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#### 1 Miscellaneous

# 1.1 Day of Date

```
// 0-based
const vector<int> T = {
  0, 3, 2, 5, 0, 3,
  5, 1, 4, 6, 2, 4
}
int day(int d, int m, int y) {
  y -= (m < 3);
  return (y + y / 4 - y / 100 + y / 400 + T[m - 1] + d) % 7;
}</pre>
```

# 1.2 Number of Days since 1-1-1

```
int rdn(int d, int m, int y) {
  if(m < 3)--y, m += 12;
  return 365 * y + y / 4 - y / 100 + y / 400 + (153 * m - 457) / 5 + d - 306;
}</pre>
```

## 1.3 Enumerate Subsets of a Bitmask

```
int x = 0;
do {
    // do stuff with the bitmask here
    x = (x + 1 + ~m) & m;
} while(x != 0);
```

# 1.4 Josephus Problem

```
ll josephus(ll n, ll k) \{ // O(k \log n) \}
  if(n == 1)
    return 0;
  if(k == 1)
    return n - 1;
  if(k > n)
    return (josephus(n - 1, k) + k) % n;
  ll cnt = n / k;
  ll res = josephus(n - cnt, k);
  res -= n % k;
  if(res < 0)
    res += n;
    res += res / (k-1);
  return res;
int josephus(int n, int k) { // O(n)
  int res = 0;
  for(int i = 1; i <= n; ++i)
    res = (res + k) \% i;
  return res + 1;
```

## 1.5 Random Primes

36671 74101 724729 825827 924997 1500005681 2010408371 2010405347

# 1.6 RNG

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```
// RNG - rand_int(min, max), inclusive
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
template<class T>
T rand_int(T mn, T mx) {
 return uniform_int_distribution<T>(mn, mx)(rng);
```

## 2 Data Structures

## 2.1 2D Segment Tree

```
struct Segtree2D
 struct Segtree
   struct node {
     int l, r, val;
     node* lc, *rc;
     node(int _l, int _r, int _val = INF) : l(_l), r(_r), val(_val),
       lc(NULL), rc(NULL) {}
   typedef node* pnode;
   pnode root;
   Segtree(int l, int r) {
     root = new node(l, r);
   void update(pnode& nw, int x, int val) {
     int l = nw - > l, r = nw - > r, mid = (l + r) / 2;
     if(l == r)
       nw->val = val;
     else {
       assert(l <= x && x <= r);
       pnode& child = x <= mid ? nw->lc : nw->rc;
       if(!child)
         child = new node(x, x, val);
       else if(child->l <= x && x <= child->r)
         update(child, x, val);
       else {
         do
           if(x \le mid)
             r = mid;
           else
             l = mid + 1;
           mid = (l + r) / 2;
         } while((x <= mid) == (child->l <= mid));</pre>
         pnode nxt = new node(l, r);
         if(child->l <= mid)</pre>
           nxt->lc = child;
         else
           nxt->rc = child;
         child = nxt;
         update(nxt, x, val);
       nw->val = min(nw->lc ? nw->lc->val : INF,
                      nw->rc ? nw->rc->val : INF);
   int query(pnode& nw, int x1, int x2) {
     if(!nw)
       return INF;
      int& l = nw->l, &r = nw->r;
     if(r < x1 \mid | x2 < l)
       return INF;
     if(x1 <= l && r <= x2)
       return nw->val;
      int ret = min(query(nw->lc, x1, x2),
                    query(nw->rc, x1, x2));
     return ret;
   void update(int x, int val) {
     assert(root->l \le x \&\& x \le root->r);
```

```
update(root, x, val);
  int query(int l, int r) {
    return query(root, l, r);
};
struct node {
  int l, r;
  Segtree y;
  node* lc, *rc;
  node(int _l, int _r) : l(_l), r(_r), y(0, MAX),
   lc(NULL), rc(NULL) {}
typedef node* pnode;
pnode root;
Segtree2D(int l, int r) {
  root = new node(l, r);
void update(pnode& nw, int x, int y, int val) {
  int& l = nw - > l, &r = nw - > r, mid = (l + r) / 2;
  if(l == r)
   nw->y.update(y, val);
  else {
    if(x \le mid) \{
      if(!nw->lc)
        nw->lc = new node(l, mid);
     update(nw->lc, x, y, val);
    } else {
      if(!nw->rc)
        nw->rc = new node(mid + 1, r);
     update(nw->rc, x, y, val);
    val = min(nw->lc ? nw->lc->y.query(y, y) : INF,
              nw->rc ? nw->rc->y.query(y, y) : INF);
    nw->y.update(y, val);
int query(pnode& nw, int x1, int x2, int y1, int y2) {
  if(!nw)
   return INF;
  int& l = nw->l, &r = nw->r;
  if(r < x1 \mid | x2 < l)
   return INF;
  if(x1 <= l && r <= x2)
   return nw->y.query(y1, y2);
  int ret = min(query(nw->lc, x1, x2, y1, y2),
                query(nw->rc, x1, x2, y1, y2));
  return ret;
void update(int x, int y, int val) {
  assert(root->l <= x && x <= root->r);
  update(root, x, y, val);
int query(int x1, int x2, int y1, int y2) {
  return query(root, x1, x2, y1, y2);
```

# 2.2 Fenwick RU-RQ

```
void updtRL(int l, int r, ll val) {
 updt(BIT1, l, val), updt(BIT1, r + 1, -val);
 updt(BIT2, l, val * (l - 1)), updt(BIT2, r + 1, -val * r);
ll query(int k) {
 return que(BIT1, k) * k - que(BIT2, k);
```

# 2.3 Heavy-Light Decomposition

```
//vertex value, klo edge value, turunin nilainya ke vertex bawahnya
class HLD {
public:
    static const int N = 100005;
    int seg[N*4], in[N], out[N], sz[N], dep[N], par[N], root[N], idx[N], val[N], t, n;
    vector<int> edge[N];
    //idx -> actual index, in -> visited time
    HLD():t(0) {}
    HLD(int n):n(n) {
        root[1] = par[1] = 1;
    void upd(int id, int l, int r, int x, int v) {
        if(l == r) +
            seg[id] = val[x] = v;
            return ;
        int m = l + r \gg 1;
        if(in[x] <= m) upd(id<<1, l, m, x, v);</pre>
        else upd(id<<1|1, m + 1, r, x, v);
        seg[id] = seg[id << 1] ^ seg[id << 1|1];
    int que(int id, int l, int r, int tl, int tr) {
        if(r < tl || l > tr) return 0;
        if(tl <= l && r <= tr) return seg[id];</pre>
        int m = l + r >> 1:
        return que(id<<1, l, m, tl, tr) ^ que(id<<1|1, m + 1, r, tl, tr);
    void build(int id, int l, int r) {
        if(l == r) {
           seg[id] = val[idx[l]];
           return ;
        int m = l + r \gg 1;
       build(id<<1, l, m);
       build(id<<1|1, m + 1, r);
       seg[id] = (seg[id<<1] ^ seg[id<<1|1]);
    void dfs(int u = 1, int p = 1, int d = 0) {
       par[u] = p; dep[u] = d; sz[u] = 1;
        int mx = -1;
        for(int &v : edge[u]) {
            if(v == p) continue;
            dfs(v, u, d + 1);
           sz[u] += sz[v];
            if(mx < sz[v]) {
               mx = sz[v];
                swap(v, edge[u][0]);
   void dfsHLD(int u = 1, int p = 1) {
        idx[++t] = u; in[u] = t;
        for(int &v : edge[u]) {
            if(v == p) continue;
            root[v] = (v == edge[u][0] ? root[u] : v);
            dfsHLD(v, u);
        // out[u] = t;
    int lca(int x, int y) {
        int res = 0;
       while(root[x] != root[y]) {
            if(dep[root[x]] < dep[root[y]]) swap(x, y);</pre>
            res ^= que(1, 1, n, in[root[x]], in[x]);
            x = par[root[x]];
        if(dep[x] > dep[y]) swap(x, y);
        res ^= que(1, 1, n, in[x], in[y]);
```

```
return res;
    void reset() {
        for(int i = 1 ; i <= n ; ++i) edge[i].clear();</pre>
// hld.build(1, 1, n);
```

#### 2.4 Li-Chao Tree

```
// max li-chao tree
// works for the range [0, MAX - 1]
// if min li-chao tree:
// replace every call to max() with min() and every > with <
// also replace -INF with INF
struct Func {
  ll m, c;
  ll operator()(ll x) {
    return x * m + c:
};
const int MAX = 1e9 + 1;
const ll INF = 1e18:
const Func NIL = {0, -INF};
struct Node {
 Func f;
  Node* lc;
  Node* rc;
  Node(): f(NIL), lc(nullptr), rc(nullptr) {}
  Node(const Node& n) : f(n.f), lc(nullptr), rc(nullptr) {}
Node* root = new Node;
void insert(Func f, Node* cur = root, int l = 0, int r = MAX - 1) {
  int m = l + (r - l) / 2;
  bool left = f(l) > cur->f(l);
  bool mid = f(m) > cur -> f(m);
  if(mid)
    swap(f, cur->f);
  if(l != r) {
    if(left != mid) {
      if(!cur->lc)
        cur->lc = new Node(*cur);
      insert(f, cur->lc, l, m);
    } else {
      if(!cur->rc)
        cur->rc = new Node(*cur);
      insert(f, cur->rc, m + 1, r);
Il query(ll \times Node \times cur = root, int l = 0, int r = MAX - 1) {
  if(!cur)
    return -INF;
  if(l == r)
    return cur->f(x);
  int m = l + (r - l) / 2;
```

```
if(x <= m)
  return max(cur->f(x), query(x, cur->lc, l, m));
else
  return max(cur->f(x), query(x, cur->rc, m + 1, r));
}
```

#### 2.5 STL PBDS

## 2.6 Treap

```
struct tNode{
    int key, prior;
   tNode *l, *r;
   int sz;
   tNode() {}
   tNode(int key) : key(key), prior(rand()), l(NULL), r(NULL), sz(1) {}
typedef tNode* pNode;
int cnt(pNode t) { return t ? t->sz : 0; }
void upd(pNode t) { if(t) t->sz = 1 + cnt(t->l) + cnt(t->r); }
void split(pNode t, int key, pNode &l, pNode &r){
   if(!t) l = r = NULL;
   else if(t->key <= key) {
        split(t->r, key, t->r, r);
   } else{
       split(t->l, key, l, t->l);
   upd(t);
void ins(pNode &t, pNode it) {
   if(!t) t = it;
   else if(it->prior > t->prior) {
       split(t, it->key, it->l, it->r);
       t = it;
    } else ins(t->key <= it->key ? t->r : t->l, it);
   upd(t);
void merge(pNode &t, pNode l, pNode r) {
   if(!l || !r) t = l ? l : r;
   else if(l->prior > r->prior) {
       merge(l->r, l->r, r);
       t = l;
    } else {
       merge(r->l, l, r->l);
       t = r;
   upd(t);
void erase(pNode &t, int key) {
   if(t->key == key) {
       pNode th = t;
       merge(t, t->l, t->r);
       delete th;
   } else erase(key < t->key ? t->l : t->r, key);
   upd(t);
```

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#### 2.8 Mo's on Tree

```
\begin{split} ST(u) &\leq ST(v) \\ P &= LCA(u,v) \\ \text{If } P &= u, \text{ query } [ST(u),ST(v)] \\ \text{Else query } [EN(u),ST(v)] + [ST(P),ST(P)] \end{split}
```

#### 2.9 Link-Cut Tree

```
// Represents a forest of unrooted trees. You can add and remove edges
// (as long as the result is still a forest), and check whether two
// nodes are in the same tree.
// Complexity: log(n)
struct Node { // Splay tree. Root's pp contains tree's parent.
  Node* p = 0, *pp = 0, *c[2];
  int sz = 0;
  Node() {
    c[0] = c[1] = 0;
    fix();
  void fix() {
    sz = 1;
    if(c[0]) c[0] -> p = this, sz += c[0] -> sz;
    if(c[1]) c[1] -> p = this, sz += c[1] -> sz;
    // (+ update sum of subtree elements etc. if wanted)
  int up() {
    return p ? p->c[1] == this : -1;
  void rot(int i, int b) {
    int h = i ^ b;
    Node* x = c[i], *y = (b == 2 ? x : x->c[h]), *z = (b ? y : x);
    if(y->p = p) p->c[up()] = y;
    c[i] = z - c[i ^ 1];
    if(b < 2) x \rightarrow c[h] = y \rightarrow c[h \land 1], z \rightarrow c[h \land 1] = b ? x : this;
    y - c[i ^ 1] = b ? this : x;
    fix();
    x \rightarrow fix();
    y->fix();
    if(p) p \rightarrow fix();
    swap(pp, y->pp);
  // Splay this up to the root. Always finishes without flip set.
  void splay() {
    while(p) {
       int c1 = up(), c2 = p - up();
      if(c2 == -1) p -> rot(c1, 2);
      else p->p->rot(c2, c1 != c2);
```

Page

```
struct LinkCut {
 vector<Node> node;
 LinkCut(int N) : node(N + 1) {}
 void link(int u, int v) { // add an edge u --> v
   assert(!connected(u, v));
   access(&node[u]);
   access(&node[v]);
   node[u].c[0] = &node[v];
   node[v].p = &node[u];
   node[u].fix();
 void cut(int u, int v) { // remove an edge u --> v
   assert(connected(u, v));
   Node* x = &node[v], *top = &node[u];
   access(top);
   top->c[0] = top->c[0]->p = 0;
   top->fix();
 bool connected(int u, int v) { // are u, v in the same tree?
   return root(u) == root(v);
  int root(int u) { // find the root id of node u
   Node* x = & node[u];
   access(x);
   for(; x \rightarrow c[0]; x = x \rightarrow c[0]);
   x->splay();
   return (int)((vector<Node>::iterator)x - node.begin());
  // Move u to root aux tree. Return the root of the root aux tree.
 Node* access(Node* u) {
   u->splay();
   Node* last = u;
   if(Node*\& x = u->c[1]) {
     x->pp = u;
     x->p = 0;
     \times = 0:
     u->fix();
    for(Node * pp; (pp = u->pp) && (last = pp);) {
     pp->splay();
     if(pp->c[1]) pp->c[1]->p = 0, pp->c[1]->pp = pp;
     pp - > c[1] = u;
     u->p = pp;
     u->pp = 0;
     pp->fix();
     u->splay();
   return last;
 int depth(int u) {
   access(&node[u]);
   return node[u].sz - 1;
 Node* lca(int u, int v) {
   access(&node[u]);
   return access(&node[v]);
};
```

#### 2.10 LineContainer

```
struct Line {
   mutable ll k, m, p;
   bool operator<(const Line& o) const { return k < o.k; }
   bool operator<(ll x) const { return p < x; }
};
// get maximum
struct LineContainer : multiset<Line, less<>> {
   // (for doubles, use inf = 1/.0, div(a,b) = a/b)
   static const ll inf = LLONG_MAX;
```

```
ll div(ll a, ll b) { // floored division
        return a / b - ((a ^ b) < 0 && a % b); }
    bool isect(iterator x, iterator y) {
        if (y == end()) return x \rightarrow p = inf, 0;
        if (x-)k == y-)k x-)p = x-)m > y-)m ? inf : -inf;
       else x->p = div(y->m - x->m, x->k - y->k);
        return x->p >= y->p;
    void add(ll k, ll m) {
        auto z = insert(\{k, m, 0\}), y = z++, x = y;
        while (isect(y, z)) z = erase(z);
        if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(y));
        while ((y = x) != begin() && (--x)->p >= y->p)
            isect(x, erase(y));
    ll querv(ll x) {
        assert(!empty());
        auto l = *lower_bound(x);
        return l.k * x + l.m;
};
```

#### 2.11 Wavelet Tree

```
class wavelet_tree {
public:
    int low, high;
    wavelet_tree* l, *r;
    vector<int> freq;
    wavelet_tree(int* from, int* to, int x, int y) {
        low = x, high = y;
        if(from >= to) return;
        if(high == low) {
            freq.reserve(to - from + 1);
            freq.push_back(0);
            for (auto it = from; it != to; it++)
                freq.push_back(freq.back() + 1);
        int mid = (low + high) / 2;
        auto lessThanMid = [mid](int x) {
            return x <= mid;
        freq.reserve(to - from + 1);
        freq.push_back(0);
        for (auto it = from; it != to; it++)
            freq.push_back(freq.back() + lessThanMid(*it));
        auto pivot = stable_partition(from, to, lessThanMid);
        l = new wavelet_tree(from, pivot, low, mid);
        r = new wavelet_tree(pivot, to, mid + 1, high);
    int kOrLess(int l, int r, int k) {
        if (l > r or k < low) return 0;</pre>
        if (high \leq k) return r - l + 1;
        int LtCount = freq[l - 1];
        int RtCount = freq[r];
        return (this->l->kOrLess(LtCount + 1, RtCount, k) +
            this->r->kOrLess(l - LtCount, r - RtCount, k));
};
```

# 3 Dynamic Programming

## 3.1 DP Convex Hull

```
/* dp[i] = min k<i {dp[k] + x[i]*m[k]}
Make sure gradient (m[i]) is either non-increasing if min,
or non-decreasing if max. x[i] must be non-decreasing. just sort */</pre>
```

```
int y[N], m[N];
// while this is true, pop back from dq. a=new line, b=last, c=2nd last
bool cekx(int a, int b, int c) {
 // if not enough, change to cross mul
 // if cross mul, beware of negative denominator, and overflow
 return (double)(y[b] - y[a]) / (m[a] - m[b]) <= (double)(y[c] - y[b]) /
         (m[b] - m[c]);
```

## 3.2 DP Knuth-Yao

```
// dp[i][j] = min{k} dp[i][k]+dp[k][j]+cost[i][j]
for(int k = 0; k \le n; k++) {
 for(int i = 0; i + k <= n; i++) {
   if(k < 2)
     dp[i][i + k] = 0, opt[i][i + k] = i;
   else {
     int sta = opt[i][i + k - 1];
      int end = opt[i + 1][i + k];
     for(int j = sta; j <= end; j++) {</pre>
       if(dp[i][j] + dp[j][i + k] + cost[i][i + k] < dp[i][i + k]) {
         dp[i][i + k] = dp[i][j] + dp[j][i + k] + cost[i][i + k];
         opt[i][i + k] = j;
```

# Geometry

# 4.1 Geometry Template

```
O. Basic Rule
    0.1. Everything is in double
    0.2. Every comparison use EPS
    0.3. Every degree in rad
1. General Double Operation
    1.1. const double EPS=1E-9
    1.2. const double PI=acos(-1.0)
    1.3. const double INFD=1E9
    1.3. between_d(double x,double l,double r)
        check whether x is between l and r inclusive with EPS
    1.4. same_d(double x,double y)
        check whether x=y with EPS
    1.5. dabs(double x)
        absolute value of x
2. Point
    2.1. struct point
        2.1.1. double x,y
            cartesian coordinate of the point
        2.1.2. point()
            default constructor
        2.1.3. point(double _x,double _y)
            constructor, set the point to (x,y)
        2.1.4. bool operator< (point other)
            regular pair<double,double> operator < with EPS
        2.1.5. bool operator== (point other)
            regular pair<double,double> operator == with EPS
        length of hypotenuse of point P to (0,0)
    2.3. e_dist(point P1,point P2)
        euclidean distance from P1 to P2
    2.4. m_dist(point P1,point P2)
       manhattan distance from P1 to P2
```

```
2.5. point rotate(point P, point O, double angle)
        rotate point P from the origin O by angle ccw
3. Vector
    3.1. struct vec
        3.1.1. double x,y
            x and y magnitude of the vector
        3.1.3. vec(double _x,double _y)
            constructor, set the vector to (_x,_y)
        3.1.4. vec(point A, point B)
            constructor, set the vector to vector AB (A->B)
/*General Double Operation*/
const double PI = acos(-1.0);
const double INFD = 1E9;
double between_d(double x, double l, double r) {
 return (min(l, r) \le x + EPS \&\& x \le max(l, r) + EPS);
double same_d(double x, double y) {
 return between_d(x, y, y);
double dabs(double x) {
 if(x < EPS)
   return -x;
 return x;
/*Point*/
struct point {
 double x, y;
  point() {
   x = y = 0.0;
  point(double _x, double _y) {
   \times = \_ \times ;
   y = _y;
  bool operator< (point other) {
   if(x < other.x + EPS)
     return true;
    if(x + EPS > other.x)
     return false;
    return y < other.y + EPS;</pre>
  bool operator == (point other) {
    return same_d(x, other.x) && same_d(y, other.y);
double e_dist(point P1, point P2) {
 return hypot(P1.x - P2.x, P1.y - P2.y);
double m_dist(point P1, point P2) {
 return dabs(P1.x - P2.x) + dabs(P1.y - P2.y);
double pointBetween(point P, point L, point R) {
 return (e_dist(L, P) + e_dist(P, R) == e_dist(L, R));
bool collinear(point P, point L,
               point R) { //newly added(luis), cek 3 poin segaris
  return P.x * (L.y - R.y) + L.x * (R.y - P.y) + R.x * (P.y - L.y) ==
         0; // bole gnti "dabs(x)<"EPS
/*Vector*/
struct vec {
 double x, y;
 vec() {
   \times = \vee = 0.0;
  vec(double _x, double _y) {
```

```
Jangan Salah B
Bina Nusantara
           \times = \_ \times ;
           y = _y;
         vec(point A) {
          \times = A.\times;
           y = A.y;
, Baca Soai 1
ara Universit
         vec(point A, point B) {
           x = B.x - A.x;
           y = B.y - A.y;
       vec scale(vec v, double s) {
        return vec(v.x * s, v.y * s);
       vec flip(vec v) {
        return vec(-v.x, -v.y);
       double dot(vec u, vec v) {
        return (u.x * v.x + u.y * v.y);
       double cross(vec u, vec v) {
        return (u.x * v.y - u.y * v.x);
       double norm_sq(vec v) {
        return (v.x * v.x + v.y * v.y);
       point translate(point P, vec v) {
        return point(P.x + v.x, P.y + v.y);
       point rotate(point P, point O, double angle) {
        vec v(0);
         P = translate(P, flip(v));
         return translate(point(P.x * cos(angle) - P.y * sin(angle),
                                  P.x * sin(angle) + P.y * cos(angle)), v);
       point mid(point P, point Q) {
        return point((P.x + Q.x) / 2, (P.y + Q.y) / 2);
       double angle(point A, point O, point B) {
         vec OA(O, A), OB(O, B);
         return acos(dot(OA, OB) / sqrt(norm_sq(OA) * norm_sq(OB)));
       int orientation(point P, point Q, point R) {
        vec PQ(P, Q), PR(P, R);
         double c = cross(PQ, PR);
        if(c < -EPS)
          return -1;
         if(c > EPS)
           return 1;
         return 0;
       /*Line*/
       struct line {
         double a, b, c;
         line() {
           a = b = c = 0.0;
         line(double _a, double _b, double _c) {
           a = _a;
           b = _b;
           C = C;
         line(point P1, point P2) {
           if(P1 < P2)swap(P1, P2);
           if(same_d(P1.x, P2.x))a = 1.0, b = 0.0, c = -P1.x;
             a = -(P1.y - P2.y) / (P1.x - P2.x), b = 1.0, c = -(a * P1.x) - P1.y;
         line(point P, double slope) {
           if(same_d(slope, INFD))a = 1.0, b = 0.0, c = -P.x;
```

```
else a = -slope, b = 1.0, c = -(a * P.x) - P.y;
 bool operator == (line other) {
   return same_d(a, other.a) && same_d(b, other.b) && same_d(c, other.c);
 double slope() {
   if(same_d(b, 0.0))
     return INFD;
   return -(a / b);
bool paralel(line L1, line L2) {
 return same_d(L1.a, L2.a) && same_d(L1.b, L2.b);
bool intersection(line L1, line L2, point& P) {
 if(paralel(L1, L2))
   return false;
 P.x = (L2.b * L1.c - L1.b * L2.c) / (L2.a * L1.b - L1.a * L2.b);
 if(same_d(L1.b, 0.0))
   P.y = -(L2.a * P.x + L2.c);
 else
   P.y = -(L1.a * P.x + L1.c);
 return true;
double pointToLine(point P, point A, point B, point& C) {
 vec AP(A, P), AB(A, B);
 double u = dot(AP, AB) / norm sq(AB);
 C = translate(A, scale(AB, u));
 return e_dist(P, C);
double lineToLine(line L1, line L2) {
 if(!paralel(L1, L2))
   return 0.0:
 return dabs(L2.c - L1.c) / sqrt(L1.a * L1.a + L1.b * L1.b);
/*Line Segment*/
struct segment -
 point P, Q;
 line L;
 segment() {
   point T1;
   P = 0 = T1;
   line T2;
   L = T2;
 segment(point _P, point _Q) {
   P = P;
   if(Q < P)swap(P, Q);
   line T(P, Q);
   L = T;
  bool operator == (segment other) {
    return P == other.P && Q == other.Q;
bool onSegment(point P, segment S) {
 if(orientation(S.P, S.Q, P) != 0)return false;
 return between_d(P.x, S.P.x, S.Q.x) && between_d(P.y, S.P.y, S.Q.y);
bool s_intersection(segment S1, segment S2) {
 double o1 = orientation(S1.P, S1.Q, S2.P);
 double o2 = orientation(S1.P, S1.Q, S2.Q);
 double o3 = orientation(S2.P, S2.Q, S1.P);
 double o4 = orientation(S2.P, S2.Q, S1.Q);
 if(o1 != o2 && o3 != o4)return true;
  if(o1 == 0 && onSegment(S2.P, S1))return true;
  if(o2 == 0 && onSegment(S2.Q, S1))return true;
  if(o3 == 0 && onSegment(S1.P, S2))return true;
  if(o4 == 0 && onSegment(S1.Q, S2))return true;
 return false;
```

```
double pointToSegment(point P, point A, point B, point& C) {
 vec AP(A, P), AB(A, B);
 double u = dot(AP, AB) / norm_sq(AB);
 if(u < EPS) {
  C = A;
   return e_dist(P, A);
 if(u + EPS > 1.0) {
   C = B;
   return e_dist(P, B);
 return pointToLine(P, A, B, C);
double segmentToSegment(segment S1, segment S2) {
 if(s intersection(S1, S2))return 0.0;
 double ret = INFD;
 point dummy;
 ret = min(ret, pointToSegment(S1.P, S2.P, S2.Q, dummy));
 ret = min(ret, pointToSegment(S1.Q, S2.P, S2.Q, dummy));
 ret = min(ret, pointToSegment(S2.P, S1.P, S1.O, dummy));
 ret = min(ret, pointToSegment(S2.Q, S1.P, S1.Q, dummy));
 return ret:
/*Circle*/
struct circle {
 point P;
 double r:
 circle() {
   point P1;
   P = P1:
   r = 0.0;
 circle(point _P, double _r) {
   P = P;
   r = _r;
 circle(point P1, point P2) {
   P = mid(P1, P2);
   r = e_dist(P, P1);
 circle(point P1, point P2, point P3) {
   vector<point> T;
   T.clear();
   T.pb(P1);
   T.pb(P2);
   T.pb(P3);
   sort(T.begin(), T.end());
   P1 = T[0];
   P2 = T[1];
   P3 = T[2];
   point M1, M2;
   M1 = mid(P1, P2);
   M2 = mid(P2, P3);
   point Q2, Q3;
   Q2 = rotate(P2, P1, PI / 2);
   Q3 = rotate(P3, P2, PI / 2);
   vec P1Q2(P1, Q2), P2Q3(P2, Q3);
   point M3, M4;
   M3 = translate(M1, P1Q2);
   M4 = translate(M2, P2Q3);
   line L1(M1, M3), L2(M2, M4);
   intersection(L1, L2, P);
   r = e_dist(P, P1);
 bool operator == (circle other) {
   return (P == other.P && same_d(r, other.r));
bool insideCircle(point P, circle C) {
 return e_dist(P, C.P) <= C.r + EPS;</pre>
```

```
bool c_intersection(circle C1, circle C2, point& P1, point& P2) {
 double d = e_dist(C1.P, C2.P);
  if(d > C1.r + C2.r) {
    return false; //d+EPS kalo butuh
  if(d < dabs(C1.r - C2.r) + EPS)
   return false:
  double x1 = C1.P.x, y1 = C1.P.y, r1 = C1.r, x2 = C2.P.x, y2 = C2.P.y, r2 = C2.r;
  double a = (r1 * r1 - r2 * r2 + d * d) / (2 * d), h = sqrt(r1 * r1 - a * a);
  point T(x1 + a * (x2 - x1) / d, y1 + a * (y2 - y1) / d);
  P1 = point(T.x - h * (y2 - y1) / d, T.y + h * (x2 - x1) / d);
  P2 = point(T.x + h * (y2 - y1) / d, T.y - h * (x2 - x1) / d);
  return true;
bool lc intersection(line L, circle O, point& P1, point& P2) {
  double a = L.a, b = L.b, c = L.c, x = 0.P.x, y = 0.P.y, r = 0.r;
  double A = a * a + b * b, B = 2 * a * b * y - 2 * a * c - 2 * b * b * x,
         C = b * b * x * x + b * b * y * y - 2 * b * c * y + c * c - b * b * r * r;
  double D = B * B - 4 * A * C;
  point T1, T2;
  if(same_d(b, 0.0)) {
    T1.x = c / a;
    if(dabs(x - T1.x) + EPS > r)
     return false;
    if(same_d(T1.x - r - x, 0.0) | | same_d(T1.x + r - x, 0.0)) {
     P1 = P2 = point(T1.x, y);
      return true;
    double dx = dabs(T1.x - x), dy = sqrt(r * r - dx * dx);
   P1 = point(T1.x, y - dy);
    P2 = point(T1.x, y + dy);
    return true;
  if(same_d(D, 0.0)) {
    T1.x = -B / (2 * A);
    T1.v = (c - a * T1.x) / b;
   P1 = P2 = T1:
   return true;
  if(D < EPS)
   return false:
  D = sqrt(D);
  T1.x = (-B - D) / (2 * A);
  T1.y = (c - a * T1.x) / b;
 P1 = T1;
  T2.x = (-B + D) / (2 * A);
  T2.y = (c - a * T2.x) / b;
 P2 = T2:
  return true;
bool sc_intersection(segment S, circle C, point& P1, point& P2) {
 bool cek = lc_intersection(S.L, C, P1, P2);
  if(!cek)
   return false;
  double x1 = S.P.x, y1 = S.P.y, x2 = S.Q.x, y2 = S.Q.y;
  bool b1 = between_d(P1.x, x1, x2) && between_d(P1.y, y1, y2);
  bool b2 = between_d(P2.x, x1, x2) && between_d(P2.y, y1, y2);
  if(P1 == P2)
   return b1;
  if(b1 || b2) -
    if(!b1)
     P1 = P2;
    if(!b2)
     P2 = P1;
    return true;
  return false;
double t_perimeter(point A, point B, point C) {
```

```
point I;
   bool B = intersection(cast, temp, I);
   if(!B)
     continue;
   else if(I == A.P[i] || I == A.P[i + 1])
   else if(pointBetween(I, A.P[i], A.P[i + 1]) && pointBetween(I, P, Q))
     cnt++;
 return cnt % 2 == 1;
// line segment p-q intersect with line A-B.
point lineIntersectSeg(point p, point q, point A, point B) {
 double a = B.y - A.y;
 double b = A.x - B.x;
 double c = B.x * A.y - A.x * B.y;
 double u = fabs(a * p.x + b * p.y + c);
 double v = fabs(a * q.x + b * q.y + c);
 return point((p.x * v + q.x * u) / (u + v), (p.y * v + q.y * u) / (u + v));
```

return e\_dist(A, B) + e\_dist(B, C) + e\_dist(C, A);

return sqrt(s \* (s - ab) \* (s - bc) \* (s - ac));

circle t\_inCircle(point A, point B, point C) {

double ratio = e\_dist(A, B) / e\_dist(A, C);

circle t\_outCircle(point A, point B, point C) {

BC = scale(BC, ratio / (1 + ratio));

ratio = e\_dist(B, A) / e\_dist(B, C);

AC = scale(AC, ratio / (1 + ratio));

vector<point> T;

sort(T.begin(), T.end());

P = translate(B, BC); line AP1(A, P);

P = translate(A, AC);

return circle(P, r);

return circle(A, B, C);

intersection(AP1, BP2, P);

polygon(vector<point>& \_P) {

point Q(P.x, 10000); line cast(P, Q);

bool rayCast(point P, polygon& A) {

FOR(i, (int)(A.P.size()) - 1) {

line temp(A.P[i], A.P[i + 1]);

T.clear();

T.pb(A);

T.pb(B);

T.pb(C);

A = T[0];B = T[1];

C = T[2];

point P;

/\*Polygon\*/

struct polygon {

polygon() {

P = P;

int cnt = 0;

vector<point> P;

P.clear();

vec AC(A, C);

line BP2(B, P);

double ab =  $e_dist(A, B)$ , bc =  $e_dist(B, C)$ , ac =  $e_dist(C, A)$ ;

double  $r = t_area(A, B, C) / (t_perimeter(A, B, C) / 2);$ 

double t\_area(point A, point B, point C) { double  $s = t_perimeter(A, B, C) / 2;$ 

```
|// cuts polygon Q along the line formed by point a -> point b
  // (note: the last point must be the same as the first point)
  vector<point> cutPolygon(point a, point b, const vector<point>& Q) {
       vector<point> P;
       for(int i = 0; i < (int)Q.size(); i++) {
           double left1 = cross(toVec(a, b), toVec(a, Q[i]));
           double left2 = 0;
           if(i != (int)Q.size() - 1)
               left2 = cross(toVec(a, b), toVec(a, Q[i + 1]));
            if(left1 > -EPS)
               P.push_back(Q[i]);
           if(left1 * left2 < -EPS)</pre>
                P.push_back(lineIntersectSeg(Q[i], Q[i + 1], a, b));
       if(!P.empty() && !(P.back() == P.front()))
           P.push back(P.front());
      return P;
  circle minCoverCircle(polygon& A) {
      vector<point> p = A.P;
      point c;
      circle ret;
      double cr = 0.0;
       int i, j, k;
      c = p[0];
       for(i = 1; i < p.size(); i++) {
           if(e_dist(p[i], c) >= cr + EPS) {
               c = p[i], cr = 0;
                ret = circle(c, cr);
                for(j = 0; j < i; j++) {
                    if(e_dist(p[j], c) >= cr + EPS) {
                         c = mid(p[i], p[j]);
                         cr = e_dist(p[i], c);
                         ret = circle(c, cr);
                          for(k = 0; k < j; k++) {
                              if(e_dist(p[k], c) >= cr + EPS) {
                                  ret = circle(p[i], p[j], p[k]);
                                  c = ret.P;
                                   cr = ret.r;
      return ret;
   /*Geometry Algorithm*/
  double DP[110][110];
  double minCostPolygonTriangulation(polygon& A) {
      if(A.P.size() < 3)return 0;</pre>
      FOR(i, A.P.size())
            for(int j = 0, k = i; k < A.P.size(); j++, k++) {
               if(k < j + 2)DP[j][k] = 0.0;
                else {
                    DP[j][k] = INFD;
                    REP(l, j + 1, k - 1) {
                         double cost = e_dist(A.P[j], A.P[k]) + e_dist(A.P[k], A.P[l]) + e_dist(A.P[l \leftrightarrow a.P[l])) + e_dist(A.P[l)) +
                         DP[j][k] = min(DP[j][k], DP[j][l] + DP[l][k] + cost);
       return DP[0][A.P.size() - 1];
```

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```

```
typedef double TD;
                                  // for precision shits
namespace GEOM {
 typedef pair<TD, TD> Pt;
                                  // vector and points
 const TD EPS = 1e-9;
 const TD maxD = 1e9;
 TD cross(Pt a, Pt b, Pt c) {
                                  // right hand rule
   TD v1 = a.first - c.first;
                                  // (a-c) X (b-c)
   TD v2 = a.second - c.second;
   TD u1 = b.first - c.first;
   TD u2 = b.second - c.second;
   return v1 * u2 - v2 * u1;
 TD cross(Pt a, Pt b) {
                                  // a X b
   return a.first * b.second - a.second * b.first;
 TD dot(Pt a, Pt b, Pt c) {
                                  // (a-c) . (b-c)
   TD v1 = a.first - c.first;
   TD v2 = a.second - c.second;
   TD u1 = b.first - c.first;
   TD u2 = b.second - c.second;
   return v1 * u1 + v2 * u2;
 TD dot(Pt a, Pt b) {
   return a.first * b.first + a.second * b.second;
 TD dist(Pt a, Pt b) {
   return sqrt((a.first - b.first) * (a.first - b.first) +
               (a.second - b.second) * (a.second - b.second));
 TD shoelaceX2(vector<Pt>& convHull) {
   TD ret = 0:
   for(int i = 0, n = convHull.size(); i < n; i++)</pre>
     ret += cross(convHull[i], convHull[(i + 1) % n]);
   return ret;
 vector<Pt> createConvexHull(vector<Pt>& points) {
   sort(points.begin(), points.end());
   vector<Pt> ret;
   for(int i = 0; i < points.size(); i++) {</pre>
     while(ret.size() > 1 &&
           cross(points[i], ret[ret.size() - 1], ret[ret.size() - 2]) < -EPS)</pre>
       ret.pop_back();
     ret.push_back(points[i]);
   for(int i = points.size() - 2, sz = ret.size(); i \ge 0; i--) {
     while(ret.size() > sz &&
           cross(points[i], ret[ret.size() - 1], ret[ret.size() - 2]) < -EPS)</pre>
       ret.pop_back();
     if(i == 0)break;
     ret.push_back(points[i]);
   return ret;
   bool isInside(Pt pv, vector<Pt>& x) { //using winding number
     int n = x.size(), wn = 0;
     x.push_back(x[0]);
     for(int i = 0; i < n; ++i) {
       if(((x[i + 1].first \leftarrow pv.first \&\& x[i].first \rightarrow pv.first)))
            (x[i + 1].first >= pv.first && x[i].first <= pv.first)) &&
           ((x[i + 1].second \le pv.second \&\& x[i].second >= pv.second) ||
            (x[i + 1].second >= pv.second && x[i].second <= pv.second))) {
         if(cross(x[i], x[i + 1], pv) == 0) {
           x.pop_back();
           return true;
     for(int i = 0; i < n; ++i) {
       if(x[i].second <= pv.second) {</pre>
         if(x[i + 1].second > pv.second && cross(x[i], x[i + 1], pv) > 0)
```

```
} else if(x[i + 1].second <= pv.second && cross(x[i], x[i + 1], pv) < 0)
    x.pop_back();
    return wn != 0;
bool isInside(Pt pv, vector<Pt>& x) { //using winding number
 int n = x.size(), wn = 0;
  x.push_back(x[0]);
  for(int i = 0; i < n; ++i) {
   if(((x[i + 1].first \le pv.first \&\& x[i].first \ge pv.first)))
        (x[i + 1].first >= pv.first && x[i].first <= pv.first)) &&
        ((x[i + 1].second \le pv.second & x[i].second >= pv.second) | |
         (x[i + 1].second >= pv.second && x[i].second <= pv.second))) {
      if(cross(x[i], x[i + 1], pv) == 0) {
       x.pop_back();
        return true;
  for(int i = 0; i < n; ++i) {
   if(x[i].second <= pv.second) {
     if(x[i + 1].second > pv.second && cross(x[i], x[i + 1], pv) > 0)
   } else if(x[i + 1].second <= pv.second && cross(x[i], x[i + 1], pv) < 0)
 x.pop_back();
  return wn != 0;
```

## 4.3 Closest Pair of Points

```
#define fi first
#define se second
typedef pair<int, int> pii;
struct Point {
 int x, y, id;
int compareX(const void* a, const void* b) {
 Point* p1 = (Point*)a, *p2 = (Point*)b;
  return (p1->x - p2->x);
int compareY(const void* a, const void* b) {
 Point* p1 = (Point*)a, *p2 = (Point*)b;
  return (p1->y - p2->y);
double dist(Point p1, Point p2) {
 return sqrt((double)(p1.x - p2.x) * (p1.x - p2.x) +
              (double)(p1.y - p2.y) * (p1.y - p2.y)
pair<pii, double> bruteForce(Point P[], int n) {
 double min = 1e8;
  pii ret = pii(-1, -1);
  for(int i = 0; i < n; ++i)
    for(int j = i + 1; j < n; ++j)
     if(dist(P[i], P[j]) < min) {</pre>
        ret = pii(P[i].id, P[j].id);
        min = dist(P[i], P[j]);
  return pair<pii, double> (ret, min);
pair<pii, double> getmin(pair<pii, double> x, pair<pii, double> y) {
  if(x.fi.fi == -1 \&\& x.fi.se == -1) return y;
  if(y.fi.fi == -1 \&\& y.fi.se == -1) return x;
  return (x.se < y.se) ? x : y;
```

Lagi ity

```
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```

```
pair<pii, double> stripClosest(Point strip[], int size, double d) {
 double min = d;
 pii ret = pii(-1, -1);
 gsort(strip, size, sizeof(Point), compareY);
 for(int i = 0; i < size; ++i)
   for(int j = i + 1; j < size && (strip[j].y - strip[i].y) < min; ++j)
      if(dist(strip[i], strip[j]) < min) {</pre>
       ret = pii(strip[i].id, strip[j].id);
       min = dist(strip[i], strip[j]);
 return pair<pii, double>(ret, min);
pair<pii, double> closestUtil(Point P[], int n) {
 if(n <= 3)return bruteForce(P, n);</pre>
 int mid = n / 2;
 Point midPoint = P[mid];
 pair<pii, double> dl = closestUtil(P, mid);
 pair<pii, double> dr = closestUtil(P + mid, n - mid);
 pair<pii, double> d = getmin(dl, dr);
 Point strip[n];
 int j = 0;
 for(int i = 0; i < n; i++)
   if(abs(P[i].x - midPoint.x) < d.second)</pre>
     strip[j] = P[i], j++;
 return getmin(d, stripClosest(strip, j, d.second));
pair<pii, double> closest(Point P[], int n) {
 gsort(P, n, sizeof(Point), compareX);
 return closestUtil(P, n);
Point P[50005];
int main() {
 int n:
 scanf("%d", &n);
 for(int a = 0; a < n; a++) {
   scanf("%d%d", &P[a].x, &P[a].y);
   P[a].id = a;
 pair<pii, double> hasil = closest(P, n);
 if(hasil.fi.fi > hasil.fi.se)
   swap(hasil.fi.fi, hasil.fi.se);
 printf("%d %d %.6lf\n", hasil.fi.fi, hasil.fi.se, hasil.se);
 return 0;
```

# 4.4 Smallest Enclosing Circle

```
// welzl's algo to find the 2d minimum enclosing circle of a set of points
// expected O(N)
// directions: remove duplicates and shuffle points, then call welzl(points)
struct Point {
 double x;
 double y;
struct Circle {
 double x, y, r;
 Circle() {}
 Circle(double \_x, double \_y, double \_r): x(\_x), y(\_y), r(\_r) {}
};
Circle trivial(const vector<Point>& r) {
 if(r.size() == 0)return Circle(0, 0, -1);
 else if(r.size() == 1)return Circle(r[0].x, r[0].y, 0);
 else if(r.size() == 2) {
   double cx = (r[0].x + r[1].x) / 2.0, cy = (r[0].y + r[1].y) / 2.0;
   double rad = hypot(r[0].x - r[1].x, r[0].y - r[1].y) / 2.0;
    return Circle(cx, cy, rad);
  } else {
```

```
double x0 = r[0].x, x1 = r[1].x, x2 = r[2].x;
    double y0 = r[0].y, y1 = r[1].y, y2 = r[2].y;
   double d = (x0 - x2) * (y1 - y2) - (x1 - x2) * (y0 - y2);
    double cx = (((x0 - x2) * (x0 + x2) + (y0 - y2) * (y0 + y2)) / 2 *
                 (y1 - y2) - ((x1 - x2) * (x1 + x2) + (y1 - y2) * (y1 + y2)) / 2
                 * (y0 - y2)) / d;
    double cy = (((x1 - x2) * (x1 + x2) + (y1 - y2) * (y1 + y2)) / 2 *
                 (x0 - x2) - ((x0 - x2) * (x0 + x2) + (y0 - y2) * (y0 + y2)) / 2
                 *(x1 - x2)) / d;
    return Circle(cx, cy, hypot(x0 - cx, y0 - cy));
Circle welzl(const vector<Point>& p, int idx = 0, vector<Point> r = {}) {
 if(idx == (int) p.size() || r.size() == 3)return trivial(r);
 Circle d = welzl(p, idx + 1, r);
 if(hypot(p[idx].x - d.x, p[idx].y - d.y) > d.r) {
   r.push_back(p[idx]);
   d = welzl(p, idx + 1, r);
 return d;
```

# 4.5 Sutherland-Hodgman Algorithm

```
// Complexity: linear time
// Ada 2 poligon, cari poligon intersectionnya
// poly_point = hasilnya, clipper = pemotongnya
#include<bits/stdc++.h>
using namespace std;
const double EPS = 1e-9;
struct point {
 double x, y;
 point(double _x, double _y): x(_x), y(_y) {}
struct vec {
 double x, y;
 vec(double _x, double _y): x(_x), y(_y) {}
point pivot(0, 0);
vec toVec(point a, point b) {
 return vec(b.x - a.x, b.y - a.y);
double dist(point a, point b) {
 return hypot(a.x - b.x, a.y - b.y);
double cross(vec a, vec b) {
 return a.x * b.y - a.y * b.x;
bool ccw(point p, point q, point r) {
 return cross(toVec(p, q), toVec(p, r)) > 0;
bool collinear(point p, point q, point r) {
 return fabs(cross(toVec(p, q), toVec(p, r))) < EPS;</pre>
bool lies(point a, point b, point c) {
  if((c.x) = min(a.x, b.x) \&\& c.x <= max(a.x, b.x)) \&\&
      (c.y >= min(a.y, b.y) \&\& c.y <= max(a.y, b.y)))
    return true;
  else return false;
bool anglecmp(point a, point b) {
  if(collinear(pivot, a, b))
    return dist(pivot, a) < dist(pivot, b);</pre>
  double d1x = a.x - pivot.x, d1y = a.y - pivot.y;
  double d2x = b.x - pivot.x, d2y = b.y - pivot.y;
```

```
return (atan2(d1y, d1x) - atan2(d2y, d2x)) < 0;
point intersect(point s1, point e1, point s2, point e2) {
 double x1, x2, x3, x4, y1, y2, y3, y4;
 y1 = s1.y;
 x2 = e1.x;
 y2 = e1.y;
 x3 = s2.x;
 y3 = s2.y;
 x4 = e2.x;
 y4 = e2.y;
 double num1 = (x1 * y2 - y1 * x2) * (x3 - x4) - (x1 - x2) * (x3 * y4 - y3 * x4);
 double num2 = (x1 * y2 - y1 * x2) * (y3 - y4) - (y1 - y2) * (x3 * y4 - y3 * x4);
 double den = (x1 - x2) * (y3 - y4) - (y1 - y2) * (x3 - x4);
 double new_x = num1 / den;
 double new_y = num2 / den;
 return point(new_x, new_y);
void clip(vector <point>& poly_points, point point1, point point2) {
 vector <point> new points;
 new_points.clear();
 for(int i = 0; i < poly_points.size(); i++) {</pre>
   int k = (i + 1) % poly_points.size();
   double i_pos = ccw(point1, point2, poly_points[i]);
   double k_pos = ccw(point1, point2, poly_points[k]);
   if(i_pos <= 0 && k_pos <= 0)
     new_points.push_back(poly_points[k]);
   else if(i_pos > 0 && k_pos <= 0) {
     new_points.push_back(intersect(point1, point2, poly_points[i],
                                     poly_points[k]));
     new_points.push_back(poly_points[k]);
   else if(i_pos <= 0 && k_pos > 0) {
     new_points.push_back(intersect(point1, point2, poly_points[i],
                                    poly_points[k]));
   else {
 poly_points.clear();
 for(int i = 0; i < new_points.size(); i++)</pre>
   poly_points.push_back(new_points[i]);
double area(const vector <point>& P) {
 double result = 0.0;
 double x1, y1, x2, y2;
 for(int i = 0; i < P.size() - 1; i++) {
   x1 = P[i].x;
   y1 = P[i].y;
   x2 = P[i + 1].x;
   y2 = P[i + 1].y;
   result += (x1 * y2 - x2 * y1);
 return fabs(result) / 2;
void suthHodgClip(vector <point>& poly_points, vector <point> clipper_points) {
 for(int i = 0; i < clipper_points.size(); i++) {</pre>
   int k = (i + 1) % clipper_points.size();
   clip(poly_points, clipper_points[i], clipper_points[k]);
vector<point> sortku(vector<point> P) {
 int P0 = 0;
 int i;
```

```
for(i = 1; i < 3; i++) +
  if(P[i].y < P[P0].y || (P[i].y == P[P0].y && P[i].x > P[P0].x))
point temp = P[0];
P[0] = P[P0];
P[P0] = temp;
pivot = P[0];
sort(++P.begin(), P.end(), anglecmp);
reverse(++P.begin(), P.end());
return P;
clipper_points = sortku(clipper_points);
suthHodgClip(poly_points, clipper_points);
```

## 4.6 Centroid of Polygon

```
C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_{i+1})(x_i \ y_{i+1} - x_{i+1} \ y_i)
C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_{i+1})(x_i \ y_{i+1} - x_{i+1} \ y_i)
```

#### 4.7 Pick Theorem

A: Area of a simply closed lattice polygon

B: Number of lattice points on the edges

I: Number of points in the interior

 $A = I + \frac{B}{2} - 1$ 

# 5 Graphs

## 5.1 Articulation Point and Bridge

```
const int SZ = 100005;
vector<int> to[SZ];
int vis[SZ], in[SZ], lw[SZ], n, T;
set<int> ap;
set<pii> bridge;
void tarjan(int u, int p = -1) {
    vis[u] = true;
    in[u] = lw[u] = ++T;
    int child = 0;
    for(int &v : to[u]) {
        if(v == p) continue;
        if(vis[v]) {
             lw[u] = min(lw[u], in[v]);
        } else {
            ++child;
            tarjan(v, u);
            lw[u] = min(lw[u], lw[v]);
            if(lw[v] >= in[u] && p != -1) ap.insert(u);
            if(lw[v] > in[u]) bridge.insert({u, v});
  if(p == -1 \&\& child > 1)
    ap.insert(u);
void getTarjan() {
    for(int i = 1 ; i <= n ; ++i) if(!vis[i]) {
        tarjan(i);
```

# 5.2 SCC and Strong Orientation

```
#define N 10020
vector<int> adj[N];
bool vis[N], ins[N];
int disc[N], low[N], gr[N];
stack<int> st;
int id, grid;
void scc(int cur, int par) {
 disc[cur] = low[cur] = ++id;
 vis[cur] = ins[cur] = 1;
 st.push(cur);
 for(int to : adj[cur]) {
   //if (to==par) continue; // ini untuk SO(scc undirected)
   if(!vis[to])
     scc(to, cur);
   if(ins[to])
      low[cur] = min(low[cur], low[to]);
 if(low[cur] == disc[cur]) {
   grid++; // group id
   while(ins[cur]) {
     gr[st.tp] = grid;
      ins[st.tp] = 0;
     st.pop();
```

# 5.3 Centroid Decomposition

```
int build_cen(int nw) {
 com_cen(nw, 0); //fungsi untuk itung size subtree
 int siz = sz[nw] / 2;
 bool found = false;
 while(!found) {
   found = true;
   for(int i : v[nw]) {
     if(!rem[i] && sz[i] < sz[nw]) {</pre>
       if(sz[i] > siz) {
         found = false;
         nw = i;
         break;
 big
 rem[nw] = true;
 for(int i : v[nw])if(!rem[i])
     par_cen[build_cen(i)] = nw;
 return nw;
```

## 5.4 Dinic's Maximum Flow

```
// O(VE log(max_flow)) if scaling == 1
// O((V + E) sqrt(E)) if unit graph (turn scaling off)
// O((V + E) sqrt(V)) if bipartite matching (turn scaling off)
// indices are 0-based
const ll INF = 1e18;
struct Dinic {
 struct Edge {
    int v;
   Edge(int _v, ll _cap): v(_v), cap(_cap), flow(0) {}
 int n;
```

```
ll lim;
  vector<vector<int>> gr;
  vector<Edge> e;
 vector<int> idx, lv;
  bool has_path(int s, int t) {
   queue<int> q;
   q.push(s);
    lv.assign(n, -1);
    lv[s] = 0;
    while(!q.empty())
     int c = q.front();
     q.pop();
     if(c == t)
       break;
      for(auto& i : gr[c]) {
       ll cur_flow = e[i].cap - e[i].flow;
       if(lv[e[i].v] == -1 && cur_flow >= lim) {
         lv[e[i].v] = lv[c] + 1;
          q.push(e[i].v);
   return lv[t] != -1;
  ll get_flow(int s, int t, ll left) {
    if(!left || s == t)
     return left;
   while(idx[s] < (int) gr[s].size()) {</pre>
     int i = gr[s][idx[s]];
      if(lv[e[i].v] == lv[s] + 1) {
        ll add = get_flow(e[i].v, t, min(left, e[i].cap - e[i].flow));
        if(add) {
         e[i].flow += add;
         e[i ^ 1].flow -= add;
         return add;
      ++idx[s];
   return 0;
 Dinic(int vertices, bool scaling = 1) : // toggle scaling here
   n(vertices), lim(scaling ? 1 << 30 : 1), gr(n) {}
  void add_edge(int from, int to, ll cap, bool directed = 1) {
   gr[from].push_back(e.size());
   e.emplace_back(to, cap);
   gr[to].push_back(e.size());
   e.emplace_back(from, directed ? 0 : cap);
  ll get_max_flow(int s, int t) { // call this
    ll res = 0;
   while(lim) { // scaling
     while(has_path(s, t)) {
       idx.assign(n, 0);
       while(ll add = get_flow(s, t, INF))
          res += add;
      lim >>= 1;
    return res;
};
```

## 5.5 Minimum Cost Maximum Flow

```
using FlowT = ll;
using CostT = ll;
const FlowT F_INF = 1e18;
const CostT C_INF = 1e18;
const int MAX_V = 1e5 + 5;
const int MAX_E = 1e6 + 5;
namespace MCMF {
 int n, E;
 int adj[MAX_E], nxt[MAX_E], lst[MAX_V], frm[MAX_V], vis[MAX_V];
 FlowT cap[MAX_E], flw[MAX_E], totalFlow;
 CostT cst[MAX_E], dst[MAX_V], totalCost;
 void init(int _n) {
   n = _n;
fill_n(lst, n, -1), E = 0;
 void add(int u, int v, FlowT ca, CostT co) {
   adj[E] = v, cap[E] = ca, flw[E] = 0, cst[E] = +co;
   nxt[E] = lst[u], lst[u] = E++;
   adj[E] = u, cap[E] = 0, flw[E] = 0, cst[E] = -co;
   nxt[E] = lst[v], lst[v] = E++;
  int spfa(int s, int t) {
   fill_n(dst, n, C_INF), dst[s] = 0;
   queue<int> que;
   que.push(s);
   while(que.size()) {
     int u = que.front();
     que.pop();
      for(int e = lst[u]; e != -1; e = nxt[e])
       if(flw[e] < cap[e]) {</pre>
          int v = adj[e];
         if(dst[v] > dst[u] + cst[e]) {
            dst[v] = dst[u] + cst[e];
            frm[v] = e;
            if(!vis[v]) {
             vis[v] = 1;
              que.push(v);
     vis[u] = 0;
   return dst[t] < C_INF;</pre>
 pair<FlowT, CostT> solve(int s, int t) {
   totalCost = 0, totalFlow = 0;
   while(1) {
      if(!spfa(s, t))break;
     FlowT mn = F_INF;
      for(int v = t, e = frm[v]; v != s; v = adj[e ^ 1], e = frm[v])
       mn = min(mn, cap[e] - flw[e]);
      for(int v = t, e = frm[v]; v != s; v = adj[e ^ 1], e = frm[v]) {
       flw[e] += mn;
       flw[e ^ 1] -= mn;
     totalFlow += mn;
     totalCost += mn * dst[t];
   return {totalFlow, totalCost};
};
```

#### 5.6 Flows with Demands

let SO be the source and TO be the original sink 1. add 2 additional nodes, call them S1 and T1

```
2. connect S0 to nodes normally
3. connect nodes to TO normally
4. for each edge(U, V), cap = original cap - demand
5. for each node N:
   1. add an edge(S1, N), cap = sum of inward demand to N
   2. add an edge(N, T1), cap = sum of outward demand from N
6. add an edge(T0, S0), cap = INF
7. the above is not a typo!
8. run max flow normally
9. for each edge(S1, V) and (U, T1), check if flow == cap
if step #9 fails, then it is not possible to satisfy the given demand
```

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Mathematically, let d(e) be the demand of edge e. Let V be the set of every vertex in the graph.

- $c'(S_1, v) = \sum_{u \in V} d(u, v)$  for each edge (s', v).
- $c'(v, T_1) = \sum_{v \in V} d(v, w)$  for each edge (v, t').
- c'(u,v) = c(u,v) d(u,v) for each edge (u,v) in the old network.
- $c'(T_0, S_0) = \infty$

## 5.7 Hungarian

```
template <typename TD> struct Hungarian {
 TD INF = le9; //max_inf
 int n;
 vector<vector<TD> > adj; // cost[left][right]
 vector<TD> hl, hr, slk;
 vector<int> fl, fr, vl, vr, pre;
 deque<int> a:
 Hungarian(int _n) {
   n = _n;
   adj = vector<vector<TD> >(n, vector<TD> (n, 0));
  int check(int i) {
   if(vl[i] = 1, fl[i] != -1)
     return q.push_back(fl[i]), vr[fl[i]] = 1;
   while(i !=-1)
     swap(i, fr[fl[i] = pre[i]]);
   return 0;
  void bfs(int s) {
   slk.assign(n, INF);
   vl.assign(n, 0);
   vr = vl;
   q.assign(vr[s] = 1, s);
    for(TD d;;) {
      for(; !q.empty(); q.pop_front()) {
        for(int i = 0, j = q.front(); i < n; i++) {</pre>
          if(d = hl[i] + hr[j] - adj[i][j], !vl[i] && d <= slk[i]) {
            if(pre[i] = j, d)
             slk[i] = d;
           else if(!check(i))return;
     d = INF;
      for(int i = 0; i < n; i++) if(!vl[i] && d > slk[i])
         d = slk[i];
      for(int i = 0; i < n; i++) {
       if(vl[i])hl[i] += d;
       else slk[i] -= d;
       if(vr[i])hr[i] -= d;
      for(int i = 0; i < n; i++) if(!vl[i] && !slk[i] && !check(i))</pre>
          return;
 TD solve() {
```

```
fl.assign(n, -1);
    fr = fl;
   hl.assign(n, 0);
   hr = hl;
    pre.assign(n, 0);
    for(int i = 0; i < n; i++)
     hl[i] = *max_element(adj[i].begin(), adj[i].begin() + n);
    for(int i = 0; i < n; i++)</pre>
     bfs(i);
    TD ret = 0;
    for(int i = 0; i < n; i++) if(adj[i][fl[i]])</pre>
        ret += adj[i][fl[i]];
    return ret;
}; //i wiLL be matched with fl[i]
```

### 5.8 Edmonds' Blossom

```
// Maximum matching on general graphs in O(V^2 E)
// Indices are 1-based
// Stolen from ko_osaga's cheatsheet
struct Blossom {
 vector<int> vis, dad, orig, match, aux;
 vector<vector<int>> conn;
 int t, N;
 queue<int> Q;
 void augment(int u, int v) {
   int pv = v;
   do {
     pv = dad[v];
     int nv = match[pv];
     match[v] = pv;
     match[pv] = v;
     \vee = n \vee ;
   } while(u != pv);
 int lca(int v, int w) {
   while(true) {
     if(v) {
       if(aux[v] == t)return v;
       aux[v] = t;
       v = orig[dad[match[v]]];
     swap(v, w);
 void blossom(int v, int w, int a) {
   while(orig[v] != a) {
     dad[v] = w;
     w = match[v];
     if(vis[w] == 1) {
       Q.push(w);
       vis[w] = 0;
     orig[v] = orig[w] = a;
     v = dad[w];
 bool bfs(int u) {
   fill(vis.begin(), vis.end(), -1);
   iota(orig.begin(), orig.end(), 0);
   Q = queue<int>();
   Q.push(u);
   vis[u] = 0;
   while(!Q.empty()) {
     int v = Q.front();
     Q.pop();
```

```
for(int x : conn[v]) {
     if(vis[x] == -1) {
       dad[x] = v;
       vis[x] = 1;
        if(!match[x]) {
         augment(u, x);
         return 1:
       Q.push(match[x]);
       vis[match[x]] = 0;
     } else if(vis[x] == 0 && orig[v] != orig[x]) {
        int a = lca(orig[v], orig[x]);
       blossom(x, v, a);
       blossom(v, x, a);
 return false;
Blossom(int n) : // n = vertices
 vis(n + 1), dad(n + 1), orig(n + 1), match(n + 1),
 aux(n + 1), conn(n + 1), t(0), N(n) {
 for(int i = 0; i <= n; ++i) {
   conn[i].clear();
   match[i] = aux[i] = dad[i] = 0;
void add_edge(int u, int v) {
 conn[u].push_back(v);
 conn[v].push back(u);
int solve() { // call this for answer
 int ans = 0;
 vector<int> V(N - 1);
 iota(V.begin(), V.end(), 1);
 shuffle(V.begin(), V.end(), mt19937(0x94949));
  for(auto x : V) {
   if(!match[x]) {
      for(auto y : conn[x]) {
       if(!match[y]) {
         match[x] = y, match[y] = x;
         ++ans;
         break;
  for(int i = 1; i <= N; ++i) {
   if(!match[i] && bfs(i))
     ++ans;
 return ans;
```

# 5.9 Eulerian Path or Cycle

```
// finds a eulerian path / cycle
// visits each edge only once
// properties:
// - cycle: degrees are even
// - path: degrees are even OR degrees are even except for 2 vertices
// how to use: g = adjacency list g[n] = connected to n, undirected
// if there is a vertex u with an odd degree, call dfs(u)
// else call on any vertex
// ans = path result
vector<set<int>> g;
```

```
vector<int> ans;
void dfs(int u) {
  while(g[u].size()) {
    int v = *g[u].begin();
    g[u].erase(v);
    g[v].erase(u);
    dfs(v);
}
  ans.push_back(u);
}
```

# 5.10 Hierholzer's Algorithm

```
// Eulerian on Directed Graph
stack<int> path;
vector<int> euler;
inline void hierholzer() {
 path.push(0);
  int cur = 0;
 while(!path.empty()) {
   if(!adj[cur].empty()) {
     path.push(cur);
      int next = adj[cur].back();
     adj[cur].pob();
     cur = next;
   } else {
     euler.pb(cur);
     cur = path.top();
     path.pop();
 reverse(euler.begin(), euler.end());
```

## 5.11 2-SAT

```
struct TwoSAT {
 int n;
 vector<vector<int>> g, gr;
 vector<int> comp, topological_order, answer;
 vector<bool> vis;
 TwoSAT() {}
 TwoSAT(int _n) :
   n(_n), g(2 * n), gr(2 * n), comp(2 * n), answer(2 * n), vis(2 * n) {}
 void add_edge(int u, int v) {
   g[u].push_back(v);
   gr[v].push_back(u);
 // For the following three functions
 // int x, bool val: if 'val' is true, we take the variable to be x.
 // Otherwise we take it to be x's complement.
 // At least one of them is true
 void add_clause_or(int i, bool f, int j, bool p) {
   add_{edge}(i + (f ? n : 0), j + (p ? 0 : n));
   add_{edge}(j + (p ? n : 0), i + (f ? 0 : n));
 // Only one of them is true
 void add_clause_xor(int i, bool f, int j, bool p) {
   add_clause_or(i, f, j, p);
   add_clause_or(i, !f, j, !p);
 // Both of them have the same value
 void add_clause_and(int i, bool f, int j, bool p) {
   add_clause_xor(i, !f, j, p);
 // Topological sort
 void dfs(int u) {
   vis[u] = true;
   for(const auto& v : g[u])
```

```
if(!vis[v])
       dfs(v);
    topological_order.push_back(u);
 // Extracting strongly connected components
  void scc(int u, int id) {
   vis[u] = true;
   comp[u] = id;
    for(const auto& v : gr[u])
     if(!vis[v])
       scc(v, id);
  bool satisfiable() {
    fill(vis.begin(), vis.end(), false);
    for(int i = 0; i < 2 * n; i++)
     if(!vis[i])
       dfs(i);
    fill(vis.begin(), vis.end(), false);
    reverse(topological_order.begin(), topological_order.end());
    for(const auto& v : topological_order)
     if(!vis[v])
       scc(v, id++);
    // Constructing the answer
    for(int i = 0; i < n; i++) {
     if(comp[i] == comp[i + n])
       return false:
     answer[i] = (comp[i] > comp[i + n] ? 1 : 0);
    return true;
};
```

## 6 Math

#### 6.1 Extended Euclidean GCD

```
// computes x and y such that ax + by = gcd(a, b) in O(log (min(a, b)))
// returns {gcd(a, b), x, y}
tuple<int, int, int> gcd(int a, int b) {
  if(b == 0) return {a, 1, 0};
  auto [d, x1, y1] = gcd(b, a % b);
  return {d, y1, x1 - y1* (a / b)};
}
```

#### 6.2 Generalized CRT

```
template<typename T>
T extended_euclid(T a, T b, T& x, T& y) {
  if(b == 0) {
    \times = 1;
    y = 0;
    return a;
  T xx, yy, gcd;
  gcd = extended_euclid(b, a % b, xx, yy);
  x = yy;
  y = xx - (yy * (a / b));
  return gcd;
template<typename T>
T MOD(T a, T b) {
 return (a % b + b) % b;
// return x, lcm. x = a % n && x = b % m
template<tvpename T>
pair<T, T> CRT(T a, T n, T b, T m) {
T _n, _m;
```

```
T gcd = extended_euclid(n, m, _n, _m);
if(n == m) {
    if(a == b)
        return pair<T, T>(a, n);
    else
        return pair<T, T>(-1, -1);
} else if(abs(a - b) % gcd != 0)
    return pair<T, T>(-1, -1);
else {
    T lcm = m * n / gcd;
    T x = MOD(a + MOD(n * MOD(_n * ((b - a) / gcd), m / gcd), lcm), lcm);
    return pair<T, T>(x, lcm);
}
```

#### 6.3 Generalized Lucas Theorem

```
/*Special Lucas : (n,k) % p^x
 fctp[n] = Product of the integers less than or equal
 to n that are not divisible by p
 Precompute fctp*/
LL p
LL E(LL n, int m) {
 LL tot = 0;
 while(n != 0)
  tot += n / m, n /= m;
 return tot;
LL funct(LL n, LL base) {
 LL ans = fast(fctp[base], n / base, base) * fctp[n % base] % base;
 return ans:
LL F(LL n, LL base) {
 LL ans = 1;
 while(n != 0) {
  ans = (ans * funct(n, base)) % base;
   n /= p;
 return ans;
LL special_lucas(LL n, LL r, LL base) {
 p = fprime(base);
 LL pow = E(n, p) - E(n - r, p) - E(r, p);
 LL TOP = fast(p, pow, base) * F(n, base) % base;
 LL BOT = F(r, base) * F(n - r, base) % base;
 return (TOP * fast(BOT, totien(base) - 1, base)) % base;
//End of Special Lucas
```

# 6.4 Linear Diophantine

```
//FOR SOLVING MINIMUM ABS(X) + ABS(Y)
ll x, y, newX, newY, target = 0;
ll extGcd(ll a, ll b) {
    if(b == 0) {
        x = 1, y = 0;
        return a;
    }
    ll ret = extGcd(b, a % b);
    newX = y;
    newY = x - y * (a / b);
    x = newX;
    y = newY;
    return ret;
}
ll fix(ll sol, ll rt) {
    ll ret = 0;
    //CASE SOLUTION(X/Y) < TARGET
    if(sol < target)ret = -floor(abs(sol + target) / (double)rt);</pre>
```

```
if(sol > target)ret = ceil(abs(sol - target) / (double)rt);
 return ret;
ll work(ll a, ll b, ll c) {
 ll gcd = extGcd(a, b);
 ll solX = x * (c / gcd);
 ll solY = y * (c / gcd);
 a /= gcd;
 b /= gcd;
 ll fi = abs(fix(solX, b));
 ll se = abs(fix(solY, a));
 ll lo = min(fi, se);
 ll hi = max(fi, se);
 ll \ ans = abs(solX) + abs(solY);
 for(ll i = lo; i <= hi; i++) {
   ans = min(ans, abs(solX + i * b) + abs(solY - i * a));
   ans = min(ans, abs(solX - i * b) + abs(solY + i * a));
 return ans;
```

# 6.5 Modular Linear Equation

```
// finds all solutions to ax = b (mod n)
vi modular_linear_equation_solver(int a, int b, int n) {
   int x, y;
   vi ret;
   int g = extended_euclid(a, n, x, y);
   if(!(b % g)) {
      x = mod(x * (b / g), n);
      for(int i = 0; i < g; i++)
      ret.push_back(mod(x + i * (n / g), n));
   }
   return ret;
}</pre>
```

#### 6.6 Miller-Rabin and Pollard's Rho

```
namespace MillerRabin {
 const vector<ll> primes = { // deterministic up to 2^64 - 1
   2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37
  ll gcd(ll a, ll b) {
   return b ? gcd(b, a % b) : a;
  ll powa(ll x, ll y, ll p) { // (x ^ y) % p
    if(!y)return 1;
   if(y & 1)
     return ((\_int128) \times powa(x, y - 1, p)) \% p;
   ll temp = powa(x, y >> 1, p);
    return ((__int128) temp * temp) % p;
  bool miller_rabin(ll n, ll a, ll d, int s) {
    ll \times = powa(a, d, n);
    if(x == 1 || x == n - 1) return 0;
    for(int i = 0; i < s; ++i) {
     x = ((__int128) \times * x) \% n;
     if(x == n - 1) return 0;
   return 1;
  bool is_prime(ll x) { // use this
   if(x < 2) return 0;
    int r = 0;
   ll d = x - 1;
    while((d & 1) == 0) {
     d >>= 1;
      ++r;
```

```
for(auto& i : primes) {
     if(x == i) return 1;
     if(miller_rabin(x, i, d, r))return 0;
   return 1;
namespace PollardRho {
 mt19937_64 generator(chrono::steady_clock::now()
                       .time_since_epoch().count());
 uniform_int_distribution<ll> rand_ll(0, LLONG_MAX);
 ll f(ll x, ll b, ll n) { // (x^2 + b) % n}
   return (((\_int128) \times \times \times) \% n + b) \% n;
 ll rho(ll n) {
   if(n \% 2 == 0) return 2;
   ll b = rand_ll(generator);
   ll x = rand_ll(generator);
   ll \vee = x;
   while(1) {
    x = f(x, b, n);
     y = f(f(y, b, n), b, n);
     il d = MillerRabin::gcd(abs(x - y), n);
     if(d != 1)return d;
 void pollard_rho(ll n, vector<ll>& res) {
   if(n == 1)return;
   if(MillerRabin::is_prime(n)) {
     res.push_back(n);
     return;
   ll d = rho(n);
   pollard_rho(d, res);
   pollard_rho(n / d, res);
 vector<ll> factorize(ll n, bool sorted = 1) { // use this
   vector<ll> res;
   pollard_rho(n, res);
   if(sorted)
     sort(res.begin(), res.end());
   return res;
```

# 6.7 Berlekamp-Massey

```
#include <bits/stdc++.h>
using namespace std;
#define pb push_back
typedef long long ll;
#define SZ 233333
const int MOD = 1e9 + 7; //or any prime
ll qp(ll a, ll b) {
 ll \times = 1;
 a %= MOD;
 while(b) {
   if(b & 1)
     x = x * a % MOD;
   a = a * a % MOD;
   b >>= 1;
 return x;
namespace linear_seq {
 vector<int> BM(vector<int> x) {
    //ls: (shortest) relation sequence (after filling zeroes) so far
    //cur: current relation sequence
```

```
vector<int> ls, cur;
    //lf: the position of ls (t')
    //ld: delta of ls (v')
    int lf = -1, ld = -1;
    for(int i = 0; i < int(x.size()); ++i) {
     ll t = 0;
     //evaluate at position i
      for(int j = 0; j < int(cur.size()); ++j)</pre>
       t = (t + x[i - j - 1] * (ll)cur[j]) % MOD;
      if((t - x[i]) \% MOD == 0) {
        continue; //good so far
      //first non-zero position
      if(!cur.size()) {
       cur.resize(i + 1);
        lf = i;
        ld = (t - x[i]) \% MOD;
        continue;
      //cur=cur-c/ld*(x[i]-t)
      ll k = -(x[i] - t) * ap(ld, MOD - 2) % MOD/*1/ld*/;
     vector<int> c(i - lf - 1); //add zeroes in front
     c.pb(k);
      for(int j = 0; j < int(ls.size()); ++j)</pre>
       c.pb(-ls[j]*k \% MOD);
      if(c.size() < cur.size())</pre>
       c.resize(cur.size());
      for(int j = 0; j < int(cur.size()); ++j)</pre>
       c[j] = (c[j] + cur[j]) \% MOD;
      //if cur is better than ls, change ls to cur
      if(i - lf + (int)ls.size() >= (int)cur.size())
        ls = cur, lf = i, ld = (t - x[i]) % MOD;
     cur = c;
    for(int i = 0; i < int(cur.size()); ++i)</pre>
     cur[i] = (cur[i] % MOD + MOD) % MOD;
    return cur;
 int m; //length of recurrence
//a: first terms
//h: relation
 ll a[SZ], h[SZ], t_[SZ], s[SZ], t[SZ];
//calculate p*q mod f
  void mull(ll* p, ll* q) {
    for(int i = 0; i < m + m; ++i)
     t_[i] = 0;
    for(int i = 0; i < m; ++i) if(p[i])</pre>
        for(int j = 0; j < m; ++j)
          t_{[i + j]} = (t_{[i + j]} + p[i] * q[j]) % MOD;
    for(int i = m + m - 1; i >= m; --i) if(t_[i])
        //miuns t_{i-m}(x^m-\sum_{j=0}^{m-1} x^{m-j-1}h_{j})
        for(int j = m - 1; ~j; --j)
          t_{i} - j - 1 = (t_{i} - j - 1) + t_{i} * h_{i} % MOD;
    for(int i = 0; i < m; ++i)
     p[i] = t_[i];
  ll calc(ll K) {
    for(int i = m; \sim i; --i)
     s[i] = t[i] = 0;
   s[0] = 1;
   if(m != 1)
     t[1] = 1;
   else
     t[0] = h[0];
    //binary-exponentiation
    while(K) {
     if(K & 1)
       mull(s, t);
      mull(t, t);
     K >> = 1;
```

```
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```

```
ll su = 0;
   for(int i = 0; i < m; ++i)
     su = (su + s[i] * a[i]) % MOD;
   return (su % MOD + MOD) % MOD;
 int work(vector<int> x, ll n) {
   if(n < int(x.size()))</pre>
     return x[n];
   vector < int > v = BM(x);
   m = v.size();
     return 0;
   for(int i = 0; i < m; ++i)
    h[i] = v[i], a[i] = x[i];
   return calc(n);
using linear_seq::work;
const vector<int> sequence = {
 0, 2, 2, 28, 60, 836, 2766
 cout << work(sequence, 7) << '\n';</pre>
```

#### 6.8 Catalan

$$C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{(n+1)!n!} = \prod_{k=2}^n \frac{n+k}{k}$$
 for  $n \ge 0$ 

# 6.9 Fast Fourier Transform

```
using ld = double; // change to long double if reach 10^18
using cd = complex<ld>;
const ld PI = acos(-(ld)1);
void fft(vector<cd>& a, int sign = 1) {
 int n = a.size();
 ld theta = sign * 2 * PI / n;
 for(int i = 0, j = 1; j < n - 1; j++) {
   for(int k = n >> 1; k > (i ^= k); k >>= 1);
   if(j < i)
     swap(a[i], a[j]);
 for(int m, mh = 1; (m = mh << 1) <= n; mh = m) {
   int irev = 0;
   for(int i = 0; i < n; i += m) {
     cd w = exp(cd(0, theta * irev));
     for(int k = n >> 2; k > (irev ^= k); k >>= 1);
      for(int j = i; j < mh + i; j++) {
       int k = j + mh;
       cd \times = a[j] - a[k];
       a[j] += a[k];
       a[k] = w * x;
 if(sign == -1) for(cd& i : a)
     i /= n;
vector<ll> multiply(vector<ll> const& a, vector<ll> const& b) {
 vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());
 int n = 1;
 while(n < a.size() + b.size())</pre>
   n <<= 1;
```

```
fa.resize(n);
fb.resize(n);
fft(fa);
fft(fb);
for(int i = 0; i < n; i++)
    fa[i] *= fb[i];
fft(fa, -1);
vector<ll> res(n);
for(int i = 0; i < n; i++)
    res[i] = round(fa[i].real());
return res;
}</pre>
```

#### 6.10 Centroid

$$C_x = \frac{1}{6A} \sum_{i=0}^{n-1} (x_i + x_i(i+1))(x_i y_i(i+1) - x_i(i+1)y_i)$$

$$C_y = \frac{1}{6A} \sum_{i=0}^{n-1} (y_i + y_i(i+1))(x_i y_i(i+1) - x_i(i+1)y_i)$$

#### 6.11 Number Theoretic Transform

```
namespace FFT {
 /* ---- Adjust the constants here ---- */
 const int LN = 24; \frac{1}{23}
 const int N = 1 << LN;</pre>
  typedef long long LL; // 2**23 * 119 + 1. 998244353
^{\prime}/ ^{\circ}MOD^{\circ} must be of the form 2**^{\circ}LN^{\circ} * k + 1, where k odd.
  const LL MOD = 9223372036737335297; // 2**24 * 54975513881 + 1.
  const LL PRIMITIVE_ROOT = 3; // Primitive root modulo `MOD`.
  /* ---- End of constants ---- */
 LL root[N];
  inline LL power(LL x, LL y) {
   LL ret = 1;
    for(; y; y >>= 1) {
     if(y & 1)
        ret = (__int128) ret * x % MOD;
      x = (\_int128) \times \times \times \% MOD;
    return ret;
  inline void init_fft() {
    const LL UNITY = power(PRIMITIVE_ROOT, MOD - 1 >> LN);
    root[0] = 1;
    for(int i = 1; i < N; i++)
     root[i] = (__int128) UNITY * root[i - 1] % MOD;
    return;
// n = 2^k is the length of polynom
  inline void fft(int n, vector<LL>& a, bool invert) {
    for(int i = 1, j = 0; i < n; ++i) {
      int bit = n >> 1;
      for(; j >= bit; bit >>= 1)
        j -= bit;
      j += bit;
      if(i < j)
        swap(a[i], a[j]);
    for(int len = 2; len <= n; len <<= 1) {
      LL wlen = (invert ? root[N - N / len] : root[N / len]);
      for(int i = 0; i < n; i += len) {
        LL w = 1;
        for(int j = 0; j<len >> 1; j++) {
          LL u = a[i + j];
          LL v = (__int128) a[i + j + len / 2] * w % MOD;
          a[i + j] = ((\_int128) u + v) % MOD;
```

```
a[i + j + len / 2] = ((\_int128) u - v + MOD) % MOD;
         w = (__int128) w * wlen % MOD;
   if(invert) {
     LL inv = power(n, MOD - 2);
     for(int i = 0; i < n; i++)
       a[i] = (\_int128) a[i] * inv % MOD;
   return;
 inline vector<LL> multiply(vector<LL> a, vector<LL> b) {
   int len = 1 << 32 - __builtin_clz(a.size() + b.size() - 2);</pre>
   a.resize(len, 0);
   b.resize(len, 0);
   fft(len, a, false);
   fft(len, b, false);
   c.resize(len);
   for(int i = 0; i < len; ++i)
     c[i] = (\_int128) a[i] * b[i] % MOD;
   fft(len, c, true);
   return c;
//FFT::init_fft(); wajib di panggil init di awal
```

## 6.12 Derangement

$$D(i) = (i-1) * (D(i-1) + D(i-2))$$
$$D(0) = 1, D(1) = 0$$

## 6.13 Gauss-Jordan

```
// Gauss-Jordan elimination with full pivoting.
    (1) solving systems of linear equations (AX=B)
    (2) inverting matrices (AX=I)
    (3) computing determinants of square matrices
  Running time: O(n^3)
             a[][] = an nxn matrix
             b[][] = an nxm matrix
                    = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
             returns determinant of a[][]
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan(VVT& a, VVT& b) {
 const int n = a.size();
 const int m = b[0].size();
 VI irow(n), icol(n), ipiv(n);
 T det = 1;
 for(int i = 0; i < n; i++) {
   int pj = -1, pk = -1;
   for(int j = 0; j < n; j++) if(!ipiv[j])</pre>
       for(int k = 0; k < n; k++) if(!ipiv[k])</pre>
            if(pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) {
             pj = j;
              pk = k;
```

```
if(fabs(a[pj][pk]) < EPS) {
     cerr << "Matrix is singular." << endl;</pre>
     exit(0);
    ipiv[pk]++;
   swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
   if(pj != pk)
     det *= -1;
    irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
   det *= a[pk][pk];
   a[pk][pk] = 1.0;
    for(int p = 0; p < n; p++)
     a[pk][p] *= c;
    for(int p = 0; p < m; p++)
     b[pk][p] *= c;
    for(int p = 0; p < n; p++) if(p != pk) {
       c = a[p][pk];
       a[p][pk] = 0;
        for(int q = 0; q < n; q++)
         a[p][q] -= a[pk][q] * c;
        for(int q = 0; q < m; q++)
         b[p][q] -= b[pk][q] * c;
  for(int p = n - 1; p \ge 0; p--) if(irow[p] != icol[p]) {
      for(int k = 0; k < n; k++)
       swap(a[k][irow[p]], a[k][icol[p]]);
 return det;
int main() {
 const int n = 4;
 const int m = 2;
 double A[n][n] = { {1, 2, 3, 4}, {1, 0, 1, 0}, {5, 3, 2, 4}, {6, 1, 4, 6} };
 double B[n][m] = \{ \{1, 2\}, \{4, 3\}, \{5, 6\}, \{8, 7\} \};
 VVT a(n), b(n);
  for(int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);
 double det = GaussJordan(a, b);
  // expected: 60
 cout << "Determinant: " << det << endl;</pre>
  // expected: -0.233333 0.166667 0.133333 0.0666667
               0.166667 0.166667 0.333333 -0.333333
               0.233333 0.833333 -0.133333 -0.0666667
               0.05 -0.75 -0.1 0.2
 cout << "Inverse: " << endl;</pre>
  for(int i = 0; i < n; i++)
   for(int j = 0; j < n; j++)
     cout << a[i][j] << ' ';
   cout << endl;</pre>
  // expected: 1.63333 1.3
               -0.166667 0.5
               2.36667 1.7
               -1.85 -1.35
  cout << "Solution: " << endl;
  for(int i = 0; i < n; i++)
   for(int j = 0; j < m; j++)</pre>
     cout << b[i][j] << ' ';
   cout << endl;</pre>
```

#### 6.14 Power

$$\sum_{k=1}^{n} k^4 = \frac{1}{30} (6n^5 + 15n^4 + 10n^3 - n) = \frac{1}{30} n(n+1)(2n+1)(3n^2 + 3n + 1)$$

$$\sum_{k=1}^{n} k^5 = \frac{1}{12} (2n^6 + 6n^5 + 5n^4 - n^2) = \frac{1}{12} n^2 (n+1)^2 (2n^2 + 2n - 1)$$

$$\sum_{k=1}^{n} k^6 = \frac{1}{42} (6n^7 + 21n^6 + 21n^5 - 7n^3 + n) = \frac{1}{42} n(n+1)(2n+1)(3n^4 + 6n^3 - 3n + 1)$$

## 6.15 Stirling

$$S(m,n) = \frac{1}{n!} \sum_{i=0}^{n} (-1)^{i} \binom{n}{i} (n-i)^{m}$$

#### 6.16 Bernoulli Number

$$\sum_{k=1}^{n} k^{m} = \frac{1}{m+1} \sum_{i=0}^{m} {m+1 \choose i} B_{i}^{+} n^{m+1-i} = m! \sum_{i=0}^{m} \frac{B_{i}^{+} n^{m+1-i}}{i!(m+1-i)!} B_{n}^{+} = 1 - \sum_{i=0}^{n-1} {n \choose i} \frac{B_{i}^{+}}{n-i+1}, \quad B_{0}^{+} = 1$$

#### 6.17 Forbenius Number

(X \* Y) - (X + Y) and total count is(X - 1) \* (Y - 1) / 2

## 6.18 Stars and Bars with Upper Bound

$$P = (1 - X^{r_1+1}) \dots (1 - X^{r_n+1}) = \sum_i c_i X^{e_i}$$
$$Ans = \sum_i c_i {N - e_i + n - 1 \choose n - 1}$$

## 6.19 Arithmetic Sequences

$$U_n = a + (n-1)a_1 + \frac{(n-1)(n-2)}{1 \times 2}a_2 + \dots + \frac{(n-1)(n-2)(n-3)\dots}{1 \times 2 \times 3 \times \dots}a_r$$
$$S_n = n \times a + \frac{n(n-1)}{1 \times 2}a_1 + \frac{n(n-1)(n-2)}{1 \times 2 \times 3}a_2 + \dots + \frac{n(n-1)(n-2)(n-3)\dots}{1 \times 2 \times 3\dots}a_r$$

#### 6.20 FWHT

```
// Desc : Transform a polynom to obtain a_i * b_j * x^(i XOR j) or combinations
// Time : O(N \log N) with N = 2^K
// OP => c00 c01 c10 c11 | c00 c01 c10 c11 inv
// XOR => +1 +1 +1 -1 \mid +1 +1 +1 -1 \mid div the inverse with size = n
// AND => 1 +1 0 1 1 -1 0 1 no comment
// OR => 1 0 +1 1 1 0 -1 1 no comment
typedef vector<long long> vec;
void FWHT(vec& a) {
 int n = a.size();
 for(int lvl = 1; 2 * lvl <= n; lvl <<= 1) {
   for(int i = 0; i < n; i += 2 * lvl) {
     for(int j = 0; j < lvl; j++) { // do not forget to modulo</pre>
       long long u = a[i + j], v = a[i + lvl + j];
       a[i + j] = u + v; // c00 * u + c01 * v
       a[i + v] + j] = u - v; // c10 * u + c11 * v
  // you can convolve as usual
```

## 6.21 Division-Polynom

```
const int M=530010+5;
inline int fastex(int x, int y) {
    int ret = 1;
        if(y & 1) ret = 1ll * ret * x % MOD;
        x = 111 * x * x * x % MOD; y >>= 1;
    return ret;
int rev[M], w[M], g[M],h[M], f[M], l[M];
inline void NTT(int *a, int N) {
    for(int i = 0; i < N; ++ i) {
        if(rev[i] > i) {
            swap(a[rev[i]], a[i]);
    for(int d = 1, t = (N >> 1); d < N; d <<= 1, t >>= 1) {
        for(int i = 0; i < N; i += (d << 1)) {
            for(int j = 0; j < d; ++ j) {
                int tmp = 111 * w[t * j] * a[i + j + d] % MOD;
                a[i + j + d] = a[i + j] - tmp + MOD; if(a[i + j + d] >= MOD) a[i + j + \leftarrow
                a[i + j] = a[i + j] + tmp; if(a[i + j] >= MOD) a[i + j] -= MOD;
inline void get_mul(int *f, int *g, int n, int m, int is_inv) {
   static int a[M], b[M];
    int N = 1, L = 0;
    for(; N < (n + m); N <<= 1, ++ L);
    for(int i = 1; i < N; ++ i) {
       rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (L - 1));
   w[0] = 1; w[1] = fastex(3, (MOD - 1) / N);
    for(int i = 2; i < N; ++ i) {
       w[i] = 1ll * w[i - 1] * w[1] % MOD;
    for(int i = 0; i < N; ++ i) {
       a[i] = b[i] = 0;
    for(int i = 0; i < n; ++ i) {
       a[i] = f[i];
    for(int i = 0; i < m; ++ i) {</pre>
       b[i] = g[i];
   NTT(a, N), NTT(b, N);
    for(int i = 0; i < N; ++ i) {
        if(is_inv) {
           a[i] = 111 * b[i] * (211 - 111 * a[i] * b[i] % MOD + MOD) % MOD;
       else -
            a[i] = 1ll * a[i] * b[i] % MOD;
   w[1] = fastex(w[1], MOD - 2);
    for(int i = 2; i < N; ++ i) {
       w[i] = 111 * w[i - 1] * w[1] % MOD;
   NTT(a, N);
    int inv = fastex(N, MOD - 2);
    for(int i = 0; i < n; ++ i) {
       a[i] = 1ll * a[i] * inv % MOD;
    for(int i = 0; i < n; ++ i) {
        if(is_inv) g[i] = a[i];
        else f[i] = a[i];
```

```
inline void get_inv(int *f, int *g, int n) {
   if(n == 1)
       g[0] = fastex(f[0], MOD - 2);
        return;
   get_inv(f, g, (n + 1) / 2);
   get_mul(f, g, n, n, 1);
inline void get_deri(int *f, int *g, int n) {
   for(int i = 1; i < n; ++ i) {
       g[i - 1] = 111 * f[i] * i % MOD;
   g[n - 1] = 0;
inline void get_inte(int *f, int *g, int n) {
   for(int i = 1; i < n; ++ i)
       g[i] = 1ll * f[i - 1] * fastex(i, MOD - 2) % MOD;
   g[0] = 0;
inline void get_ln(int *f, int *g, int n) {
   static int a[M], b[M];
   for(int i = 0; i < n; ++ i) a[i] = b[i] = 0;</pre>
   get_deri(f, a, n); get_inv(f, b, n);
   get_mul(a, b, n, n, 0);
   get_inte(a, g, n);
inline void get_exp(int *f, int *g, int n) {
   static int a[M], b[M];
    for(int i = 0; i < n; ++ i) a[i] = b[i] = 0;
   if(n == 1) {
       g[0] = 1;
       return;
   get_exp(f, g, (n + 1) / 2);
   get_ln(g, a, n);
    for(int i = 0; i < n; ++ i) {
       b[i] = (f[i] - a[i] + MOD);
       if(b[i] >= MOD) b[i] -= MOD;
   b[0] ++; if(b[0] >= MOD) b[0] -= MOD;
   get_mul(g, b, n, n, 0);
inline void get_pow(int *f, int *g, int n, int k, int k1) {
   static int a[M], b[M];
    int t = -1;
    for(int i = 0; i < n; ++ i) {
        if(f[i] != 0) {
           t = i;
           break;
   if(t == -1) {
       for(int i = 0; i < n; ++ i) {
           g[i] = 0;
       return;
   int inv = fastex(f[t], MOD - 2), pp = fastex(f[t], k1);
    for(int i = 0; i < n; ++ i) {
       a[i] = b[i] = 0;
    for(int i = 0; i < n - t; ++ i) {
       b[i] = 1ll * f[i + t] * inv % MOD;
   get_ln(b, a, n);
   for(int i = 0; i < n; ++ i) {
       a[i] = 1ll * a[i] * k % MOD;
   get_exp(a, g, n);
```

```
int lim = min(1ll * t * k, 1ll * n);
for(int i = n - 1; i >= lim; -- i) {
    g[i] = 111 * g[i - 111 * t * k] * pp % MOD;
for(int i = 0; i < lim; ++ i) {
   g[i] = 0;
```

#### 6.22 Primitive-Root

```
//cari g terkecil dimana g^k = 1 mod p dan k=phi(p)
//cari faktor dari phi(p) cek setiap angka dari 2 sampai p apakah semua fastex(res, ↔
    phi(p)/faktor) mod p !=1
```

# 7 Strings

#### 7.1 Aho-Corasick

```
const int K = 26:
struct Vertex {
public:
    int go[K], next[K], p = -1, link = -1, exit_link;
    bool leaf = false;
    char pch;
    vector<int> idx;
    Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
        fill(begin(next), end(next), -1);
        fill(begin(go), end(go), -1);
        exit_link = -1;
class Aho {
public:
    vector<Vertex> t = vector<Vertex>(1);
    vector<vector<int>> occ;
    vector<string> pat;
    string txt;
    void add_string(int num, string &s) {
        int \vee = 0;
        for(char ch : s) {
            int c = ch - 'a';
            if(t[v].next[c] == -1) {
                t[v].next[c] = t.size();
                t.emplace_back(v, ch);
            v = t[v].next[c];
        t[v].leaf = true;
        t[v].idx.pb(num);
    int get_link(int v) {
        if(t[v].link == -1) {
            if(v == 0 || t[v].p == 0) t[v].link = 0;
            else t[v].link = go(get_link(t[v].p), t[v].pch);
        return t[v].link;
    int go(int v, char ch) {
        int c = ch - 'a';
        if(t[v].go[c] == -1) {
            if(t[v].next[c] != -1) t[v].go[c] = t[v].next[c];
            else t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
        return t[v].go[c];
```

```
int find_exit(int v){
       if(t[v].exit_link != -1) return t[v].exit_link;
        if(v == 0) return 0;
        int nxt = get_link(v);
       if(t[nxt].idx.size()) return nxt;
       return t[v].exit_link = find_exit(nxt);
   void add_occur(int v, int i){
       for(int &x : t[v].idx){
           occ[x].pb(i - pat[x].length() + 1);
       if(v == 0) return ;
       add_occur(find_exit(v), i);
};
```

## 7.2 Eertree

/\*

```
Eertree - keep track of all palindromes and its occurences
   This code refers to problem Longest Palindromic Substring
https://www.spoj.com/problems/LPS/
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
struct node {
 int next[26];
 int sufflink:
 int len, cnt;
const int N = 1e5 + 69;
int n;
string s;
node tree[N]:
int idx, suff;
int ans = 0;
void init_eertree() {
 idx = suff = 2;
 tree[1].len = -1, tree[1].sufflink = 1;
 tree[2].len = 0, tree[2].sufflink = 1;
bool add_letter(int x) {
  int cur = suff, curlen = 0;
  int nw = s[x] - 'a';
  while(1) {
    curlen = tree[cur].len;
    if(x - curlen - 1 >= 0 \&\& s[x - curlen - 1] == s[x])
      break:
    cur = tree[cur].sufflink;
  if(tree[cur].next[nw]) {
    suff = tree[cur].next[nw];
    return 0;
  tree[cur].next[nw] = suff = ++idx;
  tree[idx].len = tree[cur].len + 2;
  ans = max(ans, tree[idx].len);
  if(tree[idx].len == 1) -
    tree[idx].sufflink = 2;
    tree[idx].cnt = 1;
    return 1;
  while(1) {
    cur = tree[cur].sufflink;
    curlen = tree[cur].len;
    if(x - curlen - 1 >= 0 \&\& s[x - curlen - 1] == s[x]) {
```

```
tree[idx].sufflink = tree[cur].next[nw];
     break;
 tree[idx].cnt = tree[tree[idx].sufflink].cnt + 1;
 return 1;
int main() {
 ios::sync_with_stdio(0);
 cin.tie(0);
 cin >> n >> s;
 init_eertree();
 for(int i = 0; i < n; i++)
   add_letter(i);
 cout << ans << '\n';
 return 0;
```

# 7.3 Manacher's Algorithm

```
void oddManacher(vector<int> &dl, string &s){
    int n = s.length(), l = 0, r = -1;
    d1 = vector < int > (n, 1);
    for(int i = 0; i < n; ++i){
        if(i \le r){
            int idx = l + r - i;
            d1[i] = min(d1[idx], r - i + 1);
        while(i + d1[i] < n && i - d1[i] >= 0 && s[i + d1[i]] == s[i - d1[i]]) ++d1[i↔
        if(i + d1[i] - 1 > r){
            r = i + d1[i] - 1;
            l = i - d1[i] + 1;
void evenManacher(vector<int> &d2, string &s){
    int n = s.length(), l = 0, r = -1;
    d2 = vector < int > (n, 0);
    for(int i = 0 ; i < n ; ++i){}
        if(i \le r){
            int idx = l + r - i;
            d2[i] = min(d2[idx], r - i + 1);
        while(i + d2[i] < n && i - d2[i] - 1 >= 0 && s[i + d2[i]] == s[i - d2[i] - 1])\leftrightarrow
              ++d2[i];
        if(i + d2[i] - 1 > r){
            r = i + d2[i] - 1;
            l = i - d2[i];
```

# 7.4 Suffix Array

```
const int VAL = 200005; // max(MXVAL, SZ)
const int SZ = 200005; // s.length()
const int LG = 20;
vector<int> pos[SZ], c[LG], p, pn;
map<int, int> nv;
int n, s[SZ], a[SZ];
int cnt[VAL];
vector<int> bldSA() {
 for(int i = 0; i < LG; ++i) c[i] = vector < int > (n << 2, 0);
    pn = vector < int > (n << 2, 0); p = vector < int > (n << 2, 0);
    for(int i = 0; i < n; ++i) c[0][i] = s[i];
    for(int x = 1, add = 1; add < n; add <<= 1, x += 1) {
        memset(cnt, 0, sizeof(cnt));
        for(int i = 0; i < n; ++i) ++cnt[c[x - 1][i + add]];
```

```
for(int i = 1 ; i < VAL ; ++i) cnt[i] += cnt[i - 1];</pre>
        for(int i = n - 1; i >= 0; --i) p[--cnt[c[x - 1][i + add]]] = i;
       memset(cnt, 0, sizeof(cnt));
        for(int i = 0; i < n; ++i) ++cnt[c[x - 1][i]];
        for(int i = 1; i < VAL; ++i) cnt[i] += cnt[i - 1];
        for(int i = n - 1; i \ge 0; --i) pn[--cnt[c[x - 1][p[i]]]] = p[i];
        c[x][pn[0]] = 1;
        for(int i = 1 ; i < n ; ++i) {
            c[x][pn[i]] = c[x][pn[i - 1]] + (c[x - 1][pn[i]] != c[x - 1][pn[i - 1]] ||
                          c[x - 1][pn[i] + add] != c[x - 1][pn[i - 1] + add]);
    return pn;
vector<int> kasai(string &txt, vector<int> &sa) {
    int n = txt.size();
    vector<int> lcp(n, 0), invSuff(n, 0);
    for (int i=0; i < n; i++)
        invSuff[sa[i]] = i;
    int k = 0;
    for (int i = 0; i < n; i++) {
        if (invSuff[i] == n-1){
            k = 0; continue;
        int j = sa[invSuff[i]+1];
       while (i + k < n \&\& j + k < n \&\& txt[i + k] == txt[j + k])
        lcp[invSuff[i]] = k;
        if (k > 0) k--;
   return lcp;
bool check(int i, int j) {
   int len = j - i;
  for(int x = LG - 1 ; x >= 0 ; --x) {
        if(len < (1<<x)) continue;</pre>
        if(c[x][i] == c[x][j]) {
            i += (1 << x); i += (1 << x);
            len -= (1 << x);
   return !len;
7.5 Suffix Automaton
```

```
struct state {
 int len, link;
 map<char, int>next; //use array if TLE
const int MAXLEN = 100005;
state st[MAXLEN * 2];
int sz, last;
void sa_init()
 sz = last = 0;
 st[0].len = 0;
 st[0].link = -1;
 st[0].next.clear();
 ++SZ;
void sa_extend(char c) {
 int cur = sz++;
 st[cur].len = st[last].len + 1;
 st[cur].next.clear();
 int p;
 for(p = last; p != -1 && !st[p].next.count(c); p = st[p].link)
   st[p].next[c] = cur;
 if(p == -1)st[cur].link = 0;
 else {
   int q = st[p].next[c];
```

```
if(st[p].len + 1 == st[q].len)st[cur].link = q;
    else {
      int clone = sz++;
      st[clone].len = st[p].len + 1;
      st[clone].next = st[q].next;
      st[clone].link = st[q].link;
      for(; p != -1 \&\& st[p].next[c] == q; p = st[p].link)
        st[p].next[c] = clone;
      st[q].link = st[cur].link = clone;
  last = cur;
// forwarding
for(int i = 0; i < m; i++) {
 while(cur >= 0 && st[cur].next.count(pa[i]) == 0) {
    cur = st[cur].link;
    if(cur != -1)len = st[cur].len;
  if(st[cur].next.count(pa[i])) {
    cur = st[cur].next[pa[i]];
  } else len = cur = 0;
// shortening abc -> bc
if(l == m) {
  if(l <= st[st[cur].link].len)cur = st[cur].link;</pre>
// finding lowest and highest length
int lo = st[st[cur].link].len + 1;
int hi = st[cur].len;
//Finding number of distinct substrings
//answer = distsub(0)
LL d[MAXLEN * 2];
LL distsub(int ver) {
 LL tp = 1;
  if(d[ver])return d[ver];
  for(map<char, int>::iterator it = st[ver].next.begin();
      it != st[ver].next.end(); it++)
    tp += distsub(it->second);
  d[ver] = tp;
  return d[ver];
//Total Length of all distinct substrings
//call distsub first before call lesub
LL ans[MAXLEN * 2];
LL lesub(int ver) {
 LL tp = 0;
  if(ans[ver])return ans[ver];
  for(map<char, int>::iterator it = st[ver].next.begin();
      it != st[ver].next.end(); it++)
    tp += lesub(it->second) + d[it->second];
  ans[ver] = tp;
  return ans[ver];
//find the k-th lexicographical substring
void kthsub(int ver, int K, string& ret)
  for(map<char, int>::iterator it = st[ver].next.begin();
      it != st[ver].next.end(); it++) {
    int v = it->second;
    if(K \le d[v]) {
      K--;
      if(K == 0) {
        ret.push_back(it->first);
        return;
      } else {
        ret.push_back(it->first);
        kthsub(v, K, ret);
        return;
```

```
} else
     K -= d[v];
// Smallest Cyclic Shift to obtain lexicographical smallest of All possible
int main() {
 string S;
 sa_init();
 cin >> S; //input
 tp = 0;
 t = S.length();
 S += S;
 for(int a = 0; a < S.size(); a++)
  sa_extend(S[a]);
 minshift(0);
//the function
int tp, t;
void minshift(int ver) {
 for(map<char, int>::iterator it = st[ver].next.begin();
     it != st[ver].next.end(); it++) {
   tp++;
   if(tp == t) {
     cout << st[ver].len - t + 1 << endl;
     break:
   minshift(it->second);
   break;
// LONGEST COMMON SUBSTRING OF TWO STRINGS
string lcs(string s, string t) {
 sa_init();
 for(int i = 0; i < (int)s.length(); ++i)</pre>
   sa_extend(s[i]);
 int v = 0, l = 0,
     best = 0, bestpos = 0;
 for(int i = 0; i < (int)t.length(); ++i) {</pre>
   while(v && ! st[v].next.count(t[i])) {
     v = st[v].link;
     l = st[v].length;
   if(st[v].next.count(t[i])) {
     v = st[v].next[t[i]];
     ++l;
   if(l > best)
     best = l, bestpos = i;
```

## OEIS

#### 8.1 A000127

return t.substr(bestpos - best + 1, best);

```
Maximal number of regions obtained by joining n points around a circle by
straight lines
f(n) = (n^4 - 6*n^3 + 23*n^2 - 18*n + 24) / 24
1, 2, 4, 8, 16, 31, 57, 99, 163, 256, 386, 562, 794, 1093, 1471, 1941, 2517,
3214, 4048, 5036, 6196, 7547, 9109, 10903, 12951, 15276, 17902, 20854, 24158,
27841, 31931, 36457, 41449, 46938, 52956, 59536, 66712, 74519, 82993, 92171,
102091, 112792, 124314
```

```
Number of graphs with n nodes and n edges.
0, 0, 1, 2, 6, 21, 65, 221, 771, 2769, 10250, 39243, 154658, 628635, 2632420,
11353457, 50411413, 230341716, 1082481189, 5228952960, 25945377057,
132140242356, 690238318754
```

#### 8.3 A018819

```
Binary partition function: number of partitions of n into powers of 2
f(2m+1) = f(2m); f(2m) = f(2m-1) + f(m)
1, 1, 2, 2, 4, 4, 6, 6, 10, 10, 14, 14, 20, 20, 26, 26, 36, 36, 46, 46, 60,
60, 74, 74, 94, 94, 114, 114, 140, 140, 166, 166, 202, 202, 238, 238, 284,
284, 330, 330, 390, 390, 450, 450, 524, 524, 598, 598, 692, 692, 786, 786,
900, 900, 1014, 1014, 1154, 1154, 1294, 1294
```

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#### 8.4 A092098

```
3-Portolan numbers: number of regions formed by n-secting the angles of
an equilateral triangle.
long long solve(long long n) {
    long long res = (n \% 2 == 1 ? 3*n*n - 3*n + 1 : 3*n*n - 6*n + 6);
    const int bats = n/2 - 1;
    for (long long i=1; i<=bats; i++) for (long long j=1; j<=bats; j++) {
        long long num = i * (n-j) * n;
        long long denum = (n-i) * j + i * (n-j);
        res -= 6 * (num % denum == 0 && num / denum <= bats);
    } return res;
1, 6, 19, 30, 61, 78, 127, 150, 217, 246, 331, 366, 469, 510, 625, 678, 817,
870, 1027, 1080, 1261, 1326, 1519, 1566, 1801, 1878, 2107, 2190, 2437, 2520,
2791, 2886, 3169, 3270, 3559, 3678, 3997, 4110, 4447, 4548, 4921, 5034, 5419,
5550, 5899, 6078, 6487
```