

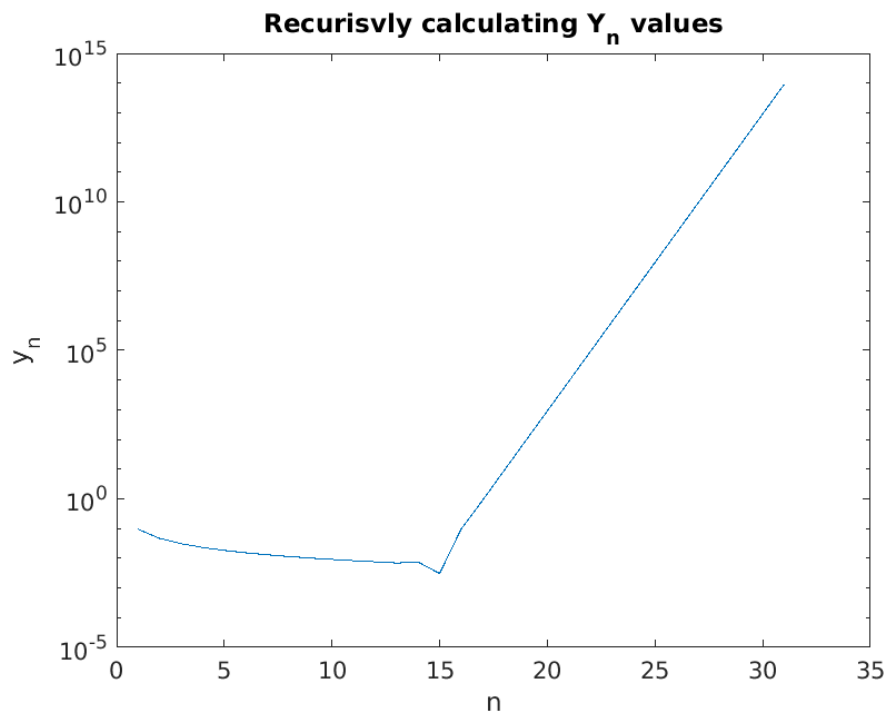
# CX 4640 Assignment 1

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1. Example 1.6 on page 13 of our textbook discusses a recursive formula for computing the integral  $y_n$  for  $n = 1, 2, \dots, 30$ .
  - i. Compute the thirty values of  $y_n$  using the recursive formula and verify the exponential growth of the errors. Plot the computed values in a graph with a logarithmic y-axis scale

n	yn
0	9.53E-02
1	4.69E-02
2	3.10E-02
3	2.32E-02
4	1.85E-02
5	1.54E-02
10	8.33E-03
15	9.74E-02
20	-9.17E+03
25	9.17E+08
30	-9.17E+13



As we can see from the table and from plotting the magnitude of the  $Y_n$  values, we can identify an exponential error revealing itself around  $n = 15$ .

- ii. In our first lecture, we mentioned it was possible to approximate the area under a curve (an integral) by the sum of the areas of a number of rectangles. Implement such a method (using Matlab or a language of your choice) and compute  $y_n$  for several values of  $n$  with your method. You should show the results when a very large number of rectangles is used, as well as when a modest (smaller) number of rectangles is used. Show the key parts of your code in your submitted solutions.

When only using a small number of rectangles ( $n=10$ ), significant portions of the curve are estimated. At large  $n$ , this causes the integral to be severely underestimated. Using larger numbers of  $n$  fixes this discretization error by using smaller rectangles.

n	$y_n$ ( $n = 10$ )	$y_n$ ( $n = 1e6$ )
0	$9.57661e-2$	$9.53106e-2$
5	$1.11786e-2$	$1.53483e-2$
10	$4.52449e-3$	$8.32342e-3$
15	$2.26379e-3$	$5.70796e-3$
20	$1.22994e-3$	$4.34249e-3$

```

num_rect = 1e6;
minX = 0;
maxX = 1;
step = (maxX-minX) / num_rect;
steps = minX:step:maxX
n=0:30;
riemann = zeros(1,length(n));
for i=n
    func = getIntegrand(i);
    for x=steps(1:end-1)
        riemann(i+1) = riemann(i+1) + (func(x));
    end
end
riemann = riemann .* step;

```

An implementation a Riemann summation approximation for estimating  $y_n$ .

- iii. What would you consider to be the most accurate way to compute the integrals  $y_n$ ? Implement your method and compare the results to those in part (b).

To compute accurately the integrals  $y_n$ , we can manipulate the recursive formula  $y_n = \frac{1}{n} - 10y_{n-1}$ , into  $y_{n-1} = \frac{1}{10}(\frac{1}{n} - y_n)$ . After each recursive iteration, any error that exists in our initial estimation of  $y_n$  is reduced by about a factor of 10. Using this method we can start at some estimation of an arbitrary  $y_n$  value, and recursively calculate  $y_{n-1}$ .

*%An alternative recursive algorithm for computing  $y_n$ .  
 %Given an estimation for  $y_{30}$  (in this %case 0), the values  
 %for  $y_0$ – $y_{29}$  very quickly converge to an accurate approximation.  
 %The error in the initial estimation is divided by 10 each iteration.*

```

y_alt = zeros(1,30);
for i=length(y_alt):-1:2
    y_alt(i-1) = 0.1 * ((1/(i)) - y_alt(i));
end

```

```

y0_alt = 0.1 * (1 - y_alt(1))
y_alt = [y0_alt , y_alt];

```

Compared with the Riemann sum method in part b, this method does not require hundreds or thousands of iterations to get a good estimate. Comparing this method and a large Riemann Sum to Matlab's integral function, the alternative recursive method approaches 5 digits of accuracy within 6 iterations ( $y_{24}$ ), even with a relatively bad starting approximation of  $y_{30} = 0$ . To increase the accuracy of values like  $y_{30}$ , a large  $n$  value could be used.

n	Matlab Integral $Y_n$	Alternative Recursive $Y_n$	Reimann Sum (n=1e4)
0	9.53102E-02	9.53102E-02	9.53106E-02
5	1.53529E-02	1.53529E-02	1.53484E-02
10	8.32797E-03	8.32797E-03	8.32342E-03
15	5.71251E-03	5.71251E-03	5.70797E-03
20	4.34704E-03	4.34704E-03	4.34249E-03
24	3.64916E-03	3.64916E-03	3.64462E-03
25	3.50835E-03	3.50838E-03	3.50381E-03
26	3.37800E-03	3.37771E-03	3.37346E-03
27	3.25699E-03	3.25993E-03	3.25245E-03
28	3.14435E-03	3.11494E-03	3.13981E-03
29	3.03924E-03	3.33333E-03	3.03470E-03
30	2.94093E-03	0	2.93639E-03

2. Perhaps a surprising property of finite precision floating-point arithmetic is that it is not associative, due to roundoff errors.
- i. Find three numbers  $a$ ,  $b$ , and  $c$  that can be represented in IEEE double precision such that

$$(a + b) + c \neq a + (b + c)$$

Explain how you found these numbers, and show using Matlab that equality does not hold.

To find three numbers  $a$ ,  $b$ , and  $c$  which are non-associative, all we need to do is find numbers such that  $a + b = b$ , and then set  $c = -b$ . If pick  $b = 2^{1023}$ , then the smallest precision number would be  $2^{1023} * 2^{-52} = 2^{971}$ . Any number less than  $2^{971}$  can not be recorded, as the mantissa does not have enough bits to record to a smaller precision. Setting  $a = 2^{970Flo}$ , and  $c = -2^{1023}$ , we find that  $(a + b) + c = 0$ , but  $a + (b + c) = 2^{970}$ . Using Matlab, we can confirm this.

```
a = 2^970;
b = 2^1023;
c = -2^1023;

d = (a+b) + c
e = a + (b+c)
```

```
d = 0
e = 9.9792e+291
```

- ii. Associativity does not hold either for finite precision multiplication. Again using IEEE double precision, explain *how common* you think it is to find that

$$(a * b) * c \neq a * (b * c)$$

for arbitrary values of  $a$ ,  $b$ , and  $c$ .

In IEEE double precision, non-associativity can happen relatively often. There are the situations where we might get an overflow or an underflow, for example when  $a = 2^{-1023}$ ,  $b = 2^{1023}$ ,  $c = 2^2$ . We know  $a * b * c = 2^2$ , however, when  $b * c$  is done first, we get an overflow error, resulting in  $(a * b) * c = 4$ ,  $a * (b * c) = Inf$ . Another occurrence of these errors would come from when a rounding error occurs in  $(a * b)$  and not  $(b * c)$ . A good example of this is  $a = 0.1$ ,  $b = 0.2$ ,  $c = 0.3$ . Using Matlab, we can see that  $(0.1 * 0.2) * 0.3 \neq 0.1 * (0.2 * 0.3)$ . Generalizing, we can create a Matlab script which randomly generates a mantissa as well as an exponent.

```
a=2^-1023;
b=2^1023;
c=2^2;
(a*b)*c
a*(b*c)

a=0.1;
b=0.2;
c=0.3;
(a*b)*c == a*(b*c)

tmp = 0
range = -128:128;
for e1=range
    for e2=range
        for e3=range
            m1 = rand(1,1);
            m2 = rand(1,1);
            m3 = rand(1,1);
            a=m1*2^e1;
            b=m2*2^e2;
            c=m3*2^e3;
            tmp = tmp + ((a*b)*c ~= a*(b*c));
        end
    end
end
tmp / (length(range))^3
```

```
ans = 4
ans = Inf
```

```
ans = logical
0
```

```
tmp = 0
```

```
ans = 0.3484
```

Sampling random values from exponents ranging from -128 to 128, we approximate that with 3 random IEEE double precision values  $a, b, c$ , about 35% of the time  $(a * b) * c \neq a * (b * c)$ .