## Lecture 15

## Construction of homotopy theories

From now on, A is a an Eilenberg-Zilber catigory with the property that, for each a  $\in$  ob (A), there are only finitely many  $b \in cb(A)$  with  $Hom_{A_+}(b,a) \neq \beta$  (E=) with  $\partial ha$  finite-for all a).

In particular, there is a weak Lacturization system on A which consists of Munomorphisms and trivial fibration.

Let \* be a terminal object in A.

We lix me and do all an interval T

We fix one and for all an interval I on A: a presheaf I on A equipped with the global actions e is I which are disjoint i.e. such that

 $\emptyset \longrightarrow e$   $\downarrow \qquad \downarrow \qquad \text{is cartesian}$   $e \longrightarrow I$ 

or, equivalently the induced maps  $C \sqcup C \longrightarrow I$  is a managementary.

Notation  $\{0\} = \text{image of d}^\circ: e \rightarrow I$  $\{1\} = \text{image of d}^\prime: e \rightarrow I$ 

de: \* = 109 => I for e = 0,1.

\* 1 \* = 2I = 10} U11 = I

We consider a set S of monomorphisms in A and we make the following assumption:

1) for any a E cb (A) the product Ix ha is finite (i.e. has finitely many non degenerate sections). 2) for any K <>> L in S, L is finite.

Exercise: the assignment X - 1 xX prearres the property of being finite -

Examples:

$$A = \Delta' \quad S = \emptyset$$

2) 
$$A = \Delta$$
,  $J = J$ ,  $S = \{\Lambda_k^n \hookrightarrow \Delta^n \mid n \geqslant 2, o < k < n\}$ 

We define  $\Lambda_{\bar{1}}(5)$  as the following at of mays in  $\hat{A}$ :

$$V_{1}^{J}(2) = V_{1}^{J} \cap V_{2}^{J}(2)$$

 $\Lambda_{I}^{\prime\prime} = \left\{ I \times \partial h_{\alpha} \cup \left\{ \epsilon \right\} \times h_{\alpha} \hookrightarrow I \times h_{\alpha} \mid \alpha \in \partial b (A), \epsilon \in \left\{ 0, 1 \right\} \right\}$   $\Lambda_{I}^{\prime\prime\prime} (S) = \left\{ I \times K \cup \partial J \times L \hookrightarrow I^{\prime\prime} \times L \mid K \hookrightarrow L \in S, n \geqslant 0 \right\}$ 

 $\Lambda_{I}^{"}(\Gamma) = \left\{ I_{\times}^{"} K \cup \partial I_{\times}^{"} L \longrightarrow I_{\times}^{"} L \mid K \hookrightarrow L \in S, n \geqslant 0 \right\}$ with  $I_{-1 \times 1 \times 1}^{"} = U_{-1 \times 1 \times 1 \times 1 \times 1}^{"} = U_{$ 

Definition. An (I,5)-anodyne extension is an element of the smallest saturated closs of maps in A containing  $\Lambda_I(S)$ .

An (I,5)-fibration is a morphism with the right lifting property

An (I,S)-fibration is a morphism with the right lifting property with respect to (I,S)-ansolyne extensions. A presheef X is (I,S)-fibrant (or simply fibrant) if X -> \* is an (I,S)-fibration.

Remark: one can apply the small esject argument. Therefore,

(1,5)-anodyne extensions and (I,5)-fibrations form a weak factorization system in A. In particular any morphism  $f: X \longrightarrow Y$  can be factored into an (I,5)-anodyne extension  $i: X \longrightarrow Z$  followed by an (I,5)-fibration  $p: Z \longrightarrow Y$ .

Proposition.

Let K - L be a monomorphism in Â.

1) for 8=0,1 the induced map

IxKUledxL ~ IxL

is (I,S) - anodyne

2) if ever Kes Lis (I,5)-anodyne, so is

IxKudIxL => IxL (when dI=10fulnfs]

In fact, the class of (I,5) - anodyne extensions is the smallest saturated class of maps containing 5 and with properties 1) and 2) above -

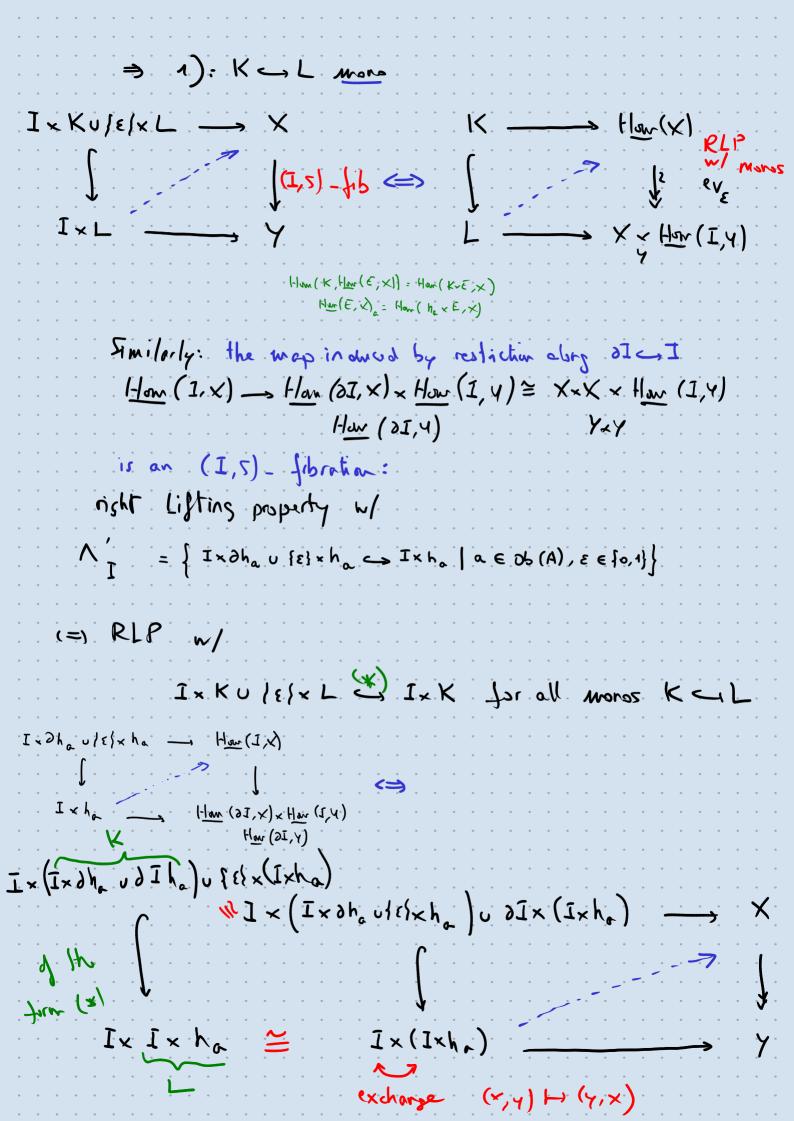
Proof. A map is an (I,5)-anodyne extension iff it has LLP w/ (I,5) - Sibration.

II p:X -> Y is an (I,5) - fibration, then, for e=0,1

ex: Hom (I, X) ~>> How (IE), X) x How (I,Y) = Xx How (I,Y)

How (IES, Y)

is a trival dibration (Los RLP N) dha (acob(A)).



It remains to check REP W/  $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \mid K \hookrightarrow L \in S, n \geqslant 0 \right\}$ For  $K \hookrightarrow L \in \Lambda_{I}^{"}(S)$  we have  $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \in \Lambda_{I}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \in \Lambda_{I}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow \Lambda_{X}^{"}(S) \right\}$   $\Lambda_{I}^{"}(S) = \left\{ I_{X}^{"} K \cup \partial I_{X}^{"} L \hookrightarrow$ 

 $\frac{1}{\sqrt{16}} \int_{-\infty}^{\infty} \frac{1}{\sqrt{16}} \int_{-\infty}^{$ 

The end of the proof is on exercise!

Example:  $A = \emptyset$ ,  $I = \emptyset'$ ,  $S = \emptyset$  (I,S)-anodyne extension = anodyne extensions

Because the class of anodyne extensions is the smallest saturated one such that

is anodyle for E=0,1 and for any mon Keyl

(1,5)-fibrations = Kan fibrations

Fibrant presheaves = Kan complixes = os-graupoids

Example:  $A = \Delta$ , I = J,  $S = \{N_k = \Delta^n \mid n \geq 2, o < k < n\}$ 

The (J, S) - fibrations exactly are those maps p: X -> Y such that

eve: How (J,X)~ Xx How (J,Y) E=0,1 (\*)

Hom (D, X) =>> Hom (V, X) x Hom (D, A) (=) b jour

are timal fibration

It is clear that:

- inner anodyne maps all are (J,J)-anodyne
- Jx KulslxL cs JxL is (J, T) anodyne for any mono K ~ L

Furthermore, for any wono Kes Land any innur and dyna map A = 13

BxKUAxL > BxL is Inner anodyne

=> 1 ( / 1/2 -> D^) = (inner an odyne maps).

Therefore, a simplicial set is (J,5)-fibrant iff it is an ac-category:

 $\times$   $\infty$  - cotigony (=)  $\times \rightarrow \times$  inner this. (=)  $\times \rightarrow \times$  inner this  $+ e_{\Sigma} : fin(J, X) \rightarrow X$  is a fir. distantian

Similarly, if p: X -> Y is a morphism between &- categories, then p is an (I, s)-tibration iff p is an isotibration:

Prod: if pis on isofibration we have a trival fib.

 $\underline{Ham}(J, \times) = h(J, \times) \xrightarrow{\sim} h(\{\xi\}, \times) \times h(J, Y) = \times \times \underline{Ham}(J, Y)$   $h(\{\xi\}, Y)$ 

and p is an inner tib. by definition

Conversely, if p is an (J, s)-fibration, then

it is an inner fibration and for any invehible moghism

b, -> b, in Y and any object as in X noth p(as) = bs

Exercise: Prove the following assertions for A=D, I=J

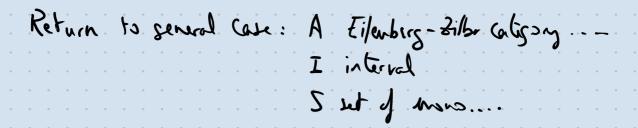
S={N\_k \leftarrow D' | n > 2, o < k < n}

- 1) for any (J, S)-anodyne map  $K \hookrightarrow L$  and any monomorphism  $U \hookrightarrow V$  the induced map  $V \times K \cup U \times L \hookrightarrow V \times L$  is (J, S)-anodyne.
- 2) a morphism p: X -> Yis an (J, r)-fibration

  ill Jor any (J, r)-anodyne map K=1 L

  Han (L, X) => Han (K, X) x Har (L, Y) is a foir. fib.

  Han (K, Y)



Definition.

Let f,g.X-17 be maphime in A.
An I-handopy from f to g is a marphism

h: IXX - Y such that the following diagram communtes

 $1 \times \times \xrightarrow{h} Y$ 111×× =× - g

Notations: h : X = { E} x X - Y

We denote by [X,Y] the quotient of Hang (X, Y) by
the smallest equivalence relation is such that

Ing whenever there exists an I-homotopy from I tog.

Given a May J: X -> Y, we write [f] for its hometopy class (or I-homotopy class) in [X,4].

1~ g (=) J~> o J d, , J, , ..., dn: X -> Y  $\left(h_{1}^{(i)}-\int_{i-1}^{i}\text{ and }h_{0}^{(i)}=\int_{i}^{i}\right)\text{ or }\left(h_{0}^{(i)}=\int_{i-1}^{i}\text{ and }h_{1}^{(i)}=\int_{i}^{i}\right)$ 

such that 1 = 1 and 1 = 9.

Remark: N is compatible with composition:

gofrgod' — frl'

abboirted to I:

objects: presheaves a A

Marghima Jm X to Y: [x, y]

Composition law: [Es] o[d] = [god]

Definition.

An I-handopy equivalure is a morphism  $J:X \to Y$  in A such that [f] is an isomorphism in the homotopy category.

Applying the Youch lume to the homotopy category me get:

Proposition. For a given morphism 1: X -> 4 in A, the following proporties are equivalent:

1) I is an I-homotopy equivalence

2) there exists a morphism  $g: Y \rightarrow X$  in  $\hat{A}$  mith  $g \rightarrow \hat{A} \sim 1_X$  and  $\hat{A} \circ \hat{g} \sim 1_Y$ .

1) for any preshed Ton A the induced may

[T, X] = [T, Y] is bijective

4) for any presheaf W on A the induced may

[Y, W] (X, W) is bijective.

Examples:

A = A, T = A',  $S = \emptyset$ 

D'-honstopies are 1-simplices in Hom.

For Y an  $\infty$ -groupoid a  $\Delta'$ -honotopy John A to g in Han(X,Y) is a normarphism in the  $\infty$ -groupoid Fun(X,Y) = Han(X,Y).

frg (=) fond gove isomorphic in Fun(xy).

An equivalence of &- caligorico X I, Y is a functor (between &- caligories)  $d: X \Rightarrow Y$  such that there exists a functor  $g: Y \to X$  as well a objectivite invertible material transformation fos - 14 and 1x - gof Observe that for an w-caligory Y and a simplicial sut X an objectuite invertible natural transformation 4-15 1- simplex in Fun(x,y)=k(x,y) in Fun(xy) is a , Fun (x,y) = Fun(x,y) anodyne J

it Jollows that:

fir iromorphic to s in Fun (x,y) (=) Ing (J-honotopy)

Therefore a morphism betreen ou-catigories is an equivalence of xo-catigories if it is a J-homotopy equivalence

A		Mo	s J	shi	Yhr	b	ke kn	رم	in w-grapside it an ignimation
y	1	. ~	· -	Ċ	ati	g s	vi o	J.	if it is a J-hanstysy equivalence
	•		۰	•	•	•	• •	•	ED THE STATE OF ALL STATES
				•	0		• •		if it is a D'howstopy equivalence

Back to the servial situation A, I, S Definition

A morphism X t, Y in A is called an (I,S)\_ weak (homotopy) equivalence or, simply a

weak equivalence, if, for any (I,5)-fibrant

preshed W the induced map  $[X, W] \xrightarrow{4} [X, W]$ 

is a bijection.

Proposition.

Any homotopy equivalence is a break equivalence

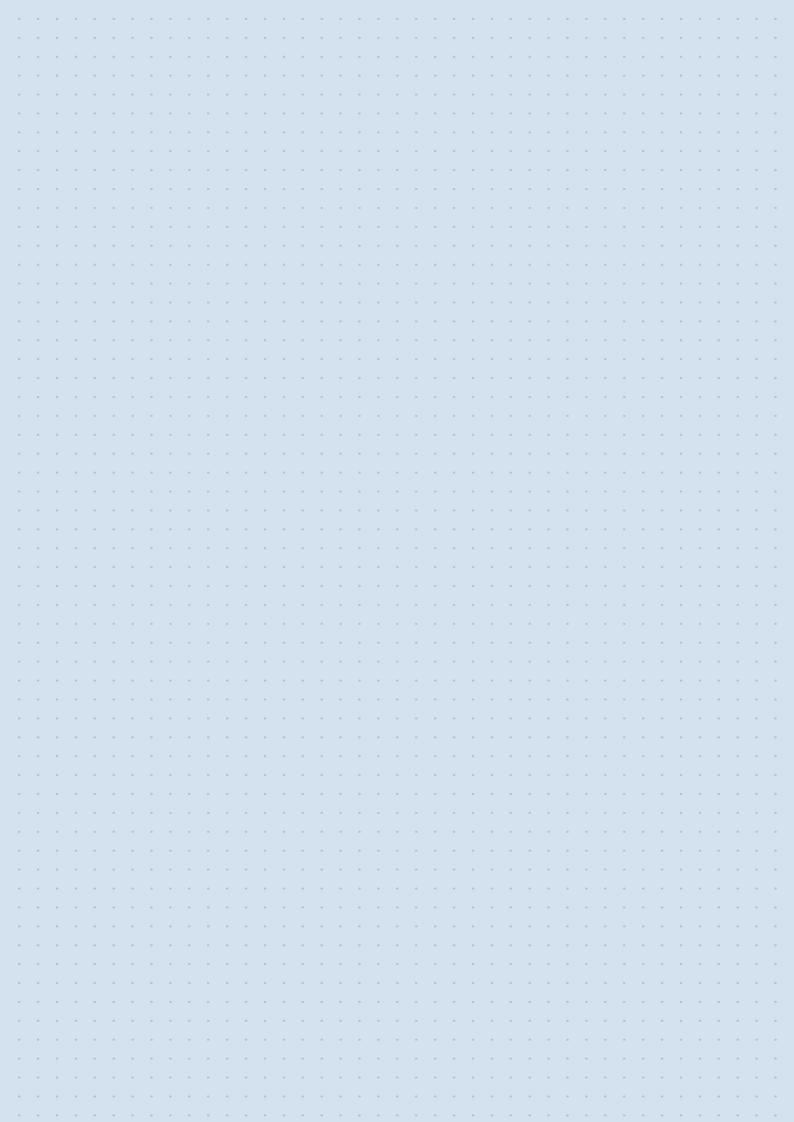
Proposition

Any weak equivalence between fibrant objects is a homotopy equivalence.

(-Sallows Jon Youda Cuma).

We will see later on that any (I,5)-anodyne extension is a weak equivalence -

weak fact syst ((I,s)-anad.), ((1,s)-fib.). car be made — functional (I,r)-and (X,e.) (I,s)-hb. (Anall object argument). Up to week equivalence, every preshed is fibrant  $\times \xrightarrow{f} \lambda$ ix Jn.e. Jiy  $R(x) \rightarrow R(y)$ R (H) weak equivalences mant to be invalid



for rulable der of statement of
propositions P P(x) ~ P(y)

( x, y objects in a 0 - 6tzzy C P: C -, D  $\Rightarrow P(x) \cong P(y)$ 

1-lating of an-lategories. C - D equiv of an lot's.

 $P'_{s} \text{ with } \Rightarrow P(c) \wedge P(b)$ 

Eventially: me mill have on oo-category of ou-categorico.

We will read a dictionary needs a bt of howster?

Heavy.

(1-6t. of a-6t's ( ) w-6t's) w-6t's)

Observation: C 00- Category 1. wtegsy. ho(c)

Conservative C - ho (c)

Constructives in C checking involibitity in ho(C) [ 1. (atessy.) W= weak handay equiv. C= Top or C = Chair complexes W = quai-12msphas W S C1 ~ C[W"] C fen His K(7) iso.  $\subset \xrightarrow{\gamma} \subset [W^{-1}]$ Afen J 3 ho(C(w')) D Let. in C there are (co) limb of interest In practice ho (C[W-1]) (co) limits do est exist us brighang.

Ch (AP) [diz.] × - U.V Z(Unv) - Z(U) Unv - U [ Justury] I pohout I Z(X) = complex of Z(v) -1 Z(x) singular chains in x

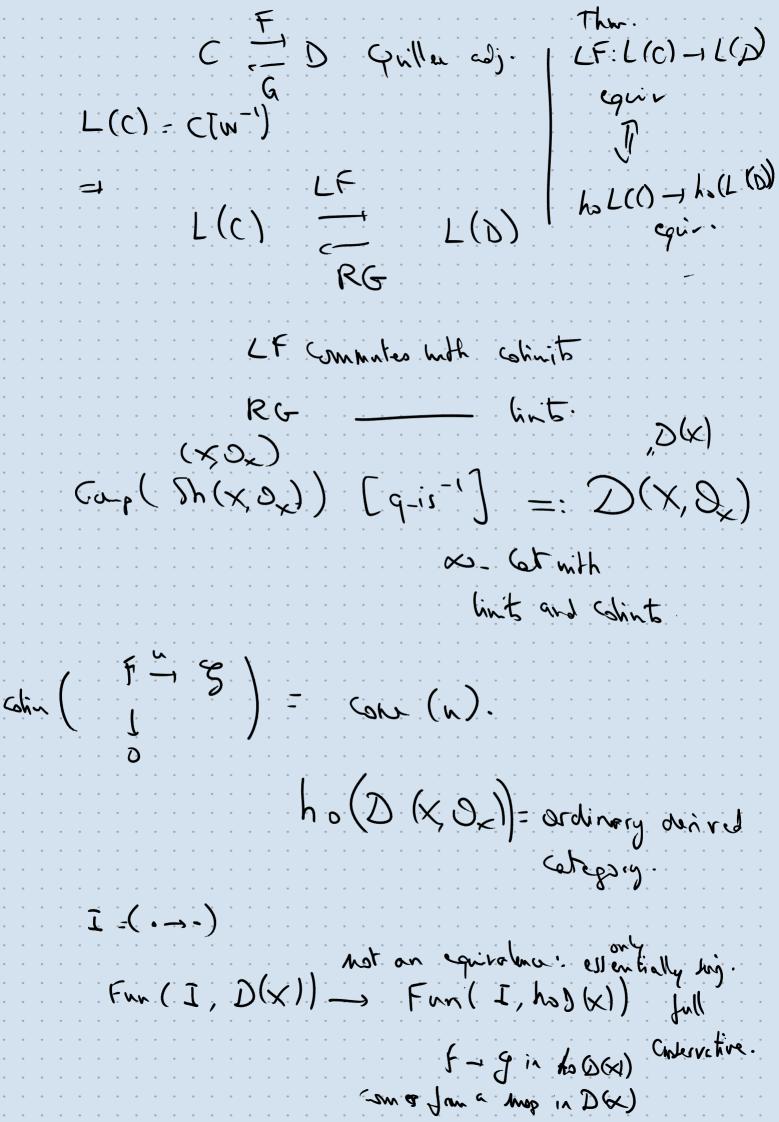


Fig in 
$$D(x)$$
 can  $(n) = \lim_{x \to \infty} (x - x)$ 
 $D(x)$  stable: how finite  $(x) = \lim_{x \to \infty} (x)$ 

and  $Cslin(x) = E(x)$ 
 $E: D(x) = D(x)$  is an equivariance  $(x - x) = X[1] = E(x)$ 

Chance  $(x - x) = X[1] = E(x)$ 

Chose realize: Short exact Lequence:

0-M'M-M (-'M -- C M' JM

C=) 1 Jv both cartesian

O - M" and wasterian

 $M \xrightarrow{\sim} N$ (+u,v) (+) 0-1M-1PBN2 - 1 onl prohab that exact

MUN ~ MXN

A abelian. F: A -> D D stable: Les finte Det: Fexaet of it preserves (co)- limits + 0 objects short exact Legranes

(=) Ignores which are both bull back and purhout E:D-D is on equir.

Douth to structure heart A

Cet Jeur a category. use that to prove theorems on category theory. ω- (at: (ω, 1) - catespies. language  $1 - \omega tegries \sim (\omega, n) - \omega tegries$ . In:  $(\omega, n) - \omega t \in (\omega, n_{11}) - \omega t$ A small tun(tun(AoP, S), C) = Fun (A, C) universal poty of S. S = { Top [W-1] = {CW-complexes { [htpq.cquir-1] C mull chimts D(Ab) = (onp(Ab) [qis") 4=h. pt Z D (Db) Commutes with copy, 5 hay. harslugg.

X has alp. struct. J: X ~ Y inv. => ]! als. Hond. a Y so that d is an isomorphism W. Gt. introduced - Boardman. Vogt. 1970's - operado used to what. P(r) x X" = X  $P(n) \sim pt$  $\times_{o} \rightarrow \times_{i} \rightarrow$ My = How (M, IR) H\*-d(X) Pon coré duality H\* (x) ~ X capact din d 1. Hy (A) - M. (X) dx: 1-1de (x) -> Har (1)

War (x) - Har (4)

1 ty g = submerin (4 (x) = --> ) 11.(3) Correspondens H\*(Xx4) > H\*(X) & H (4) Kunnth James 2x - d Rhan Couplex F(ABB) = F(A) OF(B) ABA -A F(AOA) = F(A) O+ (A) -1 F(A) Floer homology. -> Symplectic resmety.

