CS 340

#10/11: Deadlock and Producer-Consumer

Computer Systems | Sept. 27, 2022 · Wade Fagen-Ulmschneider

Solving Deadlock

On Thursday, we explored the four necessary conditions for deadlock. In the context of the dining philosophers problem, how do we remove each of the four?

- 1. Mutual Exclusion
- 2. Circular Wait
- 3. Hold and Wait
- 4. No Preemption

Deadlock Solution Considerations

- 1. Fairness:
- 2. Livelock:

```
10/producer-consumer-2.c
   #define THINGS_MAX 10
   #define THREAD_CT 5
   int things[THINGS_MAX];
   int things_ct = 0;
11
12
13
```

```
void *producer(void *vptr) {
                                    38
                                       void *consumer(void *vptr) {
                                    39
26
    while (1) {
                                        while (1) {
28
     assert(things_ct <
                                         assert(things_ct > 0);
                      THINGS_MAX):
                                    42
29
                                    43
                                          // Consume a thing:
30
     // Produce a thing:
                                    44
                                          things_ct--;
31
     things[things_ct] =
                                    45
                                          int value =
                    rand() % 100;
                                                   things[things_ct];
32
     printf("Produced [%d]: %d
                                         printf("Consumed [%d]: %d
        -> ", things_ct,
                                                     <- ", things_ct,
              things[things_ct]);
                                                           value):
33
     things_ct++:
                                         print_things_as_list();
34
     print_things_as_list();
                                    48
35
36
                                    49
```

```
52
   int main() {
53
     int i:
54
55
     // Create `thread_ct` threads of each producer and consumer:
     pthread_t tid_consumer[THREAD_CT];
56
57
     pthread_t tid_producer[THREAD_CT];
58
     for (i = 0; i < THREAD_CT; i++) {</pre>
       pthread_create(&tid_consumer[i], NULL, producer, NULL);
59
60
       pthread_create(&tid_producer[i], NULL, consumer, NULL);
61
62
63
     // Join threads:
64
     for (i = 0; i < THREAD_CT; i++) {</pre>
65
       pthread_join(tid_consumer[i], NULL);
66
       pthread_join(tid_producer[i], NULL);
67
68 }
```

Synchronization Primitives

In programming, a key synchronization primitive has evolved to become common features of many programming languages.

Primitive:

- Allow asynchronous execution until a ______.
- In JavaScript and Python, this is _____.

```
10/barrier.c
   typedef pthread_t * promise_t;
10
   void *idle_task(void *vptr) {
    printf("idle_task is running.\n");
12
13
     printf("idle_task has finished.\n");
14
     return NULL;
15
16
17
   promise_t async_task(void*(* task)(void *), void *arg) {
     pthread_t *tidptr = malloc(sizeof(pthread_t));
     pthread_create(tidptr, NULL, task, arg);
19
     return tidptr;
20
21
22
   void *async_wait(promise_t promise) {
     void *result;
25
     pthread_join(*promise, &result);
26
     return result:
27 }
28
29
   int main() {
     promise_t p = async_task(idle_task, NULL);
     printf("main thread is running at the same time...\n");
31
32
33
     printf("main thread has running again...\n");
34
     return 0;
35 }
```

General Pattern:

- An "async" call returns a ______, not a return value.
- The "await" call and then .

Multi-Threaded Uses

There are several different reasons you will come across the use of multi-threaded applications:

(1): Concurrent Compute

(2): User Interactivity

(3): Responsiveness in Requests

(4): **NOT** for Security

(5): **NOT** for Isolation