

## § 9.4 Closures of Relations (关系的闭包)

### 1. Introduction

#### (1) Example 0 (page 597)

A computer network has data centers in Boston, Chicago, Denver, Detroit, New York and San Diego. There are direct, one-way telephone lines from Boston to Chicago, from Boston to Detroit, from Chicago to Detroit, from Detroit to Denver, and from New York to San Diego.

Let  $R$  be the relation containing  $(a,b)$  if there is a telephone line from the data center  $a$  to that in  $b$ . How can we determine if there is some (possibly indirect) link composed one or more telephone lines from one center to another?

Solution:

We can find all pairs of data centers that have a link by constructing the smallest transitive relation that contains  $R$ .

This relation is called the transitive closure (传递闭包) of  $R$

#### (2) The Closure of Relation $R$ with Respect to Property $P$

Let  $R$  be a relation on  $A$ .  $R$  may or may not have some property  $P$ , such as reflexive, symmetric, or transitivity.

If there is a relation  $S$  with property  $P$  containing  $R$  such that  $S$  is a subset of every relation with property  $P$  containing  $R$ , then  $S$  is called the closure of  $R$  with respect to  $P$ .

### 2. Closures

#### (1) Reflexive Closure

The relation  $R=\{(1,1), (1,2), (2,1), (3,2)\}$  on the set  $A=\{1,2,3\}$  is not reflexive.

How can we produce a reflexive relation containing  $R$  that is as small as possible?

Answer:

by adding  $(2,2), (3,3)$  to relation  $R$ .

The new relation is called the reflexive closure of  $R$ .

#### (2) Result 1

Let  $\Delta=\{(a,a) \mid a \in A\}$ . It is called the diagonal relation (对角线的关系).

The reflexive closure of relation  $R$  on set  $A$ -  
-----  $R \cup \Delta$

#### Example 1 (page 598)

What is the reflexive closure of the relation  $R=\{(a,b) \mid a < b\}$  on the set of integers?

Solution:

$R \cup \Delta = \{(a,b) \mid a \leq b\}$

#### (3) Symmetric Closure

The relation  $\{(1,1), (1,2), (2,2), (2,3), (3,1), (3,2)\}$  on  $\{1,2,3\}$  is not symmetric.

How can we produce a symmetric relation that is as small as possible and contains  $R$ ?

Answer:

by adding  $(2,1)$  and  $(1,3)$ .

The new relation is called the symmetric closure of  $R$ .

#### (4) Result 2

Let  $R^{-1} = \{ (b,a) \mid (a,b) \in R \}$ .

The symmetric closure of relation  $R$  is  
 $R \cup R^{-1}$

Example 2 (page 598)

What is the symmetric closure of the relation  $R = \{ (a,b) \mid a > b \}$  on the set of positive integers?

Answer:

$R \cup R^{-1} = \{ (a,b) \mid a \neq b \}$

$$S = R \cup R^{-1}$$

- 1) Prove  $S$  is symmetric  
for any  $(a,b) \in S$ ,  $(a,b) \in R \cup R^{-1}$  there are two cases  
i)  $(a,b) \in R$ , then  $(b,a) \in R^{-1}$ ,  $(b,a) \in S$   
ii)
- 2)  $R \subseteq S$   
obvious

$$S = R \cup R^{-1}$$

3) For any  $T$ ,  $T$  is symmetric,  $R \subseteq T$ , we prove  $S \subseteq T$ .

$(a,b) \in S$ ,

$(a,b) \in R \cup R^{-1} \{ S = R \cup R^{-1} \}$

i)  $(a,b) \in R$

$(a,b) \in T \quad \{ R \subseteq T \}$

ii)  $(a,b) \in R^{-1}$

$(b,a) \in R \quad \{ \text{def of } R^{-1} \}$

$(b,a) \in T \quad \{ R \subseteq T \}$

$(a,b) \in T \quad \{ T \text{ is symmetric} \}$

$(a,b) \in T$

$S \subseteq T$

#### (5) Transitive Closure

Consider the relation  $R = \{ (1,3), (1,4), (2,1), (3,2) \}$  on the set  $\{1,2,3,4\}$ .

The relation is not transitive.

Add  $(1,2)$ ,  $(2,3)$ ,  $(2,4)$ , and  $(3,1)$ .

Still not transitive.

Why? ----- $(3,1)$  in

----- $(1,4)$  in

----- $(3,4)$  not in

transitive closure-----complicated

### 3. Paths in directed graphs

#### (1) Definition 1 (path)

A path from  $a$  to  $b$  in the directed graph  $G$  is a sequence of edges  $(x_0, x_1), (x_1, x_2), (x_2, x_3), \dots, (x_{n-1}, x_n)$  in  $G$ , where  $n$  is a nonnegative integer, and  $x_0 = a, x_n = b$ .

This path is denoted by  $x_0, x_1, \dots, x_n$  and has length  $n$ .

A path of length  $n \geq 1$  that begins and ends at the same vertex is called a circuit or cycle.

#### (2) Example 3 (page 599)

Which of the following are paths in the directed graph shown in Figure 1:

$a, b, e, d$ ;

$a, e, c, d, b$ ;

$b, a, c, b, a, a, b$ ;

$d, c$ ;

$c, b, a$ ;

$e, b, a, b, a, b, e$ ?

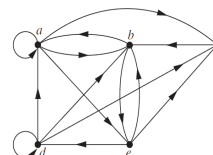


FIGURE 1 A Directed Graph.

What are the lengths of those that are paths?

Which of the paths in this list are circuits?

(3) The term *path* also applies to relations

There is a path from  $a$  to  $b$  in  $R$  if there is a sequence of elements  $a, x_1, x_2, \dots, x_{n-1}, b$  with  $(a, x_1) \in R, (x_1, x_2) \in R, \dots$ , and  $(x_{n-1}, b) \in R$ .

(4) Theorem 1 (page 600)

Let  $R$  be a relation on a set  $A$ . There is a path of length  $n$ , where  $n$  is a positive integer, from  $a$  to  $b$  if and only if  $(a, b) \in R^n$ .

Can prove formally by induction (see textbook).

#### 4. Transitive closure (传递闭包)

(1) Definition 2 (page 600)

Let  $R$  be a relation on a set  $A$ . The connectivity relation  $R^*$  consists of the pairs  $(a, b)$  such that there is a path of length at least one from  $a$  to  $b$  in  $R$ .

In other words, using theorem 1,

$$R^* = \bigcup_{n=1}^{\infty} R^n$$

(2) Example 4 (page 600)

Let  $R$  be the relation on the set of all people in the world that contains  $(a, b)$  if  $a$  has met  $b$ . What is  $R^n$ , where  $n$  is a positive integer greater than one? What is  $R^*$ ?

Answer:

(a) The relation  $R^2$  contains  $(a, b)$  if there is a person  $c$  such that  $(a, c) \in R$  and  $(c, b) \in R$ , that is, if there is a person  $c$  such that  $a$  has met  $c$  and  $c$  has met  $b$ .

(b) Similarly,  $R^n$  consists of those pairs  $(a, b)$  such that there are people  $x_1, x_2, \dots, x_{n-1}$  such that  $a$  has met  $x_1$ ,  $x_1$  has met  $x_2$ ,  $\dots$ , and  $x_{n-1}$  has met  $b$ .

(c) The relation  $R^*$  contains  $(a, b)$  if there is a sequence of people, starting with  $a$  and ending with  $b$ , such that each person in the sequence has met the next person in the sequence.

(3) Theorem 2 (page 601)

The transitive closure of a relation  $R$  equals the connectivity relation  $R^*$ .

Proof:

1)  $R^*$  contains  $R$  by definition.

2) We show that  $R^*$  is transitive. If  $(a, b) \in R^*$  and  $(b, c) \in R^*$ , then there are paths from  $a$  to  $b$  and from  $b$  to  $c$  in  $R$ . We obtain a path from  $a$  to  $c$  by starting with the path from  $a$  to  $b$  and following it with the path from  $b$  to  $c$ . Hence,  $(a, c) \in R^*$ , namely  $R^*$  is transitive.

3) Now suppose that  $S$  is a transitive relation containing  $R$ ,  $R \subseteq S$ . Because  $S$  is transitive,  $S^n \subseteq S$  (by Theorem 1 of Section 9.1). It follows that  $S^* \subseteq S$ . From  $R \subseteq S$ ,  $R^* \subseteq S^*$ , because any path in  $R$  is also a path in  $S$ . Consequently,  $R^* \subseteq S^* \subseteq S$ .

#### (4) Lemma 1 (page 601)

Let  $A$  be a set with  $n$  elements, and  $R$  be a relation on  $A$ . If there is a path of length at least one in  $R$  from  $a$  to  $b$ , then there is a path with length not exceeding  $n$ .

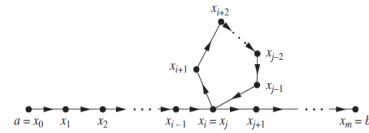


FIGURE 2 Producing a Path with Length Not Exceeding  $n$ .

Moreover, when  $a \neq b$ , if there is a path of at least one in  $R$  from  $a$  to  $b$ , then there is such a path with length not exceeding  $n-1$ .

From Lemma 1, we see that the transitive closure of  $R$  is the union of  $R, R^2, R^3, \dots, R^n$ .

$$R^* = R \cup R^2 \cup R^3 \cup \dots \cup R^n$$