## Homework 1

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1. Prove or disprove that  $\log(x^2+1)$  and  $\log(x)$  are of the same order  $(\Theta)$ .

*Proof.* Assume  $\log(x^2+1)$  and  $\log(x)$  are of the same order. Then  $\log(x^2+1) \in O(\log(x))$  and  $\log(x) \in O(\log(x^2+1))$ .

To show  $\log(x^2+1)\in O(\log(x))$ , we can start by rewriting  $\log(x^2+1)=\log(x^2(1+\frac{1}{x^2}))$ . And since  $\frac{1}{x^2}\leq 1$  for  $x\geq 1$ , we know that  $1+\frac{1}{x^2}\leq 2$ . So we have that  $\log(x^2+1)\leq \log(2x^2)=\log(x^2)+\log(2)=2\log(x)+\log(2)$ . And for  $x\geq 2$ ,  $\log(2)\leq \log(x)$ , so  $\log(x^2+1)\leq 2\log(x)+\log(2)\leq 2\log(x)+\log(x)=3\log(x)$  in this domain. So for  $x\geq 2$  and c=3 we have that  $\log(x^2+1)\in O(\log(x))$ .

To show  $\log(x) \in O(\log(x^2+1))$ , we simply follow the idea that  $\log(n) \leq \log(m)$  when  $n \leq m$  since log is always increasing. Considering x and  $x^2+1$ , we know that  $x \leq x^2+1$  for  $x \in \mathbb{R}$ . So,  $\log(x) \leq \log(x^2+1)$  for all x. Taking c=1 leaves  $\log(x) \in O(\log(x^2+1))$ .

- $\begin{aligned} \text{2.} &\quad \text{a. } \lg^k(n) \in O(n^\epsilon), \, \lg^k(n) \not\in \Omega(n^\epsilon), \, \lg^k(n) \not\in \Theta(n^\epsilon) \\ &\quad \text{b. } n^k \in O(c^n), \, n^k \not\in \Omega(c^n), \, n^k \not\in \Theta(c^n) \\ &\quad \text{f. } \lg(n!) \in O(\lg(n^n)), \, \lg(n!) \not\in \Omega(\lg(n^n)), \, \lg(n!) \not\in \Theta(\lg(n^n)) \end{aligned}$
- 3. (n+1)!: This is equivalent to (n+1)n! n!: Factorials grow faster than exponents  $2^{2^{n+1}}$ : This is equivalent to  $2^{2\cdot 2^n}=4^{2^n}$   $2^{2^n}$   $n\cdot 2^n$ : Multiply  $2^n$  by an increasing amount  $2^n$   $(3/2)^n$ :  $2^n=(4/3)^n\cdot (3/2)^n$   $n^3$   $n^2$   $n\lg(n)$   $\lg(n!)$   $(\lg(n))^{\lg(n)}, \, n^{\lg(\lg(n))}$ : The two logarithms causes these to grow much slower n  $(\lg(n))!$   $4^{\lg(n)}$   $2^{\lg(n)}$   $1^{g^2}(n)$   $n^{1/\lg(n)}$

4. Closed form solution for Tower of Hanoi problem with forwards substitution.

$$H_{i} = \begin{cases} 2H_{i-1} + 1 & i > 1\\ 1 & i = 1 \end{cases}$$

$$H_{2} = 2 * H_{1} + 1 = 2 * 1 + 1 = 2^{1} + 2^{0}$$

$$H_{3} = 2 * H_{2} + 1 = 2(2^{1} + 2^{0}) + 2^{0} = 2^{2} + 2^{1} + 2^{0}$$

$$H_{4} = 2 * H_{3} + 1 = 2(2^{2} + 2^{1} + 2^{0}) + 2^{0} = 2^{3} + 2^{2} + 2^{1} + 2^{0}$$

$$H_{n} = \sum_{i=0}^{n-1} 2^{i} = \frac{1 * 2^{n-1+1} - 1}{2 - 1} = \frac{2^{n} - 1}{1}$$

$$= 2^{n} - 1$$

5. Prove for a positive integer n that  $f_0f_1 + f_1f_2 + \ldots + f_{2n-1}f_{2n} = (f_{2n})^2$ 

Base case: 
$$n = 1$$
  
 $f_0 f_1 + f_1 f_2 = 0 * 1 + 1 * 1 = 1 = 1^2 = (f_2)^2$ 

Inductive Step: assume that 
$$f_0 f_1 + f_1 f_2 + \ldots + f_{2n-1} f_{2n} = (f_{2n})^2$$
  

$$f_0 f_1 + \ldots + f_{2n-1} f_{2n} + f_{2n} f_{2n+1} + f_{2n+1} f_{2n+2} = (f_0 f_1 + \ldots + f_{2n-1} f_{2n}) + f_{2n} f_{2n+1} + f_{2n+1} f_{2n+2}$$

$$= (f_{2n})^2 + f_{2n} f_{2n+1} + f_{2n+1} f_{2n+2}$$

$$= f_{2n} (f_{2n} + f_{2n+1}) + f_{2n+1} f_{2n+2}$$

$$= f_{2n} f_{2n+2} + f_{2n+1} f_{2n+2}$$

$$= f_{2n+2} f_{2n+2}$$

$$= (f_{2n+2})^2$$