

# Intermediate Microeconomics Lecture 12

## Game Theory

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# Strategic Interactions

- ▶ In a competitive market, consumers and firms are assumed to be price takers.
- ▶ In a monopoly market, there is only one firm in the market.
- ▶ But many markets/social activities involve only a small number of participants and economic agents can interact strategically in a variety of ways
  - ▶ sports games
  - ▶ international relations
  - ▶ several firms compete in a certain market
    - ▶ Oligopolies
    - ▶ Cartels

## Strategic Interactions (cont.)

- ▶ The outcome to a participant depends not only on his own choice but also on what others do.
- ▶ When participants make decisions, they strategically take into account their opponents' behavior.
- ▶ Game theory provides a rigorous mathematical tool to study these situations.

# Elements of a Game

- ▶ Players (博弈参与人)
  - ▶ decision making entities
  - ▶ firms, governments, individuals
- ▶ Strategies(策略) (or choices or actions) for each player
  - ▶ output, prices, advertising budget
- ▶ Payoffs (得益) to players as a function of strategy profiles
  - ▶ depend not only on own strategies, but on other's strategies as well

# Examples of Simple Games

- ▶ Matching pennies (猜硬币博弈)
  - ▶ players: Player A, B
  - ▶ actions: Heads (正面) or Tails (反面)
  - ▶ payoffs:
    - ▶ player B gets \$1 and player A loses \$1 if pennies match
    - ▶ player B loses a dollar and player A gains \$1 if pennies don't match

# Examples of Simple Games (cont.)

- ▶ Prisoner's dilemma (囚徒困境)
  - ▶ players: Prisoner A, B
  - ▶ actions: deny or confess
  - ▶ payoffs:
    - ▶ 0 if you confess while other denies
    - ▶ -1 if you both deny
    - ▶ -3 if you both confess
    - ▶ -6 if you deny and the other guy confess

# Important Terminology

- ▶ Decision nodes
  - ▶ player make decisions at various points in a game
- ▶ Actions
  - ▶ the set of choices available each decision node in a game
- ▶ Pure Strategy (纯策略)
  - ▶ a rule that tells the player what actions to take at each of his information set in the game
- ▶ Mixed strategy (混合策略)
  - ▶ consists of a probability distribution on the set of pure strategies

# Important Terminology (cont.)

- ▶ Common knowledge(公共知识)
  - ▶ if every player knows it, every player knows that every other player knows it, every player knows that every other player knows that every other player knows it, and so on
  - ▶ we usually assume that the complete description of the game is common knowledge
- ▶ Simultaneous move (同时行动)
  - ▶ if players in a game have to make their decision at the same times
- ▶ Sequential move (序贯行动)
  - ▶ if players make their decisions in a particular sequence, one after another
- ▶ Perfect information
  - ▶ every player at every decision node knows the decision taken previously by every other player



# Classification of Games

- ▶ These are several ways of classifying games
  - ▶ by the number of players
    - ▶ 2-person v.s.  $n$ -person
  - ▶ by the number of strategies available to each of the players
    - ▶ finite games v.s. infinite games
  - ▶ by the nature of payoffs
    - ▶ zero-sum v.s. non-zero sum
  - ▶ the nature of preplay negotiation
    - ▶ cooperative games v.s. non-cooperative games
  - ▶ Interactions over time
    - ▶ dynamic game v.s. static games

# Classification of Games (cont.)

- ▶ These are several ways of classifying games
  - ▶ by the nature of the states
    - ▶ stochastic v.s. deterministic game
  - ▶ perfect information vs. imperfect information game
    - ▶ Perfect information (完美信息博弈): at each move in the game, the player with the move knows the full history of the play of the game thus far
  - ▶ complete information v.s. incomplete information game
    - ▶ Incomplete information (非完全信息) : at least one player is uncertain about another player's payoff function (Asymmetric or private information)

# Classification of Games (cont.)

	Complete	Incomplete
Static	Nash Equilibrium	Bayesian NE
Dynamic	Subgame Perfect NE	Perfect Bayesian NE

# Ways to Describe a Game

- ▶ There are two ways to describe a game
  - ▶ the “Normal” or matrix form (标准型或矩阵型)
    - ▶ there are only 2 (sometimes 3) players
    - ▶ there are a finite number of strategies
    - ▶ actions approximately simultaneous
  - ▶ the “Extensive” form (扩展型)
    - ▶ if actions are sequential

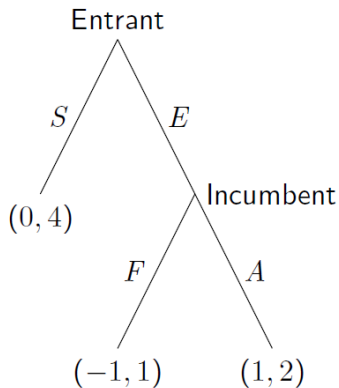
# Ways to Describe a Game (cont.)

## ► Matching Pennies Game

		Player $B$	
		Heads	Tails
Player $A$	Heads	$(-1,1)$	$(1,-1)$
	Tails	$(1,-1)$	$(-1,1)$

## Ways to Describe a Game (cont.)

- ▶ Chain store game



# Dominant Strategies

- ▶ We can describe the prisoners' dilemma in its normal form

		Prisoner <i>B</i>	
		Confess	Deny
Prisoner <i>A</i>	Confess	$(-3, -3)$	$(0, -6)$
	Deny	$(-6, 0)$	$(-1, -1)$

- ▶ If you were *A*, what would you choose?
  - ▶ If *B* chooses to confess, then you should confess because  $-3 > -6$
  - ▶ If *B* chooses to deny, then you should confess because  $0 > -1$
  - ▶ In sum, regardless of *B*'s choice, *A* should (will) choose to confess

## Dominant Strategies (cont.)

- ▶ The same logic applies to  $B$ 's decision: regardless of  $A$ 's choice, confession is better for  $B$
- ▶ Thus in this game, we know both players will choose to confess, yielding a 3-year sentence in prison for each player
- ▶ In games, if a strategy is always strictly optimal for a player regardless of other players' choices, we call this strategy a strictly dominant strategy (严格占优策略) for this player
  - ▶ strategy  $s_i$  (strictly) dominates strategy  $s'_i$  if, for all possible strategy combinations of opponents,  $s_i$  yields a (strictly) higher payoff than  $s'_i$  to player  $i$
- ▶ Confession is the strictly dominant strategy for both players in the prisoners' dilemma



# Dominant Strategies (cont.)

- ▶ If a player has a strictly dominant strategy, this player will always play this strategy.
- ▶ If all players have a strictly dominant strategy, then the game is dominance-solvable, because we know every player will simply play their strictly dominant strategy.
- ▶ Strictly dominant strategy is natural and intuitive.
- ▶ But unfortunately, in many games, players do not have a strictly dominant strategy.

# Nash Equilibrium

- ▶ Battle of the sexes:

		Wife	
		Soccer	Opera
Husband	Soccer	(2,1)	(0,0)
	Opera	(0,0)	(1,2)

- ▶ A couple wishes to go out for entertainment.
- ▶ The husband (row player) prefers soccer to opera while the wife (column player) prefers opera to soccer.
- ▶ They both prefer going together to going alone.

## Nash Equilibrium (cont.)

- ▶ A Nash equilibrium of a  $N$  player normal form game is a strategy profile  $(s_1, \dots, s_n)$  such that for each player  $i$ ,  $s_i$  is optimal given other players' strategy  $(s_1, \dots, s_{i-1}, s_{i+1}, \dots, s_n)$ 
  - ▶ no player can improve her payoff by unilaterally deviating
- ▶ Technically, a Nash equilibrium is a strategy profile  $(s_1^*, \dots, s_n^*)$  such that for every player  $i$ :

$$u_i(s_1^*, \dots, s_{i-1}^*, s_i^*, s_{i+1}^*, \dots, s_n^*) \geq u_i(s_1^*, \dots, s_{i-1}^*, s_i, s_{i+1}^*, \dots, s_n^*) \quad \forall s_i \in S_i$$

where  $u_i$  is the payoff function of player  $i$

- ▶ For two player game, this boils down to a pair of strategies  $(s_1, s_2)$  such that
  - ▶ given  $s_2$ ,  $s_1$  is optimal for player 1
  - ▶ given  $s_1$ ,  $s_2$  is optimal for player 2.

## Nash Equilibrium (cont.)

- ▶ In the battle of the sexes, there are four possible strategy profiles:
  - ▶  $(S, S)$ , Nash equilibrium
    - ▶ given that the husband goes to soccer, it is optimal for the wife to go to soccer
    - ▶ given that the wife goes to soccer, it is optimal for the husband to go to soccer
  - ▶  $(O, O)$ , Nash equilibrium
  - ▶  $(S, O)$ , not Nash equilibrium.
    - ▶ given that the husband goes to soccer, it is not optimal for the wife to go to opera
  - ▶  $(O, S)$ , not Nash equilibrium.
- ▶ The definition of Nash equilibrium does not say anything about which one should/will be played.

## Nash Equilibrium (cont.)

- ▶ Consider again the prisoners' dilemma

		Prisoner <i>B</i>	
		Confess	Deny
Prisoner <i>A</i>	Confess	$(-3,-3)$	$(0,-6)$
	Deny	$(-6,0)$	$(-1,-1)$

- ▶ (Confess, Confess) is the only Nash equilibrium of this game.

## Nash Equilibrium (cont.)

- ▶ The matching pennies game:

		Player $B$	
		Heads	Tails
Player $A$	Heads	$(-1, 1)$	$(1, -1)$
	Tails	$(1, -1)$	$(-1, 1)$

- ▶ There is no Nash equilibrium in this game

# Nash Equilibrium (cont.)

- ▶ Underlying assumptions of Nash equilibrium:
  - ▶ Players have predictions about other players' play and they optimally respond to their predictions.
  - ▶ Players' predictions are correct.
- ▶ Dominant Strategies
  - ▶ "I'm doing the best I can no matter what you do."
  - ▶ "You're doing the best you can no matter what I do."
- ▶ Nash Equilibrium
  - ▶ "I'm doing the best I can given what you are doing"
  - ▶ "You're doing the best you can given what I am doing."

## Nash Equilibrium (cont.)

- ▶ Iterated elimination of strictly dominated strategies:
  - ▶ eliminate all strategies which are dominated, relative to opponent's strategies which have not yet been eliminated
- ▶ If iterated elimination of strictly dominated strategies yields a unique strategy profile, then this strategy profile is the unique Nash equilibrium



## Nash Equilibrium (cont.)

- ▶ Example of Iterated Elimination:

		Player II	
		Left	Right
Player I	Top	1,2	4,1
	Middle	3,2	2,1
	Bottom	2,1	1,3

- ▶ Bottom is dominated by Middle (for Player I)
- ▶ Right is dominated by Left (for Player II)
- ▶ Top is dominated by Middle (for Player I)

# Mixed Strategies

- ▶ To deal with non-existence of Nash equilibrium, we extend players' strategy spaces.
- ▶ We assume players in the matching pennies games can randomize between  $H$  and  $T$ .
- ▶ We use  $p \circ H + (1 - p) \circ T$  to denote the strategy which plays  $H$  with probability  $p$  and  $T$  with probability  $1 - p$ , where  $p \in [0, 1]$
- ▶ This kind of strategies is called mixed strategies (混合策略)
  - ▶ If player  $i$  have  $K$  pure strategies available. Then a mixed strategy for player  $i$  is a probability distribution over those  $K$  strategies

## Mixed Strategies (cont.)

- ▶ The strategies we studied previously, e.g.  $H$  and  $T$  in the matching pennies game, are called pure strategies (纯策略).
  - ▶ just special cases of mixed strategies, e.g.  $H$  is equivalent to  $1 \circ H + 0 \circ T$
- ▶ There are many interpretations/justifications for mixed strategies.
  - ▶ Literally, randomization in players' brain.
  - ▶ Large population: fraction  $p$  of people play  $H$  and fraction  $1 - p$  play  $T$

## Mixed Strategies (cont.)

- ▶ Use expected payoffs to evaluate payoffs from mixed strategies
- ▶ Player 1 adopts a mixed strategy  $\sigma_1$ 
  - ▶ plays  $s_1$  with probability  $p$  and  $s'_1$  with probability  $1 - p$
- ▶ Player 2 adopts a mixed strategy  $\sigma_2$ 
  - ▶ plays  $s_2$  with probability  $q$  and  $s'_2$  with probability  $1 - q$
- ▶ Then:
  - ▶  $(s_1, s_2)$  is played with probability  $p \times q$
  - ▶  $(s_1, s'_2)$  is played with probability  $p \times (1 - q)$
  - ▶  $(s'_1, s_2)$  is played with probability  $(1 - p) \times q$
  - ▶  $(s'_1, s'_2)$  is played with probability  $(1 - p) \times (1 - q)$

## Mixed Strategies (cont.)

- ▶ Hence the expected utilities are given by

$$\begin{aligned}u_1(\sigma_1, \sigma_2) &= pq u_1(s_1, s_2) + p(1 - q) u_1(s_1, s'_2) \\ &\quad + (1 - p) q u_1(s'_1, s_2) + (1 - p)(1 - q) u_1(s'_1, s'_2)\end{aligned}$$

$$\begin{aligned}u_2(\sigma_1, \sigma_2) &= pq u_2(s_1, s_2) + p(1 - q) u_2(s_1, s'_2) \\ &\quad + (1 - p) q u_2(s'_1, s_2) + (1 - p)(1 - q) u_2(s'_1, s'_2)\end{aligned}$$

- ▶ For example, in matching pennies game, assume  $\sigma_1 = \frac{1}{2} \circ H + \frac{1}{2} \circ T$  and  $\sigma_2 = H$ . Then

$$\begin{aligned}u_1(\sigma_1, \sigma_2) &= \frac{1}{2} u_1(H, H) + \frac{1}{2} u_1(T, H) \\ &= \frac{1}{2} \times (-1) + \frac{1}{2} \times 1 = 0\end{aligned}$$

## Mixed Strategies (cont.)

- ▶ A Nash equilibrium in mixed strategies is similarly defined.
- ▶  $(\sigma_1, \sigma_2)$  is a Nash equilibrium if
  - ▶ given  $\sigma_2$ ,  $\sigma_1$  is optimal for player 1
  - ▶ given  $\sigma_1$ ,  $\sigma_2$  is optimal for player 2.
- ▶ In general, it is more difficult to find equilibria in mixed strategies than in pure strategies.
- ▶ For  $2 \times 2$  games (2 players, 2 pure strategies for each) the best response correspondence (最优反应曲线) is a powerful tool to help us analyze

## Mixed Strategies (cont.)

- ▶ Let's consider the matching pennies game again.
- ▶ Assume player 1 plays  $p \circ H + (1 - p) \circ T$  and player 2 plays  $q \circ H + (1 - q) \circ T$
- ▶ Then the payoff to player 2 is:

$$\begin{aligned}u_2(p, q) &= pq + p(1 - q)(-1) + (1 - p)q(-1) + (1 - p)(1 - q) \\&= 2(2p - 1)q + (1 - 2p)\end{aligned}$$

- ▶ if  $p < 1/2$ , player 2's optimal choice is  $q^* = 0$  (i.e.  $T$ )
- ▶ if  $p > 1/2$ , player 1's optimal choice is  $q^* = 1$  (i.e.  $H$ )
- ▶ if  $p = 1/2$ , any choice of  $q^* \in [0, 1]$  is optimal

## Mixed Strategies (cont.)

- ▶ Player 2's optimal choice can be expressed as

$$BR_2(p) = \begin{cases} \{0\} & \text{if } p < 1/2 \\ [0, 1] & \text{if } p = 1/2 \\ \{1\} & \text{if } p > 1/2 \end{cases}$$

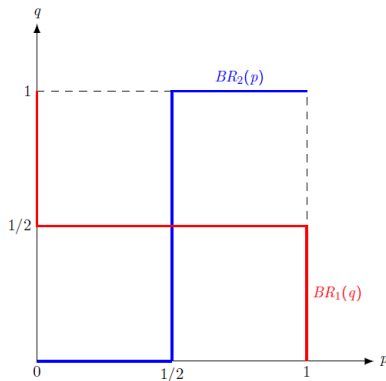
- ▶ Similarly, we can calculate player 1's optimal choice

$$BR_1(q) = \begin{cases} \{1\} & \text{if } q < 1/2 \\ [0, 1] & \text{if } q = 1/2 \\ \{0\} & \text{if } q > 1/2 \end{cases}$$



## Mixed Strategies (cont.)

- We can plot it in the  $p - q$  plane



- The intersections of the two best response correspondences are Nash equilibria

## Mixed Strategies (cont.)

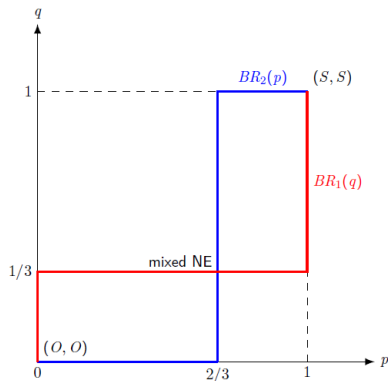
- ▶ Battle of the sexes:

		Wife	
		Soccer	Opera
Husband	Soccer	(2,1)	(0,0)
	Opera	(0,0)	(1,2)

- ▶ Any mixed Nash equilibrium?

## Mixed Strategies (cont.)

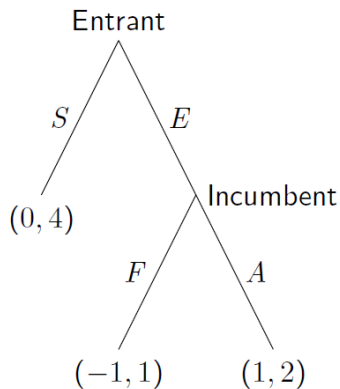
- ▶ One mixed strategy Nash equilibrium: player 1 plays  $\frac{2}{3} \circ S + \frac{1}{3} \circ O$  and player 2 plays  $\frac{1}{3} \circ S + \frac{2}{3} \circ O$



# Games with Sequential Moves

- ▶ Consider the following chain store game.
- ▶ There is an entrant who chooses to either enter a market ( $E$ ) or stay out ( $S$ )
- ▶ There is an incumbent who chooses to either fight ( $F$ ) or accommodate ( $A$ ) if the entrant enters
- ▶ Coca-Cola is an entrant while Pepsi is the incumbent in the Russian market around 1990.

## Games with Sequential Moves (cont.)



## Games with Sequential Moves (cont.)

- ▶ There are four possible strategy profiles. Which profile(s) do you think plausible?
- ▶  $(E, F)$  is not
  - ▶ if the entrant enters, then it is optimal for the incumbent not to fight
- ▶  $(S, A)$  is not
  - ▶ if the incumbent chooses to accommodate after entry, the entrant should enter instead of staying out

## Games with Sequential Moves (cont.)

- ▶  $(S, F)$  seems plausible:
  - ▶ given that the incumbent will fight after entry, it is optimal for the entrant to stay out
  - ▶ given that the entrant stays out, the choice of the incumbent is irrelevant (thus trivially optimal)
- ▶  $(E, A)$  seems plausible too:
  - ▶ given that the incumbent will accommodate, it is optimal for the entrant to enter
  - ▶ given that the entrant enters, it is optimal for the incumbent to accommodate
- ▶ Therefore, these two profiles are Nash equilibria

## Games with Sequential Moves (cont.)

- ▶ But in  $(S, F)$ , the incumbent supposedly plans to fight in case of the entrant entering.
- ▶ It is like a threat: if you enter, I will fight with you!
- ▶ But is this threat credible?
- ▶ Probably not.
- ▶ In contrast,  $(E, A)$  does not suffer from this problem.
  - ▶ it is indeed the incumbent's best response to accommodate after entry.
- ▶ But to rule out  $(S, F)$ , we need a stronger solution concept that incorporates the fact that actions are taken in sequence



# Subgame Perfect Equilibrium

- ▶ A Nash equilibrium is said to be a subgame perfect Nash equilibrium (子博弈精炼纳什均衡) if and only if it is a Nash equilibrium in every subgame of the game
- ▶ A subgame(子博弈) in an extensive-form game has the following properties:
  - ▶ it begins at a node of the tree corresponding to an information set reduced to a singleton (the set contains only one set)(单结)
  - ▶ it encompasses all parts of the tree following the starting node
  - ▶ it never divides an information set

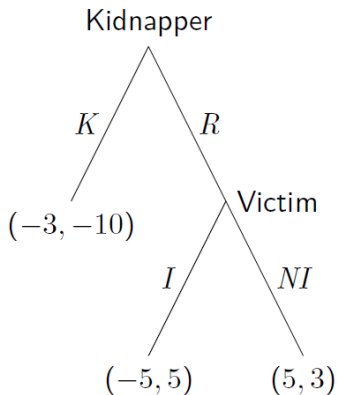
# Backward Induction

- ▶ To get an SPNE for finite horizon games, we could work back from terminal nodes.
  - ▶ go to final “decision node” and assign action to the player that maximizes his payoff
  - ▶ reduce game by trimming tree at this node and making terminal payoffs at this node, the payoffs when the player whose turn it was takes best action
  - ▶ keep working backwards
- ▶ This procedure is called backward induction solution (倒推归纳法)
- ▶  $(E, A)$  is the only backward induction solution to the entrant & incumbent game.

## Another Example: Kidnapping

- ▶ The kidnapper has two choices after receiving the ransom:
  - ▶ release ( $R$ )
  - ▶ or kill the victim ( $K$ ).
- ▶ After release, the victim has two choices
  - ▶ identify the kidnapper ( $I$ )
  - ▶ or refrain ( $NI$ )

## Another Example: Kidnapping (cont.)



## Another Example: Kidnapping (cont.)

- ▶ If the victim had been released, he would choose to identify the kidnapper
- ▶ Thus, the kidnapper would get -5 if he releases the victim
- ▶ It is then optimal for the kidnapper to kill the victim in the first place
- ▶ This is the only backward induction solution of this game
- ▶ The outcome from  $(R, NI)$  is socially optimal, but it is not consistent with backward induction
  - ▶ the victim can not commit to playing  $NI$