



TIC 2151 THEORY OF COMPUTATION
TRIMESTER 2 2018/2019

GROUP ASSIGNMENT

Prepared By

Tutorial Section: TT01 / Group 1

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1.0 Introduction

The objective of this assignment is to test the understanding of student on how to convert the concepts of the subjects 'Theory of Computation'(TIC2151) which is lead by Dr.Nbhan.D.Salih into a computer program. In this case, the knowledge of **visual basic programming** will be applied to presenting the **Graphical User Interface (GUI) program** to the users.

In this report, there are two parts of the program will be discussed which are stated below. In this case, the design flow chart and manual to operate the program will be explained as detail.

- **Part 1:** Conversion between Regular Grammar & Non-Deterministic Finite Automaton.
- **Part 2:** Convert the Context-Free Grammar(CFG) into the Chomsky Normal Form(CNF) & produce a CYK chart based on previous CNF.

A **regular grammar** is a formal grammar that is either left-regular or right-regular and every regular grammar describes a regular language. Regular grammar is a more powerful way to describe languages than finite automata. An NFA or Non-Deterministic Finite Automaton is a finite state machine which does not need to have each of its transitions uniquely determined by its source state and input symbol, as well as not needing the reading of an input symbol for each state transition.

A **context-free grammar** is an even more powerful method to describe languages in a way that it can describe certain features that have a recursive structure which is crucial in a certain application. Generally, context-free grammar often can have multiple simplified forms and one of the simplest and most useful forms is called the Chomsky normal form. A context-free grammar is successfully converted into the Chomsky normal form when every rule of the form is in; $A \rightarrow BC$, $A \rightarrow a$ where a is any terminal and A , B , and C are any variables except that B and C will not be the start variable. To determine if a word generated is part of a grammar, given a CNF grammar.

2.0 Design Flowcharts

PART I :

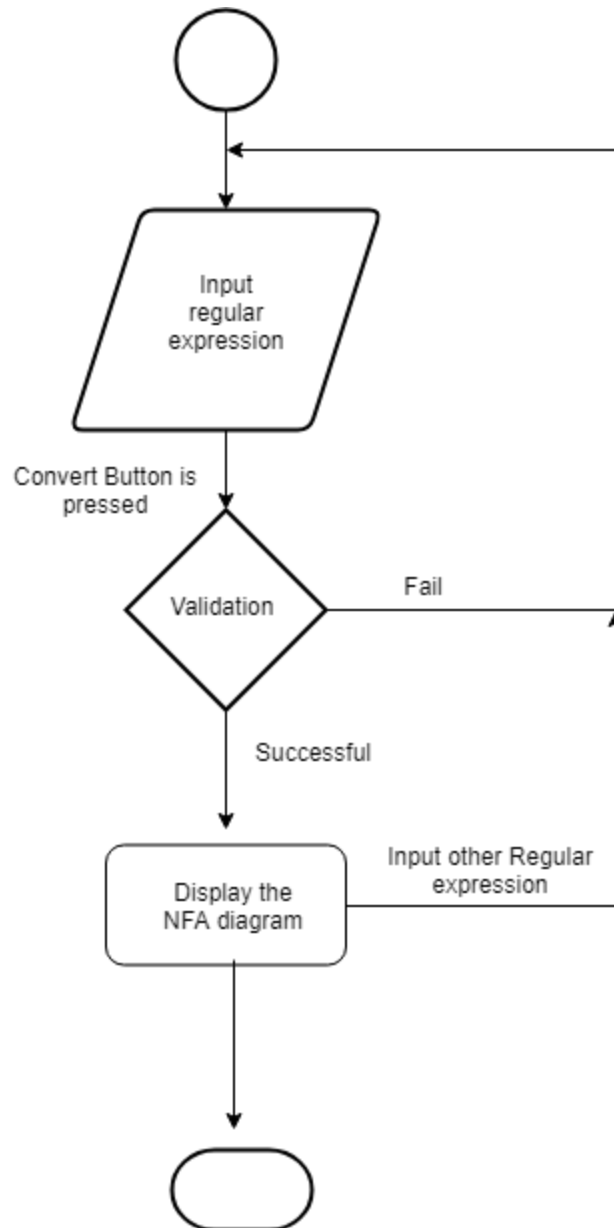


Figure 1: Flow Chart of Part 1 Feature

PART II :

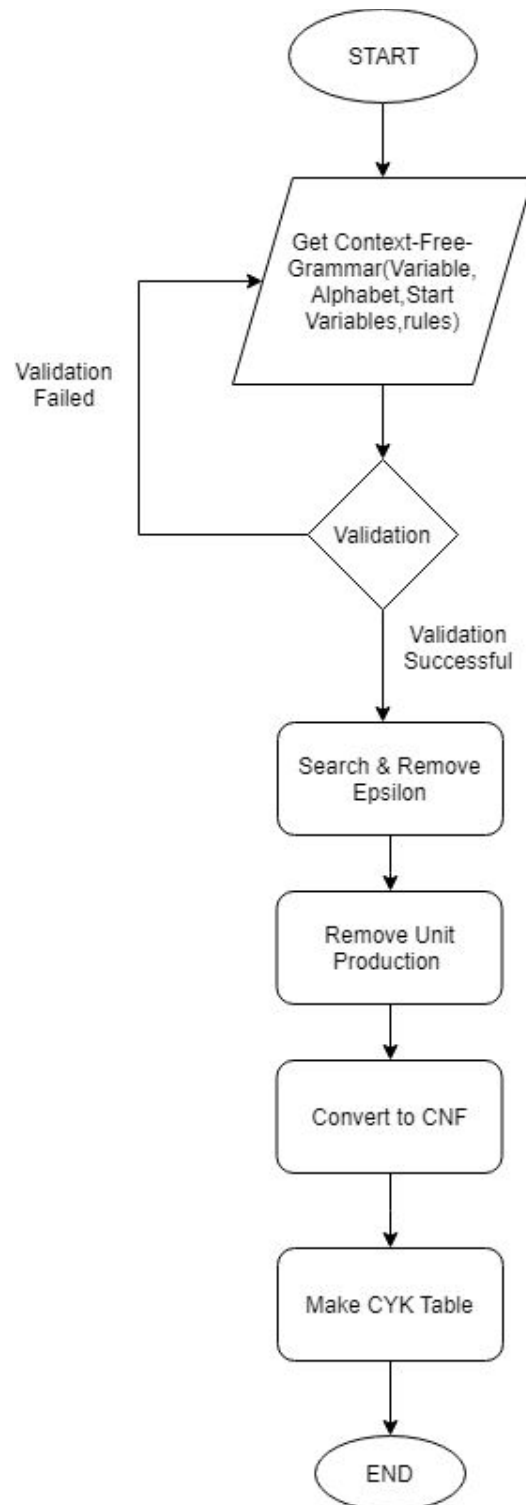
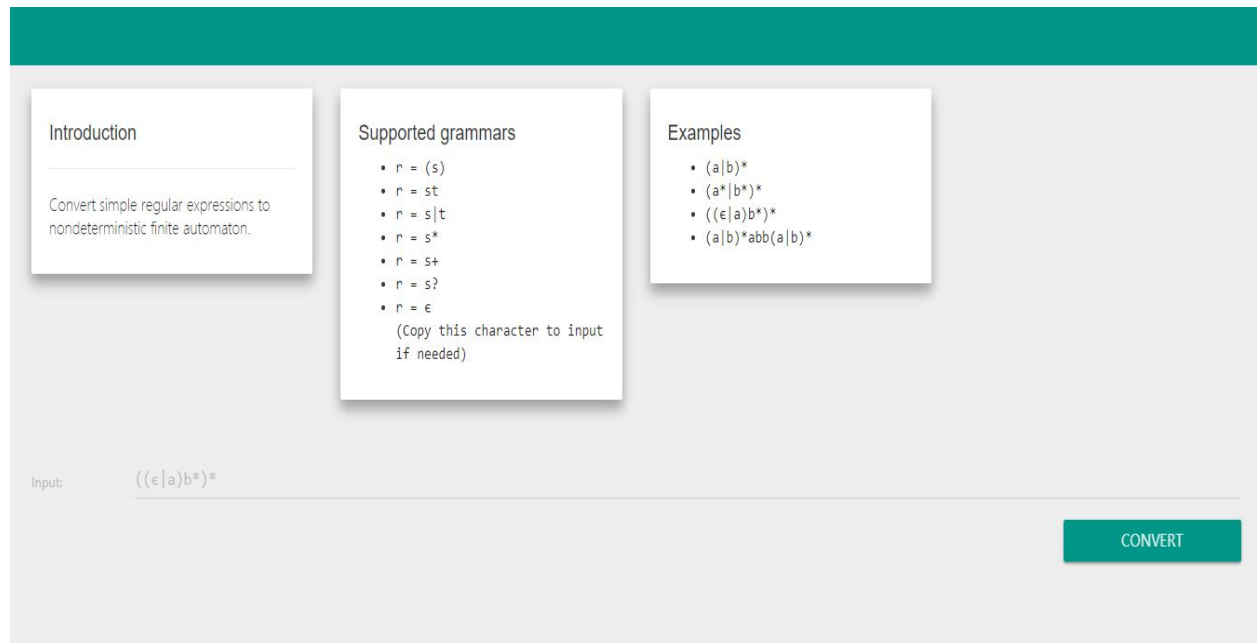


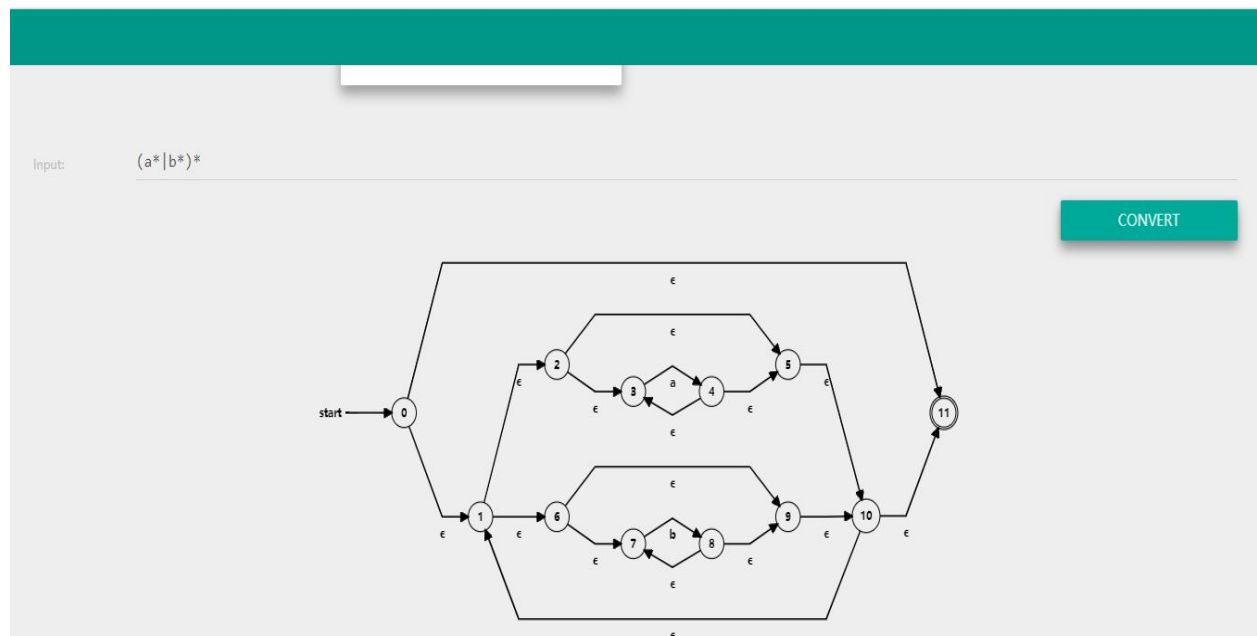
Figure 2: Flow Chart of Part 2 Feature

3.0 Screenshots

PART I:



(Screenshot 1: Main Page of Part 1 Feature)



(Screenshot 2: Result Displayed of Part 1 Feature)

PART II:

PHASE 1 - CFG & CNF CONVERSION

Variable, V = Alphabet, E = Start Variable, S = Number of Rules, R =

Submit Details

Final Chomsky Normal Form
*please submit above details at first

Documentation

(Screenshot 1: The CFG Input Layout)

PHASE 1 - RULE INSERTION

RULE INSERTION

Variable, V = S,A,B
Alphabet, A = a,b
Start Variable, S = S

Enter the Rule 1 : Enter

Rules:

(Screenshot 2: Rule Insertion Layout)

RSION

PHASE 1- RULE INSERTION

RULE INSERTION

Variable, V = S,A,B

Alphabet, A = a,b

Start Variable, S = S

Enter Rule No. 3

Rules:

S->ASA|aB

A->B|S

B->b|e

Step4: Final Chomsky Form

System.Windows.Forms.Label, Text: S->AF|EB|a|SA|AS

A->b|AF|EB|a|SA|AS

B->b

E->a

F->SA

(Screenshot 3: CFG - CNF Conversion)

PHASE 1 - CFG & CNF CONVERSION

S,A,B Alphabet, E = a,b Start Variable, S = S Number of Rules, R = 3

Submit Details

PHASE 2 - ENTER THE STRING

Enter String => aabbbaa | Submit

B->B
E->a
F->SA

tion Phase 2

(Screenshot 4: Phase2 Enter the string)

4.0 Manual with Examples



Figure 3: Main Menu

- Click to open the program 'TIC2151-Assignment.exe', and you will see the interface like the image above.
- In the main menu, the user is required to choose and click respective buttons for entering the feature of part 1, part 2 & also the documentation.

PART I:

At first, the user will direct into the page that illustrates as below after button 'Part 1' is clicked. The user needs to input the regular expression into an input text box and then press the button to conduct conversion into NFA diagram.

Introduction

Convert simple regular expressions to nondeterministic finite automaton.

Supported grammars

- $r = (s)$
- $r = st$
- $r = s|t$
- $r = s^*$
- $r = s+$
- $r = s?$
- $r = \epsilon$

(Copy this character to input if needed)

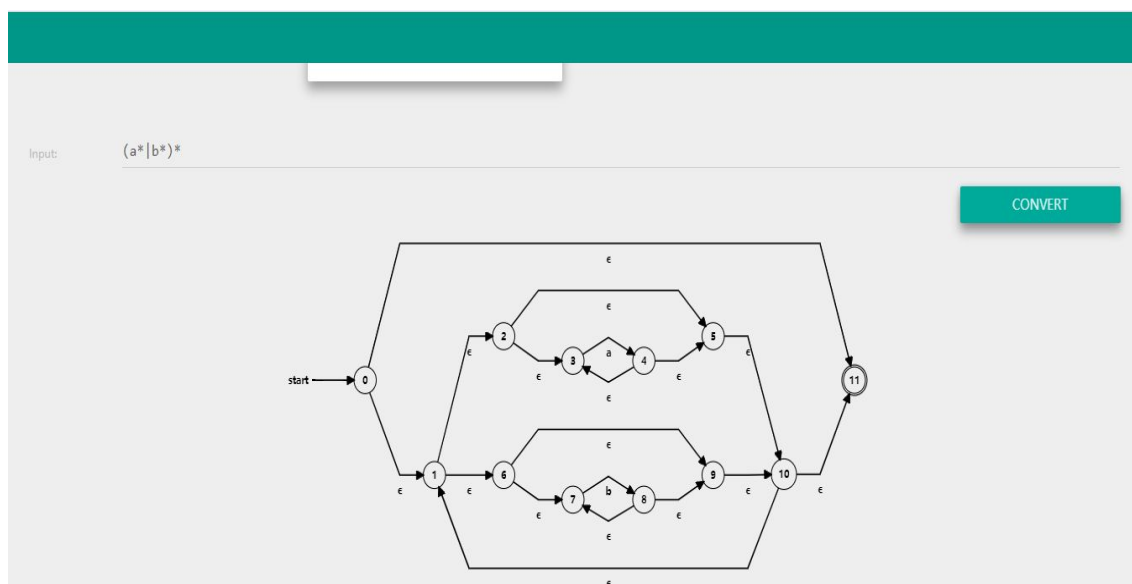
Examples

- $(a|b)^*$
- $(a^*|b^*)^*$
- $((\epsilon|a)b^*)^*$
- $(a|b)^*abb(a|b)^*$

Input:

CONVERT

After the button 'Convert' is pressed, the regular expression entered will undergo validation on checking the regular expression entered whether correct. If the validation is correct, the NFA diagram will be displayed.



If validation is failed, an error message will be displayed to alert the user to enter again.



Convert simple regular expressions to nondeterministic finite automaton.

- $r = s|t$
- $r = s^*$
- $r = s^+$
- $r = s^?$
- $r = \epsilon$

(Copy this character to input if needed)

• $((\epsilon|a)b^*)^*$

• $(a|b)^*abb(a|b)^*$

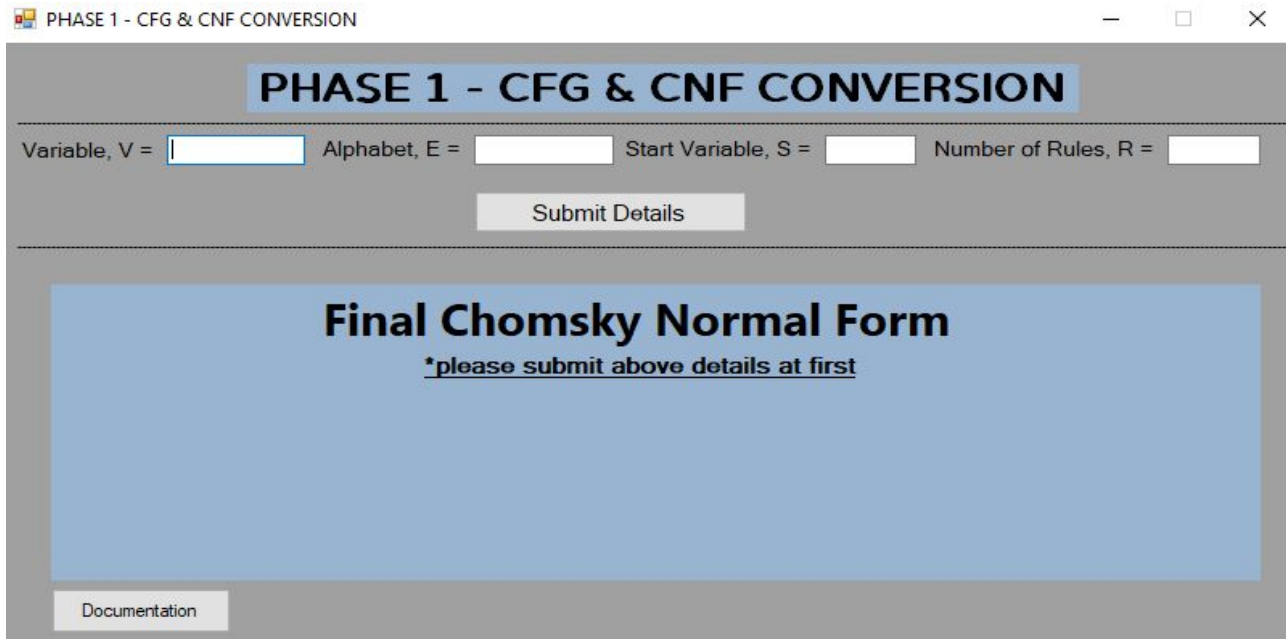
Input:

CONVERT

Invalid Grammar
Error: empty input at 0.

PART II:

Step 1: First of all, the user will see the main interface like the image below



PHASE 1 - CFG & CNF CONVERSION

Variable, V = Alphabet, E = Start Variable, S = Number of Rules, R =

Submit Details

Final Chomsky Normal Form
*please submit above details at first

Documentation

Figure 3: Interface of Phase 1- CFG->CNF conversion

Step 2: In order to input the Context-Free_Grammar correctly, user need to input of variable, alphabet, start variable and number of rules, therefore you need to type all of them into respective textbox, if user need to type more than one letter inside the textbox, make sure separate each of the letters with comma, image below can be as correct guideline, e.g, if you want to input a CF grammar

$S \Rightarrow ASA \mid aB$

$A \Rightarrow B \mid S$

$B \Rightarrow b \mid e$

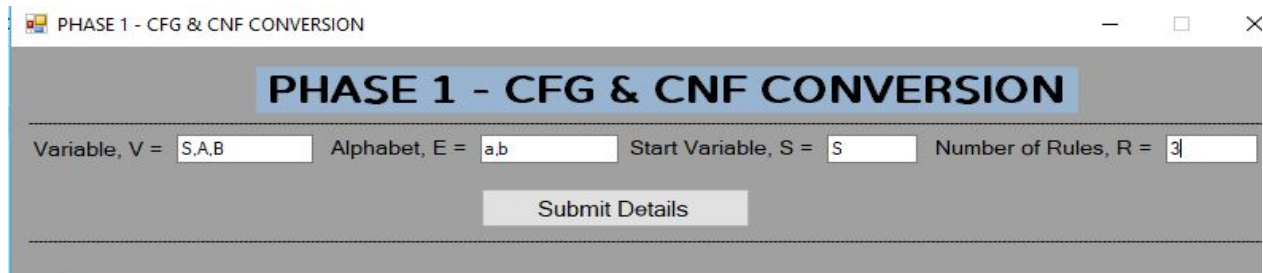


Figure 4: Procedure to Input The Data Accurately

Step 3:- Now you can click the button ‘Submit Details’, you will see a new window form pop out, this form is used type down the rules, the interface will look like this:

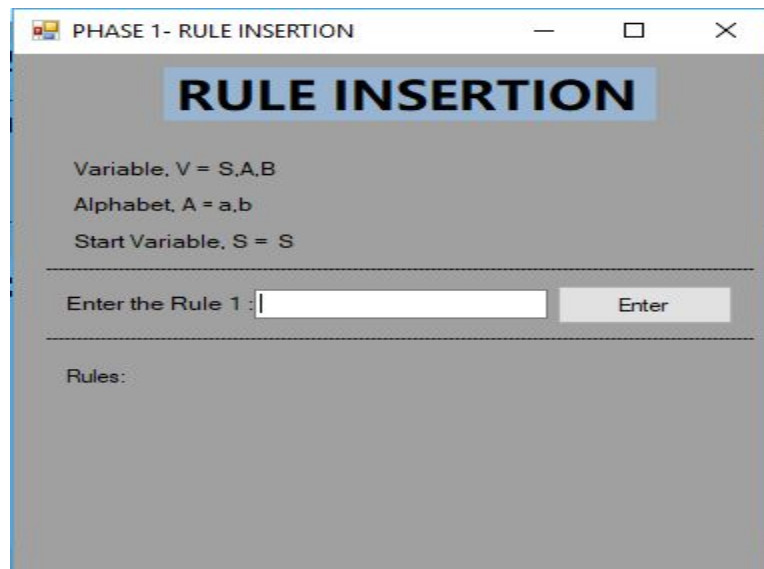


Figure 5: The interface for Rule Insertion

Step 4: It's time to enter the rule, for the sake of illustration, we will still use the example grammar after all the rules have entered, the program will carry out the transform automatically, each step of the conversion will be displayed at a pop-out dialogue as the following images shown.

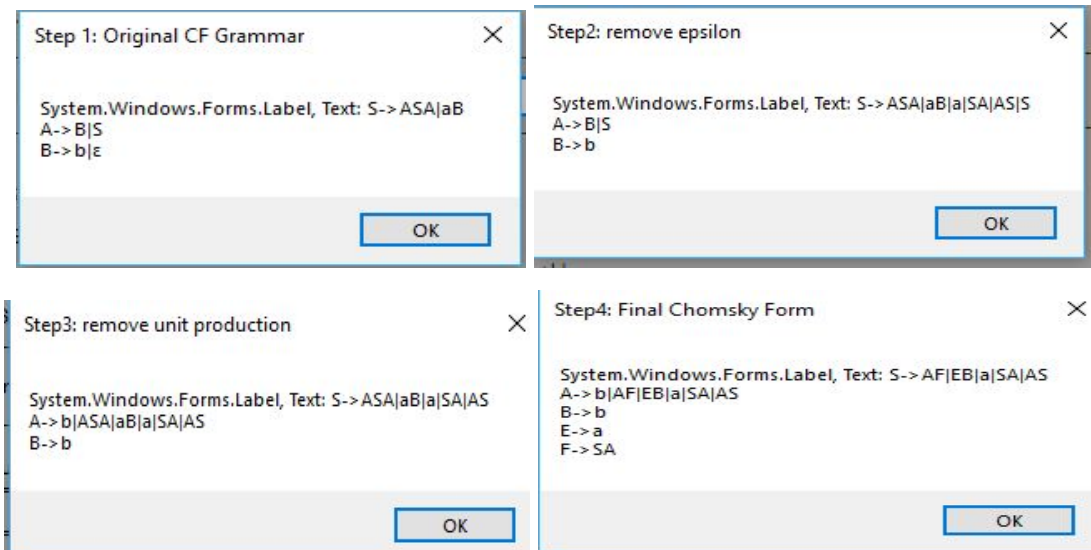


Figure 6: Four steps of conversion

Step 5: Now you got the final CNF of the grammar as the following images shown.

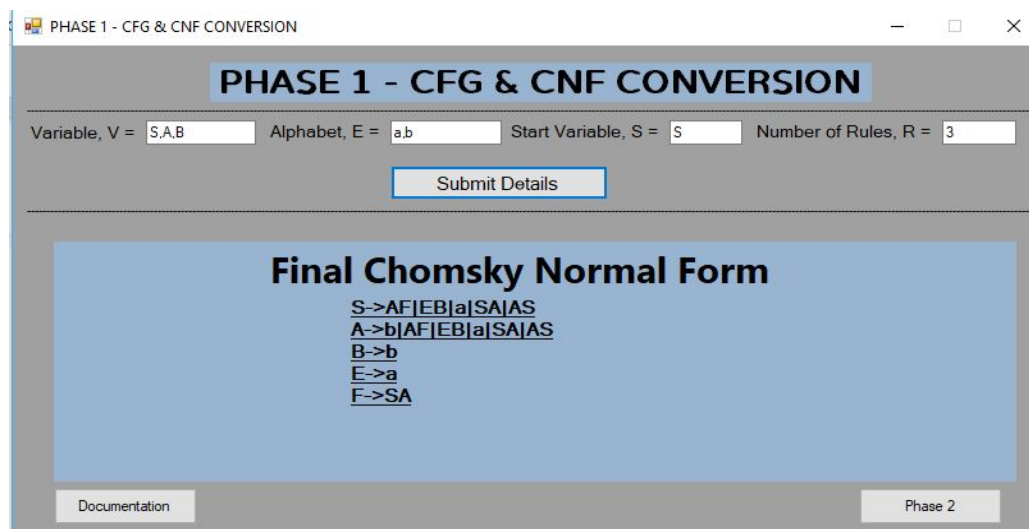


Figure 7: Output Displayed for final CNF

Step 6: Here you can click for Phase 2 which is used to make a CYK table based on the CNF you get when you enter the string, you can only type in the letter which you typed in to the textbox of Alphabet, e.g:

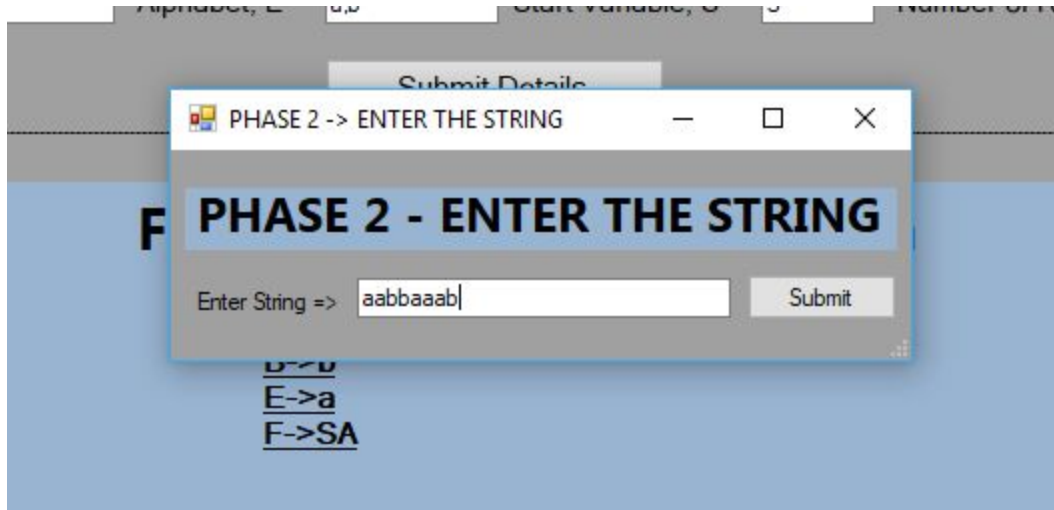


Figure 8: The string can be any combination of the letters, while the letters have to be the one which you entered in the textbox of ‘Alphabet’

Step 7: Now you can get the CYK table by simply submit the string you entered.

a	a	b	b	a	a	a	b
SAFSA	SAFSA		SA	SAFSA	SAFSA	SAFSA	
SAFSAFSA	SAFSAF	SA	SASASASA	SAFSAFSA	SAFSAFSA		
SAFSAFSA	SAFSASAF	SASASASA	SASASASA	SAFSAFSA			
SAFSAFSA	SAFSAFSA	SASASASA	SASASASA				
SAFSAFSA	SAFSAFSA	SASASASA					

Figure 9: The CYK Table based on your CNF & the string.

5.0 Important Codes

search_epsilon function:

```
Sub searchEpsilon()  
    current_node = head_node  
    current_row = head_node  
    Dim a As Integer = 0  
    Dim last_node As Node  
  
    While Not current_row Is Nothing  
        While Not current_node.child Is Nothing  
            last_node = current_node  
            current_node = current_node.child  
            If current_node.info = "ε" Then  
                variable_with_epsilon(a) = current_row.info  
                a += 1  
  
                If Not current_node.child Is Nothing Then  
                    Dim temp As Node  
                    temp = current_node.child  
                    last_node.child = temp  
                    current_node = last_node.child  
                Else  
                    current_node = Nothing  
                    last_node.child = Nothing  
                End If  
                Exit While  
            End If  
        End While  
        current_row = current_row.pnext  
        current_node = current_row  
    End While  
End Sub
```


remove_epsilon function:

```
Sub removeEpsilon()  
    current_node = head_node  
    current_row = head_node  
    Dim lastNode As Node  
    epsilon_exist = False  
    Dim checkNode As Node  
  
    For i = 0 To variable_with_epsilon.Length - 1  
        While Not current_row Is Nothing  
            node_exist = False  
            While Not current_node.child Is Nothing  
                current_node = current_node.child  
                lastNode = current_node  
                Dim eachChar() As Char = current_node.info.ToCharArray()  
  
                For j = 0 To eachChar.Length - 1  
                    For m = 0 To variable_with_epsilon.Length - 1  
                        If eachChar(j) = variable_with_epsilon(m) Then  
                            Dim newV As String = ""  
                            checkNode = current_row  
  
                            If eachChar.Length > 1 Then  
                                For k = 0 To eachChar.Length - 1  
                                    If j <> k Then  
                                        newV = newV + eachChar(k)  
                                    End If  
                                Next  
                            Else  
                                newV = "ε"  
                                'If current_row.info <> start_variable Then  
                                epsilon_exist = True  
                            End If  
  
                            While Not checkNode.child Is Nothing  
                                checkNode = checkNode.child  
                                If checkNode.info = newV Then  
                                    node_exist = True  
                                End If  
                            End While  
  
                            If node_exist = False Then  
                                While Not lastNode.child Is Nothing  
                                    lastNode = lastNode.child  
                                End While  
                            End If  
                        End If  
                    Next m  
                Next j  
            End While  
            current_row = lastNode  
        End While  
    Next i  
End Sub
```

```
                                lastNode.child = New Node
                                lastNode = lastNode.child
                                lastNode.info = newV
                            End If
                        End If
                    Next
                Next
            End While
current_row = current_row.pnext
current_node = current_row
        End While
    Next
End Sub
```

rule_insertion function:

```
'Make LinkedList for each string stored'
PARTII.current_node = PARTII.head_node
    If PARTII.head_node.info = first_letter Then 'to check if the Variable is
the start node
        For j = 0 To temp.Length - 1
            PARTII.current_node.child = New Node
            PARTII.current_node = PARTII.current_node.child
            PARTII.current_node.info = temp(j)
            'MsgBox(PARTII.current_node.info.ToString)
        Next
    Else
        While Not PARTII.current_node.pnext Is Nothing
            PARTII.current_node = PARTII.current_node.pnext
        End While

        PARTII.current_node.pnext = New Node
        PARTII.current_node = PARTII.current_node.pnext
        PARTII.current_node.info = first_letter

        For j = 0 To temp.Length - 1
            PARTII.current_node.child = New Node
            PARTII.current_node = PARTII.current_node.child
            PARTII.current_node.info = temp(j)
        Next
    End If
```

Thank You