

Lecture 3

A small category

C locally small category with small colimits

$$u: A \rightarrow C$$

$$u^*: C \rightarrow \hat{A} = \text{Fun}(A^{\text{op}}, \text{Set})$$

$$u^*(y)_a = \text{Hom}_C(u(a), y)$$

$$u_!: \hat{A} \rightarrow C$$

$$u_!(x) = \varinjlim_{(a,s) \in A_x} u(a)$$

$$\text{Hom}_C(u_!(x), y) \cong \text{Hom}_{\hat{A}}(x, u^*(y))$$

$$\text{Want: } u_!(h_a) \cong u(a)$$

Lemma. I category with terminal object ω for all i in I $F: I \rightarrow E$ functor

$$\text{Hom}_I(i, \omega) = *$$

Then F has a colimit in E and $\varinjlim F \cong F(\omega)$

proof:

We have a cocone: $F \rightarrow F(\omega)$

$$\begin{array}{ccc} i \in \text{Ob}(I) & p_i: i \rightarrow \omega & \\ \downarrow u_i & \nearrow d_i & \\ F(i) & \xrightarrow{F(p_i)} & F(\omega) \\ \downarrow F(p_i) & \nearrow F(d_i) & \\ F(i) & \xrightarrow{F(p_i)} & F(\omega) \end{array}$$

If $F \xrightarrow{q} X$ is a cocone in E

$$\begin{array}{ccc} i & F(i) & \xrightarrow{q_i} X \\ \downarrow F(p_i) & \downarrow F(p_i) & \nearrow q_i \\ F(i) & \xrightarrow{F(p_i)} & F(\omega) \end{array} \quad \text{commutes}$$

$$\text{Cocones}(F, X) \xrightarrow{\cong} \text{Hom}(F(\omega), X)$$

$$q \mapsto q_\omega$$

Observation:

 A/h_a has a final object, namely $(a, 1_a)$

$$\text{Therefore } u(a) \cong \varinjlim_{(a', s) \in A/h_a} u(a') = u_!(h_a)$$

We may actually define $u_!$ so that

$$u(a) = u_!(h_a)$$

The rest of the proof follows right away from the Yoneda Lemma.

Remark: Let $F: \hat{A} \rightarrow C$ a functor. (A, C as above)

$$\text{let } u: A \rightarrow C \quad u(a) = F(h_a)$$

$$\text{There is: } u_! \rightarrow F$$

$$u_!(x) = \varinjlim_{(a,s) \in A/x} u(a) \rightarrow F(x) \cong F\left(\varinjlim_{(a,s) \in A/x} h_a\right)$$

univ. cocone. F(Canonical Cocone)

$$u(a) = F(h_a)$$

If F commutes with colimits, then $u_! \cong F$
 $\Rightarrow F$ has a right adjoint u^*

requires all smallness hypothesis

Remark: if $F: C \rightarrow D$ has a right adjoint

$$G: D \rightarrow C$$

$$\text{Hom}_D(F(x), y) \cong \text{Hom}_C(x, G(y))$$

$\Rightarrow F$ commutes with all colimits. (exercise!)

Example: $\mathbf{1}$ terminal category
 $[0]$

$e = \{0\}$
 with only the identity

$$e^{op} = e$$

$$\text{Fun}(e, \text{Set}) = \hat{e} \xrightarrow{\cong} \text{Set}$$

$$x \mapsto x_0$$

$$x \in \text{Set}$$

$e/x = X$ seen as a category with only identities.

$$\lim_{(o,s) \in A/x} h_a = \coprod_{x \in X} 1x \cong X$$

$$A = e$$

$$\text{Set} \xrightarrow{F} C$$

F commutes with

colimits (\Rightarrow) F commutes with small sums.

Remark: Let A be small

C locally small with small colimits.

$\text{Fun}_!(\hat{A}, C) =$ full subcategory of $\text{Fun}(\hat{A}, C)$ spanned by colimit preserving functors.

$$\text{Then } \text{Fun}_!(\hat{A}, C) \rightarrow \text{Fun}(A, C)$$

$$F \mapsto F \circ h$$

is an equivalence of categories.

h : Yoneda embed. (exercise).

Example: A, B small categories.

$$u: A \rightarrow B \text{ functor}$$

$$\text{Then } u^*: \hat{B} \rightarrow \hat{A} \text{ has both a}$$

$$\text{left adjoint } u_!: \hat{A} \rightarrow \hat{B}$$

$$\text{and a right adjoint } u_*: \hat{A} \rightarrow \hat{B}$$

(co)limits are computed levelwise in \hat{A} and \hat{B}

$$F: I \rightarrow \hat{A} \quad \left(\lim_{i \in I} F(i)(a) \right) = \left(\lim_{i \in I} F(i)(a) \right)$$

$$u^*: \hat{B} \rightarrow \hat{A} \quad u^*(Y)_a = Y_{u(a)}$$

commutes with small colimits and limits

\Rightarrow it has a right adjoint u_*

$$u_*(X)_b = \text{Hom}_{\hat{A}}(u^*(h_b), X)$$

$$A \xrightarrow{u} B \xrightarrow{h} \hat{B}$$

$(hu)_!: \hat{A} \rightarrow \hat{B}$ is left adjoint to

$$(hu)^*: \hat{B} \rightarrow \hat{A}$$

$$Y \mapsto (a \mapsto \text{Hom}_{\hat{B}}(h_{u(a)}, Y))$$

$$\text{Hom}_{\hat{B}}(h_{u(a)}, Y) \cong Y_{u(a)} \quad (\text{Yoneda})$$

$$\Rightarrow (hu)^* = u^* \quad \text{Define } u_! = (hu)_!$$

$u_!: \hat{A} \rightarrow \hat{B}$ is the left Kan extension of u along h

$u_*: \hat{A} \rightarrow \hat{B}$ " right " " " "

What is a good language for category theory?

in the category of categories,

- finite products exist.

$$A \times B \quad \text{ob}(A \times B) = \text{ob}(A) \times \text{ob}(B)$$

$$\text{Hom}_{A \times B}((a_0, b_0), (a_1, b_1)) = \text{Hom}_A(a_0, a_1) \times \text{Hom}_B(b_0, b_1)$$

- Internal Hom $\underline{\text{Hom}}: \text{Cat}^{\text{op}} \times \text{Cat} \rightarrow \text{Cat}$

$$\text{Hom}_{\text{Cat}}(A \times B, C) \cong \text{Hom}_{\text{Cat}}(A, \underline{\text{Hom}}(B, C))$$

$$\text{Here } \underline{\text{Hom}}(B, C) = \text{Fun}(B, C)$$

$$\text{Hom}_{\text{Cat}}(A \times B, C) \cong \text{Hom}_{\text{Cat}}(A, \text{Fun}(B, C))$$

The concept of category is an example of algebraic structure.

Example: algebraic structure of commutative rings.

R comm. ring with unit.

$$\begin{array}{llll} R \text{ set} & R \times R \rightarrow R & R \times R \rightarrow R & 0 \in R \\ & (x, y) \mapsto x+y & (x, y) \mapsto xy & 1 \in R \\ & & & + \text{ axioms } \dots \end{array}$$

What do these axioms mean: for any polynomial $p \in \mathbb{Z}[T_1, \dots, T_n]$ there is a function

$R^n \rightarrow R$. This is compatible
 $x = (x_1, \dots, x_n) \mapsto p(x)$ with composition of
 polynomials

$$\text{Poly} \subseteq_{\text{full}} \text{CRing} = \{\text{commutative rings}\}$$

$$\text{ob}(\text{Poly}) = \{\underline{n} = \mathbb{Z}[T_1, \dots, T_n] \mid n \geq 0\}$$

Poly has finite products $\coprod_{\{1, \dots, n\}} \mathbb{Z}[T] \cong \underline{n}$
 $\text{Poly}^{\text{op}} \xrightarrow{\quad} \text{products}$

R Comm. ring

$R: \text{Poly}^{\text{op}} \rightarrow \text{Set}$ finite product preserving
 $\underline{n} \mapsto R^n$ functor

$\underline{m} \rightarrow \underline{n} (=)$ m polynomials $\leadsto m$ maps
 $R^n \rightarrow R$
 $\leadsto R^n \rightarrow R^m$

Equivalence of categories:

$\hat{\text{Poly}} \cong \{\text{finite product preserving functors } \text{Poly}^{\text{op}} \rightarrow \text{Set}\} \cong \{\underline{\mathbb{Z}} \subset \text{CRings}\}$
 $X \mapsto X(\underline{1})$

Rem: This is an instance of a Lawvere theory

observation: an element of R is a morphism $\mathbb{Z}[T] \rightarrow R$

For categories, it is similar, but a category does not consist of a set with a structure.
It is a graph with a structure.

C category

Graph:

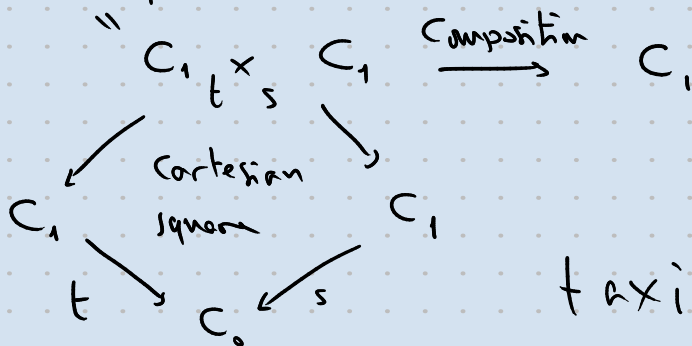
$$C_1 = \{\text{set of maps in } C\} = \text{Arr}(C)$$

$$C_0 = \text{Ob}(C)$$

$$C_1 \begin{array}{c} \xrightarrow{s} \\ \xrightarrow{t} \end{array} C_0 \quad \begin{array}{l} \text{source} \\ \text{target} \end{array}$$

$$C_0 \rightarrow C_1 \quad \text{identity} \\ x \mapsto 1_x$$

$$\{x \xrightarrow{f} y \xrightarrow{g} z\}$$



What is the analog of "Poly" for categories.

Simplicial sets

Observe: any partially ordered set E determines a category with

objects: elements of E

morphisms: $\text{Hom}(x, y) = \begin{cases} * & \text{if } x \leq y \\ \emptyset & \text{else} \end{cases}$

E, F partially ordered sets

$f: E \rightarrow F$ non decreasing $f(x) \leq f(y)$

for all $x \leq y$

iff $f: E \rightarrow F$ is a functor.

$n \geq -1$

$[n] = \{0, \dots, n\}$ with the canonical total order, seen as category

$[-1] = \emptyset$

$[0] = \text{terminal category}$

$[1]$ represents arrows.

$\text{Arr}(C) = \text{Hom}_{\text{Cat}}([1], C)$

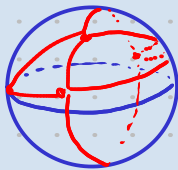
$\text{Fun}([1], C)$ is the category of arrows in C

Definition. The category of simplices is the category Δ with objects $[n], n \geq 0$ and $\text{Hom}_{\Delta}([m], [n]) = \text{Hom}_{\text{Cat}}([m], [n])$.

Definition. A simplicial set is a functor from Δ^{op} to sets.

We denote by $\mathbf{sSet} = \text{Fun}(\Delta^{\text{op}}, \text{Set}) = \hat{\Delta}$ the category of (small) simplicial sets.

Historically: simplicial sets were introduced in topology.



$u: \Delta \rightarrow \text{Top} = \{\text{topological spaces}\}$

$$[n] \mapsto \Delta_{\text{top}}^n = \{(x_0, \dots, x_n) \in [0, 1]^{n+1} \mid \sum_{i=0}^n x_i = 1\} \subseteq \mathbb{R}^{n+1}$$

$f: [m] \rightarrow [n]$ non-decreasing comp. f.

$$\leadsto f_x: \Delta_{\text{top}}^m \rightarrow \Delta_{\text{top}}^n \quad f_x(x_0, \dots, x_m) = (y_0, \dots, y_n)$$

$$\text{with } y_j = \sum_{i \in f^{-1}(j)} x_i$$

$$u^*: \text{Sing}: \text{Top} \rightarrow \text{SSet}$$

$$Y \mapsto \text{Sing}(Y)$$

$$\text{Sing}(Y)_n = C(\Delta^n_{\text{top}}, Y)$$

$$\Delta^0_{\text{top}} = \{1\} \subset \mathbb{R}$$

$$\Delta^1_{\text{top}} = \bullet \text{---} \bullet \subseteq \mathbb{R}^2$$

$$\Delta^2_{\text{top}} = \begin{array}{c} \triangle \\ \text{filled with blue wavy lines} \end{array} \subseteq \mathbb{R}^3$$

$$\Delta^3_{\text{top}} = \begin{array}{c} \triangle \\ \text{filled with red wavy lines} \end{array} \subseteq \mathbb{R}^4$$

⋮

$$\gamma: \Delta^1_{\text{top}} \rightarrow X \quad \text{path}$$

$$I = [0, 1]$$

γ, γ' homotopic.

$$I \cong \Delta^1_{\text{top}}$$

$$I \times I \xrightarrow{h} X$$

$$h(0, t) = \gamma(t)$$


$$h(1, t) = \gamma'(t)$$



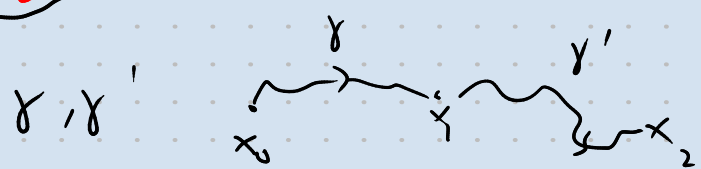
$$h(0, s) = h(0, 0)$$

$$h(1, s) = h(1, 1)$$

$$[x] = \frac{\boxed{\text{wavy line}}}{\sim} \xrightarrow{h} X$$


 \sim

$$\text{circle with wavy line} \xrightarrow{h} X$$



$$I = [0, 2] \quad \gamma' * \gamma \text{ concatenation.}$$

$$\gamma' * \gamma(t) = \begin{cases} 2t & t \leq 1/2 \\ 2t-1 & \text{else} \end{cases}$$


$$\gamma'' \sim \gamma' * \gamma$$

γ'' is a composition of γ' and γ up to homotopy.

$$\text{circle with two wavy lines} \xrightarrow{h} X$$

γ''

$$\Delta_{\text{tor}}^2 = 0 \quad \text{circle with two wavy lines and points } 0, 1, 2$$

$$\Delta_{top}^1 \hookrightarrow \Delta_{top}^2 = \Delta_{top}^2$$


$\delta_i^2: [1] \rightarrow [2]$
 unique injective map
 which avoids i

$$d_i^2 (\delta_i^2)_* : \Delta_{top}^1 \rightarrow \Delta_{top}^2$$

