CSE 331 Fall 2018

Homework 1: Q2

Name: Waiwai Kim, waiwaiki

1 Part (a) Proof Idea

We need to prove that a stable pair of schedules always exists for every n movies of platform A and platform B respectively when the lowest rating show of A is larger than the highest rating show of B. In order to prove the statement, we need to show that platform A always wins regardless of schedules.

We can do so by proof by contradiction. Assume that A does not win all match-ups. In return, this means that there exists at least a B's show that has a higher rating than a A's show. This contradicts the fact that min(A) > max(B). Thus, it must be true that A always wins all match-ups and that no change in schedule can induce A or B win more. All schedules are stable.

2 Part (b) Proof Idea

The question can be rephrased in first order logic in the following sentence. For all sets of n movies and their ratings, any pair of schedules is stable regardless of the orders. The part (b) asks us to prove that the statement is false and show that a stable pair of schedules does NOT always exist for every set of n movies and their ratings.

We can accomplish this by showing a counter-example, in which every possible pair of schedules is not stable. An example can be n = 2 in which **A** has one show that has higher rating than any **B**'s show and **A**'s second show that has the same rating as one of **B**'s shows. In this structure, we will exploit the fact that **A** wins when there is a tie and induce instability.

3 Part (b) Proof Details

As the first example of our proof, we will set an input to the counter-example. We will adopt the names of Amazing and Bombastic as they were used in the provided question. This example assumes that n = 2 and that Amazing has two movies a and b, and Bombastic has two movies c and d. In Amazing's movies, a has a rating of 50 and b has a rating of 48. In Bombastic's movies, c has a rating of 49 and d has a rating of 48. Note that movies b and d have the same rating of 48. Also note that there are four possible combinations of schedules; thus, we will go over all four cases and show that every case is NOT stable.

The number of possible schedules =
$$(n!)^2 = (2!)^2 = 4$$

Assume we start with a case shown in Table 1 in which *Amazing* wins both days. In this scenario, *Bombastic* can change its schedule as shown in 2 to win the second day. Thus, Case 1 is not stable. Conversely, *Amazing* can change its schedule to lose one day.

In response to Table 2, Amazing can change its schedule as shown in Table 3 to win both days. Thus, Case 2 is not stable.

CSE 331 Fall 2018

Table 1: Counter-example 1

	Amazing		Bombastic		***
Case 1	Movie	Rating	Movie	Rating	Winner
Day 1	a	50	c	49	Amazing
Day 2	b	48	d	48	Amazing

Table 2: Counter-example 2

	Amazing		Bombastic		
Case 2	Movie	Rating	Movie	Rating	Winner
Day 1	a	50	d	48	Amazing
Day 2	b	48	c	49	Bombastic

Thirdly, in response to Table 3, *Bombastic* can change its schedule as shown in Table 4 to win the first day. Thus, Case 3 is not stable. Conversely, *Amazing* can change its schedule to lose one day.

Table 3: Counter-example 3

	Amazing		Bombastic		
Case 3	Movie	Rating	Movie	Rating	Winner
Day 1	b	48	d	48	Amazing
Day 2	a	50	c	49	Amazing

Fourthly and lastly, in response to Table 4, Amazing can change its schedule as shown in Table 1 to win both days. Thus, Case 4 is not stable. Conversely, Bombastic can change its schedule to lose one day. In the end, we have exhausted all possible 4 cases of schedules when n = 2. It also does not matter which the initial state is; Amazing or Bombastic can always change its schedule to alter the outcomes. Thus, all schedules are unstable in the given example. In return, we have proved that a stable pair of schedule does NOT always exist for every general set of n movies and their corresponding ratings.

Table 4: Counter-example 4

	Amazing		Bombastic		
Case 4	Movie	Rating	Movie	Rating	Winner
Day 1	b	48	С	49	Bombastic
Day 2	a	50	d	48	Amazing