Higher Category Theory

Assignment 5

Exercise 1

Proof. (1) Let $\mathcal{C} = [3]$. We see that $N([3]) = \Delta_3$, which has a non-degenerate 3-simplex given by id_{Δ_3} . On the other hand, by definition all of the simplices of $Sk_2(\Delta_3)$ of dimension > 2 are degenerate, hence the canonical inclusion $Sk_2(\Delta_3) \to \Delta_3$ is not an isomorphism.

(2) Since (co)limits of presheaves are defined pointwisely and a morphism of presheaves is a monomorphism if and only if every component of the natural transformation is, we only need to check that for all $a \in \text{Ob}(A)$ the pushout squares

$$\begin{array}{ccc} X_a & \xrightarrow{f_a} & Y_a \\ i_a \downarrow & & \downarrow i'_a \\ X'_a & \xrightarrow{g_a} & Y'_a \end{array}$$

are also pullback squares, so we shall be working solely in \mathbf{Set} , allowing us to drop the a, without creating ambiguity.

Since the class of monomorphisms in **Set** is saturated, we know that i' is a monomorphism too. We will now verify that X has the universal property of the pullback by exhibiting the universal property.

Consider then $h_1\colon Z\to X',\ h_2\colon Z\to Y$ making the diagram commute. We are forced to define a candidate factorization $h\colon Z\to X$ by mapping $z\in Z$ to the unique $x\in X$ such that $h_1(z)=i(x)$, which grants us the uniqueness of an eventual factorization. By construction, h is well-defined and $h_1=i\cdot h$, so we only have to check that $h_2=f\cdot h$. Notice that $i'\cdot h_2=g\cdot h_1=g\cdot i\cdot h=i'\cdot f\cdot h$ and, by injectivity of i', we have the thesis.

Exercise 2

Proof. (1) Once more, we only need to check that for all objects $a \in Ob(A)$ the following is a coequalizer diagram.

$$X_a \times_{Y_a} X_a \xrightarrow{p_a} X_a \xrightarrow{\pi_a} \operatorname{im}(f)_a$$

Here by π we refer to the morphism we get from f by restricting the codomain. $f: X \to Y$. From now on, like in the previous exercise, we shall work in **Set** and therefore drop every a.

We begin by noticing that $\operatorname{im}(f) \cong X_{/\sim}$ under π , where $x \sim x'$ whenever f(x) = f(x'), because π is surjective by construction.

Consider then a function $g\colon X\to Z$ coequalizing p and q. All we have to do is show that, if $x\sim x'$, then g(x)=g(x'), since then g will factor through $\pi\colon X\to X_{/\sim}$ as $\tilde g\colon X_{/\sim}\to Z$, $[x]\mapsto g(x)$. By construction, $\tilde g$ will coequalize p and q, while the uniqueness of the factorization will follow from the surjectivity of π . To do this, we first characterize $X\times_Y X$ explicitly.

We claim that the pullback is given by $S = \{(x, x') \in X \times X \mid f(x) = f(x')\}$ with the obvious projection maps $\pi_1(x, x') = x$, $\pi(x, x') = x'$. Indeed, consider a pair of maps $h_1, h_2 \colon Z \to X$ such that $f \cdot h_1 = f \cdot h_2$. Then, we may construct a factorization $h \colon Z \to S$ by setting $h(z) := (h_1(z), h_2(z))$. This is well-defined since $f(h_1(z)) = (f \cdot h_1)(z) = (f \cdot h_2)(z) = f(h_2(z))$ and therefore $(h_1(z), h_2(z)) \in S$. Also, by construction $\pi_i \cdot h = h_i$ and the uniqueness of the factorization follows from the fact that these last equations (which are satisfied by all factorizations) specify both entries of a candidate h(z).

We now check that the \tilde{g} we defined earlier is actually well-defined by checking that $x \sim x'$ implies g(x) = g(x'). This follows from the fact that $x \sim x'$ means f(x) = f(x'), thus $(x, x') \in X \times_Y X$ and $g(x) = g(p(x, x')) = (g \cdot p)(x, x') = (g \cdot q)(x, x') = g(q(x, x')) = g(x')$.

(2) Suppose T to be a representable presheaf, i.e. isomorphic to \mathfrak{k}_a for some $a \in \mathrm{Ob}(\mathcal{A})$. Since \mathcal{A} is small, $\hat{\mathcal{A}}$ is locally small and therefore the hom-sets are actual sets. Writing equalities in place of natural isomorphisms, we have the following chain of identities: $\hat{\mathcal{A}}(T,Y) = \hat{\mathcal{A}}(\mathfrak{k}_a,Y) = Y_a = \bigcup_{i\in I}Y_{i,a} = \bigcup_{i\in I}\hat{\mathcal{A}}(\mathfrak{k}_a,Y_i) = \bigcup_{i\in I}\hat{\mathcal{A}}(T,Y_i)$. Here a natural transformation $s\colon T\cong \mathfrak{k}_a\to Y_i$ on the right is identified in $\bigcup_{i\in I}\hat{\mathcal{A}}(T,Y_i)$ with all other natural transformations $s'\colon T\cong \mathfrak{k}_a\to Y_j$ such that $s=s'\in Y_a$ and the equality between the two extremes is exhibited by the map sending such a natural transformation $s\colon T\to Y_i$ to the one we get by composing with the inclusion $Y_i\to Y$, which is what we get if we follow the chain of identifications.

Exercise 3

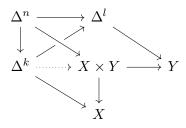
Proof. (1) Recall that the nerve functor N being a right adjoint, preserves products, and thus $\Delta^p \times \Delta^q \cong N([p] \times [q])$. For any n-simplex $s \colon \Delta^n \to \Delta^p \times \Delta^q$, under the adjunction

$$\operatorname{Hom}_{\mathbf{Cat}}([n],[p]\times[q])\cong \operatorname{Hom}_{\mathbf{sSet}}(\Delta^n,\Delta^p\times\Delta^q)$$

it corresponds to a unique $s' : [n] \to [p] \times [q]$. Suppose that s is not a monomorphism. Then s' is not either, which implies that s' factorizes through some [m] (m < n), say, into $[n] \stackrel{f'}{\to} [m] \stackrel{t'}{\to} [p] \times [q]$. Indeed, since s'(i) = s'(j) for some $0 \le i < j \le n$, it suffices to take m = n - 1 and f' to be the unique surjective map $[n] \to [n - 1]$ hitting i twice. Again f', t' correspond to some $f : [n] \to [m]$ and $t : \Delta^m \to \Delta^p \times \Delta^q$ via the adjunction

 $\tau \dashv N$, and one has $s = tf = f^*(t)$. This shows that s is degenerate. Hence the proof.

(2) We claim that if $\Delta \to X$ and $\Delta^n \to Y$ are both degenerate, then so is $\Delta^n \to X \times Y$. To see this, assume they are degenerate and then $\Delta^n \to X$ and $\Delta^n \to Y$ factorize through Δ^k , Δ^l for some $0 \le k, l < n$ respectively. Without loss of generality, one may further assume that $k \le l$, then $\Delta^n \to \Delta^k$ factorizes through Δ^l . We obtain a morphism $\Delta^k \to X \times Y$ by the universal property of products, through which $\Delta^n \to X \times Y$ factorizes, as depicted below:



Hence $\Delta^n \to X \times Y$ is degenerate, and this confirms our claim.

Therefore, if $\Delta^n \to X \times Y$ is non-degenerate, then either $\Delta^n \to X$ or $\Delta^n \to Y$ is degenerate, which implies that either $\Delta^n \to X$ or $\Delta^n \to Y$ is a monomorphism by the regularity of X and Y. We thus may assume that $\Delta^n \to X$ is monic. Then by definition, $\Delta^n([m]) \to X_m$ is an injective map of sets for all $m \ge 0$, and this in turn entails that

$$\Delta^n([m]) \to X_m \times Y_m = (X \times Y)_m$$

is an injective map of sets. Consequently $\Delta^n \to X \times Y$ is a monomorphism.

(3) Consider the diagram $F: I \to \mathbf{sSet}$ where I is finite and $X^i := F(i)$ is regular for each $i \in I$. Recall that finite limits can be exhibited by finite products and equalizers:

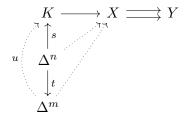
$$\lim_{I} F = \ker \left(\prod_{i \in I} X_i \Longrightarrow \prod_{i \to j} X_i \right)$$

and by (2) plus induction we know that $\prod_{i \in I} X_i$ is regular if each X_i is.

Thus the case is reduced to equalizers: in other words, it suffices to show that for any diagram $X \rightrightarrows Y$ in **sSet**, the equalizer

$$K := \ker(X \rightrightarrows Y)$$

is a regular simplicial set if X is so. To this end, suppose that an n-simplex $s \colon \Delta^n \to K$ is not a monomorphism. Then the composition $\Delta^n \to K \to X$ is not a monomorphism (since it will not be injective over some [l]) as well, and by the fact that X is regular, the composite $\Delta^n \to X$ factors through some Δ^m .



From this we can see that $\Delta^m \to X$ equalizes $X \rightrightarrows Y$, and by the universal property of equalizers, there is a unique morphism $u \colon \Delta^m \to K$. Using the universal property of equalizers again yields that ut = s, which means $s \colon \Delta^n \to K$ being degenerate. \square