Lecture 3: Borrow checker

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February 28, 2024

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Borrow Checker



```
let mut v = vec![1, 2, 3];
let x = &v[0];
v.push(4);
println!("{}", x);
```

```
let mut v = vec![1, 2, 3];
    let x = &v[0];
    v.push(4);
   println!("{}", x);
error[E0502]: cannot borrow 'v' as mutable because it is also
borrowed as immutable
 --> src/main.rs:8:5
7 | let x = &v[0];
                - immutable borrow occurs here
8 |
     v.push(4);
        ^^^^^^ mutable borrow occurs here
       println!("{}", x);
                       - immutable borrow later used here
```

```
fn sum(v: Vec<i32>) -> i32 {
    let mut result = 0;
    for i in v {
        result += i;
    result
fn main() {
    let mut v = vec![1, 2, 3];
    println!("first sum: {}", sum(v));
    v.push(4);
   println!("second sum: {}", sum(v))
```

```
error[E0382]: borrow of moved value: `v`
  --> src/main.rs:12:5
10 l
        let mut v = vec![1, 2, 3];
             ---- move occurs because `v` has type `Vec<i32>`,
      which does not implement the 'Copy' trait
11 I
        println!("first sum: {}", sum(v));
                                       - value moved here
12
        v.push(4);
         ^^^^^^ value borrowed here after move
```

- Each value in Rust has a variable that's called it's owner.
- There can be only one owner at a time.
- When the owner goes out of scope, the value will be dropped.

```
fn main() {
    let s = vec![1, 4, 8, 8];
    let u = s;
    println!("{:?}", u);
    println!("{:?}", s); // This won't compile!
}
```

```
fn om_nom_nom(s: Vec<i32>) {
    println!("I have consumed {s:?}");
}

fn main() {
    let s = vec![1, 4, 8, 8];
    om_nom_nom(s);
    println!("{s:?}");
}
```

```
fn om_nom_nom(s: Vec<i32>) {
    println!("I have consumed {s:?}");
}

fn main() {
    let s = vec![1, 4, 8, 8];
    om_nom_nom(s);
    println!("{s:?}");
}
```

- Each "owner" has the responsibility to clean up after itself.
- When you move s into om_nom_nom, it becomes the owner of s, and it will free s
 when it's no longer needed in that scope. Technically the s parameter in
 om_nom_nom become the owner.
- That means you can no longer use it in main!
- In C++, we would create a copy!

Given what we just saw, how can the following be the valid syntax?

```
fn om_nom_nom(n: u32) {
   println!("{} is a very nice number", n);
}
fn main() {
   let n: u32 = 42;
   let m = n;
    om_nom_nom(n);
    om_nom_nom(m);
   println!("{}", m + n);
}
```

- Say you have a group of lawyers that are reviewing and signing a contract over Google Docs (just pretend it's true:))
- What are some ground rules we'd need to set to avoid chaos?
- If someone modifies the contract before everyone else reviews/signs it, that's fine.
- But if someone modifies the contract while others are reviewing it, people might miss changes and think they're signing a contract that says something else.
- We should allow a single person to modify, or everyone to read, but not both.

- We can have multiple shared (immutable) references at once (with no mutable references) to a value.
- We can have only one mutable reference at once. (no shared references to it)
- This paradigm pops up a lot in systems programming, especially when you have "readers" and "writers". In fact, you've already studied it in the course of Theory and Practice of Concurrency.

- The lifetime of a value starts when it's created and ends the last time it's used
- Rust doesn't let you have a reference to a value that lasts longer than the value's lifetime
- Rust computes lifetimes at compile time using static analysis. (this is often an over-approximation!)
- Rust calls the special "drop" function on a value once its lifetime ends. (this is essentially a destructor)

```
fn main() {
    let mut x = 5;
    let y = &mut x;

    println!("y = {y}");
    x = 42; // ok
    println!("x = {x}");
}
```

```
fn main() {
    let mut x = 5;
    let y = &mut x;

    x = 42; // not ok
    println!("y = {y}");
    println!("x = {x}");
}
```

```
fn main() {
    let x1 = 42;
    let y1 = Box::new(84);
    { // starts a new scope
        let z = (x1, y1);
        // z goes out of scope, and is dropped;
        // it in turn drops the values from x1 and y1
    // x1's value is Copy, so it was not moved into z
    let x2 = x1;
    // y1's value is not Copy, so it was moved into z
    // let y2 = y1;
```

Option¹ and Result²

Let's remember their definitions:

```
enum Option<T> {
    Some(T),
    None,
}
enum Result<T, E> {
    Ok(T),
    Err(E),
}
```

¹Option documentation

²Result documentation

Matching Option:

```
let result = Some("string");
match result {
    Some(s) => println!("String inside: {s}"),
    None => println!("Ooops, no value"),
}
```

```
Useful functions .unwrap() and .expect():
    fn unwrap(self) -> T;
    fn expect(self, msg: &str) -> T;
```

```
Useful functions .unwrap() and .expect():
   let opt = Some(22022022);
    assert!(opt.is_some());
    assert!(!opt.is_none());
    assert_eq!(opt.unwrap(), 22022022);
   let x = opt.unwrap(); // Copy!
   let newest_opt: Option<i32> = None;
    // newest_opt.expect("I'll panic!");
   let new_opt = Some(Vec::<i32>::new());
    assert_eq!(new_opt.unwrap(), Vec::<i32>::new());
    // error[E0382]: use of moved value: `new_opt`
    // let x = new_opt.unwrap(); // Clone!
```

We have a magic function: fn as_ref(&self) -> Option<&T>; // &self is &Option<T> Let's solve a problem: let new_opt = Some(Vec::<i32>::new()); assert_eq!(new_opt.unwrap(), Vec::<i32>::new()); // error[E0382]: use of moved value: `new_opt` // let x = new_opt.unwrap(); // Clone! let opt_ref = Some(Vec::<i32>::new()); assert_eq!(new_opt.as_ref().unwrap(), &Vec::<i32>::new()); let x = new_opt.unwrap(); // We used reference! // There's also .as_mut() function

That means if type implements Copy, Option also implements Copy.

```
We can map Option<T> to Option<U>:
    fn map<U, F>(self, f: F) -> Option<U>;

Example:

let maybe_some_string = Some(String::from("Hello, World!"));
// `Option::map` takes self *by value*,
// consuming `maybe_some_string`
let maybe_some_len = maybe_some_string.map(|s| s.len());
assert_eq!(maybe_some_len, Some(13));
```

There's **A LOT** of different Option functions, enabling us to write beautiful functional code:

```
fn map_or<U, F>(self, default: U, f: F) -> U;
fn map_or_else<U, D, F>(self, default: D, f: F) -> U;
fn unwrap_or(self, default: T) -> T;
fn unwrap_or_else<F>(self, f: F) -> T;
fn and<U>(self, optb: Option<U>) -> Option<U>;
fn and_then<U, F>(self, f: F) -> Option<U>;
fn or(self, optb: Option<T>) -> Option<T>;
fn or_else<F>(self, f: F) -> Option<T>;
fn xor(self, optb: Option<T>) -> Option<T>;
fn zip<U>(self, other: Option<U>) -> Option<(T, U)>;
```

It's recommended for you to study the documentation and try to avoid match where possible.

Option and ownership

There's two cool methods to control ownership of the value inside:

```
fn take(&mut self) -> Option<T>;
fn replace(&mut self, value: T) -> Option<T>;
fn insert(&mut self, value: T) -> &mut T;
```

The first one takes the value out of the Option, leaving a None in its place.

The second one replaces the value inside with the given one, returning Option of the old value.

The third one inserts a value into the Option, then returns a mutable reference to it.

Option API and ownership: take

```
struct Node<T> {
    elem: T,
    next: Option<Box<Node<T>>>,
}
pub struct List<T> {
    head: Option<Box<Node<T>>>,
}
impl<T> List<T> {
    pub fn pop(&mut self) -> Option<T> {
        self.head.take().map(|node| {
            self.head = node.next;
            node.elem
        })
```

Option and optimizations

Rust guarantees to optimize the following types T such that Option<T> has the same size as T:

- Box<T>
- &T
- &mut T
- fn, extern "C" fn
- #[repr(transparent)] struct around one of the types in this list.
- num::NonZero*
- ptr::NonNull<T>

This is called the "null pointer optimization" or NPO.

Result

Functions return Result whenever errors are expected and recoverable. In the std crate, Result is most prominently used for I/O.

Results must be used! A common problem with using return values to indicate errors is that it is easy to ignore the return value, thus failing to handle the error. Result is annotated with the #[must_use] attribute, which will cause the compiler to issue a warning when a Result value is ignored.³

³The Error Model

Result API

We can match it as a regular enum:

```
let version = Ok("1.1.14");
match version {
    Ok(v) => println!("working with version: {:?}", v),
    Err(e) => println!("error: version empty"),
}
```

Result API

We have pretty the same functionality as in Option:

```
fn is_ok(&self) -> bool;
fn is_err(&self) -> bool;
fn unwrap(self) -> T;
fn unwrap_err(self) -> E;
fn expect_err(self, msg: &str) -> E;
fn expect(self, msg: &str) -> T;
fn as_ref(&self) -> Result<&T, &E>;
fn as_mut(&mut self) -> Result<&mut T, &mut E>;
fn map<U, F>(self, op: F) -> Result<U, E>;
fn map_err<F, 0>(self, op: 0) -> Result<T, F>;
// And so on
```

It's recommended for you to study the documentation and try to avoid match where possible.

Operator ?

Consider the following structure:

```
struct Info {
   name: String,
   age: i32,
}
```

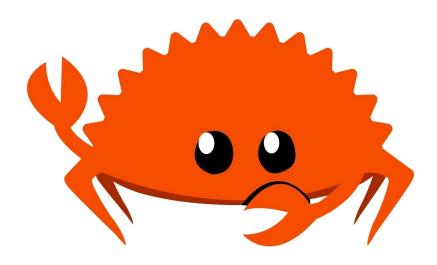
Operator ?

```
fn write_info(info: &Info) -> io::Result<()> {
    let mut file = match File::create("my_best_friends.txt") {
        Err(e) => return Err(e),
        0k(f) \Rightarrow f
    };
    if let Err(e) = file
        .write_all(format!("name: {}\n", info.name)
        .as_bytes()) {
        return Err(e)
    }
    if let Err(e) = file
        .write_all(format!("age: {}\n", info.age)
        .as_bytes()) {
        return Err(e)
    }
    0k(())
```

Operator ?

```
We can use the ? operator to make the code smaller!
fn write_info(info: &Info) -> io::Result<()> {
    let mut file = File::create("my_best_friends.txt")?;
    file.write_all(format!("name: {}\n", info.name).as_bytes())?;
    file.write_all(format!("age: {}\n", info.age).as_bytes())?;
    0k(())
Beautiful, isn't it?
We can use it for Option too!
```

Questions?



Box and Rc

We are already familiar with Box type. Let's check one advanced function:

```
fn leak<'a>(b: Box<T, A>) -> &'a mut T;
fn into_raw(b: Box<T, A>) -> *mut T;

Example:
let x = Box::new(41);
```

```
let x = Box::new(41);
let static_ref: &'static mut usize = Box::leak(x);
*static_ref += 1;
assert_eq!(*static_ref, 42);
```

Box

But stop! Rust is the safe language, no memory unsafety, no undefined behavior, what's wrong!?

Box

But stop! Rust is the safe language, no memory unsafety, no undefined behavior, what's wrong!?

In reality, when you're creating global objects or interacting with other languages, you have to leak objects. Moreover, it's safe to leak memory, just not good!

Rc is single-threaded reference-counting pointer. "Rc" stands for "Reference Counted".

```
let rc = Rc::new(());
let rc2 = rc.clone(); // Clones Rc, not what inside!
let rc3 = Rc::clone(&rc); // The same
```

Rc is dropped when all instances of Rc are dropped.

Primary functions:

```
fn get_mut(this: &mut Rc<T>) -> Option<&mut T>;
fn downgrade(this: &Rc<T>) -> Weak<T>;
fn weak_count(this: &Rc<T>) -> usize;
fn strong_count(this: &Rc<T>) -> usize;
```

References to the variable inside Rc are controlled at runtime:

```
let mut rc = Rc::new(42);
println!("{}", *rc);
*Rc::get_mut(&mut rc).unwrap() -= 41;
println!("{}", *rc);
let mut rc1 = rc.clone();
println!("{}", *rc1);
// thread 'main' panicked at 'called `Option::unwrap()`
// on a `None` value'
// *Rc::get_mut(&mut rc1).unwrap() -= 1;
```

get_mut guarantees that it will return mutable reference only if there's only one pointer. If there are more, you won't have a chance to modify Rc.

Weak

Rc is a **strong** pointer, while Weak is a **weak** pointer. Both of them have *ownership* over allocation, but only Rc have *ownership* over the value inside:

You can upgrade Weak to Rc:

```
fn upgrade(&self) -> Option<Rc<T>>;
```

```
let rc1 = Rc::new(String::from("string"));
let rc2 = rc1.clone():
let weak1 = Rc::downgrade(&rc1);
let weak2 = Rc::downgrade(&rc1);
drop(rc1); // The string is not deallocated
assert!(weak1.upgrade().is_some());
drop(weak1); // Nothing happens
drop(rc2); // The string is deallocated
assert_eq!(weak2.strong_count(), 0);
// If no strong pointers remain, this will return zero.
assert_eq!(weak2.weak_count(), 0);
assert!(weak2.upgrade().is_none());
drop(weak2); // The Rc is deallocated
```

Arc

There's also Arc - a thread-safe reference-counting pointer. Arc stands for "Atomically Reference Counted".

We'll need it to share data safely across threads in the future.