



Note by Roshan Bist
SNSC ,Mnr
IT Helloprogrammers-Google search
<https://www.helloprogrammers.com>

Unit-4

Recursion

⊗ Definition:- The process in which a function calls itself directly or indirectly is called recursion. It is a powerful technique of writing a complicated algorithm in an easy way. According to this technique a problem is defined in terms of itself. The problem is solved by dividing it into smaller problems, which are similar in nature to the original problem. These smaller problems are solved and their solutions are applied to get the final solution of our original problem.

A function will be recursive, if it contains following features;

i) Function should call itself.

ii) Function should have a stopping condition (base criteria).

Using recursive algorithm, certain problems can be solved quite easily. Examples of such problems are Towers of Hanoi (TOH), Inorder/Preorder/Postorder tree traversals, DFS of Graph etc.

⊗ Recursive vs. Iterative algorithm:-

| Property | Recursion | Iteration |
|---------------------|---|---|
| i) Definition | Function calls itself | A set of instructions repeatedly executed. |
| ii) Application | It is used for functions | It is used for loops. |
| iii) Termination | It will terminate at base case, where there will be no function call. | It will terminate when the condition for iteration is satisfied. |
| iv) Usage | It is used when code size needs to be small, and time complexity is not an issue. | It is used when time complexity needs to be balanced against an expanded code size. |
| v) Code Size | It has smaller code size | vi) It has larger code size. |
| vi) Time Complexity | It has very high (generally exponential) time complexity. | It has relatively lower time complexity (generally polynomial or logarithmic). |

⊗ Tail recursion (Definition & Example):-

A recursive function is tail recursive when recursive call is the last thing executed by the function. So, when nothing is left to do after coming back from the recursive call.

Example (Using C++):

```
#include <iostream>
using namespace std;
void printN(int n) {
    if (n < 0) {
        return;
    }
    cout << n << " ";
    printN(n-1);
}

int main() {
    printN(10);
}
```

Output:

10 9 8 7 6 5 4 3 2 1 0

The tail recursion is better than non-tail recursion. As there is no task left after the recursive call, it will be easier for the compiler to ~~organize~~ optimize the code.

A non-tail recursive function can be written as tail-recursive to optimize it. Consider the following function to calculate factorial of n.

```
long fact(int n) {
    if (n <= 1)
        return 1;
    n * fact(n-1);
}
```

This function is a non-tail recursive function because the value returned by fact(n-1) is used in fact(n), so the call to fact(n-1) is not the last thing done by fact n. The above function can be written as tail recursive function by adding some other parameters as follows:-

```
long fact(long n, long a) {
    if (n == 0)
        return a;
    return fact(n-1, a*n);
}
```


⊗ Examples of Recursive Algorithms:

➤ Factorial: Factorial of any number n is denoted as $n!$ and is equal to; $n! = n \times (n-1) \times (n-2) \times \dots \times 3 \times 2 \times 1$.

for eg; Factorial of 5 is;

$$5! = 5 \times 4 \times 3 \times 2 \times 1 \\ = 120$$

Recursive Algorithm for calculation of factorial of given integer;

Step 1: Start

Step 2: Read number n

Step 3: call factorial(n)

Step 4: Print factorial f

Step 5: Stop.

function to call

int factorial(n)

Step 1: if $n == 1$ then return 1.

Step 2: else,

$$f = n * \text{factorial}(n-1)$$

Step 3: return f ;

Now we can easily write C or C++ program on the basis of this algorithm.

2) Fibonacci Sequence: Fibonacci series are the numbers in the sequence 0, 1, 1, 2, 3, 5, 8, 13, 21, ...

By the definition, the first two numbers are 0 and 1. Each subsequent numbers in the series is equal to the sum of the previous two numbers.

Recursive Algorithm for Fibonacci Sequence:

Step 1: Start

Step 2: Read number n .

Step 3: Call fib(n).

Step 4: Print fib(n).

Step 5: Stop.

int fib(int n)

Step1: if $(n=1 || n==2)$
return 1;

Step2: else,
 $fib(n) = fib(n-1) + fib(n-2)$

Step3: return fib(n);

3) Greatest Common Divisor (GCD): [Imp]

GCD or Greatest Common Divisor of two or more integers is the largest positive integer that can divide both the number without leaving any remainder.

Example:- GCD of 20 and 8 is 4.

Recursive Algorithm for GCD:

Step1: Start

Step2: Read two numbers a and b.

Step3: Call gcd(a,b)

Step4: Print gcd

Step5: Stop.

int gcd(int a, int b)

Step1: if $(b==0)$
return a;

Step2: else,
 $gcd = gcd(b, a \% b);$

Step3: return gcd;

⊗ Tower of Hanoi:- [Imp]

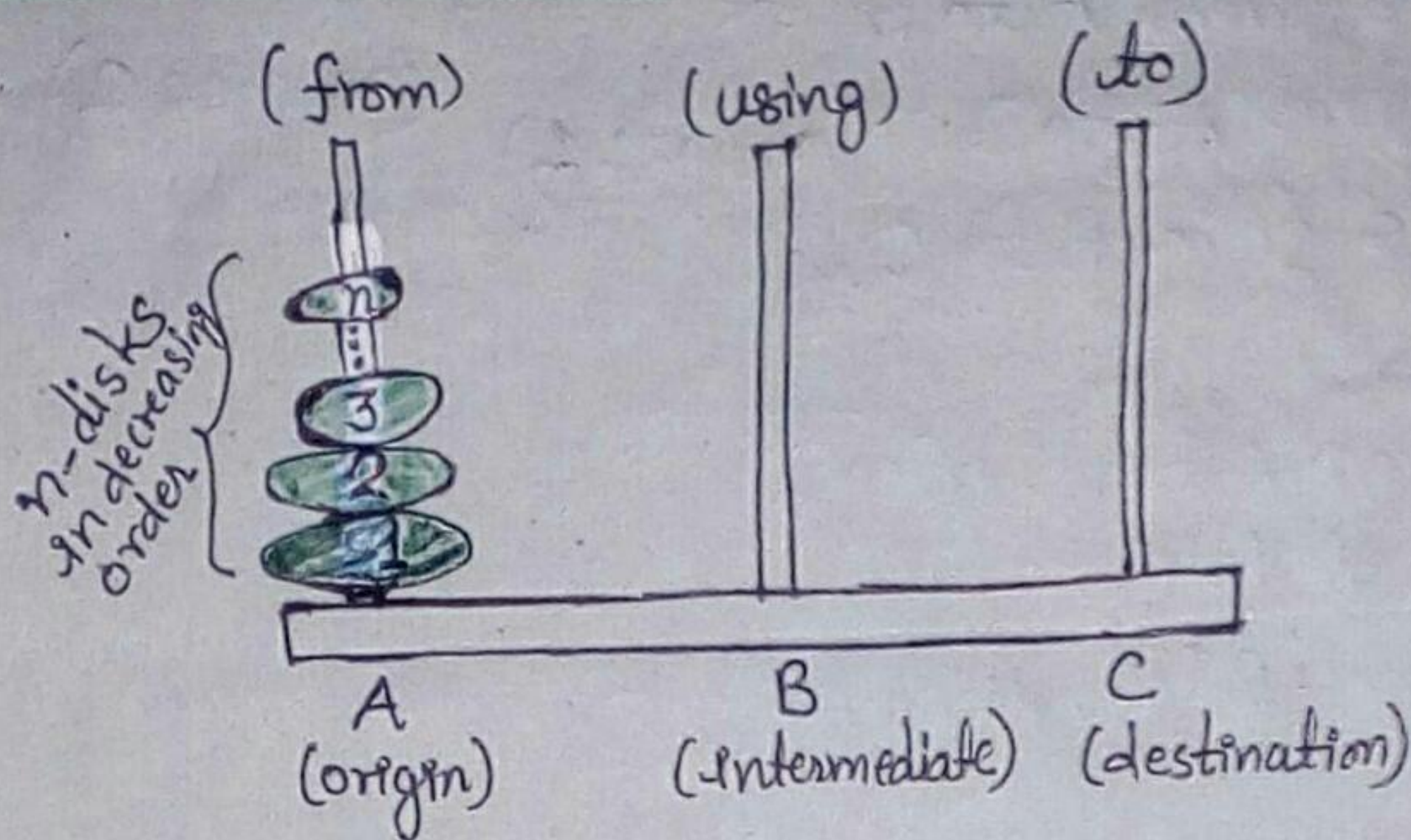


Fig. Representation of tower of Hanoi (Initial state)

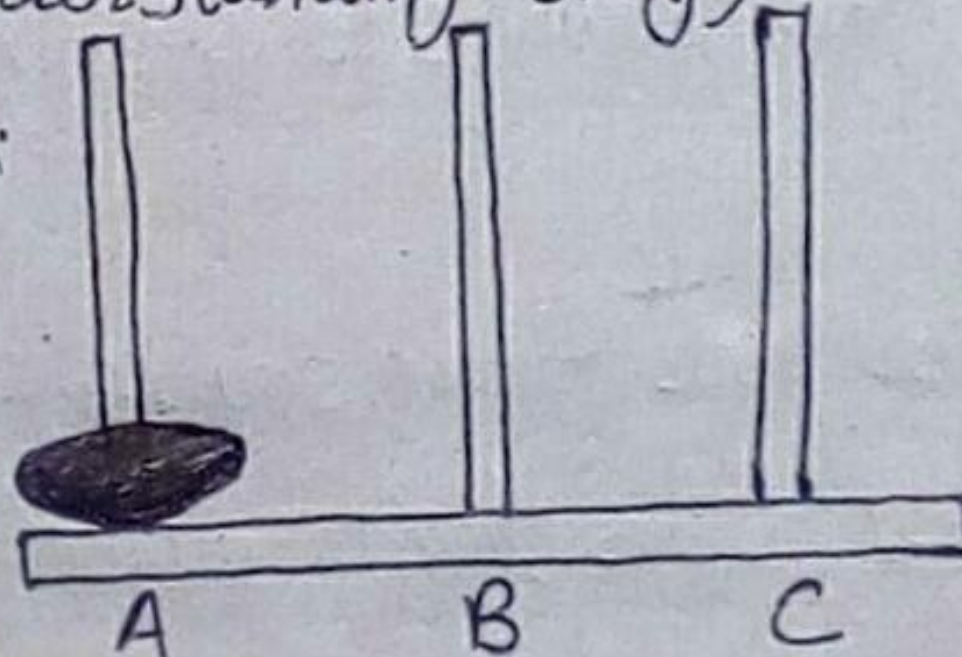
- There are three poles named as A (origin), B (intermediate) and C (destination).
- n number of different sized disks having hole at the center is stacked around the origin pole in decreasing order.
- The disks are numbered as $1, 2, 3, \dots, n$.

✓ Objective: The objective of tower of Hanoi is to transfer all disks from origin pole to destination pole using intermediate pole for temporary storage.

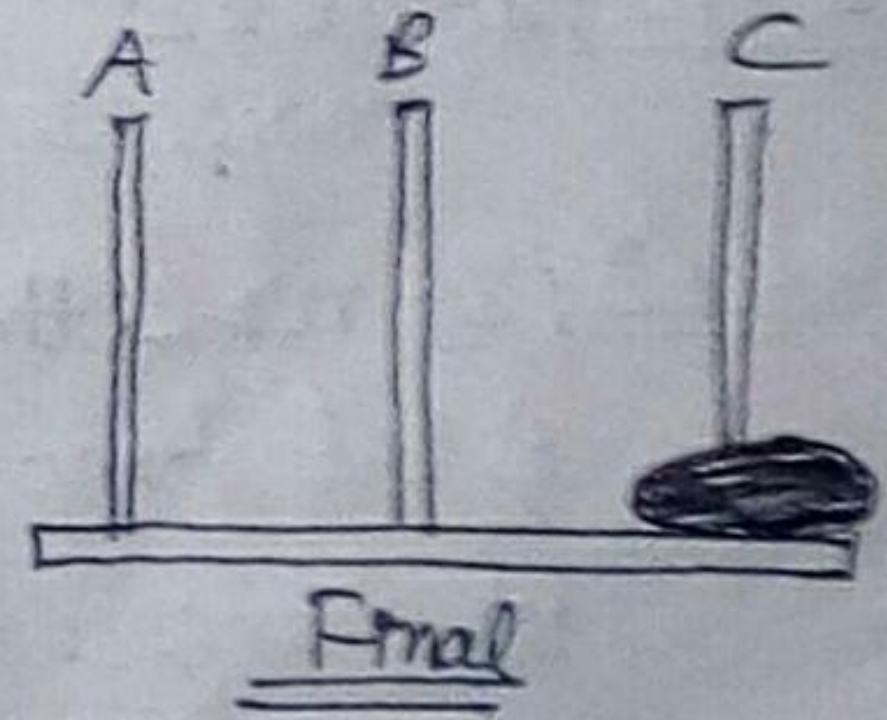
Not imp ← Cases: (for understanding only)

Case 1:

Initial



after moving



➤ If only one disk is given then move a disk from A to C directly.

Case 2: If two disks are given then;

→ Move a disk from A to B

→ Then move a disk from A to C

→ Finally move a disk from B to C.

(Same as single case)
Since single remaining after moving one

Case 3: if 3 disks are given;

- Move 2 disks from A to B ←
 - Move a disk from A to C (as case 1)
 - Move 2 disks from B to C. ←
- (as case 2)

∴ For n-disks

- Move $n-1$ disks from A to B using C.
- Move a disk from A to C
- Move $n-1$ disks from B to C using A

✓ Conditions:

- i) Move only one disk at a time.
- ii) Each disk must always be placed around one of the pole.
- iii) Never place larger disk on top of smaller disk.

✓ [Imp] Algorithm:

Algorithm to move a tower of n disks from source to destination. (where n is positive integer).

Step 1: If $n=1$;
move a single disk from source to destination.

Step 2: If $n > 1$;
i) Let temp be the other pole than source and destination.
ii) Move a tower of $(n-1)$ disks from source to temp.
iii) Move a single disk from source to destination.
iv) Move a tower of $(n-1)$ disks from temp to destination.

Step 3: Terminate.

⊗ Implementation of tower of Hanoi using C;

```
#include <stdio.h>
```

```
#include <conio.h>
```

```
void TOH (int, char, char, char); // function prototype.
```

```
void main() {
```

```
    int n;
```

```
    printf ("Enter number of disks");
```

```
    scanf ("%d", &n);
```

```
    TOH (n, 'A', 'B', 'C'); // function call
```

```
    getch();
```

```
}
```

```
void TOH (int n, char A, char B, char C) {
```

```
    if (n > 0) {
```

```
        TOH (n-1, A, C, B);
```

```
        printf ("Move disk %d from %c to %c \n",
```

```
                n, A, B);
```

```
        TOH (n-1, C, B, A);
```

```
    }
```

```
}
```

⊗ Applications of Recursion:

i) The most important data structure 'Tree' does not exist without recursion. We can solve that in iterative way also but that will be a very tough task.

ii) All puzzle games like chess, Candy Crush etc broadly uses recursion.

iii) Recursion is the backbone of AI.

iv) Many of the well known sorting algorithms (like Quick, Merge etc.) uses recursion.

v) It is the backbone for searching.

Since we are moving from A to B first in algorithm

source

intermediate
destination
temp.

⊗. Advantages of Recursion:

- i> The code becomes much easier to write.
- ii> It helps to solve some problems which are naturally recursive such as tower of Hanoi.
- iii> Reduces unnecessary calling of functions.

⊗. Disadvantages of Recursion:

- i> Recursive functions are generally slower than non-recursive functions.
- ii> It may require a lot of memory to hold intermediate results on the system stack.
- iii> ~~It is~~ It is difficult to think recursively so one must be very careful when writing recursive functions.