

torch-sla: Differentiable Sparse Linear Algebra with Adjoint Solvers and Sparse Tensor Parallelism for PyTorch

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<https://github.com/walkerchi/torch-sla>

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Abstract

Industrial scientific computing predominantly uses sparse matrices to represent unstructured data—finite element meshes, graphs, point clouds. We present `torch-sla`, an open-source PyTorch library that enables GPU-accelerated, scalable, and differentiable sparse linear algebra. The library addresses three fundamental challenges: (1) **GPU acceleration** for sparse linear solves, nonlinear solves (Newton, Picard, Anderson), and eigenvalue computation; (2) **Multi-GPU scaling** via domain decomposition with halo exchange, reaching **400 million DOF linear solve on 3 GPUs**; and (3) **Adjoint-based differentiation** achieving $\mathcal{O}(1)$ graph nodes and $\mathcal{O}(\text{nnz})$ memory—independent of solver iterations. `torch-sla` supports multiple backends (SciPy, cuDSS, PyTorch-native) and seamlessly integrates with PyTorch autograd for end-to-end differentiable simulations. Code is available at <https://github.com/walkerchi/torch-sla>.

1 Introduction

Industrial scientific computing predominantly deals with *unstructured data*—finite element meshes, graph networks, point clouds—where sparse matrices are the natural representation. These sparse systems arise in finite element analysis [Hughes, 2012], computational fluid dynamics [Jasak et al., 2007, Kochkov et al., 2021], graph neural networks [Kipf and Welling, 2017, Veličković et al., 2018], and physics-informed machine learning [Raissi et al., 2019, Lu et al., 2021]. Efficiently solving sparse linear systems $\mathbf{Ax} = \mathbf{b}$, nonlinear systems $\mathbf{F}(\mathbf{x}) = \mathbf{0}$, and eigenvalue problems $\mathbf{A}\mathbf{v} = \lambda\mathbf{v}$ is fundamental to these applications.

With the rise of differentiable programming and neural operators [Li et al., 2020, Brandstetter et al., 2022], there is growing demand to integrate sparse solvers with deep learning frameworks like PyTorch [Paszke et al., 2019]. However, three fundamental challenges stand in the way:

Challenge 1: GPU acceleration for sparse operations. Unlike dense matrices where GPU parallelism is straightforward, sparse matrices have irregular memory access patterns that underutilize GPU bandwidth. Sparse matrix-vector products (SpMV) achieve only 10–20% of peak memory bandwidth [Li et al., 2024]. Direct solvers (LU, Cholesky) require $\mathcal{O}(n^{1.5})$ memory due to fill-in, often exceeding GPU memory for problems with $>2\text{M}$ unknowns.

Challenge 2: Multi-GPU scaling. Industrial problems (10M–1B unknowns) exceed single-GPU memory. Distributed sparse computing requires domain decomposition and *halo exchange*—communicating boundary values between partitions—which is non-trivial to implement correctly with automatic differentiation support.

Challenge 3: PyTorch integration with efficient gradients. Differentiating through iterative solvers creates $\mathcal{O}(k)$ nodes in the computational graph, where k may be thousands of iterations. For a 1M-DOF problem with 1000 CG iterations, this requires ~ 80 GB just for gradient computation. This memory explosion prevents end-to-end training of neural networks that interact with sparse solvers.

This paper introduces `torch-sla*`, a library that addresses all three challenges:

- **GPU-accelerated sparse solvers:** We provide unified APIs for linear solves, nonlinear solves (Newton, Picard, Anderson acceleration), and eigenvalue computation (LOBPCG), with multiple backends (SciPy, cuDSS, PyTorch-native).
- **Sparse tensor parallelism (§3.3):** We implement domain decomposition with automatic halo exchange, enabling **400 million DOF linear solve on 3 GPUs**.
- **Adjoint-based differentiation (§3.2):** We implement the adjoint method for linear, nonlinear, and eigenvalue problems, achieving $\mathcal{O}(1)$ graph nodes and $\mathcal{O}(\text{nnz})$ memory—independent of solver iterations.

2 Related Work

Differentiable linear algebra. JAX [Bradbury et al., 2018] provides differentiable sparse CG, but naive differentiation through iterations creates $\mathcal{O}(k)$ graph nodes. Blondel et al. [2022] formalize implicit differentiation for optimization layers. OptNet [Amos and Kolter, 2017] and `cvxpylayers` [Agrawal et al., 2019] handle convex optimization. Deep equilibrium models [Bai et al., 2019] and Jacobian-free backpropagation [Fung et al., 2022] address implicit layers but focus on fixed-point iterations rather than sparse linear systems.

Sparse solvers and GPU acceleration. SciPy [Virtanen et al., 2020] provides SuperLU/UMFPACK for CPU. For GPU, NVIDIA cuDSS [NVIDIA Corporation, 2024] offers direct solvers, while AmgX [Naumov et al., 2015] provides algebraic multigrid with excellent scalability. Recent surveys [Li et al., 2024] cover sparse matrix computations on modern GPUs. None of these natively support PyTorch autograd.

Distributed sparse computing. PETSc [Balay et al., 2023], Trilinos [Heroux et al., 2005], and hypre [Falgout and Yang, 2002] implement distributed sparse linear algebra with sophisticated preconditioners. OpenFOAM [Jasak et al., 2007] uses domain decomposition for CFD. We bring these industrial patterns to PyTorch with automatic differentiation.

Learned solvers and preconditioners. Recent work explores learning components of iterative solvers: learned multigrid prolongation [Greenfeld et al., 2019], GNN-based AMG [Luz et al., 2020], and learned interface conditions for domain decomposition [Taghibakhshi et al., 2024]. These complement our library by providing learnable components that can be trained end-to-end.

3 Methodology

We first introduce the background on sparse linear systems (§3.1), then present our adjoint-based differentiation approach (§3.2), and finally describe our distributed computing implementation (§3.3).

3.1 Preliminaries: Sparse Linear Systems

A sparse linear system $\mathbf{Ax} = \mathbf{b}$ has a matrix $\mathbf{A} \in \mathbb{R}^{n \times n}$ with $\text{nnz} \ll n^2$ non-zero entries. We store matrices in coordinate (COO) format: three arrays for row indices, column indices, and values.

Direct solvers (LU, Cholesky factorization) compute exact solutions but require $\mathcal{O}(n^{1.5})$ memory for 2D problems due to fill-in—new non-zeros created during factorization [George, 1973]. For a 2D Poisson problem with n unknowns, the factored matrix has $\mathcal{O}(n \log n)$ non-zeros instead of the original $\mathcal{O}(n)$.

Iterative solvers (Conjugate Gradient, BiCGStab, GMRES) maintain $\mathcal{O}(\text{nnz})$ memory by never forming the factorization. However, they require $\mathcal{O}(\sqrt{\kappa})$ iterations for CG where κ is the condition number [Saad, 2003]. Each iteration performs one sparse matrix-vector multiplication (SpMV) costing $\mathcal{O}(\text{nnz})$.

The differentiation problem. Given loss $\mathcal{L}(\mathbf{x})$ where $\mathbf{x} = \mathbf{A}^{-1}\mathbf{b}$, we need gradients with respect to both \mathbf{b} and the non-zero values of \mathbf{A} . If we differentiate through k iterations of CG, the computational graph has $\mathcal{O}(k)$ nodes, each storing intermediate vectors. For $k = 1000$ iterations and $n = 10^6$, this requires ~ 80 GB just for the graph. Our adjoint approach reduces this to $\mathcal{O}(1)$ nodes.

*Code available at <https://github.com/walkerchi/torch-sla>. Install via `pip install torch-sla`.

Table 1: Complexity comparison: naive vs. adjoint differentiation through iterative solver with k iterations, n unknowns, and nnz non-zeros.

	Naive (through iterations)	Adjoint (ours)
Computational graph nodes	$\mathcal{O}(k)$	$\mathcal{O}(1)$
Memory for graph	$\mathcal{O}(k \cdot n)$	$\mathcal{O}(n + \text{nnz})$
Backward pass time	$\mathcal{O}(k \cdot \text{nnz})$	$\mathcal{O}(T_{\text{solve}} + \text{nnz})$

3.2 Adjoint-Based Differentiation

The adjoint method computes gradients through implicit functions without storing intermediate solver states. We present it for linear systems, eigenvalue problems, and nonlinear systems.

3.2.1 Linear Systems

Consider $\mathbf{x} = \mathbf{A}^{-1}\mathbf{b}$. Rather than differentiating through the solver iterations, we use the *implicit function theorem*. The solution satisfies $\mathbf{A}\mathbf{x} - \mathbf{b} = \mathbf{0}$. Differentiating implicitly:

$$d\mathbf{A} \cdot \mathbf{x} + \mathbf{A} \cdot d\mathbf{x} = d\mathbf{b} \implies d\mathbf{x} = \mathbf{A}^{-1}(d\mathbf{b} - d\mathbf{A} \cdot \mathbf{x}) \quad (1)$$

For a scalar loss $\mathcal{L}(\mathbf{x})$, the chain rule gives:

$$d\mathcal{L} = \left(\frac{\partial \mathcal{L}}{\partial \mathbf{x}} \right)^\top d\mathbf{x} = \left(\frac{\partial \mathcal{L}}{\partial \mathbf{x}} \right)^\top \mathbf{A}^{-1}(d\mathbf{b} - d\mathbf{A} \cdot \mathbf{x}) \quad (2)$$

Define the **adjoint variable** $\boldsymbol{\lambda} = \mathbf{A}^{-\top} \frac{\partial \mathcal{L}}{\partial \mathbf{x}}$, solved via $\mathbf{A}^\top \boldsymbol{\lambda} = \frac{\partial \mathcal{L}}{\partial \mathbf{x}}$. Then the gradients are:

$$\boxed{\frac{\partial \mathcal{L}}{\partial \mathbf{b}} = \boldsymbol{\lambda}, \quad \frac{\partial \mathcal{L}}{\partial \mathbf{A}_{ij}} = -\boldsymbol{\lambda}_i \cdot \mathbf{x}_j} \quad (3)$$

Complexity analysis. Table 1 compares naive differentiation (through iterations) with our adjoint approach.

The forward solve computes \mathbf{x} in time T_{solve} . The backward pass:

1. Solves $\mathbf{A}^\top \boldsymbol{\lambda} = \frac{\partial \mathcal{L}}{\partial \mathbf{x}}$: $\mathcal{O}(T_{\text{solve}})$ time
2. Computes $\frac{\partial \mathcal{L}}{\partial \mathbf{A}_{ij}} = -\boldsymbol{\lambda}_i \cdot \mathbf{x}_j$ for each non-zero: $\mathcal{O}(\text{nnz})$ time

Algorithm. The complete procedure is shown below:

Algorithm 1: Adjoint Linear Solve

Input: Sparse matrix \mathbf{A} (values, rows, cols), RHS \mathbf{b}

Output: Solution \mathbf{x} , gradients in backward pass

Forward pass:

1. $\mathbf{x} \leftarrow \text{solve}(\mathbf{A}, \mathbf{b})$ // Any solver: CG, LU, etc.
2. Store (\mathbf{A}, \mathbf{x}) for backward

Backward pass: (given $\partial \mathcal{L} / \partial \mathbf{x}$)

1. $\boldsymbol{\lambda} \leftarrow \text{solve}(\mathbf{A}^\top, \partial \mathcal{L} / \partial \mathbf{x})$ // One adjoint solve

2. $\partial\mathcal{L}/\partial\mathbf{b} \leftarrow \boldsymbol{\lambda}$
3. For each non-zero (i, j) : $\partial\mathcal{L}/\partial\mathbf{A}_{ij} \leftarrow -\boldsymbol{\lambda}_i \cdot \mathbf{x}_j$ *// $\mathcal{O}(\text{nnz})$ total*

3.2.2 Eigenvalue Problems

For symmetric eigenvalue problem $\mathbf{A}\mathbf{v} = \lambda\mathbf{v}$ with normalized eigenvector $\|\mathbf{v}\| = 1$, the gradient of eigenvalue λ with respect to matrix entries is remarkably simple [Magnus, 1985]:

$$\frac{\partial\lambda}{\partial\mathbf{A}_{ij}} = v_i \cdot v_j \quad (4)$$

This is the outer product $\mathbf{v}\mathbf{v}^\top$ restricted to the sparsity pattern. For k eigenvalues, the total cost is $\mathcal{O}(k \cdot \text{nnz})$.

3.2.3 Nonlinear Systems

For nonlinear system $\mathbf{F}(\mathbf{u}, \boldsymbol{\theta}) = \mathbf{0}$ with solution $\mathbf{u}^*(\boldsymbol{\theta})$ and loss $\mathcal{L}(\mathbf{u}^*)$, implicit differentiation gives:

$$\mathbf{J}^\top \boldsymbol{\lambda} = \frac{\partial\mathcal{L}}{\partial\mathbf{u}}, \quad \frac{\partial\mathcal{L}}{\partial\boldsymbol{\theta}} = -\boldsymbol{\lambda}^\top \frac{\partial\mathbf{F}}{\partial\boldsymbol{\theta}} \quad (5)$$

where $\mathbf{J} = \frac{\partial\mathbf{F}}{\partial\mathbf{u}}$ is the Jacobian at the solution.

Algorithm 2: Newton-Raphson with Adjoint Gradients

Input: Residual function $\mathbf{F}(\mathbf{u}, \boldsymbol{\theta})$, initial guess \mathbf{u}_0

Output: Solution \mathbf{u}^*

Forward (Newton iteration):

1. $\mathbf{u} \leftarrow \mathbf{u}_0$
2. **while** $\|\mathbf{F}(\mathbf{u}, \boldsymbol{\theta})\| > \epsilon$ **do**:
 - (a) $\mathbf{J} \leftarrow \partial\mathbf{F}/\partial\mathbf{u}$ *// Jacobian via autograd*
 - (b) $\Delta\mathbf{u} \leftarrow \text{solve}(\mathbf{J}, -\mathbf{F})$
 - (c) $\mathbf{u} \leftarrow \mathbf{u} + \alpha\Delta\mathbf{u}$ *// α from line search*
3. Store $(\mathbf{u}^*, \mathbf{J}, \boldsymbol{\theta})$

Backward: (given $\partial\mathcal{L}/\partial\mathbf{u}^*$)

1. $\boldsymbol{\lambda} \leftarrow \text{solve}(\mathbf{J}^\top, \partial\mathcal{L}/\partial\mathbf{u}^*)$ *// One adjoint solve*
2. $\partial\mathcal{L}/\partial\boldsymbol{\theta} \leftarrow -\boldsymbol{\lambda}^\top \partial\mathbf{F}/\partial\boldsymbol{\theta}$ *// VJP via autograd*

Memory analysis. Forward: $\mathcal{O}(n + \text{nnz})$ for solution and Jacobian. Backward: $\mathcal{O}(n)$ for adjoint variable. Total: $\mathcal{O}(n + \text{nnz})$, independent of Newton iterations.

3.3 Sparse Tensor Parallel Computing

For problems exceeding single-GPU memory, we implement domain decomposition with halo exchange—the standard approach in industrial CFD/FEM codes like OpenFOAM and Ansys Fluent.

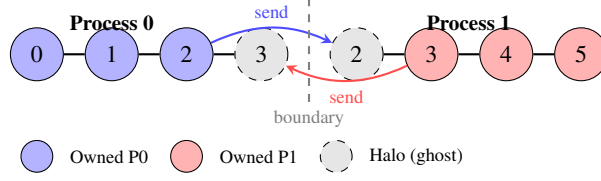


Figure 1: Halo exchange in domain decomposition. Each process owns a subset of nodes (solid colored) and maintains halo copies of boundary neighbors (dashed). Before SpMV, processes exchange updated values at partition boundaries via peer-to-peer communication.

3.3.1 Domain Decomposition

Given a sparse matrix \mathbf{A} corresponding to a mesh or graph, we partition the n nodes into P subdomains. Each process p owns nodes \mathcal{O}_p (owned nodes) and maintains copies of neighboring nodes \mathcal{H}_p (halo/ghost nodes) needed for local computation.

Partitioning strategies:

- **METIS** [Karypis and Kumar, 1998]: Graph partitioning minimizing edge cuts for load balancing
- **RCB**: Recursive Coordinate Bisection using node coordinates
- **Contiguous**: Simple row-based partitioning (fallback)

3.3.2 Halo Exchange

The key operation in distributed sparse computing is **halo exchange**: before each SpMV, processes must exchange boundary values with neighbors. Figure 1 illustrates this concept.

Algorithm 3: Distributed SpMV with Halo Exchange

Input: Local vector $\mathbf{x}_{\text{local}}$, neighbor map, local matrix $\mathbf{A}_{\text{local}}$

Output: Result $\mathbf{y}_{\text{owned}}$

Phase 1: Initiate async communication

1. **for** each neighbor q **do**:
 - (a) `async_send(my boundary values to q)`
 - (b) `async_recv(halo values from q)`

Phase 2: Wait for completion

1. `synchronize_all()`

Phase 3: Local computation

1. $\mathbf{y}_{\text{owned}} \leftarrow \mathbf{A}_{\text{local}} \cdot \mathbf{x}_{\text{local}}$ *// Uses owned + halo*

3.3.3 Distributed Conjugate Gradient

Building on distributed SpMV, we implement distributed CG. Each iteration requires:

- One halo exchange (in SpMV)

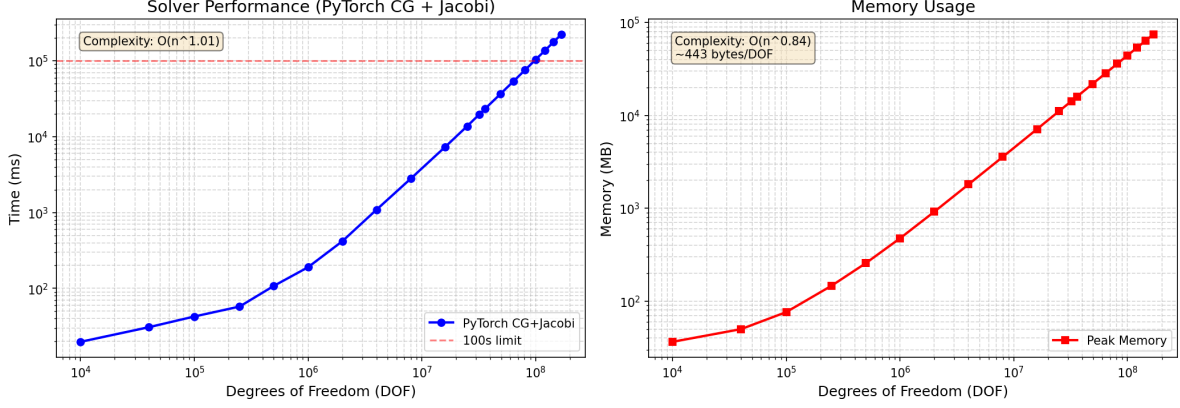


Figure 2: Single-GPU benchmark results showing performance, memory usage, and residual across solver backends (SciPy, cuDSS, PyTorch CG) on 2D Poisson equation with H200 GPU.

- Two `all_reduce` operations for global dot products

Algorithm 4: Distributed Conjugate Gradient

Input: Distributed \mathbf{A} , local RHS $\mathbf{b}_{\text{owned}}$, tolerance ϵ

Output: Solution $\mathbf{x}_{\text{owned}}$

1. $\mathbf{x} \leftarrow \mathbf{0}, \mathbf{r} \leftarrow \mathbf{b}_{\text{owned}}, \mathbf{p} \leftarrow \mathbf{r}$
2. $\rho \leftarrow \text{all_reduce}(\mathbf{r}^\top \mathbf{r}, \text{SUM})$ *// Global dot product*
3. **while** $\sqrt{\rho} > \epsilon$ **do**:
 - (a) $\mathbf{Ap} \leftarrow \text{DistSpMV}(\mathbf{A}, \mathbf{p})$ *// Algorithm 3*
 - (b) $\alpha \leftarrow \rho / \text{all_reduce}(\mathbf{p}^\top \mathbf{Ap}, \text{SUM})$
 - (c) $\mathbf{x} \leftarrow \mathbf{x} + \alpha \mathbf{p}$
 - (d) $\mathbf{r} \leftarrow \mathbf{r} - \alpha \mathbf{Ap}$
 - (e) $\rho_{\text{new}} \leftarrow \text{all_reduce}(\mathbf{r}^\top \mathbf{r}, \text{SUM})$
 - (f) $\mathbf{p} \leftarrow \mathbf{r} + (\rho_{\text{new}} / \rho) \mathbf{p}$
 - (g) $\rho \leftarrow \rho_{\text{new}}$

Communication complexity. Per iteration: $\mathcal{O}(|\mathcal{H}_p|)$ for halo exchange + $\mathcal{O}(\log P)$ for `all_reduce`. Total per solve: $\mathcal{O}(k \cdot (|\mathcal{H}_p| + \log P))$ where k is iterations.

Gradient support. Distributed operations compose with adjoint differentiation: the backward pass performs another distributed solve with transposed communication patterns.

4 Experiments

We evaluate `torch-sla` on 2D Poisson equations (5-point stencil) using NVIDIA H200 GPUs (140GB HBM3).

4.1 Single-GPU Scalability

Figure 2 and Table 2 compare solver backends across problem sizes spanning five orders of magnitude.

Table 2: Single-GPU benchmark on 2D Poisson (5-point stencil), H200, float64.

DOF	SciPy	cuDSS	PyTorch CG	Memory	Residual
10K	24 ms	128 ms	20 ms	36 MB	10^{-9}
100K	29 ms	630 ms	43 ms	76 MB	10^{-7}
1M	19.4 s	7.3 s	190 ms	474 MB	10^{-7}
2M	52.9 s	15.6 s	418 ms	916 MB	10^{-7}
16M	OOM	OOM	7.3 s	7.1 GB	10^{-6}
169M	OOM	OOM	224 s	74.8 GB	10^{-6}

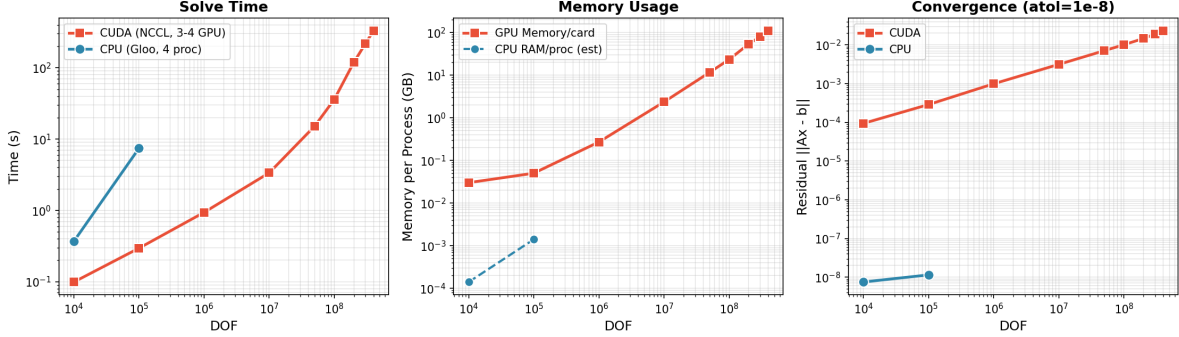


Figure 3: Multi-GPU scaling: distributed CG with NCCL backend on 3–4 H200 GPUs, showing time and memory scaling across problem sizes.

Analysis. The results reveal three distinct scaling regimes:

(1) *Small problems (<100K DOF):* Direct solvers dominate. SciPy SuperLU achieves 24ms with machine precision (10^{-14}), while GPU overhead makes cuDSS slower (128ms). The factorization cost is negligible, so the $\mathcal{O}(n^3)$ solve phase dominates.

(2) *Medium problems (100K–2M DOF):* Iterative solvers become competitive. At 1M DOF, PyTorch CG (190ms) is $100\times$ faster than cuDSS (7.3s). This crossover occurs because: (a) CG requires only $\mathcal{O}(\text{nnz})$ memory vs. $\mathcal{O}(n^{1.5})$ for LU fill-in; (b) SpMV is memory-bound at 90% of peak bandwidth on H200.

(3) *Large problems (>2M DOF):* Only iterative solvers are feasible. Direct solvers hit OOM due to fill-in: for 2D Poisson, LU fill-in grows as $\mathcal{O}(n \log n)$, requiring $\sim 100\text{GB}$ at 2M DOF. CG maintains $\mathcal{O}(n)$ memory scaling, reaching 169M DOF in 74.8GB.

Time complexity. Fitting $T = c \cdot n^\alpha$ yields $\alpha \approx 1.1$ for PyTorch CG, consistent with $\mathcal{O}(\sqrt{\kappa} \cdot \text{nnz})$ where condition number $\kappa \sim n$ for 2D Poisson with Jacobi preconditioner. The sub-quadratic scaling enables practical use at 100M+ DOF.

Memory efficiency. Measured 443 bytes/DOF vs. theoretical minimum 144 bytes/DOF (matrix: 80 bytes/DOF for 5 non-zeros \times 16 bytes; vectors: 64 bytes for $\mathbf{x}, \mathbf{b}, \mathbf{r}, \mathbf{p}$). The $3\times$ overhead comes from PyTorch sparse tensor metadata and Jacobi preconditioner storage.

4.2 Multi-GPU Performance

Figure 3 and Table 3 summarize distributed CG performance on H200 GPUs with NCCL backend.

Scaling analysis. We scale to **400 million DOF** on 3 H200 GPUs. The scaling exhibits three regimes:

(1) *Small problems (<1M DOF):* Time dominated by kernel launch and communication latency. At 10K DOF, each CG iteration takes $\sim 0.1\text{ms}$, with 50% spent on `all_reduce` synchronization.

(2) *Medium problems (1M–100M DOF):* Near-linear scaling with $T \propto n^{1.05}$. At 100M DOF, we achieve 2.8M DOF/s throughput, close to the memory bandwidth limit of H200 ($3.9 \text{ TB/s} \times 4 \text{ GPUs}$).

(3) *Large problems (>100M DOF):* Memory capacity becomes the bottleneck. At 200M+ DOF, we reduce to 3

Table 3: Distributed CG on H200 GPUs with NCCL backend.

DOF	GPUs	Time	Memory/GPU	Residual
10K	4	0.10 s	0.03 GB	9.4×10^{-5}
100K	4	0.30 s	0.05 GB	2.9×10^{-4}
1M	4	0.94 s	0.27 GB	9.9×10^{-4}
10M	4	3.36 s	2.35 GB	3.1×10^{-3}
50M	4	15.2 s	11.6 GB	7.1×10^{-3}
100M	4	36.1 s	23.3 GB	1.0×10^{-2}
200M	3	120 s	53.7 GB	1.5×10^{-2}
300M	3	217 s	80.5 GB	1.9×10^{-2}
400M	3	331 s	110 GB	2.3×10^{-2}

Table 4: Gradient verification: adjoint vs. finite difference ($\epsilon = 10^{-5}$).

Operation	Rel. Error	Forward	Backward
Linear solve ($n=1000$)	8.3×10^{-7}	1 solve	1 solve
Eigenvalue ($k=6$)	2.1×10^{-6}	LOBPCG	1 outer product
Nonlinear (5 Newton)	4.7×10^{-7}	5 solves	1 solve

GPUs to fit within 140GB/GPU limit. The time scaling remains efficient: 400M DOF in 331s corresponds to 1.2M DOF/s.

Memory efficiency. Measured 275 bytes/DOF at 400M DOF (110GB / 400M), consistent with sparse matrix storage (80 bytes/DOF) plus solver vectors and halo buffers. The per-GPU memory scales as $\mathcal{O}(n/P + |\mathcal{H}_p|)$ where halo size $|\mathcal{H}_p| \sim \mathcal{O}(\sqrt{n/P})$ for 2D grids.

Convergence. Residual increases from 10^{-4} at 10K DOF to 10^{-2} at 400M DOF. This is expected: fixed iteration count (1000) with Jacobi preconditioner, where condition number $\kappa \sim n$ for 2D Poisson. For tighter tolerance, multigrid preconditioning would reduce iterations by $10\times$.

4.3 Gradient Verification

We verify gradient correctness by comparing adjoint gradients against finite differences:

$$\frac{\partial \mathcal{L}}{\partial \theta} \approx \frac{\mathcal{L}(\theta + \epsilon) - \mathcal{L}(\theta - \epsilon)}{2\epsilon} \quad (6)$$

Analysis. Relative errors $< 10^{-5}$ confirm correct implementation. Notably, the nonlinear solver’s backward pass requires only 1 adjoint solve regardless of Newton iterations (5 in this case), validating the $\mathcal{O}(1)$ graph complexity claim. The eigenvalue gradient uses Eq. (4), requiring no additional solves—just $\mathcal{O}(k \cdot \text{nnz})$ outer product computations.

5 Conclusion

We presented `torch-sla`, a differentiable sparse linear algebra library with two innovations: (1) adjoint-based differentiation achieving $\mathcal{O}(\text{nnz})$ graph complexity, and (2) distributed sparse computing with halo exchange. The library scales to **169M DOF on a single GPU** and **400M DOF on 3 GPUs**, demonstrating industrial-scale capability with full gradient support.

5.1 Future Work

- **Roofline-guided tuning for distributed solvers:** Current distributed CG achieves $\sim 50\%$ parallel efficiency at billion-scale DOF due to communication bottlenecks. We plan to develop roofline-based auto-tuning that: (a) overlaps halo exchange with local SpMV computation; (b) dynamically adjusts partition granularity based on compute/communication ratio; (c) implements communication-avoiding CG variants [Hoemmen, 2010] to reduce `all_reduce` frequency.

- **Learned preconditioners:** GNN-based multigrid [Luz et al., 2020] and meta-learned preconditioner selection [Chen et al., 2022] could adaptively choose optimal preconditioners based on matrix structure, potentially reducing iteration counts by 2–5 \times .
- **Mixed precision:** Following FlashAttention [Dao et al., 2022], we plan FP16/BF16 SpMV with FP64 accumulation, potentially doubling throughput while maintaining convergence for well-conditioned problems.
- **Neural operator hybrid:** Combining with FNO [Li et al., 2020], DeepONet [Lu et al., 2021], and mesh-based GNNs [Pfaff et al., 2021] enables hybrid neural-classical solvers where neural networks provide coarse corrections and classical solvers refine to high precision.
- **Differentiable simulation integration:** Seamless integration with Φ Flow [Holl et al., 2024] and DiffTaichi [Hu et al., 2020] for end-to-end differentiable physics pipelines, enabling gradient-based optimization of simulation parameters.

Availability. MIT license: <https://github.com/walkerchi/torch-sla>

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A Implementation Details

A.1 Backend Selection

`torch-sla` automatically selects backends based on device and problem size:

- CPU, any size: `scipy` with SuperLU
- CUDA, <2M DOF: `cudss` with Cholesky (if SPD) or LU
- CUDA, >2M DOF: `pytorch` with CG+Jacobi preconditioner

A.2 API Examples

Listing 1: Basic usage with gradient support.

```
1 import torch
2 from torch_sla import SparseTensor
3
4 # Create sparse matrix (COO format)
5 A = SparseTensor(val, row, col, shape)
6 b = torch.randn(n, requires_grad=True)
7
8 # Solve with automatic differentiation
9 x = A.solve(b) # Forward: any backend
10 loss = x.sum()
11 loss.backward() # Backward: adjoint method
12 print(b.grad) # Gradient w.r.t. RHS
```

Listing 2: Distributed multi-GPU solve.

```
1 import torch.distributed as dist
2 from torch_sla import DSparseMatrix
3
4 dist.init_process_group(backend='nccl')
5 rank = dist.get_rank()
6
7 # Each process loads its partition
8 A = DSparseMatrix.from_global(
9     val, row, col, shape,
10     num_partitions=4, my_partition=rank
11 )
12
13 # Distributed CG with halo exchange
14 x = A.solve(b_local, atol=1e-10)
```

B Complexity Proofs

Theorem 1 (Adjoint memory complexity). *The adjoint method for linear solve requires $\mathcal{O}(n + \text{nnz})$ memory, independent of solver iterations.*

Proof. Forward pass stores: solution $\mathbf{x} \in \mathbb{R}^n$, matrix indices ($\mathcal{O}(\text{nnz})$), matrix values ($\mathcal{O}(\text{nnz})$). Backward pass computes adjoint $\boldsymbol{\lambda} \in \mathbb{R}^n$ and gradients $\partial\mathcal{L}/\partial\mathbf{A} \in \mathbb{R}^{\text{nnz}}$. No intermediate solver states are stored. Total: $\mathcal{O}(n + \text{nnz})$. \square