

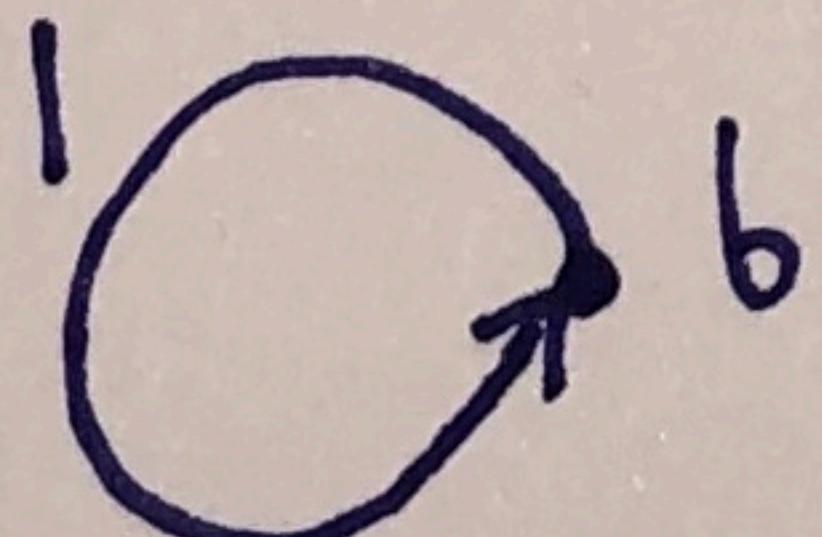
# Circle

To do algebra we need:

- terms  $x, y, \dots$
- operations  $\ast, \dots$
- relations  $xy = yx, \dots$

→ We need points and path constructors, i.e. higher inductive types HIT

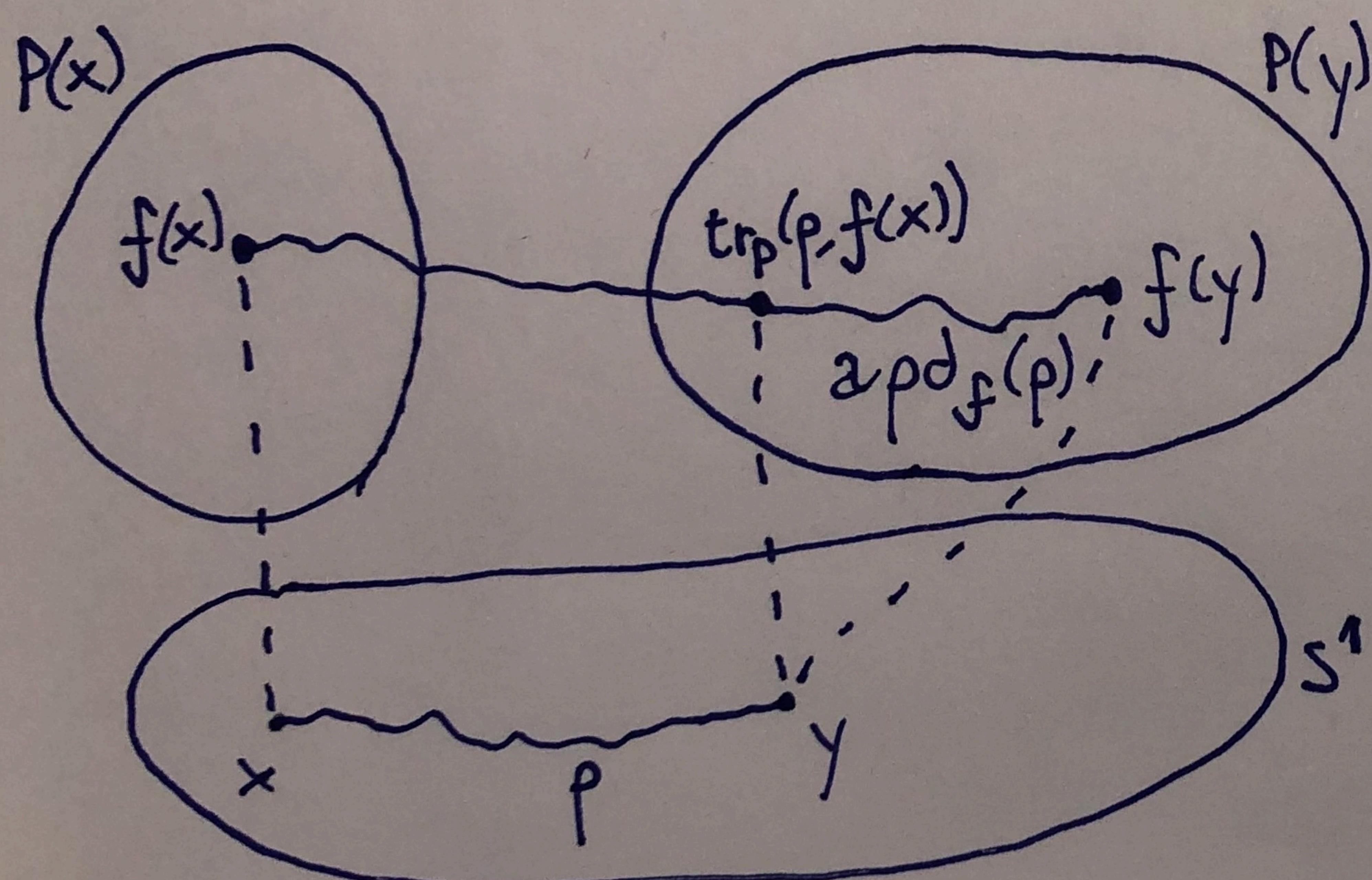
Ex The circle  $b : S^1$  (base)  $l : b = b$  (loop)



Given a dependent function on  $S^1$ ,  $f : \prod_{x:S^1} P(x)$ , we get:

$$1 - f(b) : P(b)$$

$$2 - \text{apd}_f(l) : \text{tr}_P(l, f(b)) \\ \text{apd}_f(l) = f(b)$$



To give an induction principle on  $S^1$  we need to consider the action on the loop,  $\text{apd}_f(l)$ !

Dependent action  
on generators

$$\begin{aligned} \text{dgen} : \left( \prod_{x:S^1} P(x) \right) &\rightarrow \left( \sum_{y:P(b)} \text{tr}_P(l, y) = y \right) \\ f &\mapsto (f(b), \text{apd}_f(l)) \end{aligned}$$

Induction: dgen has a section  $\text{ind}$

we get comp:  $\text{dgen} \cdot \text{ind} \cdot \text{id}$ , our computation rule

Prm  $\exists n \sum_{y:P(b)} \text{tr}_P(l, y) = y$  we have

$((y', p') = (y, p)) \simeq \text{type of pairs } \alpha, \beta \text{ s.t.}$

$$\text{tr}_P(l, y') \xrightarrow{\alpha \text{ pr}_{tr_P(l)}(\alpha)} \text{tr}_P(l, y) \\ p' \parallel \quad \beta \parallel \quad \parallel p \\ y' \xrightarrow[\alpha]{} y$$

$(y, p) \rightsquigarrow f : \prod_{x:S^1} P(x)$  from ind and  $\alpha : f(b) = y, \beta$  from comp

$$\text{tr}_P(l, f(b)) \xrightarrow{\alpha \text{ pr}_{tr_P(l)}(\alpha)} \text{tr}_P(l, y) \\ p' \parallel \quad \beta \parallel \quad \parallel p \\ f(b) \xrightarrow[\alpha]{} y$$

Thm dgen is an equivalence (dependent universal property)

Pf We need  $\text{ind} \cdot \text{dgen} \sim \text{id}$ , i.e.  $\text{ind}(\text{dgen}(f)) = f$ ,

i.e.  $g(x) = f(x)$  for  $g : \equiv \text{ind}(\text{dgen}(f))$

Idea: use  $S^1$ -induction with  $E(x) : \equiv g(x) = f(x)$

$$\text{ind} : \left( \sum_{\alpha : E(b)} \text{tr}_E(l, \alpha) = \alpha \right) \rightarrow \left( \prod_{x:S^1} E(x) \right)$$

Need:  $\alpha : g(b) = f(b)$  and  $\beta : \text{tr}_E(l, \alpha) = \alpha$

[Get a map  $(\text{apd}_g(p) \cdot r = \text{ap}_{tr_E(p)}(q) \cdot \text{apd}_f(p)) \rightarrow (\text{tr}_E(p, q) = r)$   
by path induction on  $p : x = x'$  with  $r : x' = x'$ ,  $q : x = x$

Enough:  $\alpha : g(b) = f(b)$  and  $\beta$

$$\text{tr}_P(l, g(b)) \xrightarrow{\alpha \text{ pr}_{tr_P(l)}(\alpha)} \text{tr}_P(l, f(b)) \\ \parallel \quad \beta \parallel \quad \parallel \text{apd}_f(l) \\ g(b) \xrightarrow[\alpha]{} f(b)$$

They are given by  $\text{comp}(f(b), \text{apd}_f(l))$   
by the remark

Cor  $P/S^1, y: P(b), p: \text{tr}_P(l, y) = y$

The type of triples  $f: \prod_{x:S} P(x), \alpha: f(b) = y, \beta$  is contractible,

$$\begin{array}{ccc} \text{tr}_P(l, f(b)) & \xrightarrow{\alpha \cdot \text{tr}_P(l)(\alpha)} & \text{tr}_P(l, y) \\ \text{apd}_f(l) \parallel & \beta \parallel & \parallel p \\ f(b) & \xlongequal{\alpha} & y \end{array}$$

i.e. there's a unique function with unique proofs of the desired equalities.

Thm  $\text{gen}: (S^1 \rightarrow X) \rightarrow (\sum_{x:X} x=x)$

$$f \mapsto (f(b), \text{ap}_f(l))$$

is an equivalence.

(universal property)

Rf Careful with  $\text{tr}_X(l)$ !

Strategy: consider the triangle and construct the bottom map

$$\begin{array}{ccc} (S^1 \rightarrow X) & & \\ \text{gen} \swarrow & \searrow \text{dgen} & \\ \sum_{x:X} x=x & \xrightarrow{\sim} & \sum_{x:X} \text{tr}_X(l, x) = x \end{array}$$

1- Construct  $\text{tr}_X(p, x): \text{tr}_X(p, x) = x$  by path induction

2- Define the map as  $(x, e) \mapsto (x, \text{tr}_X(l, x) \cdot e)$ .

Check commutativity: consider  $\text{tr}_X(p, f(y)) \xrightarrow{\text{tr}_X(p, f(y))} f(y)$

$$\begin{array}{c} \text{apd}_f(p) \parallel \\ f(y') \parallel \\ \text{ap}_f(p) \end{array}$$

It commutes by path induction. Now substitute  $b/y, b/y', l/p$ .

Cor  $p: x=x \text{ in } X$

$$\sum_{f: S^1 \rightarrow X} \sum_{\alpha: f(b) = x} \alpha \cdot p = \text{ap}_f(l) \cdot \alpha$$

is contractible.

$$f(b) \xlongequal{\alpha} x$$

$$\begin{array}{ccc} \text{ap}_f(l) \parallel & \beta \parallel & \parallel p \\ f(b) & \xlongequal{\alpha} & x \end{array}$$

In order to do algebra we still need an operation.

Def

We construct  $\text{mul} : S^1 \rightarrow (S^1 \rightarrow S^1)$

We want  $\text{mul}(b) = \text{id}$ , so consider the universal property wrt  
 $\text{ind} : \left( \sum_{f:S^1 \rightarrow S^1, f=f} f \right) \rightarrow (S^1 \rightarrow (S^1 \rightarrow S^1))$

$(\text{id}, \text{eq-htpy}(H)) \xrightarrow{\text{ap}} b\text{-mul} : \text{mul}(b) = \text{id} \text{ and } l\text{-mul} \text{ (plus mul)}$

$$\begin{array}{c} \text{mul}(b) \xrightarrow[b\text{-mul}]{\quad} \text{id} \\ \text{ap}_{\text{mul}}(l) // \quad \quad \quad // \text{eq-htpy}(H) \\ \text{mul}(b) \xrightarrow[l\text{-mul}]{\quad} \text{id} \end{array}$$

We have to specify  $H : \text{id} \sim \text{id}$

Use the dependent universal property on  $E(x) : \exists x = x$

$\text{ind} : \left( \sum_{p:E(b)} \text{tr}_E(l, p) = p \right) \rightarrow (\prod_{x:S}, E(x))$  (this is exactly  $\text{id} \sim \text{id}$ )

$$(l, \gamma) \mapsto H : \text{id} \sim \text{id} \text{ with } \alpha : H(b) = l \text{ and } \beta \quad \begin{array}{c} \text{tr}_E(l, H(b)) \xrightarrow[\quad]{\text{ap}_{\text{tr}_E(l)}(\alpha)} \text{tr}_E(l, l) \\ \text{ap}_H(l) // \quad \quad \quad \beta // \quad \quad \quad \gamma \\ H(b) \xrightarrow[\alpha]{\quad} l \end{array}$$

We have to specify  $\gamma$

Get the map  $(p \cdot r = q \cdot p) \rightarrow (\text{tr}_E(p, q) = r)$  by path induction on  $p : b = x$   
 with  $q : b = b$ ,  $r : x = x$

$$\text{refl}_{l, l} \mapsto \gamma$$

Rem  $H \neq \text{refl-htpy}$  for otherwise  $\text{mul}(x, y) = y$

We have  $\text{mul}(b, x) = x$  by the computation rule, while  $\text{mul}(x, b) = x$  is more involved.

$$\text{left-unit} := \text{htpy-eq}(b\text{-mul})$$

Thm  $\text{mul}(x, b) = x$

Pf We want to use  $S^1$ -induction on  $P(x) := \text{mul}(x, b) = x$

$$\left( \sum_{p: \text{mul}(b, b) = b} \text{tr}_p(l, p) = p \right) \rightarrow \left( \prod_{x: S^1} \text{mul}(x, b) = x \right)$$

We have  $\text{left-unit}(b) : \text{mul}(b, b) = b$ , so we need  $\text{tr}_p(l, \text{left-unit}(b)) = \text{left-unit}(b)$

Construct  $(\text{htpy-eq}(\lambda p \text{mul}(p))(b) \cdot r = q \cdot p) \rightarrow (\text{tr}_p(p, q) = r)$  by path induction on  $p: b = x$  with  $q: \text{mul}(b, b) = b$ ,  $r: \text{mul}(x, b) = x$

Enough:

$$\begin{array}{ccc} \text{mul}(b, b) & \xrightarrow{\text{left-unit}(b)} & b \\ \text{htpy-eq}(\lambda p \text{mul}(p))(b) & \parallel & \parallel \\ \text{mul}(b, b) & \xrightarrow{\quad \beta \quad} & b \\ & \parallel & \parallel \\ & \text{left-unit}(b) & \end{array}$$

Since  $H(b) = l$ , we get  $\beta$  from

$$\begin{array}{ccc} \text{mul}(b, b) & \xrightarrow{\text{htpy-eq}(b-\text{mul})(b)} & b \\ \text{htpy-eq}(\lambda p \text{mul}(p))(b) & \parallel & \parallel \\ \text{mul}(b, b) & \xrightarrow{\quad \text{htpy-eq}(b-\text{mul})(b) \quad} & b \\ & \parallel & \parallel \\ & H(b) & \end{array}$$

which commutes by definition  
of  $l\text{-mul} \mapsto \text{htpy-eq}(l\text{-mul})(b)$

What is the fundamental cover of  $S^1$ ? It's a dependent function  $D(\mathbb{Z}, \text{succ}) : S^1 \rightarrow \mathcal{U}$  induced by the universal property from the pair  $(\mathbb{Z} : \mathcal{U}, \text{eq-equiv}(\text{succ}) : \mathbb{Z} = \mathbb{Z})$

$$\begin{array}{ccc} D(\mathbb{Z}, \text{succ})(b) & \xrightarrow{\tilde{e}} & \mathbb{Z} \\ \text{tr}_{D(\mathbb{Z}, \text{succ})}(l) \downarrow & & \downarrow \text{succ} \\ D(\mathbb{Z}, \text{succ})(b) & \xrightarrow{\tilde{e}} & \mathbb{Z} \end{array}$$

The fiber over the base is  $\mathbb{Z}$  and going through the loop once sends a point in the fiber to the next one.

## Homotopy Pushouts

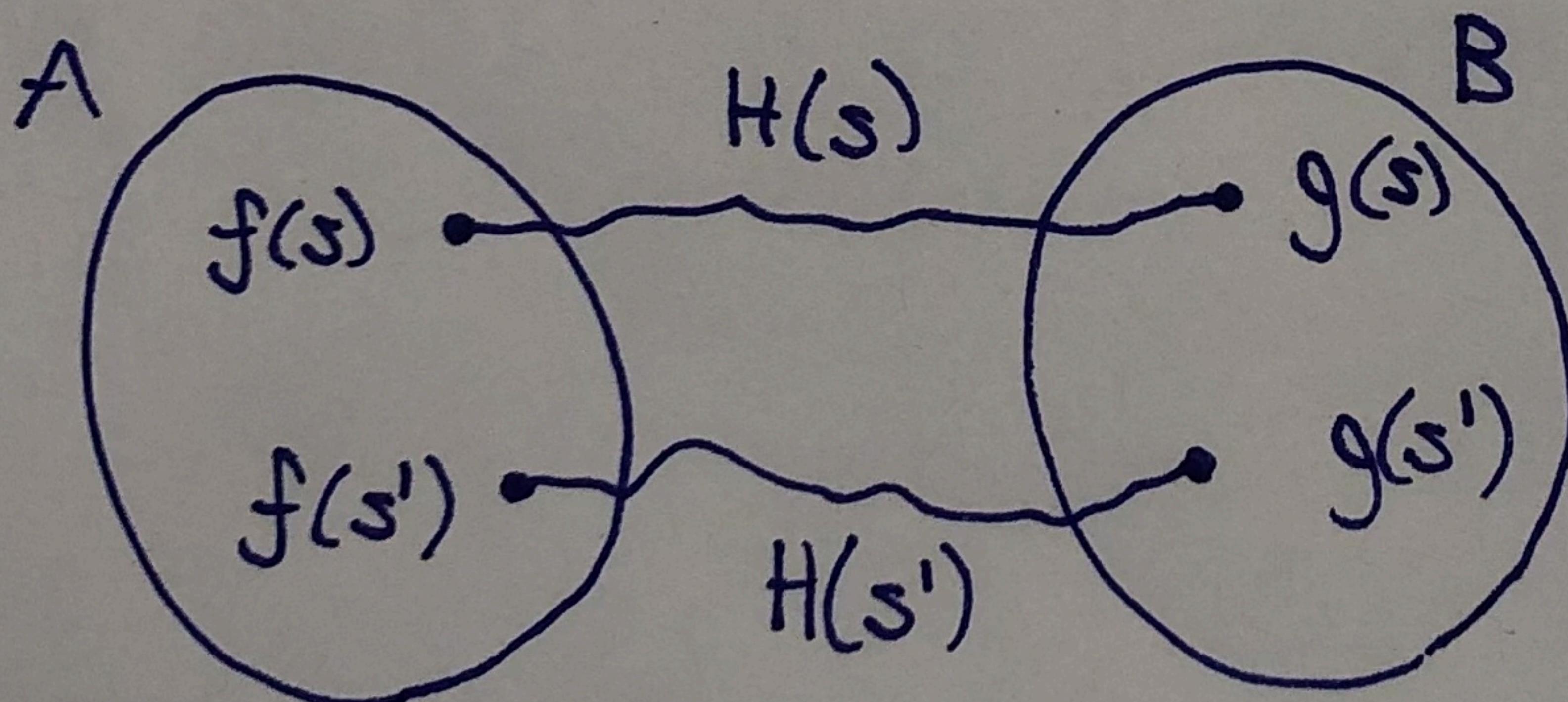
How to construct new types?

In algebraic topology we attach cells to spaces by pushout.

Def

- A span  $S$  is a triple  $(S, f, g)$
- A commutative square is a span  $S$  with a triple  $(i, j, H)$ :  $\sum_{i: A \rightarrow X} \sum_{j: B \rightarrow X} \text{if} \sim jg$
- cocone-map $_{X'}(i, j, H): (X \rightarrow X') \rightarrow \text{cocones}_S(X')$   
 $h \mapsto (hi, hj, hH)$
- A (homotopy) pushout is a commutative square s.t. this map is an equivalence  $\forall X'$

A pushout  $X$  is a disjoint union of  $A$  and  $B$  with added paths from  $f(s)$  to  $g(s)$ .



Pushouts may not exist! They do allow to construct new types.

Lemma

$$(i, j, H), (i', j', H'): \text{cocones}_S(X)$$

$$(i, j, H) = (i', j', H') \Leftrightarrow \sum_{K: i \sim i'} \sum_{L: j \sim j'} H' K f \sim L g H$$

$$\begin{array}{ccc} i f & \xrightarrow{Kf} & i' f \\ \downarrow H & \Downarrow M & \downarrow H' \\ j g & \xrightarrow{Lg} & j' g \end{array}$$

Lemma

$$\begin{array}{ccc} S & \xrightarrow{g} & B \\ f \downarrow & \nearrow H & \downarrow j \\ A & \xrightarrow{i} & X \end{array} \quad \text{pushout}$$

$$\begin{array}{ccc} S & \xrightarrow{g} & B \\ f \downarrow & \nearrow H' & \downarrow j' \\ A & \xrightarrow{i'} & X' \end{array}$$

$$\sum_{h:X \rightarrow X'} \sum_{K:h \sim i} \sum_{L:h \sim j'} H' K f \sim L g H \text{ is contractible}$$

$$\begin{array}{ccc} hif & \xrightarrow{Kf} & i'f \\ hH \downarrow & \nearrow M & \downarrow H' \\ h gj & \xrightarrow{Lg} & gj \end{array} \quad (\text{Uniqueness of factorization})$$

Thm

Two commutative squares.

If two of the following hold, so does the third:

- (i)  $(i, j, H)$  is a pushout on  $S$
- (ii)  $(i', j', H')$  //
- (iii)  $h:X \rightarrow X'$  is an equivalence

pf

$$(X' \rightarrow Y) \xrightarrow{h^*} (X \rightarrow Y)$$

cocone-map( $i', j', H'$ )  $\swarrow$   $\downarrow$  cocone-map( $i, j, H$ )

cocone  $S(Y)$

$h$  is an equivalence  $\Leftrightarrow h^*$  is  $\forall Y$

for

If the two squares are pushouts, then

\*  $\sum_{h:X \simeq X'} \sum_{K:h \sim i'} \sum_{L:h \sim j'} H' K f \sim L g H \text{ is contractible}$

for

$S$  is a span in  $\mathcal{U} \Rightarrow \sum_{X:\mathcal{U}} \sum_{c:\text{cocone}_S(X)} \prod_{Y:\mathcal{U}} \text{is-equiv}(\text{cocone-map}(c))$   
is a proposition

pf

\*  $\simeq ((X, c, u) = (X', c', u'))$

## Duality

Lemma

$$\begin{array}{ccc} \text{cocone}_S(X) & \xrightarrow{\pi_2} & X^B \\ \pi_1 \downarrow & \pi_3 \nearrow & \downarrow g^* \\ X^A & & X^S \end{array}$$

~~TAKEAWAY~~

$\pi_3 : (i, j, H). \text{eq-htpy}(H)$   
is a pullback

Pf

The gap map  $\text{cocone}_S(X) \rightarrow X^A \times_{X^S} X^B$  is induced by  $\text{eq-htpy}$  on the total spaces.

Prop

Given a commutative square, TFAE:

- (i) it's a pushout
- (ii) the commutative square with  $H' := \lambda h. \text{eq-htpy}(hH)$  is a pullback  $\forall T$

$$\begin{array}{ccc} T^X & \xrightarrow{j^*} & T^B \\ i^* \downarrow & H' \nearrow & \downarrow g^* \\ T^A & \xrightarrow{f^*} & T^S \end{array}$$

Pf

$$\begin{array}{ccc} T^X & & \text{gap}(i^*, j^*, \lambda h. \text{eq-htpy}(hH)) \\ \text{cocone-map}(i, j, H) \swarrow & \searrow & \\ \text{cocone}_S(T) & \xrightarrow{\sim} & T^A \times_{T^S} T^B \\ & & \end{array}$$

$\text{gap}(i, j, \text{eq-htpy}(H))$

Def

•  $X$  type

Its suspension is given by:

•  $\sum X$  type

•  $N, S : \sum X$  poles

• merid:  $X \rightarrow (N = S)$

s.t.

$$\begin{array}{ccc} X & \longrightarrow & 1 \\ \downarrow \text{merid} & \nearrow & \downarrow \text{const}_S \\ 1 & \xrightarrow{\text{const}_N} & \sum X \end{array}$$

is a pushout

$$\bullet \text{ Spheres: } S^0 \cong 2 \quad S^{n+1} \cong \sum S^n$$

It works because  $X^{S^1} \rightarrow X^1$  is a pullback  $\forall X$

$$2 \xrightarrow{\text{id}} (b=b)$$

$$0_2 \mapsto 1$$

$$0_1 \mapsto \text{refl}_b$$

$$\begin{array}{ccc} X^{S^1} & \rightarrow & X^1 \\ \downarrow & & \downarrow \\ X^1 & \rightarrow & X^2 \end{array} \Rightarrow S^1 \cong \sum S^0$$

Inspired by the link between  $S^1$  and suspensions, we get the following lemma.

Lemma  $(\sum X \rightarrow Y) \rightarrow (\sum_{Y, Y' : Y} X \rightarrow y = y')$  is an equivalence

$$f \mapsto (f(N), f(S), f \cdot \text{merid})$$

Ex  $(\sum X \rightarrow Y)$   
cocone-map ↗  
 $\text{cocone}_{S^1}(Y) \xrightarrow{\sim} (\sum_{Y, Y' : Y} X \rightarrow y = y')$   
 $(i, j, H) \mapsto (i(*), j(*), H)$

Rhm If the left square is a pushout, then  
the right one is iff the rectangle is  
(posting)

$$\begin{array}{ccccc} A & \xrightarrow{i} & B & \xrightarrow{k} & C \\ f \downarrow & \circ & \downarrow g & \circ & \downarrow h \\ X & \xrightarrow{j} & Y & \xrightarrow{l} & Z \end{array}$$

Def  $f: A \rightarrow B$

Its cofiber is the pushout

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow & \text{H} \nearrow & \downarrow \text{inr}_f \\ 1 & \longrightarrow & \text{cofib}(f) \\ & \text{inl}_f & \end{array}$$

Lemma  $\text{cofiber}(\text{inr}_f) \cong \sum A$

Ex  $\begin{array}{ccccc} A & \xrightarrow{f} & B & \longrightarrow & 1 \\ \downarrow & & \downarrow \text{inr}_f & & \downarrow \text{inr}_{\text{inr}f} \\ 1 & \xrightarrow{\text{inl}_f} & \text{cofib}(f) & \xrightarrow{\text{inl}_{\text{inr}f}} & \text{cofiber}(\text{inr}_f) \cong \sum A \end{array}$