

## 9.1

### DEFINITION 6

Let  $R$  be a relation from a set  $A$  to a set  $B$  and  $S$  a relation from  $B$  to a set  $C$ . The *composite* of  $R$  and  $S$  is the relation consisting of ordered pairs  $(a, c)$ , where  $a \in A$ ,  $c \in C$ , and for which there exists an element  $b \in B$  such that  $(a, b) \in R$  and  $(b, c) \in S$ . We denote the composite of  $R$  and  $S$  by  $S \circ R$ .

$R \subseteq A \times B$   $S \subseteq B \times C$   $\exists b$  that  $(a, b) \in R \wedge (b, c) \in S$   
 $\Rightarrow (a, c) \in S \circ R$

### DEFINITION 7

Let  $R$  be a relation on the set  $A$ . The powers  $R^n$ ,  $n = 1, 2, 3, \dots$ , are defined recursively by

$$R^1 = R \quad \text{and} \quad R^{n+1} = R^n \circ R. \quad (a, b) \in R \wedge (b, c) \in R^n$$

### THEOREM 1

The relation  $R$  on a set  $A$  is transitive if and only if  $R^n \subseteq R$  for  $n = 1, 2, 3, \dots$ .

传递  $\Leftrightarrow R$  的幂都是  $R$  的子集

## 9.2

### DEFINITION 1

Let  $A_1, A_2, \dots, A_n$  be sets. An  $n$ -ary relation on these sets is a subset of  $A_1 \times A_2 \times \dots \times A_n$ . The sets  $A_1, A_2, \dots, A_n$  are called the *domains* of the relation, and  $n$  is called its *degree*.

A database consists of records, which are  $n$ -tuples, made up of fields. The fields are the entries of the  $n$ -tuples. For instance, a database of student records may be made up of fields containing the name, student number, major, and GPA of the student.

TABLE 1 Students.

$S \times I \times M \times G$

Student_name	ID_number	Major	GPA
Ackermann	231455	Computer Science	3.88
Adams	888323	Physics	3.45
Chou	102147	Computer Science	3.49
Goodfriend	453876	Mathematics	3.45
Rao	678543	Mathematics	3.90
Stevens	786576	Psychology	2.99

The relational data model represents a database of records as an  $n$ -ary relation. Thus, students are represented as 4-tuples of the form (Student\_name, ID\_number, Major, GPA).

表的分解 (仅作了解)

TABLE 2 GPAs.	
Student_name	GPA
Ackermann	3.88
Adams	3.45
Chou	3.49
Goodfriend	3.45
Rao	3.90
Stevens	2.99

$S \times G$

TABLE 3 Enrollments.		
Student	Major	Course
Glauser	Biology	BI 290
Glauser	Biology	MS 475
Glauser	Biology	PY 410
Marcus	Mathematics	MS 511
Marcus	Mathematics	MS 603
Marcus	Mathematics	CS 322
Miller	Computer Science	MS 575
Miller	Computer Science	CS 455

$S \times M \times C$

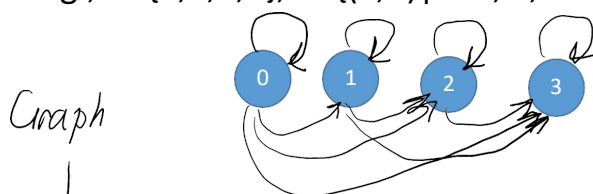
TABLE 4 Majors.	
Student	Major
Glauser	Biology
Marcus	Mathematics
Miller	Computer Science

$S \times M$

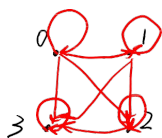
Supplementary (9.3) 仅作了解

Representation of relations on a set:

E.g.,  $A = \{0, 1, 2, 3\}$ ,  $R = \{(a, b) \mid a \leq b; a, b \in A\}$



Directed Graph



Table

a\b	0	1	2	3
0	x	x	x	x
1		x	x	x
2			x	x
3				x

(Binary) Matrix  
(For programming)

$$\begin{bmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

## 9.4 Closure of Relations

### DEFINITION

In mathematics, a subset of a given set is closed under an operation of the larger set if performing that operation on members of the subset always produces a member of that subset.

$$1+1=2 \in \mathbb{N}$$

$$2-1=-1 \notin \mathbb{N}$$

For example, the natural numbers of  $\mathbb{R}$  are closed under addition, but not under subtraction:  $1-2 \notin \mathbb{N}$ . The closure of  $\mathbb{N}$  under subtraction is  $\mathbb{Z}$ .

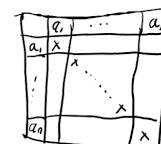
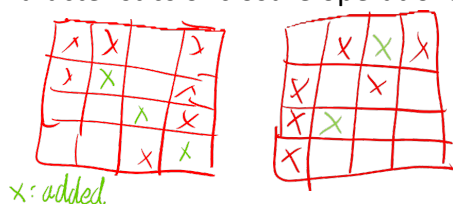
### DEFINITION

If there is a relation  $S$  with property  $P$  containing  $R$  such that  $S$  is a subset of every relation with property  $P$  containing  $R$ , then  $S$  is called the **closure** of  $R$  with respect to  $P$ .

$R \cup I$  is the reflexive closure of  $R$   
 $R \cup R^{-1}$  is symmetric closure of  $R$

Useful notations: Let  $\Delta = \{(a, a) \mid a \in A\}$  be the diagonal relation on  $A$ .

Characteristics of closure operations:



Let  $R$  be a relation that is reflexive and transitive. Prove that  $R^n = R$  for all positive integers  $n$ .

带  $n$ , 多个 item, 没有对应的定理/定义可用 考虑 数学归纳法

Proof: We use

**Mathematical Induction** to prove that  $R^n = R$  holds for  $n \in \mathbb{N}^+$

(i)  $n=1$ ,  $R=R$

(平凡情况, 视情况并入起点)

(ii)  $n=2$ , To prove that  $R^2 = R$  (归纳起点)

$\forall (a, a) \in R \Rightarrow (a, a) \in R \circ R = R^2$   
 $\forall (a, b), (b, b) \in R$

By composition  $(a, b) \in R \circ R = R^2$   
 $\therefore R \subseteq R^2$

$R$  is transitive; by theorem,  $R^2 \subseteq R$

$\therefore R^2 = R$

(iii) Assume that for  $n=i$ ,  $R^i = R$  holds (归纳递推)  
 To prove that for  $n=i+1$ ,  $R^{i+1} = R$  ( $i \geq 2$ )

$\forall (a, b) \in R$

$(b, b) \in R = R^i$  by reflexive

$\therefore$  For  $R^{i+1} = R^i \circ R$ ,  $(a, b) \in R^{i+1}$

$\therefore R \subseteq R^{i+1}$

$R$  is transitive, by theorem,

$R^{i+1} \subseteq R$

$\therefore R^{i+1} = R$  is proved.

How many triples are in the relation  $\{(a, b, c) \mid a, b, \text{ and } c \text{ are integers with } 0 < a < b < c < 5\}$ ?

0 | 1 2 3 4 | 5  
 a b c  $\Rightarrow$  1 2 3  $\nexists$   
 a b c

3 people

4 seats

unordered

$$C_4^3 = 4$$

(1, 2, 3)

(1, 2, 4)

(1, 3, 4)

(2, 3, 4)

Which 4-tuples are in the relation  $\{(a, b, c, d) \mid a, b, c, \text{ and } d \text{ are positive integers with } abcd = 6\}$ ?

factorization  $6 = 6 \times 1 \times 1 \times 1$  (i)  
 $= 2 \times 3 \times 1 \times 1$  (ii)

(i)  $C_4^1 C_3^3 = 4$   
 (6, 1, 1, 1)  
 (1, 6, 1, 1)  
 (1, 1, 6, 1)  
 (1, 1, 1, 6)

(ii)  $C_4^1 C_3^1 C_2^2 = 12$

(2, 3, 1, 1)  
 (2, 1, 3, 1)  
 (2, 1, 1, 3)

(1, 2, 3, 1)  
 (3, 2, 1, 1)  
 (1, 2, 1, 3)

(1, 1, 2, 3)  
 (3, 1, 2, 1)  
 (1, 3, 2, 1)  
 (1, 1, 3, 2)  
 (3, 1, 1, 2)  
 (1, 3, 1, 2)

Find the reflexive closure and the symmetric closure of the relation  $\{(1, 2), (2, 3), (2, 4), (3, 1)\}$  on the set  $\{1, 2, 3, 4\}$ .

	1	2	3	4
1	\	X	X	
2	X	\	X	X
3	X	X	\	
4		X		\

$$r(R) = \{(1, 2), (2, 3), (2, 4), (3, 1), (1, 1), (2, 2), (3, 3), (4, 4)\}$$

$$s(R) = \{ \dots \} = (R \cup R^{-1})$$

Let  $R$  be the symmetric relation on the Set  $A$ .  
Prove that the reflexive closure of  $R$  is symmetric.

Proof: Let  $\Delta = \{(a, a) \mid a \in A\}$  be the diagonal relation on  $A$

The reflexive closure of  $R$  is  $R \cup \Delta$ , we denote it by  $S$ .

For  $\forall (a, b) \in S$ ,

① if  $a = b$ ,  $(a, a) \in \Delta \subseteq S$   $(b, a) \in S$

② if  $a \neq b$ ,  $(a, b) \in R$

Since  $R$  is symmetric

$(b, a) \in R \subseteq S$   $(b, a) \in S$

$\therefore$  By symmetric relation definition,  $S$  is symmetric.

Let  $R$  be the relation  $\{(a, b) \mid a \neq b\}$  on the set of integers.  
What is the reflexive closure of  $R$ ?

$$\Rightarrow \{(a, b) \mid a, b \in \mathbb{Z}\} = \mathbb{Z} \times \mathbb{Z} \quad \text{+ } I_{\mathbb{Z}} / \Delta$$

Prove. Note  $S = \{(a, b) \mid a, b \in \mathbb{Z}\}$

(1)  $\forall (a, a) \in \mathbb{Z} \times \mathbb{Z}, (a, a) \in S$  ①  
 $\therefore S$  is reflexive.

(2)  $R \subseteq S = R \cup \Delta$  ( $\Delta$  is the diagonal rel.)  
an arbitrary

(3) Assume that  $T$  is ref. relation on  $\mathbb{Z}$  and  $R \subseteq T$ .

(i)  $\forall (a, a) \in S$  By reflexive,  $(a, a) \in T$

$$\begin{aligned} \forall (a, b) \in S \quad (a \neq b) \\ &= R \subseteq T \\ &\therefore (a, b) \in R \implies R \subseteq S \\ &\therefore (a, b) \in T \\ \Rightarrow S \subseteq T \end{aligned}$$