

## Higher Category Theory

### Assignment 8

#### Exercise 1

*Proof.* (1) From definition one sees that  $\text{Ob}(\mathcal{A}) = \text{Ob}(\mathcal{A}_0) \cup \text{Ob}(\mathcal{A}_1)$ . We construct the functor  $u: \mathcal{A} \rightarrow \mathcal{C} * \mathcal{D}$  as follows: on objects,

$$u(a) := \begin{cases} u_0(a), & \text{if } a \in \text{Ob}(\mathcal{A}_0) \\ u_1(a), & \text{if } a \in \text{Ob}(\mathcal{A}_1) \end{cases}$$

and on morphisms,

$$u(a \rightarrow b) := \begin{cases} u_i(a \rightarrow b), & \text{if } a, b \in \text{Ob}(\mathcal{A}_i), i = 0, 1 \\ 0, & \text{if } a \in \text{Ob}(\mathcal{A}_0) \text{ and } b \in \text{Ob}(\mathcal{A}_1). \end{cases}$$

Note that there is no  $a \rightarrow b$  with  $a \in \text{Ob}(\mathcal{A}_1)$  and  $b \in \text{Ob}(\mathcal{A}_0)$ , since otherwise applying  $q: \mathcal{A} \rightarrow [1]$  to it yields a morphism  $1 \rightarrow 0$ . From the definition it follows that the restriction of  $u$  to  $\mathcal{A}_0$  and  $\mathcal{A}_1$  are  $u_0$  and  $u_1$  respectively. Next we check that  $pu = q$ . Indeed, we have  $pu(a) = pu_i(a) = i = q(a)$  for  $a \in \text{Ob}(\mathcal{A}_i)$  ( $i = 0, 1$ ), and

$$pu(a \rightarrow b) = \begin{cases} pu_i(a \rightarrow b) = \text{id}_i = q(a \rightarrow b), & \text{if } a, b \in \text{Ob}(\mathcal{A}_i), i = 0, 1 \\ 0 \rightarrow 1 = q(a \rightarrow b), & \text{if } a \in \text{Ob}(\mathcal{A}_0) \text{ and } b \in \text{Ob}(\mathcal{A}_1). \end{cases}$$

Suppose that there is another  $u': \mathcal{A} \rightarrow \mathcal{C} * \mathcal{D}$  such that  $pu' = q$  and that  $u'$  restricts to  $u_i$  on  $\mathcal{A}_i$ . Then  $u$  and  $u'$  agree on  $\mathcal{A}_i$ , and for any  $a \rightarrow b$  in  $\mathcal{A}$  with  $a \in \text{Ob}(\mathcal{A}_0)$ ,  $b \in \text{Ob}(\mathcal{A}_1)$ ,  $u'(a \rightarrow b) = u(a \rightarrow b) - 0$  is the only morphism between  $u(a) = u'(a) \in \text{Ob}(\mathcal{C})$  and  $u(b) = u'(b) \in \text{Ob}(\mathcal{D})$ . Hence  $u = u'$ .

(2) Recall that  $N(\mathcal{C}) * N(\mathcal{D})$  is given by

$$(N(\mathcal{C}) * N(\mathcal{D}))_n = \coprod_{\substack{i+1+j=n \\ -1 \leq i, j \leq n}} N(\mathcal{C})_i \times N(\mathcal{D})_j$$

for each  $[n] \in \text{Ob}(\Delta)$ . We then define a map

$$\varphi_n: (N(\mathcal{C}) * N(\mathcal{D}))_n \rightarrow N(\mathcal{C} * \mathcal{D})_n$$

as below. Take an arbitrary  $(x, y) \in N(\mathcal{C})_i \times N(\mathcal{D})_j$  with  $-1 \leq i, j \leq n$  and  $i+1+j = n$ , where  $x$  or  $y$  may be empty. Then  $(x, y)$  corresponds to a unique  $([i] \xrightarrow{u_0} \mathcal{C}, [j] \xrightarrow{u_1} \mathcal{D})$  via

the adjunction  $\tau \dashv N$  plus the facts that the counit is an isomorphism and  $\Delta^i = N([i])$ . Moreover, let us define a functor  $q: [n] \rightarrow [1]$  by sending  $i \mapsto 0$  and  $i + 1 \mapsto 1$ . Then by (1), we get a unique functor  $u: [n] \rightarrow \mathcal{C} * \mathcal{D}$  such that  $q = pu$  and  $u|_{[i]} = u_0$ ,  $u|_{[j]} = u_1$ , where  $p: \mathcal{C} * \mathcal{D} \rightarrow [1]$  is the same as in (1). Again under the adjunction,  $u$  corresponds uniquely to a simplicial map  $\Delta^n \rightarrow N(\mathcal{C} * \mathcal{D})$  (a.k.a an element of  $N(\mathcal{C} * \mathcal{D})_n$ ), which we denote by  $\varphi_n(x, y)$ .

We claim that  $\varphi_n$  is a bijection. To this end, we construct an inverse  $\psi_n$  to  $\varphi_n$ . Take an element  $z$  in  $N(\mathcal{C} * \mathcal{D})_n$ , and it corresponds via adjunction to some  $u: [n] \rightarrow \mathcal{C} * \mathcal{D}$ . Put  $q := pu$ ,  $i := \max\{i \mid q(i) = 0\}$  and  $j := n - i - 1$ . Then we can define  $u_0: [i] \rightarrow \mathcal{C}$  by restricting  $u$  to  $[i]$ , and  $u_1: [j] \rightarrow \mathcal{D}$  by the composition  $[j] \xrightarrow{k \mapsto k+i+1} [n] \xrightarrow{u} \mathcal{C} * \mathcal{D}$ , which actually lands in  $\mathcal{D}$ . Again the pair  $(u_0, u_1)$  corresponds under adjunction to an element of  $N(\mathcal{C})_i \times N(\mathcal{D})_j$ , for which we write  $\psi_n(z)$ .

The well-definedness of  $\varphi_n$  and  $\psi_n$  lies in the adjunction bijection and the universal property of the join, which in every step of our construction provides a unique choice.

Verifying  $\psi_n$  and  $\varphi_n$  being mutually inverse is straightforward. For example, to check that  $\varphi_n \psi_n = \text{id}_{N(\mathcal{C} * \mathcal{D})_n}$ , we consider an arbitrary  $u: [n] \rightarrow \mathcal{C} * \mathcal{D}$  and  $u_0, u_1$  constructed as above. By the universal property of the join, the  $u': [n] \rightarrow \mathcal{C} * \mathcal{D}$  such that  $q = pu'$  and  $u'|_{[i]} = u_0$ ,  $u'|_{[j]} = u_1$  is unique (and thus equals to  $u$ ), which corresponds to the image under  $\varphi_n \psi_n$ . The argument for  $\psi_n \varphi_n = \text{id}_{(N(\mathcal{C}) * N(\mathcal{D}))_n}$  is similar.

In what follows we show that the bijection  $\varphi_n: (N(\mathcal{C}) * N(\mathcal{D}))_n \rightarrow N(\mathcal{C} * \mathcal{D})_n$  is natural in  $[n]$ . For this, we take a functor  $f: [m] \rightarrow [n]$  and  $-1 \leq i, j \leq n$  such that  $i + j + 1 = n$ . Then there exists a unique pair of integers  $(a, b)$  and functors  $f_a: [a] \rightarrow [i]$ ,  $f_b: [b] \rightarrow [j]$  satisfying  $a + 1 + b = m$  and  $f_a * f_b = f$ . Explicitly, one has  $a = \max\{a \mid f(a) \leq i\}$ ,  $f_a(k) = f(k)$  and  $f_b(k) = f(k + a + 1) - i - 1$ . Consider the following diagram

$$\begin{array}{ccccc}
(N(\mathcal{C}) * N(\mathcal{D}))_n & \longrightarrow & N(\mathcal{C} * \mathcal{D})_n & & (x, y) \longmapsto \varphi_n(x, y) \\
\downarrow & & \downarrow & & \downarrow \\
(N(\mathcal{C}) * N(\mathcal{D}))_m & \longrightarrow & N(\mathcal{C} * \mathcal{D})_m & & (f_a^* x, f_b^* y) \mapsto \varphi_m(f_a^* x, f_b^* y) \quad f^* \varphi_n(x, y)
\end{array}$$

Note that under the adjunction,  $f^* \varphi_n(x, y)$  corresponds to  $u \circ f$ , whereas  $f_a^* x$ ,  $f_b^* y$  corresponds to  $u_0 \circ f_a$  and  $u_1 \circ f_b$ , which correspond to some  $u': [m] \rightarrow \mathcal{C} * \mathcal{D}$ . Note that the restriction of  $u \circ f$  on  $[a]$  and  $[b]$  are respectively  $u_0 \circ f_a$  and  $u_1 \circ f_b$ , and also that  $p \circ u' = q_m = q_n \circ f = p \circ u \circ f$ , where  $q_m: [m] \rightarrow [1]$  and  $q_n: [n] \rightarrow [1]$  are given by  $[a] \mapsto 0, [a + 1] \mapsto 1$  and  $[i] \mapsto 0, [i + 1] \mapsto 1$  respectively. By the universal property of the join (1), one has  $u' = u \circ f$ . Therefore  $\varphi_m(f_a^* x, f_b^* y) = f^* \varphi_n(x, y)$ , and in conclusion,  $\varphi_n$  is natural with regard to  $[n]$ .

So far we have proved  $N(\mathcal{C} * \mathcal{D}) \cong N(\mathcal{C}) * N(\mathcal{D})$ . □

## Exercise 2

*Proof.* (1) Notice that  $N(0) = \Delta^{-1}$ . Now, applying (1.2), we see that  $\Delta^i * \Delta^{n-i-1} = N([i]) * N([n-i-1]) \cong N([i] * [n-i-1])$ , so it is enough to check that  $[n] \cong [i] * [n-i-1]$ .

In  $[i] * [n-i-1]$  there is exactly one morphism between any pair of objects coming from  $[i]$  or from  $[n-i-1]$ . Also, given an object in  $[i]$  and one in  $[n-i-1]$ , by definition of  $[i] * [n-i-1]$  there is exactly one morphism between the two of them in this category, from the former to the latter. This shows that  $[i] * [n-i-1]$  is an order and, since its set of objects has cardinality  $n+1 = (i+1) + ((n-i-1)+1)$  like the one of  $[n]$ , we get that the two categories are (uniquely) isomorphic, as desired.

(2) Note that the  $m$ -simplices of  $\Lambda_k^i$  are those non-surjective  $f: [m] \rightarrow [i]$  whose images do not contain  $[i] \setminus \{k\}$ . Since  $v(i) = 0$ , then  $v$  carries each  $m$ -simplex to  $0: [m] \rightarrow [1]$ . Therefore  $v$  sends  $\Lambda_k^i$  to  $0$  in  $\Delta^1$ .

Next we show that there exists  $\alpha: \Delta^i \rightarrow X$  extending  $u|_{\Lambda_k^i}$ . For this, we claim that  $u|_{\Lambda_k^i}: \Lambda_k^i \rightarrow X * Y$  lands in  $X$ . Indeed, in the commutative diagram below,

$$\begin{array}{ccc} \Lambda_k^i & \xrightarrow{u|_{\Lambda_k^i}} & X * Y \\ \downarrow & & \downarrow p \\ \Delta^i & \xrightarrow{v|_{\Delta^i}} & \Delta^1 \end{array}$$

since the restriction of  $v$  to  $\Lambda_k^i$  sends it to  $0$ , so does  $pu|_{\Lambda_k^i}$ . Note also that  $p: X * Y \rightarrow \Delta^0 * \Delta^0 = \Delta^1$  is given by sending each  $n$ -simplex  $(x, y) \in X_r \times Y_s$  ( $r+1+s=n$ ) to  $([r] \rightarrow [0], [s] \rightarrow [0])$ . Hence  $pu|_{\Lambda_k^i}$  sending  $\Lambda_k^i$  to  $0$  in  $\Delta^1$  means that the  $Y$ -entries of  $u|_{\Lambda_k^i}$  are all empty, i.e. it lands in  $X$ . Then by the fact that  $X$  is an  $\infty$ -category, we get a lift  $\alpha: \Delta^i \rightarrow X$  extending  $u|_{\Lambda_k^i}$ .

If there is  $\beta: \Delta^{n-i-1} \rightarrow Y$ , then by (1) we have  $\alpha * \beta: \Delta^n = \Delta^i * \Delta^{n-i-1} \rightarrow X * Y$ . To show  $p(\alpha * \beta) = v$ , we claim that  $v = v_0 * v_1$ , where  $v_0 = (\Delta^i \rightarrow \Delta^0)$  and  $v_1 = (\Delta^{n-i-1} \rightarrow \Delta^0)$  are the unique simplicial maps. Indeed, it suffices to note that  $v$  is given by  $[i] * [n-i-1] \rightarrow [0] * [0]$  since  $i = \max\{i \mid v(i) = 0\}$ . Then we conclude that  $p(\alpha * \beta) = (p_0\alpha) * (p_1\beta) = v_0 * v_1 = v$  by noting again that  $\Delta^i \rightarrow \Delta^0$  and  $\Delta^{n-i-1} \rightarrow \Delta^0$  are unique, where  $p_0: X \rightarrow \Delta^0$  and  $p_1: Y \rightarrow \Delta^0$  are the simplicial maps defining  $p$  (i.e.  $p = p_0 * p_1$ ).

(3) Let's apply the operator  $(-)^{\text{op}}$  to the commutative diagram

$$\begin{array}{ccc} \Lambda_k^n & \xrightarrow{u} & X * Y \\ \downarrow & & \downarrow p \\ \Delta^n & \xrightarrow{v} & \Delta^1 \end{array} ,$$

giving us a commutative diagram which admits a filler  $g$  by (2.2). Here we use the fact that  $(X * Y)^{\text{op}} \cong Y^{\text{op}} * X^{\text{op}}$ .

$$\begin{array}{ccc} \Lambda_{n-k}^n & \xrightarrow{u^{\text{op}}} & Y^{\text{op}} * X^{\text{op}} \\ \downarrow & \nearrow g & \downarrow p^{\text{op}} \\ \Delta^n & \xrightarrow{v^{\text{op}}} & \Delta^1 \end{array}$$

By reapplying the operator (which is an involution) we get then the desired filler  $f = g^{\text{op}}$ .

$$\begin{array}{ccc} \Lambda_k^n & \xrightarrow{u} & X * Y \\ \downarrow & \nearrow f & \downarrow p \\ \Delta^n & \xrightarrow{v} & \Delta^1 \end{array}$$

(4) Since the diagram is commutative and the map on the left is a monomorphism bijective on objects, the fact that  $v(j) = 0$  is equivalent to  $pu(j) = 0$  and therefore, by definition of  $p$  and  $i$ ,  $u(j) \in X_0$  for all  $0 \leq j \leq i$ ,  $u(j) \in Y_0$  for all  $i < j \leq n$ .

Suppose to have a lifting  $f$  already. We will start showing its uniqueness by rewriting  $\Delta^n$  as  $\Delta^i * \Delta^{n-i-1}$ . This gives us the restrictions  $v|_{\Delta^i} = v|_{v^{-1}(0)}$ ,  $v|_{\Delta^{n-i-1}} = v|_{v^{-1}(1)}$ , which map all 0-simplices respectively to 0 and 1 by our previous observation. Precomposing by the inclusion  $\Lambda_i^n \rightarrow \Delta_i^n$ , we get that  $v|_{\Delta^i} = pu|_{\Delta^i}$ ,  $v|_{\Delta^{n-i-1}} = pu|_{\Delta^{n-i-1}}$ , thus all of  $\Delta^i$  is sent to  $X$  and all of  $\Delta^{n-i-1}$  to  $Y$  under  $u$  by the description of  $p$ . This allows us to construct the following commutative diagram

$$\begin{array}{ccccc} & & & & X \sqcup Y \\ & & & \nearrow u|_{\Delta^i \sqcup \Delta^{n-i-1}} & \downarrow \\ \Delta^i \sqcup \Delta^{n-i-1} & \hookrightarrow & \Lambda_i^n & \xrightarrow{u} & X * Y \\ \downarrow & & \downarrow & \nearrow f & \downarrow p \\ \Delta^i * \Delta^{n-i-1} & \xlongequal{\quad} & \Delta^n & \xrightarrow{v} & \Delta^1 \\ \searrow & & & & \uparrow \\ & \partial \Delta^1 & & & \end{array}$$

Now, restricting our focus to the commutative diagram

$$\begin{array}{ccc}
& X \sqcup Y & \hookrightarrow X * Y \\
u|_{\Delta^i \sqcup u|_{\Delta^{n-i-1}}} \nearrow & & \nearrow f \\
\Delta^i \sqcup \Delta^{n-i-1} & \xrightarrow{\quad} & \Delta^n \\
\downarrow & \nwarrow & \downarrow v \\
\partial \Delta^1 & \xrightarrow{\quad} & \Delta^1
\end{array} \quad , \quad p$$

we see that there can be at most one  $f$  solving our initial lifting problem since all solutions must fill this diagram and we have the universal property of the join.

Remember that  $v = v|_{\Delta^i} * v|_{\Delta^{n-i-1}}: \Delta^i * \Delta^{n-i-1} \rightarrow \Delta^0 * \Delta^0$  and, by essentially the same argument,  $p = p|_X * p|_Y: X * Y \rightarrow \Delta^0 * \Delta^0$ .

Now,  $f := u|_{\Delta^i} * u|_{\Delta^{n-i-1}}: \Delta^n \cong \Delta^i * \Delta^{n-i-1} \rightarrow X * Y$  is such that  $pf = (p|_X * p|_Y) \cdot (u|_{\Delta^i} * u|_{\Delta^{n-i-1}}) = (p|_X \cdot u|_{\Delta^i}) * (p|_Y \cdot u|_{\Delta^{n-i-1}}) = pu|_{\Delta^i} * pu|_{\Delta^{n-i-1}} = v|_{\Delta^i} * v|_{\Delta^{n-i-1}} = v$  and, by construction,  $f$  coincides with  $u$  when restricted to  $\Delta^i * \partial \Delta^1$  and  $\Lambda_i^i * \Delta^{n-i-1}$  seen as subobjects of  $\Lambda_i^n$  covering it. This shows that  $u = f|_{\Lambda_i^n}$ , thus  $f$  solves the lifting problem we started from.

(5) It is enough to check that  $ho(X * Y)$  has the universal property of the join of  $ho(X)$  and  $ho(Y)$ . Let's consider then functors  $q: \mathcal{A} \rightarrow [1]$ ,  $u_0: \mathcal{A}_0 \rightarrow ho(X)$ ,  $u_1: \mathcal{A}_1 \rightarrow ho(Y)$ , and the obvious embedding  $ho(X) \sqcup ho(Y) \rightarrow ho(X * Y)$  (it's faithful because joining two  $\infty$ -categories does not produce new 2-simplices establishing homotopy relations between the morphisms in  $X$  or in  $Y$ ).  $p: ho(X * Y) \rightarrow [1]$  will be given by  $x \mapsto 0$ ,  $y \mapsto 1$ , and  $p(x \rightarrow y) = (0 \rightarrow 1)$ . Notice that there's no map  $y \rightarrow x$  since it would come from  $(y, x) \in Y_0 \times X_0$ .

$$\begin{array}{ccc}
& ho(X) \sqcup ho(Y) & \hookrightarrow ho(X * Y) \\
u_0 \sqcup u_1 \nearrow & & \nearrow f \\
\mathcal{A}_0 \sqcup \mathcal{A}_1 & \xrightarrow{\quad} & \mathcal{A} \\
\downarrow & \nwarrow & \downarrow q \\
[0] \sqcup [0] & \xrightarrow{\quad} & [1]
\end{array} \quad p$$

To construct a factorization  $f: \mathcal{A} \rightarrow ho(X * Y)$  of  $q$  making the diagram commute we are forced to start by composing  $u_0 \sqcup u_1$  with the embedding, which gives us  $a_i \mapsto u_i(a)$  for  $a_i \in \text{Ob}(\mathcal{A}_i)$ ,  $g \mapsto u_i(g)$  for  $g \in \text{Mor}(\mathcal{A}_i)$ . To extend then this functor to  $\mathcal{A}$ , we have to send maps  $a_0 \rightarrow a_1$  to the unique morphism  $f(a_0) \rightarrow f(a_1)$  given by the element  $(f(a_0), f(a_1)) \in X_0 \times Y_0 \subset (X * Y)_1$ . Notice that there are no morphisms  $a_1 \rightarrow a_0$  in  $\mathcal{A}$  by the definition of the  $\mathcal{A}_i$  since they would need to be mapped to an arrow  $1 \rightarrow 0$  under  $q$ , but it is not there.

We see that identities are trivially preserved and compositions of arrows all in  $\mathcal{A}_i$  are too since the  $u_i$  and the embedding are functors. If one composes instead an arrow  $a'_0 \rightarrow a_0$  with one  $a_0 \rightarrow a_1$  whose domain and codomain lie in different categories the result is again a map  $a'_0 \rightarrow a_1$  with domain and codomain lying in different categories and will therefore be mapped to the unique map  $f(a'_0) \rightarrow f(a_1)$ . Likewise, the composition of the maps one obtains by first applying  $f$  and then composing in  $ho(X * Y)$  is again the unique map  $f(a'_0) \rightarrow f(a_1)$ . A symmetric argument for  $a_0 \rightarrow a_1$  and  $a_1 \rightarrow a'_1$  then gives us functoriality.

By construction, the desired diagram commutes and uniqueness of factorization follows from the fact that when we were defining  $f$  we had a unique possible choice at every step.  $\square$