frac{1}{2} \sum_{i=0}^{k} ||t_i g_i||_*^2$$

Step size selection

$$f_{\text{best},k} - f^{\star} \leq \frac{d(x^{\star}, x_0)}{\sum_{i=0}^{k} t_i} + \frac{\sum_{i=0}^{k} ||t_i g_i||_*^2}{2\sum_{i=0}^{k} t_i} \leq \frac{d(x^{\star}, x_0)}{\sum_{i=0}^{k} t_i} + \frac{G^2 \sum_{i=0}^{k} t_i^2}{2\sum_{i=0}^{k} t_i}$$

• diminishing step size: $f_{\text{best},k} \to f^*$ if

$$t_i \to 0, \qquad \sum_{i=0}^{\infty} t_i = \infty$$

(see page 3.7)

• optimal step size for fixed number of iterations k, if we know that $d(x^*, x_0) \leq D$:

$$t_i = \frac{\sqrt{2D}}{\|g_i\|_* \sqrt{k+1}}, \qquad f_{\text{best},k} \le \frac{G\sqrt{2D}}{\sqrt{k+1}}$$

(see page 3.10)

Entropic mirror descent

apply mirror descent with relative entropy distance and

$$C = \{x \in \mathbf{R}^n \mid x \ge 0, \ \mathbf{1}^T x = 1\}$$

Algorithm: choose $x_0 > 0$, $\mathbf{1}^T x_0 = 1$, and repeat

$$x_{k+1} = \frac{1}{s^T x_k} (s \circ x_k)$$
 where $s = (e^{-t_k g_{k,1}}, \dots, e^{-t_k g_{k,n}})$

- g_k is any subgradient of f at x_k
- o denotes component-wise vector product

Convergence

in the analysis on page 13.26

• if we choose $x_0 = (1/n)\mathbf{1}$, then we can take $D = \log n$:

$$d(x^*, x_0) = \log n + \sum_{i=1}^n x_i^* \log x_i^* \le \log n$$

• $\phi(x) = \sum_{i} x_{i} \log x_{i}$ is 1-strongly convex for $\|\cdot\|_{1}$ on C: by Cauchy–Schwarz,

$$v^T \nabla^2 \phi(x) v = \sum_{i=1}^n \frac{v_i^2}{x_i} \ge ||v||_1^2$$
 if $x > 0$, $\mathbf{1}^T x = 1$

with optimal step size for k iterations,

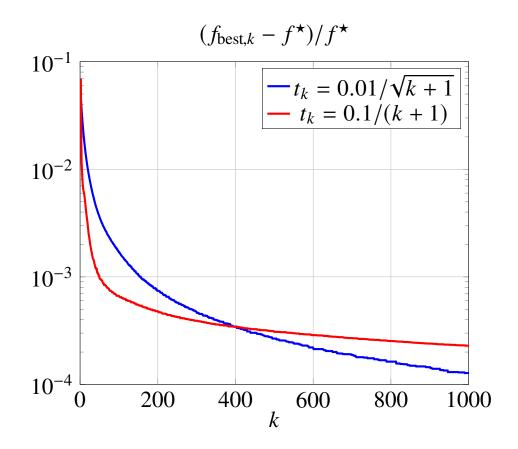
$$f_{\text{best},k} \le \frac{G\sqrt{2\log n}}{\sqrt{k+1}}$$

where G is Lipschitz constant of f for $\|\cdot\|_1$ -norm

minimize
$$||Ax - b||_1$$

subject to $x \ge 0$, $\mathbf{1}^T x = 1$

- subgradient $g = A^T \operatorname{sign}(Ax b)$, so $||g||_{\infty} \le G = \max_j \sum_i |A_{ij}|$
- example with randomly generated $A \in \mathbf{R}^{1000 \times 500}$, $b \in \mathbf{R}^{1000}$



References

Generalized distances

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