

CSE 40622 Cryptography
Writing Assignment 05 (Lecture 09-10)

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1. (20 pts, Page 4) In \mathbb{Z}_p^* with an odd prime p , prove that a QR has exactly two square roots. That is, for any $x \in \mathbb{Z}_p^*$ that is a QR, there exists exactly two distinct y 's such that $y^2 = x$.
 - First show that a QR has at least two square roots, and then show that a QR has at most two square roots.
 - Hint: recall that \mathbb{Z}_p^* has exactly $\frac{p-1}{2}$ QRs.

Answer:

Proof. Let x be a QR in \mathbb{Z}_p^* , then there exists some $y \in \mathbb{Z}_p^*$ such that $y^2 = x$. Then we also know that $(-y)^2 = x$, so x has at least two square roots. Consider two QRs x_1, x_2 that share a square root y . Then $y^2 = x_1 = x_2$, so $x_1 = x_2$. Thus, each QR has distinct square roots. Since we know that there are exactly $\frac{p-1}{2}$ QRs and $p-1$ elements in \mathbb{Z}_p^* and each QR has at least 2 distinct roots that are not shared by any other QR, then each QR must have exactly 2 square roots. \square

2. (20 pts, Page 6) Prove that, given \mathbb{Z}_p^* with $p = 2q + 1$ where p, q are both odd prime numbers, if a random number $g \in \mathbb{Z}_p^*$ satisfies $g \neq 1$ and $g^2 \bmod p \neq 1$, $\langle g \rangle$ must be a subgroup of \mathbb{Z}_p^* which contains ALL of QRs in \mathbb{Z}_p^* .
 - You need to prove all QRs of \mathbb{Z}_p^* belong to $\langle g \rangle$. You do NOT need to prove $\langle g \rangle$ contains QRs only.

Answer:

Proof. Since the order of any subgroup must divide the order of the group, we know that $\text{ord}(g) \mid \text{ord}(\mathbb{Z}_p^*)$. Since it is given that the order of g is neither 1 nor 2 and $|\mathbb{Z}_p^*| = 2q$ where q is an odd prime, then $|\langle g \rangle|$ must be equal to either q or $2q$. First, if $|\langle g \rangle| = 2q$, then $\langle g \rangle$ contains all elements of \mathbb{Z}_p^* and thus all QRs. Second, if $|\langle g \rangle| = q$, then we know $g^q = 1$. Since 1 is a QR, we know that $g^q = y^2$ for some $y \in \mathbb{Z}_p^*$. Then $g^{\frac{q}{2}} = y$ but since q is odd, $\frac{q}{2}$ is not an integer. So for this to hold, g must be a QR, meaning $g = z^2$ for some $z \in \mathbb{Z}_p^*$. So, $g^q = (z^2)^q = 1$. Then since the subgroup generated by g must contain elements expressed in powers of $g = z^2$, all elements of $\langle g \rangle$ are QRs. \square

3. (10 pts, Coming from nowhere) Prove that any generator g of \mathbb{Z}_p^* is a QNR if p is an odd prime number, WITHOUT using Corollary 1 in the note of Lecture 09-10.

Answer:

Proof. If g generates \mathbb{Z}_p^* , then $|\langle g \rangle| = p-1$. Suppose, towards a contradiction, that g is a QR so there exists some z where $z^2 = g$. Then $g^{p-1} = (z^2)^{p-1} = 1$ and every element generated by g is of the form $z^{2k} = (z^k)^2$ meaning it is a QR. But we know that g generates all of \mathbb{Z}_p^* , so it cannot generate only QRs. Here lies our contradiction, so g must be a QNR. \square

4. (10 pts, Page 6) What is the consequence if we have $p = kq + 1$ with a large positive even number k instead of 2? In other words, what do you need to do **additionally** in order to find a generator of \mathbb{Z}_p^* ?

Answer:

If k is some large positive even number instead of just 2, we have more to check when looking for a generator. Instead of just checking whether $g^{\frac{p-1}{2}} \neq 1$ and $g^{\frac{p-1}{q}} \neq 1$, we must also check this for all other prime factors of kq . That is, for every prime factor m of kq , we must ensure that $g^{\frac{p-1}{m}} \neq 1$ before we can say that g is a generator.

5. (15 pts, Coming from nowhere) In class, I said one of the countermeasures to QR/QNR attacks is to use the subset of \mathbb{Z}_p^* which contains all of its QRs and use it as \mathbb{G} in ElGamal encryption. Then, Legendre symbols do not give much information to ciphertexts and public keys since all parameters will be QRs. Why is it not possible to use a subset of \mathbb{Z}_p^* which contains all of its QNRs and use it as \mathbb{G} ?

Answer:

The subset of \mathbb{Z}_p^* containing only QNRs is not a subgroup as it is not closed under the group operator. Consider some y QNR so that y belongs to this subset. If the QNRs formed a subgroup, then y^2 would also be in this subset, but it is obvious that y^2 is a QR and thus is not contained in the QNR subset. Since the QNR subset does not form a group, it is impossible to use it as a group in ElGamal encryption.