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1

Contest

2

Data structures

3

Dynamic Programming

4

Game theory

5

Geometry

6

Graph

7

Mathematics

8

Number theory

9

Strings

10

Miscellaneous

# Contest (1)

template.cpp47 lines

```
#include <bits/stdc++.h>
using namespace std;
#define sws cin.tie(0)->sync_with_stdio(0)

#define endl '\n'
#define ll long long
#define ld long double
#define pb push_back
#define ff first
#define ss second
#define pll pair<ll, ll>
#define vll vector<ll>

#define teto(a, b) (((a)+(b)-1)/(b))
#define LSB(i) ((i) & -(i))
#define MSB(i) (63 - __builtin_clzll(i)) // for ll
#define BITS(i) __builtin_popcountll(i)

template<class A> void debug(A a) {
    cout << "container: ";
    for(auto b : a) cout << b << " ";
    cout << endl;
}

template<class... A> void dbg(A const&... a) {
    ((cout << "{ " << a << " } ", ...);
    cout << endl;
}

const ll MAX = 2e5+10;
const ll MOD = 998'244'353;
const ll INF = INT32_MAX; // INT64_MAX
const ld EPS = 1e-7;
const ld PI = acos(-1);

#include <chrono>
using namespace std::chrono;
int32_t main(){ sws;
    auto start = high_resolution_clock::now();
```

```
// function to be timed here
auto stop = high_resolution_clock::now();
auto duration = duration_cast<milliseconds>(stop - start);
cout << duration.count() << endl;
}

// add to beginning
#pragma GCC optimize("O3,unroll-loops")
#pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")

.bashrc1 lines
alias comp='g++ -std=c++17 -O2 -g3 -ggdb3 -fsanitize=address,
    undefined -Wall -Wextra -Wshadow -Wconversion -o test'

hash.sh3 lines
# Hashes a file, ignoring all whitespace and comments. Use for
# verifying that code was correctly typed. CTRL+D to send EOF
cpp -dD -P -fpreprocessed | tr -d '[:space:]' | md5sum | cut -c
    -6

troubleshoot.txt52 lines
Pre-submit:
Write a few simple test cases if sample is not enough.
Are time limits close? If so, generate max cases.
Is the memory usage fine?
Could anything overflow?
Make sure to submit the right file.

Wrong answer:
Print your solution! Print debug output, as well.
Are you clearing all data structures between test cases?
Can your algorithm handle the whole range of input?
Read the full problem statement again.
Do you handle all corner cases correctly?
Have you understood the problem correctly?
Any uninitialized variables?
Any overflows?
Confusing N and M, i and j, etc.?
Are you sure your algorithm works?
What special cases have you not thought of?
Are you sure the STL functions you use work as you think?
Add some assertions, maybe resubmit.
Create some testcases to run your algorithm on.
Go through the algorithm for a simple case.
Go through this list again.
Explain your algorithm to a teammate.
Ask the teammate to look at your code.
Go for a small walk, e.g. to the toilet.
Is your output format correct? (including whitespace)
Rewrite your solution from the start or let a teammate do it.

Runtime error:
Have you tested all corner cases locally?
Any uninitialized variables?
Are you reading or writing outside the range of any vector?
Any assertions that might fail?
Any possible division by 0? (mod 0 for example)
Any possible infinite recursion?
Invalidated pointers or iterators?
Are you using too much memory?
Debug with resubmits (e.g. remapped signals, see Various).

Time limit exceeded:
Do you have any possible infinite loops?
What is the complexity of your algorithm?
Are you copying a lot of unnecessary data? (References)
```

How big is the input and output? (consider scanf)  
Avoid vector, map. (use arrays/unordered\_map)  
What do your teammates think about your algorithm?

Memory limit exceeded:  
What is the max amount of memory your algorithm should need?  
Are you clearing all data structures between test cases?

## Data structures (2)

### 2.1 Stack

An optimization for *std::stack* is to use a *std::vector* as the container, instead of *std::deque* !

```
stack<int, vector<int>> st;
```

A stack can be used to efficiently solve the maximum rectangle in a histogram problem:

max-rectangle-histogram.cpp

**Description:** solves the problem of finding the maximum rectangle area in a grid setting (different widths, different heights)  
**Time:**  $O(nm)$ 461370, 54 lines

```
// Example Problem: You are given a map of a forest where some
    squares are empty and some squares have trees.
// What is the maximum area of a rectangular building that can
    be placed in the forest so that no trees must be cut down?

ll maxRectangleHistogram(vector<ll> x) { // O(n)

    // add an end point with heigth 0 to compute the last
        rectangles
    x.pb(0);

    ll area = 0;
    ll n = x.size();
    stack<pll, vector<pll>> st; // {maxLeft, height for this
        rectangle}

    for(ll i=0; i<n; i++) {
        ll h = x[i];
        ll maxLeft = i;

        while(!st.empty() and st.top().ss >= h) {
            auto [maxLeft2, h2] = st.top(); st.pop();

            // compute the area of the de-stacked rectangle
            area = max(area, (i-maxLeft2)*h2 );

            // extend current rectangle width with previous
            maxLeft = maxLeft2;
        }

        st.push({maxLeft, h});
    }

    return area;
}
```

```
int32_t main(){ sws;
    ll n, m; cin >> n >> m;

    vector<vector<ll>> grid(n, vector<ll>(m));

    // convert the problem into N histogram subproblems, O(n m)
    for(ll i=0; i<n; i++) {
```

```
for(ll j=0; j<m; j++) {
    char c; cin >> c;
    if (c == '*') grid[i][j] = 0;
    else if (i == 0) grid[i][j] = 1;
    else grid[i][j] = grid[i-1][j] + 1;
}

ll area = 0;
for(ll i=0; i<n; i++) {
    area = max(area, maxRectangleHistogram(grid[i]));
}

cout << area << endl;
}
```

Also can be used to solve the maximum rectangle in a grid, with some blocked spots:

max-rectangle-grid.cpp

**Description:** solves the problem of finding the maximum rectangle area in a histogram setting (same bottom, different heights).  
**Time:**  $\mathcal{O}(n)$

```
8610da, 38 lines
// Example Problem: A fence consists of n vertical boards. The
// width of each board is 1 and their heights may vary.
// You want to attach a rectangular advertisement to the fence.
// What is the maximum area of such an advertisement?

ll maxRectangleHistogram(vector<ll> x) { // O(n)

    // add an end point with heighth 0 to compute the last
    // rectangles
    x.pb(0);

    ll area = 0;
    ll n = x.size();
    stack<pll, vector<pll>> st; // {maxLeft, height for this
    // rectangle}

    for(ll i=0; i<n; i++) {
        ll h = x[i];
        ll maxLeft = i;

        while(!st.empty() and st.top().ss >= h) {
            auto [maxLeft2, h2] = st.top(); st.pop();

            // compute the area of the de-stacked rectangle
            area = max(area, (i-maxLeft2)*h2 );

            // extend current rectangle width with previous
            maxLeft = maxLeft2;
        }

        st.push({maxLeft, h});
    }

    return area;
}

int32_t main(){ sws;
ll n; cin >> n;
vector<ll> x;
for(ll i=0, a; i<n; i++) cin >> a, x.pb(a);
cout << maxRectangleHistogram(x) << endl;
}
```

2.2 List

`std::list` is a container that supports constant time insertion and removal of elements from anywhere in the container.

Adding, removing and moving the elements within the list or across several lists does not invalidate the iterators or references. An iterator is invalidated only when the corresponding element is deleted.

Element Access:  $\mathcal{O}(1)$

- `list.back()`
- `list.front()`

Modifiers:  $\mathcal{O}(1)$

- `list.insert(itr, val)` inserts val before itr and returns an itr to the inserted value
- `list.erase(itr)` erases the element referenced by itr and returns the itr for the next value (or .end())
- `list.push_back(val)`
- `list.pop_back(val)`
- `list.push_front(val)`
- `list.pop_front(val)`

2.3 Ordered Set

Policy Based Data Structures (PBDS) from gcc compiler

Ordered Multiset can be created using `ordered_set<pll>val, idx`

`order_of_key()` can search for non-existent keys!

`find_by_order()` requires existent key and return the 0-idx position of the given value. Therefore, it returns the numbers of elements that are smaller than the given value;

ordered-set.cpp

```
Description: Set with index operators, implemented by gnu pbds. Remember to compile with gcc!!
Time: O(log(N)) but with slow constant
<bits/extc++.h>, <bits/extc++.h> 8578e5, 11 lines

// 0-idx
// find_by_order(i) -> iterator to elem with index i
// order_of_key(val) -> index of key

// Ordered Set
using namespace __gnu_pbds;
template <class T> using ordered_set = tree<T, null_type, less<
T>, rb_tree_tag, tree_order_statistics_node_update>;

// Ordered Map
using namespace __gnu_pbds;
template <class K, class V> using ordered_map = tree<K, V, less
<K>, rb_tree_tag, tree_order_statistics_node_update>;
```

2.3.1 Pyramid Array min-cost

You are given an array consisting of n integers. On each move, you can swap any two adjacent values. You want to transform the array into a pyramid array. This means that the final array has to be first increasing and then decreasing. It is also allowed that the final array is only increasing or decreasing. What is the minimum number of moves needed?

pyramid-array.cpp

**Description:** algorithm to find the min-cost of sorting an array in a pyramid order  
**Time:**  $\mathcal{O}(N\log(N))$ , or  $\mathcal{O}(N\log^2(N))$  if iterating the map directly

```
1651d7, 25 lines
int32_t main() { sws;
ll n; cin >> n;
map<ll, vll> freq;
for(ll i=0; i<n; i++) {
    ll val; cin >> val;
    freq[val].pb(i);
}

ordered_set<ll> os; // os with indexes of greater processed
// elements
ll ans = 0;
// iterate from greater values to lesser one.
// for each element,
// consider inserting it to the left of all greater
// elements
// or to the right of all greater elements
for(auto itr = freq.rbegin(); itr != freq.rend(); itr++) {
    auto [val, vec] = *itr;
    for(auto idx : vec) {
        ll pos = os.order_of_key({idx});
        ll left_cost = pos;
        ll right_cost = (ll)os.size() - pos;
        ans += min(left_cost, right_cost);
    }
    for(auto idx : vec) os.insert(idx);
}

cout << ans << endl;
}
```

2.4 Interval Set

interval-set.cpp

**Description:** A set that contains closed [l, r] interval which are disjoint (no intersection). This set is ordered and each interval [l1, r1] < [l2, r2] has r1 < l2. When a new interval is added, it checks which intersections will occur and rearranges the intervals.  
**Time:**  $\mathcal{O}(\log(N))$  per insertion, slow constant

```
a3c7e0, 29 lines
// keeps track of disjoint closed intervals [l, r]
// a new interval added may replace parts of an older one
struct IntervalSet {
    using T = array<ll, 3>;
    set<T> ranges;

    void add(T arr) {
        auto [l, r, k] = arr;

        while(ranges.upper_bound({r, INF, INF}) != ranges.begin()
            ) {
            auto itr = prev(ranges.upper_bound({r, INF, INF}));
            auto [l2, r2, k2] = *itr;

            if (r2 < l) break;
            // guarantees that there is an intersection: l2 <= r
            // and r2 >= l
        }
    }
};
```

```

        ranges.erase(itr);

    if (l2 <= l-1) {
        ranges.insert({l2, l-1, k2});
    }

    if (r+1 <= r2) {
        ranges.insert({r+1, r2, k2});
    }

    ranges.insert({l, r, k});
}
};

```

## 2.5 Disjoint Set Union

There are two optional improvements:

- Tree Balancing
- Path Compression

If one improvement is used, the time complexity will become  $O(\log N)$

If both are used,  $O(\alpha) \approx O(5)$

dsu.cpp

**Description:** Disjoint Set Union with path compression and tree balancing  
**Time:**  $O(\alpha)$

424a7d, 22 lines

```

struct DSU{
    vector<ll> group, card;
    DSU (ll n){
        n += 1; // 0-idx -> 1-idx
        group = vector<ll>(n);
        iota(group.begin(), group.end(), 0);
        card = vll(n, 1);
    }
    ll find(ll i){
        return (i == group[i]) ? i : (group[i] = find(group[i]));
    }
    // returns false if a and b are already in the same
    // component
    bool join(ll a ,ll b){
        a = find(a);
        b = find(b);
        if (a == b) return false;
        if (card[a] < card[b]) swap(a, b);
        card[a] += card[b];
        group[b] = a;
        return true;
    }
};

```

## 2.6 Trie

Also called a **digital tree** or **prefix tree**.

trie.cpp

**Description:** Creates a trie by pre-allocating the trie array, which contains the indices for the child nodes. The trie can be easily modified to support alphanumeric strings instead of binary strings.

**Time:**  $O(D)$ ,  $D$  = depth of trie

efeb7e, 40 lines

// MAX = maximum number of nodes that can be created

```

struct Trie{
    ll trie[MAX][26];

```

```

    bool isWordEnd[MAX];
    ll nxt = 1, wordsCnt = 0;

    void add(string s){ // O(Depth)
        ll node = 0;
        for(auto c: s) {
            if(trie[node][c-'a'] == 0) { // create new node
                trie[node][c-'a'] = nxt++;
            }
            node = trie[node][c-'a'];
        }
        if(!isWordEnd[node]){
            isWordEnd[node] = true;
            wordsCnt++;
        }
    }

    bool find(string s, bool remove=false){ // O(Depth)
        ll node = 0;
        for(auto c: s) {
            if(trie[node][c-'a'] == 0) {
                return false;
            }
            else {
                node = trie[node][c-'a'];
            }
        }

        if(remove and isWordEnd[node]){
            isWordEnd[node] = false;
            wordsCnt--;
        }

        return isWordEnd[node];
    }
};

```

## 2.7 Segment Trees

Each node of the segment tree represents the cumulative value of a range.

**Observation:** For some problems, such as range distinct values query, considerer offline approach, ordering the queries by L for example.

### 2.7.1 Recursive SegTree

seg-recursive-sum.cpp

**Description:** Basic Recursive Segment Tree for points increase and range sum query. When initializing, choose an appropriate value for n.

**Time:**  $O(N \log N)$  to build,  $O(\log N)$  to increase or query

156cd2, 71 lines

```

// [0, n] segtree for range sum query, point increase
// 0 or 1-idx
ll L=0, R;
struct Segtree {

```

```

    struct Node {
        // null element:
        ll ps = 0;
    };

```

```

    vector<Node> tree;
    vector<ll> v;

```

```

    Segtree(ll n) {
        R = n;
        v.assign(n+1, 0);
    }

```

```

        tree.assign(4*(n+1), Node{});
    }

    Node merge(Node a, Node b) {
        return Node {
            // merge operation:
            a.ps + b.ps
        };
    }

    void build(ll l=L, ll r=R, ll i=1 ) {
        if (l == r) {
            tree[i] = Node {
                // leaf element:
                v[l]
            };
        }
        else {
            ll mid = (l+r)/2;
            build(l, mid, 2*i);
            build(mid+1, r, 2*i+1);
            tree[i] = merge(tree[2*i], tree[2*i+1]);
        }
    }

    void increase(ll idx=1, ll val=0, ll l=L, ll r=R, ll i=1 )
    {
        if (l == r) {
            // increase operation:
            tree[i].ps += val;
        }
        else {
            ll mid = (l+r)/2;
            if (idx <= mid) increase(idx, val, l, mid, 2*i);
            else increase(idx, val, mid+1, r, 2*i+1);
            tree[i] = merge(tree[2*i], tree[2*i+1]);
        }
    }

    Node query(ll left=L, ll right=R, ll l=L, ll r=R, ll i=1) {
        // left/right are the range limits for the query
        // l / r are the internal variables of the tree
        if (right < l or r < left){
            // null element:
            return Node{};
        }
        else if (left <= l and r <= right) return tree[i];
        else{
            ll mid = (l+r)/2;
            return merge(
                query(left, right, l, mid, 2*i),
                query(left, right, mid+1, r, 2*i+1)
            );
        }
    }
};

```

seg-recursive-minmax.cpp

**Description:** Basic Recursive Segment Tree for point update, range min/-max query When initializing, choose an appropriate value for n.

**Time:**  $O(N \log N)$  to build,  $O(\log N)$  to update or query

fb9b46, 71 lines

```

// [0, n] segtree for point update, range min/max query
// 0 or 1-idx
ll L=0, R;
struct Segtree {

```

```

    struct Node {
        // null element:

```

```
ll mn = INF, mx = -INF;
};

vector<Node> tree;
vector<ll> v;

Segtree(ll n) {
R = n;
v.assign(n+1, 0);
tree.assign(4*(n+1), Node{});
}

Node merge(Node a, Node b) {
return Node {
// merge operation:
min(a.mn, b.mn),
max(a.mx, b.mx)
};
}

void build(ll l=L, ll r=R, ll i=1 ) {
if (l == r) {
tree[i] = Node {
// leaf element:
v[l],
v[l]
};
}
else {
ll mid = (l+r)/2;
build(l, mid, 2*i);
build(mid+1, r, 2*i+1);
tree[i] = merge(tree[2*i], tree[2*i+1]);
}
}

void update(ll idx=1, ll val=0, ll l=L, ll r=R, ll i=1 ) {
if (l == r) {
// increase operation:
tree[i].mn = tree[i].mx = val;
}
else {
ll mid = (l+r)/2;
if (idx <= mid) update(idx, val, l, mid, 2*i);
else update(idx, val, mid+1, r, 2*i+1);
tree[i] = merge(tree[2*i], tree[2*i+1]);
}
}

Node query(ll left=L, ll right=R, ll l=L, ll r=R, ll i=1) {
if (right < l or r < left){
// null element:
return Node{};
}
else if (left <= l and r <= right) return tree[i];
else{
ll mid = (l+r)/2;
return merge(
query(left, right, l, mid, 2*i),
query(left, right, mid+1, r, 2*i+1)
);
}
}
};
```

2.7.2 Inverted Segtree

Instead of keeping the prefix sum for all the children in each node, store only the delta encoding value.

Therefore, to check a value in a certain position, iterate and sum all delta values from root to leaf.

seg-inverted.cpp

**Description:** Basic Inverted Segment Tree for point query stored value, range increase When initializing, choose an appropriate value for n.

**Time:**  $\mathcal{O}(N \log N)$  to build,  $\mathcal{O}(\log N)$  to range increase or point query

```
// [0, n] segtree for point query stored value, range increase
// 0 or 1-idx
ll L=0, R;
struct Segtree {

struct Node {
// null element:
ll ps = 0;
};

vector<Node> tree;
vector<ll> v;

Segtree(ll n) {
R = n;
v.assign(n+1, 0);
tree.assign(4*(n+1), Node{});
}

Node merge(Node a, Node b) {
return Node {
// merge operation:
a.ps + b.ps
};
}

void build(ll l=L, ll r=R, ll i=1 ) {
if (l == r) {
tree[i] = Node {
// leaf element:
v[l]
};
}
else {
ll mid = (l+r)/2;
build(l, mid, 2*i);
build(mid+1, r, 2*i+1);
tree[i] = Node{};
}
}

void increase(ll left, ll right, ll val=0, ll l=L, ll r=R, ll i=1 ) {
if (right < l or r < left) {
return;
}
else if (left <= l and r <= right) {
// increase operation
tree[i].ps += val;
}
else {
ll mid = (l+r)/2;
increase(left, right, val, l, mid, 2*i);
increase(left, right, val, mid+1, r, 2*i+1);
}
}

Node query(ll idx, ll l=L, ll r=R, ll i=1) {
if (l == r) {
return tree[i];
}
}
```

```
else {
ll mid = (l+r)/2;
if (idx <= mid)
return merge(tree[i], query(idx, l, mid, 2*i));
else
return merge(tree[i], query(idx, mid+1, r, 2*i+1));
}
};
```

2.7.3 PA Segtree

seg-pa.cpp

**Description:** Seg with PA (Progressao Aritmetica / Arithmetic Progression) When initializing the segmente tree, remeber to choose the range limits (L, R) and call build()

**Time:**  $\mathcal{O}(N \log N)$  to build,  $\mathcal{O}(\log N)$  to increase or query

```
// [0, n] segtree for range sum query, point increase
ll L=0, R;
struct SegtreePA {

struct Node {
// null element:
ll ps = 0;
};

vector<Node> tree;
vector<ll> v;
vector<pll> lazy; // {x, y} of {x*i + y}
// x = razao da PA, y = constante

SegtreePA(ll n) {
R = n;
v.assign(n+1, 0);
tree.assign(4*(n+1), Node{});
lazy.assign(4*(n+1), pll());
}

Node merge(Node a, Node b) {
return Node {
// merge operaton:
a.ps + b.ps
};
}

inline pll sum(pll a, pll b) {
return {a.ff+b.ff, a.ss+b.ss};
}

void build(ll l=L, ll r=R, ll i=1) {
if (l == r) {
tree[i] = Node {
// leaf element:
v[l]
};
}
else {
ll mid = (l+r)/2;
build(l, mid, 2*i);
build(mid+1, r, 2*i+1);
tree[i] = merge(tree[2*i], tree[2*i+1]);
}
lazy[i] = {0, 0};
}

void prop(ll l=L, ll r=R, ll i=1) {
auto [x, y] = lazy[i];
if (x == 0 and y == 0) return;
}
```

```

11 len = r-l+1;

// (l_val + r_val) * len / 2
Node val{ ((y + y + x*(len-1))*len) / 2 };
tree[i] = merge(tree[i], val);

if (l != r) {
    11 mid = (l+r)/2;
    lazy[2*i] = sum(lazy[2*i], lazy[i]);
    lazy[2*i+1] = sum(lazy[2*i+1], {x, y + x*(mid-l+1)});
}

lazy[i] = {0, 0};
}

// left/right are the range limits for the query
// l / r are the internal variables of the tree
void increase(11 left, 11 right, 11 x, 11 y, 11 l=L, 11 r=R,
, 11 i=1 ) {
    prop(1, r, i);
    if (right < 1 or r < left) return;
    else if (left <= 1 and r <= right) {
        lazy[i] = {x, y};
        prop(1, r, i);
    }
    else{
        11 mid = (l+r)/2;
        increase(left, right, x, y, 1, mid, 2*i);
        11 ny = y + max( x*( mid-max(left, 1) + 1), 0LL);
        increase(left, right, x, ny, mid+1, r, 2*i+1);
        tree[i] = merge(tree[2*i], tree[2*i+1]);
    }
}

Node query(11 left=L, 11 right=R, 11 l=L, 11 r=R, 11 i=1) {
    prop(1, r, i);
    if (right < 1 or r < left) {
        // null element:
        return Node{};
    }
    else if (left <= 1 and r <= right) return tree[i];
    else{
        11 mid = (l+r)/2;
        return merge(
            query(left, right, l, mid, 2*i),
            query(left, right, mid+1, r, 2*i+1)
        );
    }
}
};

```

## Dynamic Programming (3)

### 3.1 Divide-Conquer Optimization

Some dynamic programming problems have a recurrence of this form:

$$dp(i, j) = \min_{1 \leq k \leq j} dp(i-1, k-1) + C(k, j)$$

where  $C(k, j)$  is a cost function and  $dp(i, j) = 0$  when  $j \leq 0$  (using 1-idx).

Say  $0 \leq i < m$  and  $1 \leq j \leq n$ , and evaluating  $C$  takes  $O(1)$  time. Then the straightforward evaluation of the above recurrence is  $O(mn^2)$ . There are  $m \times n$  states, and  $n$  transitions for each state.

Let  $opt(i, j)$  be the value of  $k$  that minimizes the above expression. Assuming that the cost function satisfies the quadrangle inequality, we can show that  $opt(i, j-1) \leq opt(i, j)$  for all  $i, j$ . This is known as the monotonicity condition. Then, we can apply divide and conquer DP. The optimal "splitting point" for a fixed  $i$  increases as  $j$  increases.

This lets us solve for all states more efficiently. Say we compute  $opt(i, j)$  for some fixed  $i$  and  $j$ . Then for any  $j' < j$  we know that  $opt(i, j') \leq opt(i, j)$ . This means when computing  $opt(i, j')$ , we don't have to consider as many splitting points!

To minimize the runtime, we apply the idea behind divide and conquer. First, compute  $opt(i, n/2)$ . Then, compute  $opt(i, n/4)$ , knowing that it is less than or equal to  $opt(i, n/2)$ ; and  $opt(i, 3n/4)$  knowing that it is greater than or equal to  $opt(i, n/2)$ .

By recursively keeping track of the lower and upper bounds on  $opt$ , we reach a  $O(n \log n)$  runtime per  $i$ . Each possible value of  $opt(i, j)$  only appears in  $\log n$  different nodes.

#### divide-conquer-dp.cpp

**Description:** Optimize an  $O(mn^2)$  dp to  $O(mn \log(n))$  using divide and conquer. cost function must have quadrangle inequality ("wider is worse")  
**Time:**  $O(mn \log(n))$

a2746c, 48 lines

```

// m partitions, n elements
11 divideConquerDP(11 m, vector<11> &vec) {
    // vec indexed with 1-idx, vec[0] = 0
    11 n = vec.size() - 1;

    vector<11> ps(n+1, 0);
    for(11 i=1; i<=n; i++) {
        ps[i] = ps[i-1] + vec[i];
    }
    auto cost = [&](11 l, 11 r) {
        11 sum = ps[r] - ps[l-1];
        return sum * sum;
    };

    vector<11> cur(n+1, 0), nxt(n+1);

    // O(n log(n))
    function<void(11, 11, 11, 11)> compute = [&](11 l, 11 r, 11
        optl, 11 optr) {
        if (r < 1) return;

        11 mid = (l+r)/2;

        11 best = INF;
        11 opt = -1;
        for(11 k=optl; k<=min(mid, optr); k++) {
            11 val = cur[k-1] + cost(k, mid);
            if (val < best) {
                best = val;
                opt = k;
            }
        }
    };
}

```

```

}
nxt[mid] = best;

compute(l, mid-1, optl, opt);
compute(mid+1, r, opt, optr);
};

for(11 i=1; i<=n; i++) // 1 partition
    cur[i] = cost(1, i);

for(11 i=1; i<=m; i++) { // m partitions
    nxt[0] = 0;
    compute(1, n, 1, n);
    swap(cur, nxt);
}

return cur[n];
}

```

### 3.2 Knuth Optimization

Knuth's optimization, also known as the Knuth-Yao Speedup, is a special case of dynamic programming on ranges, that can optimize the time complexity of solutions by a linear factor, from  $O(n^3)$  for standard range DP to  $O(n^2)$ .

#### 3.2.1 Conditions

The Speedup is applied for transitions of the form:

$$dp(i, j) = \min_{i \leq k < j} [dp(i, k) + dp(k+1, j) + C(i, j)].$$

Similar to divide and conquer DP, let  $opt(i, j)$  be the value of  $k$  that minimizes the expression in the transition ( $opt$  is referred to as the "optimal splitting point" further in this article). The optimization requires that the following holds:

$$opt(i, j-1) \leq opt(i, j) \leq opt(i+1, j).$$

We can show that it is true when the cost function  $C$  satisfies the following conditions for  $a \leq b \leq c \leq d$ :

$$C(b, c) \leq C(a, d);$$

$$C(a, c) + C(b, d) \leq C(a, d) + C(b, c) \text{ (the quadrangle inequality [QI])}.$$

A common cost function that satisfies the above condition is the **sum of the values in a subarray**.

#### knuth.cpp

**Description:** Optimize  $O(n^3)$  to  $O(n^2)$  dp with transitions of finding a optimal division point  $k$  for  $[l, r]$ .

**Time:**  $O(n^2)$

8863a7, 37 lines

```

// dp[l][r] (inclusive) -> min cost
// opt[l][r] (inclusive) -> optimal spliting point k in l<=k<=r
11 dp[MAX][MAX], opt[MAX][MAX];

11 knuth(vector<11> &vec) {
    // vec indexed with 1-idx, vec[0] = 0
    11 n = vec.size() - 1;
}

```



```
vector<ll> ps(n+1, 0);
for(ll i=1; i<=n; i++) {
    ps[i] = ps[i-1] + vec[i];
}
auto C = [&](ll l, ll r) {
    return ps[r] - ps[l-1];
};

for(ll i=1; i<=n; i++) {
    opt[i][i] = i;
}

for(ll l=n-1; l>=1; l--) {
    for(ll r=l+1; r<=n; r++) {
        ll mn = INF;
        ll cost = C(l, r);
        for(ll k=opt[l][r-1]; k<=min(r-1, opt[l+1][r]); k++) {
            ll aux = dp[l][k] + dp[k+1][r] + cost;
            if (aux <= mn) {
                mn = aux;
                opt[l][r] = k;
            }
        }
        dp[l][r] = mn;
    }
}

return dp[1][n];
}
```

3.3 Slope Optimizations

3.3.1 Convex Hull Trick

If multiple transitions of the DP can be seen as first degree polynomials (lines). CHT can be used to optimized it

Some valid functions:

$ax + b$

$cx^2 + ax + b$  (ignore  $cx^2$  if c is independent)

cht-dynamic.cpp

**Description:** Dynamic version of CHT, thefore, one can insert lines in any order. There is no line removal operator

**Time:**  $\mathcal{O}(\log N)$  per query and per insertion

```
707da4, 51 lines

// Convex Hull Trick Dinamico
//
// Para float, use LLINF = 1/.0, div(a, b) = a/b
//
// update(x) atualiza o ponto de intersecao da reta x
// overlap(x) verifica se a reta x sobrepoe a proxima
// add(a, b) adiciona reta da forma ax + b
// query(x) computa maximo de ax + b para entre as retas
// se quiser computar o minimo, eh soh fazer (-a)x + (-b)
//
// O(log(n)) amortizado por insercao
// O(log(n)) por query

struct Line {
    mutable ll a, b, p;
    bool operator<(const Line& o) const { return a < o.a; }
    bool operator<(ll x) const { return p < x; }
};

struct DynamicCHT : multiset<Line, less<>> {
```

```
ll div(ll a, ll b) {
    return a / b - ((a ^ b) < 0 and a % b);
}

void update(iterator x) {
    if (next(x) == end()) x->p = LLINF;
    else if (x->a == next(x)->a) x->p = x->b >= next(x)->b ?
        LLINF : -LLINF;
    else x->p = div(next(x)->b - x->b, x->a - next(x)->a);
}

bool overlap(iterator x) {
    update(x);
    if (next(x) == end()) return 0;
    if (x->a == next(x)->a) return x->b >= next(x)->b;
    return x->p >= next(x)->p;
}

void add(ll a, ll b) {
    auto x = insert({a, b, 0});
    while (overlap(x)) erase(next(x)), update(x);
    if (x != begin() and !overlap(prev(x))) x = prev(x), update
        (x);
    while (x != begin() and overlap(prev(x)))
        x = prev(x), erase(next(x)), update(x);
}

ll query(ll x) {
    assert(!empty());
    auto l = *lower_bound(x);
    return *l->a < x ? l->b;
```

3.3.2 Li-chao Tree

Works for any type of function that has the **transcending property**:

Given two functions  $f(x), g(x)$  of that type, if  $f(t)$  is greater than/smaller than  $g(t)$  for some  $x=t$ , then  $f(x)$  will be greater than/smaller than  $g(x)$  for  $x \neq t$ . In other words, once  $f(x)$  “win/lose”  $g(x)$ ,  $f(x)$  will continue to “win/lose”  $g(x)$ .

3.3.3 Slope Trick

You are given an array of  $n$  integers. You want to modify the array so that it is non-decreasing, i.e., every element is at least as large as the previous element. On each move, you can increase or decrease the value of any element by one. What is the minimum number of moves required?

**Observation:** It is also possible to solve the problem of modifying the array to stricly increasing.

slope-trick.cpp

**Description:** Using Slope trick, compute the min cost to modify arry to be non-decreasing

**Time:**  $\mathcal{O}(n \log(n))$

```
25b6fc, 58 lines

// funcao f_i(x) = custo de deixar todo mundo ate i
// nao decrescente e <= x

// os pontos em changepoints sao os pontos da
// piecewise linear function convexa
```

```
// eu calculo g_i(x) = custo de deixar todo mundo ate i
// nao decrescente e v[i] == x

// entao f_i(x) = min(g_i(t) pra t <= x)

// podemos escrever g_i(x) = f_{i-1}(x) + |x-v[i]|
// entao a gente ta somando as funcoes e gerando outra convexa

// a resposta vai armazenar o custo (coord y) do opt
// e o topo do change-points vai ser o opt atual

// se opt < v[i] entao a gente calcula o g_i e o novo opt
// vai ser v[i]

// se opt > v[i] entao o slope entre opt e anterior opt vai
ficar reto
// (este anterior opt podendo ser o v[i] que vai ser inserido),
// entao basta retirar o ultimo opt e teremos de novo a
resposta
// neste caso devemos aumentar o custo do opt, que vai ser por
// (opt - v[i]) (so olhar a funcao em V do || e a convexa do f_{i-1})

// o v[i] vai ser inserido varias vezes no change-points
// pra denotar a inclinacao no slope dele

int32_t main() { sws;

    ll n; cin >> n;

    vector<ll> v(n);
    for(ll i = 0; i < n; i++) {
        cin >> v[i];
        // to change the problem
        // from increasing to non-decreasing
        // v[i] -= i;
    }

    priority_queue<ll> change_points;
    change_points.push(-INF);
    ll ans = 0;

    for(ll i = 0; i < n; i++) {
        ll opt = change_points.top();
        change_points.push(v[i]);

        if(opt > v[i]) {
            ans += opt - v[i];
            change_points.push(v[i]);
            change_points.pop();
        }

        cout << ans << endl;
    }
}
```

3.4 SOS DP

**Sum over Subsets DP (SOS DP)** computes how many elements there are for each mask which are a subset of this mask.

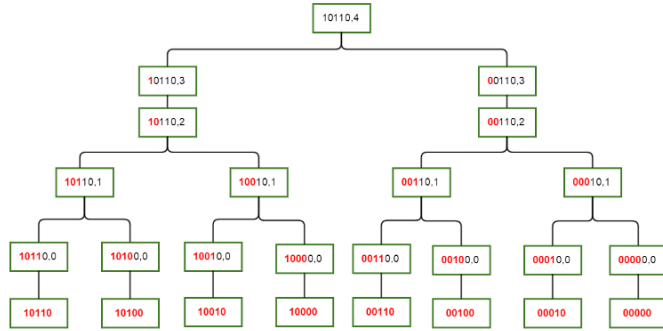
This can be modified for other operations in which the subset contributes for the mask . *Example:*

**1001** if a subset of **1101**;

0001 if a subset of 1101;

1100 if a subset of 1101;

1101 if a subset of 1101;



sos-dp.cpp

**Description:** Efficiently compute a bitmask dp, in which a subset of this bitmask contributes for the value of this bitmask.

**Time:**  $\mathcal{O}(2^N N)$ , N = number of bits

19e50a, 35 lines

```
// problem: Given a list of n integers, your task is to
// calculate for each element x:
// the number of elements y such that x | y = x
// the number of elements y such that x & y = x
// the number of elements y such that x & y != 0
```

```
const ll LOGMAX = 20;
ll dp[1 << LOGMAX];
ll dp2[1 << LOGMAX];

int32_t main(){ sws;
    ll n; cin >> n;

    vector<ll> a(n);
    for(auto &val : a) cin >> val;

    ll full = (1LL << LOGMAX) - 1;

    for(auto val : a) dp2[full^val] += 1;
    for(auto val : a) dp[val] += 1;

    for(ll b=0; b<LOGMAX; b++) {
        for(ll mask=0; mask<(1LL<<LOGMAX); mask++) {
            if (mask & (1LL << b)) {
                dp[mask] += dp[mask ^ (1LL << b)];
                dp2[mask] += dp2[mask ^ (1LL << b)];
            }
        }
    }

    for(auto val : a) {
        cout << dp[val] << " ";
        cout << dp2[full ^ val] << " ";
        cout << n - dp[full^val] << endl;
    }
}
```

## Game theory (4)

### 4.1 Classic Game

- There are n piles (heaps), each one with  $x_i$  stones.
- Each turn, a players must remove t stones (non-zero) from a pile, turning  $x_i$  into  $y_i$ .
- The game ends when it's impossible to make any more moves and the player without moves left lose.

### 4.2 Bouton's Theorem

Let  $s$  be the xor-sum value of all the piles sizes, a state  $s = 0$  is a losing position and a state  $s! = 0$  is a winning position

#### 4.2.1 Proof

All wining positions will have at least one valid move to turn the game into a losing position.

All losing positions will only have moves that turns the game into winning positions (except the base case when there are no piles left and the player already lost)

### 4.3 DAG Representation

Consider all game positions or states of the game as **Vertices** of a graph

Valid moves are the transition between states, therefore, the directed **Edges** of the graph

If a state has no outgoing edges, it's a dead end and a losing state (degenerated state).

If a state has only edges to winning states, therefore it is a losing state.

if a state has at least one edge that is a losing state, it is a winning state.

### 4.4 Sprague-Grundy Theorem

Let's consider a state  $u$  of a two-player impartial game and let  $v_i$  be the states reachable from it.

To this state, we can assign a fully equivalent game of Nim with one pile of size  $x$ . The number  $x$  is called the **Grundy value** or **nim-value** of the state  $u$ .

If **all transitions** lead to a *winning state*, the current state must be a *losing state* with number 0.

If **at least one transition** lead to a *losing state*, the current state must be a *winning state* with number  $\neq 0$ .

The **MEX** operator satisfies both condition above and can be used to calculate the nim-value of a state:

$nimber_u = \text{MEX of all } nimber_{v_i}$

Viewing the game as a DAG, we can gradually calculate the Grundy values starting from vertices without outgoing edges (number=0).

Note that the MEX operator **guarantees** that all nim-values smaller than the considered number can be reached, which is essentially the nim game with a single heap with pile size = number.

There are only two operations that are used when considering a Sprague-Grundy game:

#### 4.4.1 Composition

*XOR operator to compose sub-games into a single composite game*

When a game is played with multiple sub-games (as nim is played with multiple piles), you are actually choosing one sub-game and making a valid move there (choosing a pile and subtracting a value from it).

The final result/winner will depend on all the sub-games played. Because you need to play all games.

To compute the final result, one can simply consider the XOR of the numbers of all sub-games.

#### 4.4.2 Decomposition

*MEX operator to compute the number of a state that has multiple transitions to other states*

A state with number  $x$  can be transitioned (decomposed) into all states with number  $y < x$

Nevertheless a state may reach several states, only a single one will be used during the game. This shows the difference between **states** and **sub-games**: All sub-games must be played by the players, but the states of a sub-game may be ignored.

To compute the mex of a set efficiently:

mex.cpp

**Description:** Compute MEX efficiently by keeping track of the frequency of all existent elements and also the missing ones

**Time:**  $\mathcal{O}(\log N)$  per addition/removal,  $\mathcal{O}(1)$  to get mex value,  $\mathcal{O}(N \log(N))$  to initialize

d6f2b9, 27 lines

```
struct MEX {
    map<ll, ll> freq;
    set<ll> missing;

    // initialize set with values up to {max_valid_value}
    inclusive
    MEX(ll max_valid_value) { // O(n log(n))
        for(ll i=0; i<=max_valid_value; i++)
            missing.insert(i);
    }

    ll get() { // O(1)
        if (missing.empty()) return 0;
        return *missing.begin();
    }
}
```



```
void remove(ll val) { // O(log(n))
    freq[val]--;
    if (freq[val] == 0)
        missing.insert(val);
}

void add(ll val) { // O(log(n))
    freq[val]++;
    if (missing.count(val))
        missing.erase(val);
}
};
```

4.5 Variations and Extensions

4.5.1 Nim with Increases

Consider a modification of the classical nim game: a player can now add stones to a chosen pile instead of removing.

Note that this extra rule needs to have a restriction to keep the game acyclic (finite game).

**Lemma:** This move is not used in a winnig strategy and can be ignored.

**Proof:** If a player adds  $t$  stones in a pile, the next player just needs to remove  $t$  stones from this pile.

Considering that the game is finite and this ends sooner or later.

**Example:** If the set of possible outcomes for a state is 0, 1, 2, 7, 8, 9. The nimber is 3, because the MEX is 3, which is the smallest nim-value you can't transition into and also you can transition to all smaller nim-values.

Note that 7, 8, 9 transitions can be ignored, because you can simply revert the play by subtracting the same amount.

4.6 Misère Game

In this version, the player who takes the last object loses. To consider this version, simply swap the winning and losing player of the normal version.

4.7 Staircase Nim

4.7.1 Description

In Staircase Nim, there is a staircase with  $n$  steps, indexed from 0 to  $n-1$ . In each step, there are zero or more coins. Two players play in turns. In his/her move, a player can choose a step ( $i > 0$ ) and move one or more coins to step below it ( $i-1$ ). The player who is unable to make a move lose the game. That means the game ends when all the coins are in step 0.

4.7.2 Strategy

We can divide the steps into two types, odd steps, and even steps.

Now let's think what will happen if a player A move  $x$  coins from an even step(non-zero) to an odd step. Player B can always move these same  $x$  coins to another even position and **the state of odd positions aren't affected**

But if player A moves a coin from an odd step to an even step, similar logic won't work. Due to the degenerated case, there is a situation when  $x$  coins are moved from stair 1 to 0, and player B can't move these coins from stair 0 to -1 (not a valid move).

From this argument, we can agree that coins in even steps are useless, they don't interfere to decide if a game state is winning or losing.

Therefore, the staircase nim can be visualized as a simple nim game with only the odd steps.

When stones are sent from an odd step to an even step, it is the same as removing stones from a pile in a classic nim game.

And when stones are sent from even steps to odd ones, it is the same as the increasing variation described before.

4.8 Grundy's Game

Initially there is only one pile with  $x$  stones. Each turn, a player must divide a pile into two non-zero piles with different sizes. The player who can't do any more moves loses.

4.8.1 Degenerate (Base) States

$x = 1$  (nim-val = 0) (losing)

$x = 2$  (nim-val = 0) (losing)

4.8.2 Other States

nim-val = MEX (all transitions)

Examples

**x = 3:**

```
{2, 1} -> {0} xor {0} -> 0
nim-val = MEX({0}) = 1
```

**x = 4:**

```
{3, 1} -> {1} xor {0} -> 1
nim-val = MEX({1}) = 0
```

**x = 5:**

```
{4, 1} -> {0} xor {0} -> 0
{3, 2} -> {1} xor {0} -> 1
nim-val = MEX({0, 1}) = 2
```

**x = 6:**

```
{5, 1} -> {2} xor {0} -> 2
{4, 2} -> {0} xor {0} -> 0
```

```
nim-val = MEX({0, 2}) = 1
```

**Important observation:** All numbers for ( $n \geq 2000$ ) are non-zero. (missing proof here and testing for values above  $1e6$ ).

4.9 Insta-Winning States

Classic nim game: if **all** piles become 0, you lose. (no more moves)

Modified nim game: if **any** pile becomes 0, you lose.

To adapt to this version of nim game, we create insta-winning states, which represents states that have a transition to any empty pile (will instantly win). Insta-winning states must have an specific number so they don't conflict with other nimbers when computing. A possible solution is nimber=INF, because no other nimber will be high enough to cause conflict.

Because of this adaptation, we can now ignore states with empty piles, and consider them with (*nullvalue* = -1). And the (*number* = 0) now represents the states that only have transitions to insta-winning states.

After this, beside winning states and losing states, we have added two new categories of states (insta-winning and empty-pile). Notice that:

```
empty-pile <- insta-winning <- number(0)
```

Therefore, we have returned to the classical nim game and can proceed normally.

OBS: *Empty piles* (wasn't empty before) (*number* = -1) is different from *Non-existent piles* (never existed) (number = 0)

Usage Example:  
<https://codeforces.com/gym/101908/problem/B>

4.10 References

[https://cp-algorithms.com/game\\_theory/sprague-grundy-nim.html](https://cp-algorithms.com/game_theory/sprague-grundy-nim.html)

<https://codeforces.com/blog/entry/66040>

<https://brilliant.org/wiki/nim/>

Geometry (5)

5.1 Point Struct

```
point.cpp
Description: Point struct for point operations, supports floating points and integers
Time:  $\mathcal{O}(1)$ 
Te11ab, 43 lines

const ld EPS = 1e-9;

// T can be int, long long, float, double, long double
```

```
template<class T> bool eq(T a, T b) {
    if (is_integral<T>::value) return a == b;
    else return abs(a-b) <= EPS;
}

template<class T> struct P {
    T x, y;
    ll id; // (optional)

    P(T xx=0, T yy=0): x(xx), y(yy) {}

    P operator +(P const& o) const { return { x+o.x, y+o.y }; }
    P operator -(P const& o) const { return { x-o.x, y-o.y }; }
    P operator *(T const& t) const { return { x*t, y*t }; }
    P operator /(T const& t) const { return { x/t, y/t }; }
    T operator *(P const& o) const { return x*o.x + y*o.y; }
    T operator ^(P const& o) const { return x*o.y - y*o.x; }

    bool operator <(P const& o) const { // enables sorting, set
        , etc
        return (eq(x, o.x) ? y < o.y : x < o.x);
    }

    bool operator ==(P const& o) const {
        return eq(x, o.x) and eq(y, o.y);
    }

    bool operator !=(P const& o) const {
        return !(*this == o);
    }

    friend istream& operator >>(istream& in, P &p) {
        return in >> p.x >> p.y;
    }

    friend ostream& operator <<(ostream& out, P const& p) {
        return out << p.x << ' ' << p.y;
    }
};
using point = P<ll>;
// using point = P<ld>;
```

5.2 Line Struct

line.cpp

Description: Line struct for line operations

Time:  $\mathcal{O}(1)$

5f33bd, 29 lines

```
template<class T> struct L {
    point p1, p2;
    T a, b, c; // ax+by+c = 0;

    // y-y1 = ((y2-y1)/(x2-x1))(x-x1)
    L(point pp1=0, point pp2=0) : p1(pp1), p2(pp2) {
        a = p1.y - p2.y;
        b = p2.x - p1.x;
        c = p1 ^ p2;
    }

    T eval(point p) {
        return a*p.x + b*p.y + c;
    }

    bool inside(point p) { // reta
        return eq(eval(p), T(0));
    }

    point normal() {
        return point(a, b);
    }
}
```

```
bool insideSeg(point p) { // segmento [p1, p2]
    return ( ((p1-p) ^ (p2-p)) == 0 and ((p1-p) * (p2-p))
        <= 0 );
}

};
using line = L<ll>;
// using line = L<ld>;

5.3 Manhattan Minimum Spanning Tree

Also called the rectilinear or L1 Minimum Spanning Tree
problem.

manhattanMST.cpp
Description: Given N points, returns up to 4*N edges, which are guaran-
teed to contain a minimum spanning tree for the graph with edge weights
w(p, q) = -p.x - q.x - + -p.y - q.y-. Edges are in the form (distance,
src, dst). Use a standard MST algorithm on the result to find the final MST.
Time:  $\mathcal{O}(N \log N)$ 
00e093, 25 lines

// requires point struct, at least the constructor and operator -

vector<array<ll, 3>> manhattanMST(vector<point> ps) {
    vector<ll> id(size(ps));
    iota(id.begin(), id.end(), 0);
    vector<array<ll, 3>> edges;
    for(ll k=0; k<4; k++) {
        sort(id.begin(), id.end(), [&](ll i, ll j) {
            return (ps[i]-ps[j]).x < (ps[j]-ps[i]).y;
        });

        map<ll, ll> sweep;
        for (ll i : id) {
            for (auto it = sweep.lower_bound(-ps[i].y); it !=
                sweep.end(); sweep.erase(it++)) {
                ll j = it->ss;
                point d = ps[i] - ps[j];
                if (d.y > d.x) break;
                edges.pb({d.y + d.x, i, j});
            }
            sweep[-ps[i].y] = i;
        }
        for (point& p : ps) if (k & 1) p.x = -p.x; else swap(p.
            x, p.y);
    }
    return edges;
}
```

Graph (6)

6.1 Fundamentals

dfs.cpp

Description: Simple DFS template for anonymous function

Time:  $\mathcal{O}(V + E)$

af867f, 19 lines

```
int32_t main(){ sws;
    // compute cardinality of each subtree
    vector<vll> g(n);
    vector<ll> card(n);
    vector<bool> vis(n); // redundant here

    function<ll (ll, ll)> dfs = [&](ll u, ll p) -> ll {
        if (vis[u]) return;
        vis[u] = 1;
```

```
card[u] += 1;
    for(auto v : g[u]) if (v != p) {
        card[u] += dfs(v, u);
    }
    return card[u];
};

dfs(1, -1);
}

bfs.cpp
Description: Simple BFS template
Time:  $\mathcal{O}(V + E)$ 
7bed46, 34 lines

vector<vll> g(n);
vector<ll> d(n);
vector<bool> vis(n);

void bfs(ll src, ll sink) {

    queue<ll> q;
    q.push(src);
    d[src] = 0;
    vis[src] = 1;

    while(!q.empty()) {
        auto u = q.front(); q.pop();

        // add here a special break condition if needed, ex:
        if (u == sink) break;

        for(auto v : g[u]) {

            // each v is added to queue only once
            // due to checking visited inside for(auto v : g[u]
                ])
            // and setting vis[v] = 1 before pushing to queue
            if (!vis[v]) {
                vis[v] = 1;
                d[v] = d[u] + 1;
                q.push(v);
            }

            else { // already added to queue, but there may be
                a shorter path
                d[v] = min(d[v], d[u] + 1);
            }
        }
    }
}
```

6.2 Network flow

In optimization theory, maximum flow problems involve finding a feasible flow through a flow network that obtains the maximum possible flow rate.

dinic.cpp

Description: Run several bfs to compute the residual graph until a max flow configuration is discovered

Time: General Case,  $\mathcal{O}(V^2E)$ ; Unit Capacity,  $\mathcal{O}\left((V + E)\sqrt{E}\right)$ ; Bipartite and unit capacity,  $\mathcal{O}\left((V + E)\sqrt{V}\right)$

dea1b7, 86 lines

```
// remember to duplicate vertices for the bipartite graph
// N = number of nodes, including sink and source
const ll N = 700;
```

```

struct Dinic {
    struct Edge {
        ll from, to, flow, cap;
    };
    vector<Edge> edges;

    vector<ll> g[N];
    ll ne = 0, lvl[N], vis[N], pass;
    ll qu[N], px[N], qt;

    ll run(ll s, ll sink, ll minE) {
        if (s == sink) return minE;
        ll ans = 0;
        for (; px[s] < (int)g[s].size(); px[s]++){
            ll e = g[s][px[s]];
            auto &v = edges[e], &rev = edges[e^1];
            if (lvl[v.to] != lvl[s]+1 || v.flow >= v.cap)
                continue;
            ll tmp = run(v.to, sink, min(minE, v.cap - v.flow));
            ;
            v.flow += tmp, rev.flow -= tmp;
            ans += tmp, minE -= tmp;
            if (minE == 0) break;
        }
        return ans;
    }

    bool bfs(ll source, ll sink) {
        qt = 0;
        qu[qt++] = source;
        lvl[source] = 1;
        vis[source] = ++pass;
        for (ll i=0; i<qt; i++) {
            ll u = qu[i];
            px[u] = 0;
            if (u == sink) return 1;
            for (auto& ed : g[u]) {
                auto v = edges[ed];
                if (v.flow >= v.cap || vis[v.to] == pass)
                    continue;
                vis[v.to] = pass;
                lvl[v.to] = lvl[u]+1;
                qu[qt++] = v.to;
            }
        }
        return false;
    }

    ll flow(ll source, ll sink) { // max_flow
        reset_flow();
        ll ans = 0;
        while (bfs(source, sink))
            ans += run(source, sink, LLINF);
        return ans;
    }

    void addEdge(ll u, ll v, ll c, ll rc = 0) { // c = capacity
        , rc = retro-capacity;
        Edge e = {u, v, 0, c};
        edges.pb(e);
        g[u].pb(ne++);
        e = {v, u, 0, rc};
        edges.pb(e);
        g[v].pb(ne++);
    }

    void reset_flow() {
        for (ll i=0; i<ne; i++) edges[i].flow = 0;
        memset(lvl, 0, sizeof(lvl));

```

```

        memset(vis, 0, sizeof(vis));
        memset(qu, 0, sizeof(qu));
        memset(px, 0, sizeof(px));
        qt = 0; pass = 0;
    }

    // cut set cost = minimum cost = max flow
    // cut set is the set of edges that, if removed,
    // will disrupt flow from source to sink and make it 0.
    vector<pll> cut() {
        vector<pll> cuts;
        for (auto [from, to, flow, cap]: edges)
            if (flow == cap and vis[from] == pass and vis[to] <
                pass and cap > 0)
                cuts.pb({from, to});
        return cuts;
    }
};

dinitz.cpp
Description: This second version may be slower due to dynamic allocation,
queue, etc but it's more readable, more memory efficient
Time: General Case,  $\mathcal{O}(V^2E)$ ; Unit Capacity,  $\mathcal{O}((V+E)\sqrt{E})$ ; Bipartite
and unit capacity,  $\mathcal{O}((V+E)\sqrt{V})$ 
49775d, 88 lines

struct Dinitz {
    struct Edge { // u -> v
        ll u, v, cap, flow=0; // u is redundant, but nice for
            some problems
    };

    vector<Edge> edges;
    vector<vector<ll>> g;
    vector<ll> level, ptr;

    // n need to be big enough for all nodes, including src/
    sink
    ll n, src, sink;
    Dinitz(ll nn, ll s = -1, ll t = -1) : n(nn+10) {
        src = (s == -1 ? n-2 : s);
        sink = (t == -1 ? n-1 : t);
        g.resize(n);
    }

    void addEdge(ll u, ll v, ll cap, ll rcap = 0) { // rcap =
        retrocapacity for bidirectional edges
        g[u].push_back( (ll)edges.size() );
        edges.push_back({u, v, cap});
        g[v].push_back( (ll)edges.size() );
        edges.push_back({v, u, rcap});
    }

    bool bfs() {
        level.assign(n, -1); // not vis
        level[src] = 0;
        queue<ll> q;
        q.push(src);
        while (!q.empty()) {
            ll u = q.front(); q.pop();
            for (auto eid : g[u]) {
                auto e = edges[eid];
                if (e.flow >= e.cap or level[e.v] != -1)
                    continue;
                level[e.v] = level[u] + 1;
                q.push(e.v);
            }
        }
        return level[sink] != -1;

```

```

    }

    ll dfs(ll u, ll f) {
        if (f == 0 or u == sink) return f;
        for (ll &i = ptr[u]; i < (ll)g[u].size(); i++) {
            ll eid = g[u][i];
            auto &e = edges[eid];
            if (e.flow >= e.cap or level[u]+1 != level[e.v])
                continue;
            ll newf = dfs(e.v, min(f, e.cap - e.flow));
            if (newf == 0) continue;
            e.flow += newf;
            edges[eid^1].flow -= newf;
            return newf;
        }
        return 0;
    }

    ll max_flow = 0;
    ll flow(bool reset_flow = true) {
        if (reset_flow) {
            max_flow = 0;
            for (ll u=0; u<n; u++) {
                for (auto eid : g[u]) {
                    auto &e = edges[eid];
                    e.flow = 0;
                }
            }
        }
        while (bfs()) {
            ptr.assign(n, 0);
            while (ll newf = dfs(src, INF))
                max_flow += newf;
        }
        return max_flow;
    }

    // minimum cut set cost = minimum cost = max flow
    // minimum cut set is the minimum set of edges that, if
    removed,
    // will disrupt flow from source to sink and make it 0.
    vector<pll> cut() {
        vector<pll> cuts;
        for (auto [u, v, cap, flow]: edges) {
            if (level[u] != -1 and level[v] == -1) {
                cuts.pb({u, v});
            }
        }
        return cuts;
    }
};

```

### 6.2.1 Matching with Flow

By modeling a bipartite graph, with some Vertices (that will choose a match) to be on the L graph and some Vertices (that will be chosen) on the R. Set the correct capacities for these edges and also of matching is the perfect matching, and source. After this involving from the bipartite graph and Algorithm, one will generate a possible matching (if it is possible).

A maximum matching has the maximum cardinality. A perfect matching is a maximum matching. But the opposite is not necessarily true.

It's possible to access `dinic.edges`, which is a vector that contains all edges and also its respective attributes, like the *flow* passing through each edge. Remember to consider that negative flow exist for reverse edges.

### 6.2.2 Minimum Cut

In computer science and optimization theory, the max-flow min-cut theorem states that, in a flow network, the maximum amount of flow passing from the source to the sink is equal to the total weight of the edges in a minimum cut, i.e., the smallest total weight of the edges which if removed would disconnect the source from the sink.

Let's define an s-t cut  $C = (S\text{-component}, T\text{-component})$  as a partition of  $V \in G$  such that source  $s \in S\text{-component}$  and sink  $t \in T\text{-component}$ . Let's also define a cut-set of  $C$  to be the set  $(u, v) \in E \mid u \in S\text{-component}, v \in T\text{-component}$  such that if all edges in the cut-set of  $C$  are removed, the Max Flow from  $s$  to  $t$  is 0 (i.e.,  $s$  and  $t$  are disconnected). The cost of an s-t cut  $C$  is defined by the sum of the capacities of the edges in the cut-set of  $C$ .

The by-product of computing Max Flow is Min Cut! After Max Flow algorithm stops, we run graph traversal (DFS/BFS) from source  $s$  again. All reachable vertices from source  $s$  using positive weighted edges in the residual graph belong to the S-component. All other unreachable vertices belong to the T-component. All edges connecting the S-component to the T-component belong to the cut-set of  $C$ . The Min Cut value is equal to the Max Flow value. This is the minimum over all possible s-t cuts values.

### 6.2.3 Minimum Vertex Cover

The **Konig's Theorem** describes an equivalence between the maximum matching problem and the minimum vertex cover problem in bipartite graphs.

Therefore, the value for the maximum flow in a bipartite graph is the same value as the number of nodes in a minimum vertex cover.

To retrieve the set of vertices of the minimum vertex cover:

- Give orientation to the edges, matched edges start from the right side of the graph to the left side, and free edges start from the left side of the graph to the right side.
- Run DFS from unmatched nodes of the left side, in this traversal some nodes will become visited, others will stay unvisited.

- The MVC nodes are the visited nodes from the right side, and unvisited nodes from the left side.

$$MVC = Visited_{Right} \cup Unvisited_{Left}$$

#### min-vertex-cover.cpp

**Description:** computes the min vertex cover for a bipartite graph matched with dinitz

**Time:**  $\mathcal{O}(E \log(E))$

963b5e, 55 lines

```
// a vertex cover is a set of vertices that contains
// at least one endpoint for each edge in the bipartite match
// A vertex cover in minimum if no other vertex cover has fewer
// vertices.
// only for bipartite graphs
vector<ll> minVertexCover(Dinitz &dinitz) {
    ll n = dinitz.n;

    vector<vector<ll>> g(n);
    set<ll> left, right; // unique
    vector<bool> matched(n);

    for(auto e : dinitz.edges) {
        if (e.u == dinitz.src or e.u == dinitz.sink) continue;
        if (e.v == dinitz.src or e.v == dinitz.sink) continue;
        if (e.cap > 0) { // not retro edge

            left.insert(e.u);
            right.insert(e.v);

            if (e.flow == e.cap) {
                // orient matched edges from right to left
                g[e.v].pb(e.u);
                matched[e.u] = 1;
                matched[e.v] = 1;
            }
            else {
                // orient non-matched edges from left to right
                g[e.u].pb(e.v);
            }
        }
    }

    vector<bool> vis(n, 0);
    function<void (ll)> dfs = [&](ll u) {
        vis[u] = 1;
        for(auto v : g[u])
            if (!vis[v])
                dfs(v);
    };

    for(auto l : left) if (!matched[l]) {
        dfs(l);
    }

    vector<ll> ans;
    for(auto l : left) if (!vis[l]) {
        ans.pb(l);
    }
    for(auto r : right) if (vis[r]) {
        ans.pb(r);
    }

    // remember, right nodes ids are dislocated by an offset
    return ans;
}
```

### 6.2.4 Maximum Independent Set

A **Independent Set** is a subset of nodes, in which all pairs  $u, v$  in the subset are not adjacent (There is no direct edge between nodes  $u$  and  $v$ ).

A **Maximum Independent Set** is a *Independent Set* with maximum cardinality;

The **Maximum Independent Set** is complementary to the **Minimum Vertex Cover**.

$$MaxIS = all_{Vertices} \setminus MVC$$

Therefore, to acquire the **Maximum Independent Set**, run the MVC algorithm and subtract them from the set of vertices and it will end up with the maxIS.

## 6.3 Minimum Cost Matching

### 6.3.1 Minimum Cost with Dinitz

#### min-cost-dinitz.cpp

**Description:** change bfs to spfa to attribute a weight for the edges

**Time:** SPFA is  $\mathcal{O}(E)$  at average and  $\mathcal{O}(VE)$  in the worst case

c51a5b, 88 lines

```
struct Dinitz {
    struct Edge { // u -> v
        ll u, v, cost, cap, flow=0;
    };

    vector<Edge> edges;
    vector<vector<ll>> g;
    vector<ll> dist, ptr; // uses dist instead of level

    // n need to be big enough for all nodes, including src/sink
    ll n, src, sink;
    Dinitz(ll nn, ll s = -1, ll t = -1) : n(nn+10) {
        src = (s == -1 ? n-2 : s);
        sink = (t == -1 ? n-1 : t);
        g.resize(n);
    }

    void addEdge(ll u, ll v, ll cost, ll cap, ll rcap = 0) { //
        rcap = retrocapacity for bidirectional edges
        g[u].push_back((ll)edges.size());
        edges.push_back({u, v, cost, cap});
        g[v].push_back((ll)edges.size());
        edges.push_back({v, u, -cost, rcap});
    }

    bool spfa() {
        dist.assign(n, INF);
        vector<bool> inqueue(n, false);

        queue<ll> q; q.push(src);
        dist[src] = 0;
        inqueue[src] = true;

        while (!q.empty()) {
            ll u = q.front(); q.pop();
            inqueue[u] = false;

            for (auto eid : g[u]) {
                auto const& e = edges[eid];
                if (e.flow >= e.cap) continue;
```

```

        if (dist[e.u] + e.cost < dist[e.v]) {
            dist[e.v] = dist[e.u] + e.cost;
            if (!inqueue[e.v]) {
                q.push(e.v);
                inqueue[e.v] = true;
            }
        }
    }
    return dist[sink] != INF;
}

11 min_cost = 0;
11 dfs(11 u, 11 f) {
    if (f == 0 or u == sink) return f;
    for (11 &i = ptr[u]; i < (11)g[u].size(); i++) {
        11 eid = g[u][i++];
        auto &e = edges[eid];
        if (e.flow >= e.cap or (dist[e.u] + e.cost) != dist[e.v]) continue;
        11 newf = dfs(e.v, min(f, e.cap - e.flow));
        if (newf == 0) continue;
        e.flow += newf;
        edges[eid^1].flow -= newf;
        min_cost += e.cost * newf;
        return newf;
    }
    return 0;
}

11 max_flow = 0;
pair<11, 11> flow(bool reset_flow_cost = true) {
    if (reset_flow_cost) {
        max_flow = 0;
        min_cost = 0;
        for (11 u=0; u<n; u++) {
            for(auto eid : g[u]) {
                auto &e = edges[eid];
                e.flow = 0;
            }
        }
    }
    while (spfa()) {
        ptr.assign(n, 0);
        while (11 newf = dfs(src, INF))
            max_flow += newf;
    }
    return {min_cost, max_flow};
}
};

```

There are  $n$  jobs and  $n$  workers. Each worker specifies the amount of money they expect for a particular job. Each worker can be assigned to only one job. The objective is to assign jobs to workers in a way that minimizes the total cost.

## hungarian dijkstra

Given an  $n \times n$  matrix  $A$ , the task is to select one number from each row such that exactly one number is chosen from each column, and the sum of the selected numbers is minimized.

Given an  $n \times n$  matrix  $A$ , the task is to find a permutation  $p$  of length  $n$  such that the value  $\sum A[i][p[i]]$  is minimized.

Consider a complete bipartite graph with  $n$  vertices per part, where each edge is assigned a weight. The objective is to find a perfect matching with the minimum total weight.

It is important to note that all the above scenarios are "square" problems, meaning both dimensions are always equal to  $n$ . In practice, similar "rectangular" formulations are often encountered, where  $n$  is not equal to  $m$ , and the task is to select  $\min(n, m)$  elements. However, it can be observed that a "rectangular" problem can always be transformed into a "square" problem by adding rows or columns with zero or infinite values, respectively.

We also note that by analogy with the search for a minimum solution, one can also pose the problem of finding a maximum solution. However, these two problems are equivalent to each other: it is enough to multiply all the weights by  $-1$ .

### hungarian.cpp

**Description:** Solves the assignment problem

**Time:**  $\mathcal{O}(n^3)$

06d970, 72 lines

```

// Hungaro
//
// Resolve o problema de assignment (matriz n x n)
// Colocar os valores da matriz em 'a' (pode < 0)
// assignment() retorna um par com o valor do
// assignment minimo, e a coluna escolhida por cada linha
// 0-idx
//
//  $\mathcal{O}(n^3)$ 

```

```

template<typename T> struct Hungarian {
    int n;
    vector<vector<T>>> a;
    vector<T> u, v;
    vector<int> p, way;
    T inf;

    Hungarian(int n_) : n(n_), u(n+1), v(n+1), p(n+1), way(n+1) {
        a = vector<vector<T>>>(n, vector<T>(n));
        inf = numeric_limits<T>::max();
    }

    pair<T, vector<int>>> assignment() {
        for (int i = 1; i <= n; i++) {
            p[0] = i;
            int j0 = 0;
            vector<T> minv(n+1, inf);
            vector<int> used(n+1, 0);
            do {
                used[j0] = true;
                int i0 = p[j0], j1 = -1;
                T delta = inf;
                for (int j = 1; j <= n; j++) if (!used[j]) {

```

```

                    T cur = a[i0-1][j-1] - u[i0] - v[j];
                    if (cur < minv[j]) minv[j] = cur, way[j] = j0;
                    if (minv[j] < delta) delta = minv[j], j1 = j;
                }
            } for (int j = 0; j <= n; j++)
                if (used[j]) u[p[j]] += delta, v[j] -= delta;
                else minv[j] -= delta;
            j0 = j1;
        } while (p[j0] != 0);
        do {
            int j1 = way[j0];
            p[j0] = p[j1];
            j0 = j1;
        } while (j0);
    }
    vector<int> ans(n);
    for (int j = 1; j <= n; j++) ans[p[j]-1] = j-1;
    return make_pair(-v[0], ans);
}
};

int32_t main(){ sws;
    11 n; cin >> n;
    Hungarian<11> h(n);

    for(11 i=0; i<n; i++) {
        for(11 j=0; j<n; j++) {
            cin >> h.a[i][j];
        }
    }

    auto [cost, match] = h.assignment();

    cout << cost << endl;

    for(11 i=0; i<n; i++) {
        cout << i+1 << " " << match[i]+1 << endl;
    }
}

```

## 6.4 Coloring

### 6.4.1 k-Coloring

**TODO:** Add this blog

<https://codeforces.com/blog/entry/57496>

[https://en.wikipedia.org/wiki/Graph\\_coloring](https://en.wikipedia.org/wiki/Graph_coloring)

<https://open.kattis.com/problems/coloring>

## 6.5 Shortest Paths

For weighted directed graphs

### 6.5.1 Dijkstra

Single Source and there **cannot** be any negative weighted edges.

#### dijkstra.cpp

**Description:** By keeping track of the distances sorted using a priority queue of candidates. if an edge can reduce the current min distance, insert into the priority queue. ONLY when the vertice is dequeued and its cost is  $\leq d[u]$ , it is in fact a part of a shortest path

**Time:**  $\mathcal{O}((V + E) \log V)$

76e3a5, 21 lines

```

priority_queue<p11, vector<p11>, greater<p11>> pq; // min pq
vector<vector<p11>>> g(MAX);
vector<11> d(MAX, INF);

void dijkstra(11 start){

```



```
    pq.push({0, start});
    d[start] = 0;

    while(!pq.empty()){

        auto [cost, u] = pq.top(); pq.pop();
        if (cost > d[u]) continue;

        for (auto [v, w] : g[u]) {
            if (d[u] + w < d[v]) {
                d[v] = d[u] + w;
                pq.push({d[v], v});
            }
        }
    }
}
```

By inverting the sorting order, Dijkstra can be modified for the opposite operation: *longest paths*.

Furthermore, Dijkstra be extended to keep track of more information, such as:

- how many minimum-price routes are there? (modulo  $10^9 + 7$ )
- what is the minimum number of flights in a minimum-price route?
- what is the maximum number of flights in a minimum-price route?

```
extendedDijkstra.cpp
Description: Also counts the numbers of shortest paths, the minimum and maximum number of edges transversed in any shortest path.
Time:  $\mathcal{O}((V + E) \log V)$ 
f93094, 32 lines

priority_queue<pll, vector<pll>, greater<pll>> pq; // min pq
vector<vector<pll>> g(MAX);
vector<ll> d(MAX, INF), ways(MAX, 0), mx(MAX, -INF), mn(MAX, INF);
// INF = INT64_MAX

void dijkstra(ll start){
    pq.push({0, start});
    ways[start] = 1;
    d[start] = mn[start] = mx[start] = 0;

    while( !pq.empty() ){
        auto [p1, u] = pq.top(); pq.pop();
        if (p1 > d[u]) continue;
        for(auto [v, p2] : g[u]){
            // reset info, shorter path found, previous ones are discarded
            if (d[u] + p2 < d[v]){
                d[v] = d[u] + p2;
                ways[v] = ways[u];
                mx[v] = mx[u]+1;
                mn[v] = mn[u]+1;

                pq.push({d[v], v});
            }
            // same distance, different path, update info
            else if (d[u] + p2 == d[v]) {
                ways[v] = (ways[v] + ways[u]) % MOD;
                mn[v] = min(mn[v], mn[u]+1);
                mx[v] = max(mx[v], mx[u]+1);
            }
        }
    }
}
```

```
    }
    }
}

6.5.2 Bellman-Ford
Single Source and it supports negative edges

Conjecture: After at most n-1 (Vertices-1) iterations, all shortest paths will be found.

bellman-ford.cpp
Description: n-1 iterations is sufficient to find all shortest paths
Time:  $\mathcal{O}(V * E) \rightarrow \mathcal{O}(N^2)$ 
d749f1, 15 lines

using T = array<ll, 3>;
vector<T> edges;
vector<ll> d(MAX, INF);
// INF = 0x3f3f3f3f3f3f3f3f, to avoid overflow

void BellmanFord(ll src, ll n) {
    d[src] = 0;
    for(ll i=0; i<n-1; i++) { // n-1 iterations
        for(auto [u, v, w] : edges) {
            if (d[u] + w < d[v]) {
                d[v] = d[u] + w;
            }
        }
    }
}

By iterating once more, one can check if the last iteration reduced once again any distance. If so, it means that there must be a negative cycle, because the shortest distance should have been found before elseway.

To retrieve the negative cycle itself, one can keep track of the last vertice that reaches a considered vertice

bellman-ford-cycle.cpp
Description: By using the property that n-1 iterations is sufficient to find all shortest paths in a graph that doesn't have negative cycles. Iterate n times and retrieve the path using a vector of parents
Time:  $\mathcal{O}(V * E) \rightarrow \mathcal{O}(N^2)$ 
0506b5, 35 lines

using T = array<ll, 3>;
vector<T> edges;
vector<ll> d(MAX, INF), p(MAX, -1);
vector<ll> cycle;
// INF = 0x3f3f3f3f3f3f3f3f, to avoid overflow

void BellmanFordCycle(ll src, ll n) {
    d[src] = 0;
    ll x = -1; // possible node inside a negative cycle
    for(ll i=0; i<n; i++) { // n iterations
        x = -1;
        for(auto [u, v, w] : edges) {
            if (d[u] + w < d[v]) {
                d[v] = d[u] + w;
                p[v] = u;
                x = v;
            }
        }
    }

    if (x != -1) {
        // set x to a node, contained in a cycle in p[]
    }
}
```

```
for(ll i=0; i<n; i++) x = p[x];

ll tmp = x;
do {
    cycle.pb(tmp);
    tmp = p[tmp];
}
while (tmp != x);
cycle.pb(x);

reverse(cycle.begin(), cycle.end());
}
}
```

```
6.5.3 Floyd Warshall
All-Pair Shortest Paths

floyd-warshall.cpp
Description: By using an auxiliar vertice, check if a smaller path exists between a pair (u, v) of vertices, if so, update minimum distance.
Time:  $\mathcal{O}(V^3)$ 
fa5f60, 15 lines

// N < sqrt(1e8) = 460
ll N = 200;

// d[u][v] = INF (no edge)
vector<vll> d(N+1, vll(N+1, INF));

void floydWarshall() { // O(N^3)
    for(ll i=1; i<=N; i++) d[i][i] = 0;

    for(ll aux=1; aux<=N; aux++)
        for(ll u=1; u<=N; u++)
            for(ll v=1; v<=N; v++)
                if (d[u][aux] < INF and d[v][aux] < INF)
                    d[u][v] = min(d[u][v], d[u][aux] + d[v][aux]);
}
}
```

6.6 Undirected Graph

Bridges and Articulation Points are concepts for undirected graphs!

6.6.1 Bridges (Cut Edges)

Also called **isthmus** or **cut arc**.

A back-edge is never a bridge!

A **lowlink** for a vertice *U* is the closest vertice to the root reachable using only span edges and a *single* back-edge, starting in the subtree of *U*.

After constructing a DFS Tree, an edge (u, v) is a bridge  $\iff$  there is no back-edge from *v* (or a descendent of *v*) to *u* (or an ancestor of *u*)

To do this efficiently, it's used *tin*[*i*] (entry time of node *i*) and *low*[*i*] (minimum entry time considering all nodes that can be reached from node *i*).

In another words, a edge (u, v) is a bridge  $\iff$  the *low*[*v*]  $\geq$  *tin*[*u*].



bridges.cpp  
**Description:** Using the concepts of entry time (tin) and lowlink (low), an edge is a bridge if, and only if,  $\text{low}[v] > \text{tin}[u]$   
**Time:**  $\mathcal{O}(V + E)$

87e0d3, 25 lines

```
vector<vll> g(MAX);
ll timer = 1;
ll tin[MAX], low[MAX];
vector<pll> bridges;

void dfs(ll u, ll p = -1){
    tin[u] = low[u] = timer++;
    for(auto v : g[u]) if (v != p) {
        if (tin[v]) // v was visited ({u,v} is a back-edge)
            // considering a single back-edge:
            low[u] = min(low[u], tin[v]);
        else { // v wasn't visited ({u, v} is a span-edge)
            dfs(v, u);
            // after low[v] was computed by dfs(v, u):
            low[u] = min(low[u], low[v]);
            if (low[v] > tin[u])
                bridges.pb({u, v});
        }
    }
}

void findBridges(ll n) {
    for(ll i=1; i<=n; i++) if (!tin[i])
        dfs(i);
}
```

6.6.2 Bridge Tree

After merging *vertices* of a **2-edge connected component** into single vertices, and leaving only bridges, one can generate a Bridge Tree.

Every **2-edge connected component** has following properties:

- For each pair of vertices A, B inside the same component, there are at least 2 distinct paths from A to B (which may repeat vertices).

bridgeTree.cpp  
**Description:** After finding bridges, set an component id for each vertice, then merge vertices that are in the same 2-edge connected component  
**Time:**  $\mathcal{O}(V + E)$

6d00bd, 47 lines

```
// g: u -> {v, edge id}
vector<vector<pll>> g(MAX);
vector<vll> gc(MAX);
ll timer = 1;
ll tin[MAX], low[MAX], comp[MAX];
bool isBridge[MAX];

void dfs(ll u, ll p = -1) {
    tin[u] = low[u] = timer++;
    for(auto [v, id] : g[u]) if (v != p) {
        if (tin[v])
            low[u] = min(low[u], tin[v]);
        else {
            dfs(v, u);
            low[u] = min(low[u], low[v]);
            if (low[v] > tin[u])
                isBridge[id] = 1;
        }
    }
}
```

```
}

void dfs2(ll u, ll c, ll p = -1) {
    comp[u] = c;
    for(auto [v, id] : g[u]) if (v != p) {
        if (isBridge[id]) continue;
        if (!comp[v]) dfs2(v, c, u);
    }
}

void bridgeTree(ll n) {
    // find bridges
    for(ll i=1; i<=n; i++) if (!tin[i])
        dfs(i);

    // find components
    for(ll i=1; i<=n; i++) if (!comp[i])
        dfs2(i, i);

    // condensate into a TREE (or TREES if disconnected)
    for(ll u=1; u<=n; u++) {
        for(auto [v, id] : g[u]) {
            if (comp[u] != comp[v]) {
                gc[comp[u]].pb(comp[v]);
            }
        }
    }
}
```

6.6.3 Articulation Points

One Vertice in a graph is considered a Articulation Points or Cut Vertice if its removal in the graph will generate more disconnected components

articulation.cpp  
**Description:** if  $\text{low}[v] \geq \text{tin}[u]$ , u is an articulation points The root is a corner case  
**Time:**  $\mathcal{O}(V + E)$

8707a0, 29 lines

```
vector<vll> g(MAX);
ll timer = 1;
ll low[MAX], tin[MAX], isAP[MAX];
// when vertex i is removed from graph
// isAP[i] is the quantity of new disjoint components created
// isAP[i] >= 1 {i is a Articulation Point}
void dfs(ll u, ll p = -1) {
    low[u] = tin[u] = timer++;

    for(auto v : g[u]) if (v != p) {
        if (tin[v]) // visited
            low[u] = min(low[u], tin[v]);
        else { // not visited
            dfs(v, u);
            low[u] = min(low[u], low[v]);

            if (low[v] >= tin[u])
                isAP[u]++;
        }
    }

    // corner case: root
    if (p == -1 and isAP[u]) isAP[u]--;
}

void findAP(ll n) {
    for(ll i=1; i<=n; i++) if (!tin[i])
        dfs(i);
}
```

6.6.4 Block Cut Tree

After merging *edges* of a **2-vertex connected component** into single vertices, one can obtain a block cut tree.

2-vertex connected components are also called as biconnected component

Every bridge by itself is a biconnected component

Each edge in the block-cut tree connects exactly an Articulation Point and a biconnected component (bipartite graph)

Each biconnected component has the following properties:

- For each pair of edges, there is a cycle that contains both edges
- For each pair of vertices A, B inside the same connected component, there are at least 2 distinct paths from A to B (which do not repeat vertices).

blockCutTree.cpp  
**Description:** After Merging 2-Vertex Connected Components, one can generate a block cut tree  
**Time:**  $\mathcal{O}(V + E)$

f752d5, 100 lines

```
// Block-Cut Tree (bruno monteiro)
//
// Cria a block-cut tree, uma arvore com os blocos
// e os pontos de articulacao
// Blocos sao as componentes 2-vertice-conexas maximais
// Uma 2-coloracao da arvore eh tal que uma cor sao
// os componentes, e a outra cor sao os pontos de articulacao
//
// Funciona para grafo nao conexo
//
// isAP[i] responde o numero de novas componentes conexas
// criadas apos a remocao de i do grafo g
// Se isAP[i] >= 1, i eh ponto de articulacao
//
// Para todo i < blocks.size()
// blocks[i] eh uma componente 2-vertce-conexa maximal
// blockEdges[i] sao as arestas do bloco i
//
// tree eh a arvore block-cut-tree
// tree[i] eh um vertice da arvore que corresponde ao bloco i
//
// comp[i] responde a qual vertice da arvore vertice i pertence
//
// Arvore tem no maximo 2n vertices
//
// O(n+m)

// 0-idx graph!!!
vector<vll> g(MAX), tree, blocks;
vector<vector<pll>> blockEdges;
stack<ll> st; // st for vertices,
stack<pll> st2; // st2 for edges
vector<ll> low, tin, comp, isAP;
ll timer = 1;

void dfs(ll u, ll p = -1) {
    low[u] = tin[u] = timer++;

    for(auto [v, id] : g[u]) if (v != p) {
        if (tin[v])
            low[u] = min(low[u], tin[v]);
        else {
            dfs(v, u);
            low[u] = min(low[u], low[v]);
            if (low[v] >= tin[u])
                isAP[u]++;
            if (low[v] > tin[u])
                blockEdges[id].pb({u, v});
        }
    }

    st.push(u);
}
```

```

// add only back-edges to stack
if (p != -1) st2.push({u, p});
for(auto v : g[u]) if (v != p) {
    if (tin[v] != -1) // visited
        st2.push({u, v});
}

for(auto v : g[u]) if (v != p) {
    if (tin[v] != -1) // visited
        low[u] = min(low[u], tin[v]);
    else { // not visited
        dfs(v, u);
        low[u] = min(low[u], low[v]);

        if (low[v] >= tin[u]) {
            isAP[u] += 1;

            blocks.pb(vll(1, u));
            while(blocks.back().back() != v)
                blocks.back().pb(st2.top()), st2.pop();

            blockEdges.pb(vector<pll>(1, st2.top())), st2.pop();
            while(blockEdges.back().back() != pair<ll, ll>(v, u))
                blockEdges.back().pb(st2.top()), st2.pop();
        }
    }
}

// corner case: root
if (p == -1 and isAP[u]) isAP[u]--;
}

void blockCutTree(ll n) {

    // initialize vectors and reset
    tree.clear(), blocks.clear(), blockEdges.clear();
    st = stack<ll>(), st2 = stack<pll>();
    tin.assign(n, -1);
    low.assign(n, 0), comp.assign(n, 0), isAP.assign(n, 0);
    timer = 1;

    // find Articulation Points
    for(ll i=0; i<n; i++) if (tin[i] == -1)
        dfs(i);

    // set component id for APs
    tree.assign(blocks.size(), vll());
    for(ll i=0; i<n; i++) if (isAP[i])
        comp[i] = tree.size(), tree.pb(vll());

    // set component id for non-APs and construct tree
    for(ll u=0; u<(ll)blocks.size(); u++) {
        for(auto v : blocks[u]) {
            if (!isAP[v])
                comp[v] = u;
            else
                tree[u].pb(comp[v]), tree[comp[v]].pb(u);
        }
    }
}

```

### 6.6.5 Strong Orientation

A **strong orientation** of an undirected graph is an assignment of a direction to each edge that makes it a strongly connected graph. That is, after the orientation we should be able to visit any vertex from any vertex by following the directed edges.

Of course, this cannot be done to every graph. Consider a **bridge** in a graph. We have to assign a direction to it and by doing so we make this bridge "crossable" in only one direction. That means we can't go from one of the bridge's ends to the other, so we can't make the graph strongly connected.

Now consider a DFS through a bridgeless connected graph. Clearly, we will visit each vertex. And since there are no bridges, we can remove any DFS tree edge and still be able to go from below the edge to above the edge by using a path that contains at least one back edge. From this follows that from any vertex we can go to the root of the DFS tree. Also, from the root of the DFS tree we can visit any vertex we choose. We found a strong orientation!

In other words, to strongly orient a bridgeless connected graph, run a DFS on it and let the DFS tree edges point away from the DFS root and all other edges from the descendant to the ancestor in the DFS tree.

The result that bridgeless connected graphs are exactly the graphs that have strong orientations is called **Robbins' theorem**.

### Acyclic Graph Orientation

**Problem:** Given an undirected graph, your task is to choose a direction for each edge so that the resulting directed graph is acyclic.

**Solution:** Do a dfs tree, every span-edge is oriented according to the dfs transversal, and every back-edge is oriented contrary to the dfs transversal

### 6.6.6 Minimum Spanning Tree

A minimum spanning tree (MST) or minimum weight spanning tree is a subset of the edges of a connected, edge-weighted undirected graph that connects all the vertices together, without any cycles and with the minimum possible total edge weight. That is, it is a spanning tree whose sum of edge weights is as small as possible.

#### kruskal.cpp

**Description:** Sort all edges in crescent order by weight, include all edges which joins two disconnected trees. In case of tie, choose whichever. Dont include edges that will join a already connected part of the tree.  
**Time:**  $\mathcal{O}(E \log E\alpha)$

```

// use DSU struct
struct DSU{};

```

```

set<array<ll, 3>> edges;

int32_t main(){ sws;
    ll n, m; cin >> n >> m;
    DSU dsu(n+1);
    for(ll i=0; i<m; i++) {
        ll u, v, w; cin >> u >> v >> w;
        edges.insert({w, u, v});
    }
    ll minCost = 0;
    for(auto [w, u, v] : edges) {
        if (dsu.find(u) != dsu.find(v)) {
            dsu.join(u, v);
            minCost += w;
        }
    }
    cout << minCost << endl;
}

```

## 6.7 Directed Graph

### 6.7.1 Topological Sort

Sort a directed graph with no cycles (DAG) in an order which each source of an edge is visited before the sink of this edge.

Cannot have cycles, because it would create a contradiction of which vertices whould come before.

It can be done with a DFS, appending in the reverse order of transversal. Also a stack can be used to reverse order

#### toposort.cpp

**Description:** Using DFS pos order transversal and inverting the order, one can obtain the topological order

**Time:**  $\mathcal{O}(V + E)$  75f781, 17 lines

```

vector<vll> g(MAX, vll());
vector<bool> vis;
vll topological;

void dfs(ll u) {
    vis[u] = 1;
    for(auto v : g[u]) if (!vis[v]) dfs(v);
    topological.pb(u);
}

```

```

// 1 - indexed
void topological_sort(ll n) {
    vis.assign(n+1, 0);
    topological.clear();
    for(ll i=1; i<=n; i++) if (!vis[i]) dfs(i);
    reverse(topological.begin(), topological.end());
}

```

### 6.7.2 Kosaraju

A Strongly Connected Component is a maximal subgraph in which every vertex is reachable from any vertex inside this same subgraph.

A important *property* is that the inverted graph or transposed graph has the same SCCs as the original graph.

kosaraju.cpp

**Description:** By using the fact that the inverted graph has the same SCCs, just do a DFS twice to find all SCCs. A condensated graph can be created if wished. The condensated graph is a DAG!!  
**Time:**  $\mathcal{O}(V + E)$

381904, 45 lines

```
struct Kosaraju {
    ll n;
    vector<vll> g, gi, gc;
    vector<bool> vis;
    vector<ll> comp;
    stack<ll, vll> st;

    void dfs(ll u) { // g
        vis[u] = 1;
        for(auto v : g[u]) if (!vis[v]) dfs(v);
        st.push(u);
    }

    void dfs2(ll u, ll c) { // gi
        comp[u] = c;
        for(auto v : gi[u]) if (comp[v] == -1) dfs2(v, c);
    }

    Kosaraju(vector<vll> &g_)
        : g(g_), n(g_.size()-1) { // 1-idx

        gi.assign(n+1, vll());
        for(ll i=1; i<=n; i++) {
            for(auto j : g[i])
                gi[j].pb(i);
        }

        gc.assign(n+1, vll());
        vis.assign(n+1, 0);
        comp.assign(n+1, -1);
        st = stack<ll, vll>();

        for(ll i=1; i<=n; i++) if (!vis[i]) dfs(i);

        while(!st.empty()) {
            auto u = st.top(); st.pop();
            if (comp[u] == -1) dfs2(u, u);
        }

        for(ll i=1; i<=n; i++)
            for(auto v : g[u])
                gc[comp[u]].pb(comp[v]);
    }
};
```

6.7.3 2-SAT

SAT (Boolean satisfiability problem) is NP-Complete.

2-SAT is a restriction of the SAT problem, in 2-SAT every clause has exactly two variables:  $(X_1 \vee X_2) \wedge (X_2 \vee X_3)$

Every restriction or implication are represented in the graph as directed edges.

The algorithm uses kosaraju to check if any  $(X$  and  $\neg X)$  are in the same Strongly Connected Component (which implies that the problem is impossible).

If it doesn't, there is at least one solution, which can be generated using the topological sort of the same kosaraju (opting for the variables that appers latter in the sorted order)

2sat.cpp

**Description:** Kosaraju to find if there are SCCs. If there are not cycles, use toposort to choose states  
**Time:**  $\mathcal{O}(V + E)$

87417c, 83 lines

```
// 0-idx graph !!!!
struct TwoSat {
    ll N; // needs to be the twice of the number of variables
    // node with idx 2x => variable x
    // node with idx 2x+1 => variable !x

    vector<vll> g, gi;
    // g = graph; gi = transposed graph (all edges are inverted)

    TwoSat(ll n) { // number of variables (add +1 faor 1-idx)
        N = 2*n;
        g.assign(N, vll());
        gi.assign(N, vll());
    }

    ll idx; // component idx
    vector<ll> comp, order; // topological order (reversed)
    vector<bool> vis, chosen;
    // chosen[x] == 0 -> x was assigned
    // chosen[x] == 1 -> !x was assigned

    // dfs and dfs2 are part of kosaraju algorithm
    void dfs(ll u) {
        vis[u] = 1;
        for (ll v : g[u]) if (!vis[v]) dfs(v);
        order.pb(u);
    }

    void dfs2(ll u, ll c) {
        comp[u] = c;
        for (ll v : gi[u]) if (comp[v] == -1) dfs2(v, c);
    }

    bool solve() {
        vis.assign(N, 0);
        order = vector<ll>();
        for (ll i = 0; i < N; i++) if (!vis[i]) dfs(i);

        comp.assign(N, -1); // comp = 0 can exist
        idx = 1;
        for(ll i=(ll)order.size()-1; i>=0; i--) {
            ll u = order[i];
            if (comp[u] == -1) dfs2(u, idx++);
        }

        chosen.assign(N/2, 0);
        for (ll i = 0; i < N; i += 2) {
            // x and !x in the same component => contradiction
            if (comp[i] == comp[i+1]) return false;
            chosen[i/2] = comp[i] < comp[i+1]; // choose latter node
        }
        return true;
    }

    // a (with flagA) implies => b (with flagB)
    void add(ll a, bool fa, ll b, bool fb) {
        // {fa == 0} => a
        // {fa == 1} => !a
    }
};
```

```
a = 2*a + fa;
b = 2*b + fb;
g[a].pb(b);
gi[b].pb(a);
}

// force a state for a certain variable (must be true)
void force(ll a, bool fa) {
    add(a, fa^1, a, fa);
}

// xor operation: one must exist, and only one can exist
void exclusive(ll a, bool fa, ll b, bool fb) {
    add(a, fa^0, b, fb^1);
    add(a, fa^1, b, fb^0);
    add(b, fb^0, a, fa^1);
    add(b, fb^1, a, fa^0);
}

// nand operation: no more than one can exist
void nand(ll a, bool fa, ll b, bool fb) {
    add(a, fa^0, b, fb^1);
    add(b, fb^0, a, fa^1);
}

};
```

6.8 Trees

lca.cpp

**Description:** Solves LCA for trees  
**Time:**  $\mathcal{O}(N \log(N))$  to build,  $\mathcal{O}(\log(N))$  per query

7afc1a, 54 lines

```
struct BinaryLifting {
    ll n, logN = 20; // ~1e6
    vector<vll> g;
    vector<ll> depth;
    vector<vll> up;

    BinaryLifting(vector<vll> &g_)
        : g(g_), n(g_.size() + 1) { // 1-idx
            depth.assign(n, 0);

            while((1 << logN) < n) logN++;
            up.assign(n, vll(logN, 0));
            build();
        }

    void build(ll u = 1, ll p = -1) {
        for(ll i=1; i<logN; i++) {
            up[u][i] = up[ up[u][i-1] ][i-1];
        }

        for(auto v : g[u]) if (v != p) {
            up[v][0] = u;
            depth[v] = depth[u] + 1;
            build(v, u);
        }
    }

    ll go(ll u, ll dist) { // O(log(n))
        for(ll i=logN-1; i>=0; i--) { // bigger jumps first
            if (dist & (1LL << i)) {
                u = up[u][i];
            }
        }
        return u;
    }

    ll lca(ll a, ll b) { // O(log(n))
        if (depth[a] < depth[b]) swap(a, b);
```

```
a = go(a, depth[a] - depth[b]);
if (a == b) return a;

for(ll i=logN-1; i>=0; i--) {
    if (up[a][i] != up[b][i]) {
        a = up[a][i];
        b = up[b][i];
    }
}
return up[a][0];
}

ll lca(ll a, ll b, ll root) { // lca(a, b) when tree is
    rooted at 'root'
    return lca(a, b)^lca(b, root)^lca(a, root); //magic
}

};
```

queryTree.cpp  
**Description:** Binary Lifting for min, max weight present in a simple path  
**Time:**  $\mathcal{O}(N \log(N))$  to build;  $\mathcal{O}(\log(N))$  per query

75ba37, 67 lines

```
struct BinaryLifting {
    ll n, logN = 20; // ~1e6
    vector<vp11> g;
    vector<ll> depth;
    vector<v11> up, mx, mn;

    BinaryLifting(vector<vp11> &g_)
        : g(g_), n(g_.size() + 1) { // 1-idx
            depth.assign(n, 0);

            while((1 << logN) < n) logN++;
            up.assign(n, v11(logN, 0));
            mx.assign(n, v11(logN, -INF));
            mn.assign(n, v11(logN, INF));
            build();
        }

    void build(ll u = 1, ll p = -1) {

        for(ll i=1; i<logN; i++) {
            mx[u][i] = max(mx[u][i-1], mx[ up[u][i-1] ][i-1]);
            mn[u][i] = min(mn[u][i-1], mn[ up[u][i-1] ][i-1]);
            up[u][i] = up[ up[u][i-1] ][i-1];
        }

        for(auto [v, w] : g[u]) if (v != p) {
            mx[v][0] = mn[v][0] = w;
            up[v][0] = u;
            depth[v] = depth[u] + 1;
            build(v, u);
        }
    }

    array<ll, 3> go(ll u, ll dist) { //  $\mathcal{O}(\log(n))$ 
        ll mxval = -INF, mnval = INF;
        for(ll i=logN-1; i>=0; i--) { // bigger jumps first
            if (dist & (1LL << i)) {
                mxval = max(mxval, mx[u][i]);
                mnval = min(mnval, mn[u][i]);
                u = up[u][i];
            }
        }
        return {u, mxval, mnval};
    }

    array<ll, 3> query(ll u, ll v) { //  $\mathcal{O}(\log(n))$ 
        if (depth[u] < depth[v]) swap(u, v);
```

```
auto [a, mxval, mnval] = go(u, depth[u] - depth[v]);
ll b = v;

if (a == b) return {a, mxval, mnval};

for(ll i=logN-1; i>=0; i--) {
    if (up[a][i] != up[b][i]) {
        mxval = max({mxval, mx[a][i], mx[b][i]});
        mnval = min({mnval, mn[a][i], mn[b][i]});
        a = up[a][i];
        b = up[b][i];
    }
}

mxval = max({mxval, mx[a][0], mx[b][0]});
mnval = min({mnval, mn[a][0], mn[b][0]});
return {up[a][0], mxval, mnval};
}

};
```

Mathematics (7)

7.1 Modular Arithmetic

modular.cpp  
**Description:** mint struct for modular arithmetic operations  
**Time:**  $\mathcal{O}(1)$  for most operations,  $\mathcal{O}(\log(n))$  for division and exponentiation

```
// supports operations between int/ll and mint,
// and it will return a mint object independently of the order
// of operations
template <ll P> struct Z {
    ll val;

    Z(ll a = 0) {
        val = a % P;
        if (val < 0) val += P;
    }

    Z& operator +=(Z rhs) {
        val += rhs.val;
        if (val >= P) val -= P;
        return *this;
    }

    friend Z operator +(Z lhs, Z rhs) { return lhs += rhs; }

    Z& operator --(Z rhs) {
        val += P - rhs.val;
        if (val >= P) val -= P;
        return *this;
    }

    friend Z operator -(Z lhs, Z rhs) { return lhs -= rhs; }

    Z& operator *=(Z rhs) {
        val = (val * rhs.val) % P;
        return *this;
    }

    friend Z operator *(Z lhs, Z rhs) { return lhs *= rhs; }

    Z fexp(Z x, ll i) {
        if (i == 0) return 1;
        if (i == 1) return x;
        Z m = fexp(x, i/2);
        m *= m;
        if (i & 1) return m * x;
        else return m;
    }
};
```

```
Z& operator /=(Z rhs) {
    return *this *= fexp(rhs, P-2);
}

friend Z operator /(Z lhs, Z rhs) { return lhs /= rhs; }

bool operator ==(Z rhs) { return val == rhs.val; }

bool operator !=(Z rhs) { return val != rhs.val; }

friend ostream& operator <<(ostream& out, Z a) { return out
    << a.val; }

friend istream& operator >>(istream& in, Z& a) {
    ll x; in >> x;
    a = Z(x);
    return in;
}

};

using mint = Z<MOD>;
```

7.1.1 Lucas's Theorem

$$\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$$

For  $p$  prime.  $n_i$  and  $m_i$  are the coefficients of the representations of  $n$  and  $m$  in base  $p$ .

Example:

$11 \text{ (in base } p=3) = 1 \cdot 3^2 + 0 \cdot 3^1 + 2 \cdot 3^0$   
 $\implies n_2 = 1, n_1 = 0, n_0 = 2$

7.2 Combinatorics

$$\binom{n}{m} = \frac{n!}{m! \cdot (n-m)!}, \quad 0 \leq m \leq n$$
  
$$0, \quad \text{otherwise}$$

7.2.1 Factorial

$n$	1	2	3	4	5	6	7	8	9	10
$n!$	1	2	6	24	120	720	5040	40320	362880	3628800
$n$	11	12	13	14	15	16	17			
$n!$	4.0e7	4.8e8	6.2e9	8.7e10	1.3e12	2.1e13	3.6e14			
$n$	20	25	30	40	50	100	150	171		
$n!$	2e18	2e25	3e32	8e47	3e64	9e157	6e262	>DBL_MAX		

7.2.2 Combinatorial Struct

```
combinatorics.cpp
Description: basic operations for combinatorics problems under a certain modulo
Time:  $\mathcal{O}(n)$  to construct,  $\mathcal{O}(1)$  operations

// remeber to import mint struct !!
struct Combinatorics {
    vector<mint> fact, ifact;

    Combinatorics(ll n) : fact(n+1), ifact(n+1) { // inclusive
        n
```

```

fact[0] = 1;
for (ll i=1; i<=n; i++) fact[i] = fact[i-1] * i;

ifact[n] = 1 / fact[n];
for (ll i=n; i>0; i--) ifact[i-1] = ifact[i] * i;
}

// "Combinacao / Binomio de Newton"
// n objects to place in k spaces
// the order doesn't matter, so we consider the re-
// orderings
// = n! / (k! * (n-k)!)
mint combination(ll n, ll k) {
    if (k < 0 or n < k) return 0;
    return fact[n] * ifact[k] * ifact[n-k];
}

// "Permutacao"
// n objects to place in n spaces
// = n!
mint permutation(ll n) {
    if (n < 0) return 0;
    return fact[n];
}

// "Permutacao com repeticao"
// n objects to place in n spaces
// some objects are equal
// therefore, we consider the possible re-orderings
// = n! / (k1! k2! k3!)
mint permutationRepetition(ll n, vector<ll> vec) {
    if (n < 0) return 0;
    mint ans = fact[n];
    for(auto val : vec) ans /= ifact[val];
    return ans;
}

// "Arranjo Simples"
// n objects to place in k spaces (k < n)
// n * (n-1) * ... * (n-k+1)
// = n! / (n-k)!
mint arrangement(ll n, ll k) {
    if (n < 0) return 0;
    return fact[n] * ifact[n-k];
}

// "Pontos e Virgulas"
// n stars to distribute among
// k distinct groups, that can contain 0, 1 or more stars
// separated by k-1 bars
// = (n+k-1)! / (n! * (k-1)!)
mint starsBars(ll n, ll k) {
    if (k == 0) {
        if (n == 0) return 1;
        else return 0;
    }
    return combination(n + k - 1, k - 1);
}

// a derangement is a permutation of the elements of a set
// in which no element appears in its original position
// In other words, a derangement is a permutation that has
// no fixed points.
// derangement(n) = subfactorial(n) = !n
// !n = (n-1) * ( !(n-1) + !(n-2) ), for n >= 2
// !1 = 0, !0 = 1
vector<mint> subfact;
void computeSubfactorials(ll n) {
    subfact.assign(n+1, 0);

```

```

subfact[0] = 1;
subfact[1] = 0;

for(ll i=2; i<=n; i++) {
    subfact[i] = (i-1) * (subfact[i-1] + subfact[i-2]);
}

// remeber to compute subfactorials first !!
mint derangement(ll n) {
    if (n < 0) return 0;
    return subfact[n];
}

};
Combinatorics op(MAX); // MAX = inclusive max_value for fact[]

```

### 7.2.3 Burside Lemma

Let  $G$  be a group that acts on a set  $X$ . The Burnside Lemma states that the number of distinct orbits is equal to the average number of points fixed by an element of  $G$ .

$$T = \frac{1}{|G|} \sum_{g \in G} |\text{fix}(g)|$$

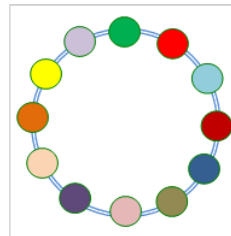
Where a orbit  $\text{orb}(x)$  is defined as

$$\text{orb}(x) = \{y \in X : \exists g \in G \text{ } gx = y\}$$

and  $\text{fix}(g)$  is the set of elements in  $X$  fixed by  $g$

$$\text{fix}(g) = \{x \in X : gx = x\}$$

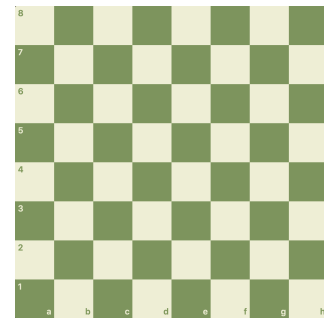
**Example1:** With  $k$  distinct types of beads how many distinct necklaces of size  $n$  can be made? Considering that two necklaces are equal if the rotation of one gives the other.



$$\frac{1}{n} \sum_{i=1}^n k^{\text{gcd}(i,n)}$$

**Example2:** Count the number of different  $n \times n$  grids whose each square is black or white.

Two grids are considered to be different if it is not possible to rotate one of them so that they look the same.



$$G(\text{Rotations}) = 0^\circ, 90^\circ, 180^\circ, 270^\circ$$

$$\begin{aligned}
 f(\text{rotation}) = & \\
 & 0^\circ : 2^{(n^2)} \\
 & 90^\circ / 270^\circ : 2^{\frac{n^2}{4}}, & n_{\text{even}} \\
 & 2^{\frac{n^2-1}{4}} \cdot 2, & n_{\text{odd}} \\
 & 180^\circ : 2^{\frac{n^2}{2}}, & n_{\text{even}} \\
 & 2^{\frac{n^2-1}{2}} \cdot 2, & n_{\text{odd}}
 \end{aligned}$$

$$\text{ans} = \frac{1}{4} (f(0^\circ) + f(90^\circ) + f(180^\circ) + f(270^\circ))$$

### 7.2.4 Interesting Recursion

$$f(a, b) = f(a-1, b) + f(a, b-1)$$

$$\Rightarrow f(a, b) = \frac{(a+b)!}{a!b!} = \binom{a+b}{a}$$

*Proof:*

$$f(a,b) = \frac{(a+b)!}{a!b!}$$
$$\implies f(a-1,b) = \frac{(a-1+b)!}{(a-1)!b!}, f(a,b-1) = \frac{(a+b-1)!}{a!(b-1)!}$$
$$\implies f(a-1,b) + f(a,b-1) = \frac{(a-1+b)!}{(a-1)!b!} + \frac{(a+b-1)!}{a!(b-1)!}$$
$$\implies f(a,b) = (a+b-1)! \cdot \left( \frac{1}{(a-1)!b!} + \frac{1}{a!(b-1)!} \right)$$
$$\implies f(a,b) = (a+b-1)! \cdot \left( \frac{a+b}{a!b!} \right)$$
$$\implies f(a,b) = \frac{(a+b)!}{a!b!} = \binom{a+b}{a}$$

7.3 FFT

FFT can be used to turn a polynomial multiplication complexity to  $O(N \log N)$ .

A **convolution** is easily computed by inverting one of the vector and doing the polynomial multiplication normally.

fft-simple.cpp  
**Description:** Computes the product between two polynomials using fft  
**Time:**  $O(N \log N)$

```
// #define ld long double
// const ld PI = acos(-1);

struct num{
    ld a {0.0}, b {0.0};
    num() {}
    num(ld na) : a{na} {}
    num(ld na, ld nb) : a{na}, b{nb} {}
    const num operator +(const num &c) const{
        return num(a + c.a, b + c.b);
    }
    const num operator -(const num &c) const{
        return num(a - c.a, b - c.b);
    }
    const num operator *(const num &c) const{
        return num(a*c.a - b*c.b, a*c.b + b*c.a);
    }
    const num operator /(const int &c) const{
        return num(a/c, b/c);
    }
};

void fft(vector<num> &a, bool invert) {
    int n = (int)a.size();
    for(int i=1,j=0;i<n;i++){
        int bit = n>>1;
        for(; j&bit; bit>>=1)
            j^=bit;
        j^=bit;
        if(i<j) swap(a[i], a[j]);
    }
    for(int len = 2; len <= n; len <= 1) {
        ld ang = 2 * PI / len * (invert ? -1 : 1);
        num wlen(cos(ang), sin(ang));
        for(int i=0; i<n; i+=len) {
            num w(1);
```

```
for (int j=0;j<len/2;j++){
    num u = a[i+j], v = a[i+j+len/2] * w;
    a[i+j] = u + v;
    a[i+j+len/2] = u - v;
    w = w * wlen;
}
}
if(invert)
    for(num &x: a)
        x = x/n;
}

vector<ll> multiply(vector<int> const& a, vector<int> const& b)
{
    vector<num> fa(a.begin(), a.end());
    vector<num> fb(b.begin(), b.end());
    int n = 1;
    while(n < int(a.size() + b.size()) )
        n <= 1;
    fa.resize(n);
    fb.resize(n);
    fft(fa, false);
    fft(fb, false);
    for(int i=0;i<n;i++)
        fa[i] = fa[i]*fb[i];
    fft(fa, true);
    vector<ll> result(n);
    for(int i=0;i<n;i++)
        result[i] = (ll) round(fa[i].a);
    // while(result.back()==0) result.pop_back();
    return result;
}
```

Number theory (8)

8.1 Sieves

These sieves are used to find all primes up to an upper bound N,  $1 \leq N \leq 10^7$ .  
8.1.1 Eratosthenes  
Eratosthenes uses less memory than the linear sieve and is almost as fast

```
eratosthenes.cpp
Description: Optimized sieve of eratosthenes
Time: O(N log log N)

// O(N log^2(N)) -> Teorema de Merten
vector<ll> primes {2, 3};
bitset<MAX> sieve; // {sieve[i] == 1} if i is prime
// MAX can be ~1e7

void eratosthenes(ll n){
    sieve.set();
    for(ll i=5, step=2; i<=n; i+=step, step = 6 - step){
        if(sieve[i]){ // i is prime
            primes.pb(i);
            for(ll j= i*i; j<=n; j += 2*i) // sieving all odd multiples of i >= i*i
                sieve[j] = false;
        }
    }
}
```

8.1.2 Linear Sieve

Due to the lp vector, one can compute the factorization of any number very quickly!

Can check primality with  $lp[i] == i$

Uses more memory, because lp is a vector of int or ll and not bits.

Proof of time complexity:

We need to prove that the algorithm sets all values  $lp[]$  correctly, and that every value will be set exactly once. Hence, the algorithm will have linear runtime, since all the remaining actions of the algorithm, obviously, work for  $O(n)$ .

Notice that every number  $i$  has exactly one representation in form:

$$i = lp[i] \cdot x,$$

where  $lp[i]$  is the minimal prime factor of  $i$ , and the number  $x$  doesn't have any prime factors less than  $lp[i]$ , i.e.

$$lp[i] \leq lp[x].$$

Now, let's compare this with the actions of our algorithm: in fact, for every  $x$  it goes through all prime numbers it could be multiplied by, i.e. all prime numbers up to  $lp[x]$  inclusive, in order to get the numbers in the form given above.

Hence, the algorithm will go through every composite number exactly once, setting the correct values  $lp[]$  there. Q.E.D.

linear-sieve.cpp  
**Description:** Linear Sieve that iterates every value once (prime) or twice (composite)  
**Time:**  $O(N)$

```
vector<ll> primes, lp(MAX);
// lp[i] = smallest prime divisor of i

void linearSieve(ll n) {
    for (ll i=2; i <= n; i++) {
        if (lp[i] == 0) { // i is prime
            lp[i] = i; // {lp[i] == i} for prime numbers
            primes.pb(i);
        }
        // visit every composite number that has primes[j] as the lp
        for (ll j = 0; i * primes[j] <= n; j++) {
            lp[i * primes[j]] = primes[j];

            if (primes[j] == lp[i])
                break;
        }
    }
}
```

8.2 Extended Euclid

Solves the  $ax + by = gcd(a, b)$  equation.



8.2.1 Inverse Multiplicative

if  $\gcd(a,b) = 1$  :

then:

$$ax + by \equiv 1$$

also, if you apply  $(\bmod b)$  to the equation:

$$\begin{aligned} ax \pmod b + by \pmod b &\equiv 1 \pmod b \\ ax &\equiv 1 \pmod b \end{aligned}$$

In other words, one can find the inverse multiplicative of any number  $a$  in modulo  $b$  if  $\gcd(a,b) = 1$

8.2.2 Diofantine Equation

$$ax \equiv c \pmod b$$

if  $g = \gcd(a,b,c) \neq 1$ , divide everything by  $g$ .

After this, if  $\gcd(a,b) = 1$ , find  $a^{-1}$ , then multiply both sides of the Diofantine equation.

$$x \equiv c * a^{-1} \pmod b$$

After this, one has simply found  $x$

extended-euclid.cpp

**Description:** Solves the  $a * x + b * y = \gcd(a,b)$  equation

**Time:**  $\mathcal{O}(\log \min(a,b))$

c3a650, 43 lines

```
// equation: a*x + b*y = gcd(a, b)
// input: (a, b)
// returns gcd of (a, b)
// also computes &x and &y, which are passed by reference

11 extendedEuclid(11 a, 11 b, 11 &x, 11 &y) {
    x = 1, y = 0;
    11 x1 = 0, y1 = 1, a1 = a, b1 = b;
    while (b1) {
        11 q = a1 / b1;
        tie(x, x1) = pll{x1, x - q * x1};
        tie(y, y1) = pll{y1, y - q * y1};
        tie(a1, b1) = pll{b1, a1 - q * b1};
    }
    return a1;
}
```

Strings (9)

9.1 Hashing

Hashing consists in generating a Polynomial for the string, therefore, assigning each distint string to a specific numeric value In practice, there will always be some collisions:

Probability of colision:  $= \frac{n^2}{2m}$

extended-euclid hashing zfunction

n = Comparissons, m = mod size

when using multiple mods, they multiply: m = m1 \* m2

hashing.cpp

**Description:** Create a numerical value for a string by using polynomial hashing

**Time:**  $\mathcal{O}(n)$  to build,  $\mathcal{O}(1)$  per query

c3a650, 43 lines

```
// s[0]*P^n + s[1]*P^(n-1) + ... + s[n]*P^0
// 0-idx
struct Hashing {
    11 n, mod;
    string s;
    vector<11> p, h; // p = P^i, h = accumulated hash sum

    const 11 P = 31; // can be 53

    Hashing(string &s_, 11 m)
        : n(s_.size()), s(s_), mod(m), p(n), h(n) {

        for(11 i=0; i<n; i++)
            p[i] = (i ? P*p[i-1] : 1) % mod;

        for(11 i=0; i<n; i++)
            h[i] = (s[i] + P*(i ? h[i-1] : 0)) % mod;
    }

    11 query(11 l, 11 r) { // [l, r] inclusive (0-idx)
        11 hash = h[r] - (l ? (p[r-l+1]*h[l-1]) % mod : 0);
        return hash < 0 ? hash + mod : hash;
    }
};

// for codeforces:
mt19937 rng(chrono::steady_clock::now().time_since_epoch().
count());

int32_t main() { sws;
    vector<11> mods = {
        1000000009,10000000021,10000000033,
        10000000087,10000000093,10000000097,
        1000000103,10000000123,10000000181,
        1000000207,10000000223,10000000241,
        10000000271,10000000289,10000000297
    };

    shuffle(mods.begin(), mods.end(), rng);

    string s; cin >> s;

    Hashing hash(s, mods[0]);
}
```

9.2 Z-Function

Suppose we are given a string  $s$  of length  $n$ . The Z-function for this string is an array of length  $n$  where the  $i$ -th element is equal to the greatest number of characters starting from the position  $i$  that coincide with the first characters of  $s$  (the prefix of  $s$ )

The first element of the Z-function,  $z[0]$ , is generally not well defined. This implementation assumes it as  $z[0] = 0$ . But it can also be interpreted as  $z[0] = n$  (all characters coincide).

Can be used to solve the following simples problems:

- Find all ocurrences of a pattern  $p$  in another string  $s$ . ( $p + '$' + s$ ) ( $z[i] == p.size()$ )
- Find all borders. A border of a string is a prefix that is also a suffix of the string but not the whole string. For example, the borders of  $ab cab ab cab$  are  $ab$  and  $ab cab$ . ( $z[8] = 2, z[5] = 5$ ) ( $z[i] = n-i$ )
- Find all period lengths of a string. A period of a string is a prefix that can be used to generate the whole string by repeating the prefix. The last repetition may be partial. For example, the periods of  $abc ab ca$  are **abc**, **abc abc** and **abc ab ca**.

It works because  $(z[i] + i \leq n)$  is the condition when the common characters of  $z[i]$  in addition to the elements already passed, exceeds or is equal to the end of the string. For example:

$abaababab$   $z[8] = 2$

**abaababa** is the period; the remaining  $(z[i]$  characters) are a prefix of the period; and when all these characters are combined, it can form the string (which has  $n$  characters).

zfunction.cpp

**Description:** For each substring starting at position  $i$ , compute the maximum match with the original prefix.  $z[0] = 0$

**Time:**  $\mathcal{O}(n)$

14b37c, 12 lines

```
vector<11> z_function(string &s) { // O(n)
    11 n = (11) s.length();
    vector<11> z(n);
    for (11 i=1, l=0, r=0; i<n; i++) {
        if (i <= r) z[i] = min(r - i + 1, z[i - l]);

        while (i + z[i] < n and s[z[i]] == s[i + z[i]]) z[i]++;

        if (r < i + z[i] - 1) l = i, r = i + z[i] - 1;
    }
    return z;
}
```

9.3 KMP

KMP stands for Knuth-Morris-Pratt and computes the prefix function.

You are given a string  $s$  of length  $n$ . The prefix function for this string is defined as an array  $\pi$  of length  $n$ , where  $\pi[i]$  is the length of the longest proper prefix of the substring  $s[0 \dots i]$  which is also a suffix of this substring. A proper prefix of a string is a prefix that is not equal to the string itself. By definition,  $\pi[0] = 0$ .

For example, prefix function of string  $"abc abcd"$  is  $[0,0,0,1,2,3,0]$ , and prefix function of string  $"aabaab"$  is  $[0,1,0,1,2,2,3]$ .

kmpp.cpp

**Description:** Computes the prefix function

**Time:**  $\mathcal{O}(n)$

```
vector<ll> kmp(string &s) { //  $\mathcal{O}(n)$ 
    ll n = (ll) s.length();
    vector<ll> pi(n, 0); //  $pi[0] = 0$ 
    for (ll i=1; i<n; i++) {
        ll j = pi[i-1];
        while (j > 0 and s[i] != s[j])
            j = pi[j-1];
        if (s[i] == s[j])
            j++;
        pi[i] = j;
    }
    return pi;
}
```

9.3.1 Patterns in a String

Given a string  $p$  (pattern) and a string  $s$ , we want to find and display the positions of all occurrences of the string  $p$  in the string  $s$ .

**Solution:** Concatenate  $p + '$' + s$ , each position where  $pi[i] == p.size() \implies$  a match of the pattern in this substring.

9.4 Suffix Array

The suffix array is the array with size  $n$ , whose values are the indexes from the longest substring (0) to the smallest substring ( $n$ ) after ordering it lexicographically. Example:

Let the given string be "banana".		
0 banana		5 a
1 anana	Sort the Suffixes	3 ana
2 nana	----->	1 anana
3 ana	alphabetically	0 banana
4 na		4 na
5 a		2 nana
So the suffix array for "banana" is {5, 3, 1, 0, 4, 2}		

Note that the length of the string  $i$  is:  $(s.size())-sa[i]$

suffix-array.cpp

**Description:** Creates the Suffix Array

**Time:**  $\mathcal{O}(N \log N)$

```
vector<ll> suffixArray(string s) {
    s += "!";
    ll n = s.size(), N = max(n, 260LL);
    vector<ll> sa(n), ra(n);
    for (ll i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];

    for (ll k = 0; k < n; k ? k *= 2 : k++) {
        vector<ll> nsa(sa), nra(n), cnt(N);

        for (ll i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n)%n, cnt[ra[i]]++;
        for (ll i = 1; i < N; i++) cnt[i] += cnt[i-1];
        for (ll i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]]] = nsa[i];

        for (ll i = 1, r = 0; i < n; i++) nra[sa[i]] = r += ra[sa[i]] !=
```

kasai.cpp

**Description:** Creates the Longest Common Prefix array (LCP)

**Time:**  $\mathcal{O}(N \log N)$

```
vector<ll> kasai(string s, vector<ll> sa) {
    ll n = s.size(), k = 0;
    vector<ll> ra(n), lcp(n);
    for (ll i = 0; i < n; i++) ra[sa[i]] = i;

    for (ll i = 0; i < n; i++, k -= !k) {
        if (ra[i] == n-1) { k = 0; continue; }
        ll j = sa[ra[i]+1];
        while (i+k < n and j+k < n and s[i+k] == s[j+k]) k++;
        lcp[ra[i]] = k;
    }
    return lcp;
}
```

kmpp suffix-array kasai manacher

**Description:** Creates the Longest Common Prefix array (LCP)

**Time:**  $\mathcal{O}(N \log N)$

```
vector<ll> kasai(string s, vector<ll> sa) {
    ll n = s.size(), k = 0;
    vector<ll> ra(n), lcp(n);
    for (ll i = 0; i < n; i++) ra[sa[i]] = i;

    for (ll i = 0; i < n; i++, k -= !k) {
        if (ra[i] == n-1) { k = 0; continue; }
        ll j = sa[ra[i]+1];
        while (i+k < n and j+k < n and s[i+k] == s[j+k]) k++;
        lcp[ra[i]] = k;
    }
    return lcp;
}
```

Problems that can be solved:

Numbers of Distinct Substrings:

- $\frac{n(n+1)}{2} - lcp[i]$  (for all  $i$ )

Longest Repeated Substring:

- biggest  $lcp[i]$ . The position can be found in  $sa[i]$

Find how many distinct substrings there are for each  $len$  in  $[1:n]$ :

- Use delta encoding and the fact that  $lcp[i]$  counts the repeated substring between  $s.substr(sa[i])$  and  $s.substr(sa[i+1])$ , which are the substrings corresponding to the common prefix.

Find the  $k$ -th distinct substring:

kmpp suffix-array kasai manacher

**Description:** Creates the Longest Common Prefix array (LCP)

**Time:**  $\mathcal{O}(N \log N)$

```
string s; cin >> s;
ll n = s.size();

auto sa = suffix_array(s);
auto lcp = kasai(s, sa);

ll k; cin >> k;

for (ll i=0; i<n; i++) {
    ll len = n-sa[i];
    if (k <= len) {
        cout << s.substr(sa[i], k) << endl;
        break;
    }
    k += lcp[i] - len;
}
```

9.5 Manacher

Manacher's Algorithm is used to find all palindromes in a string.

For each substring, centered at  $i$ , find the longest palindrome that can be formed.

Works best for odd size string, so we convert all string to odd ones by adding and extra characters between the original ones

Therefore, the value stored in the vector  $cnt$  is actually  $palindrome-len + 1$ .

manacher.cpp

**Description:** Covert String to odd length to use manacher, which computes all the maximum lengths of all palindromes in the given string

**Time:**  $\mathcal{O}(2n)$

```
struct Manacher {
    string s, t;
    vector<ll> cnt;

    // t is the transformed string of s, with odd size
    Manacher(string &s_) : s(s_) {
        t = "#";
        for (auto c : s) {
            t += c, t += "#";
        }
        count();

        // perform manacher on the odd string
        // cnt will give all the palindromes centered in i
        // for the odd string t
        void count() {
            ll n = t.size();
            string aux = "$" + t + "^";
            vector<ll> p(n + 2);
            ll l = 1, r = 1;
            for (ll i = 1; i <= n; i++) {
                p[i] = max(0LL, min(r - i, p[l + (r - i)]));
                while (aux[i - p[i]] == aux[i + p[i]]) {
                    p[i]++;
                }
                if (i + p[i] > r) {
                    l = i - p[i], r = i + p[i];
                }
            }
            cnt = vector<ll>(p.begin() + 1, p.end() - 1);
        }

        // compute a longest palindrome present in s
        string getLongest() {
            ll len = 0, pos = 0;
            for (ll i=0; i<(ll)t.size(); i++) {
                ll sz = cnt[i]-1;
                if (sz > len) {
                    len = sz;
                    pos = i;
                }
            }
            return s.substr(pos/2 - len/2, len);
        }
    };
};
```

## 9.6 Booth

An efficient algorithm which uses a modified version of KMP to compute the least amount of rotation needed to reach the **lexicographically minimal string rotation**.

A rotation of a string can be generated by moving characters one after another from beginning to end. For example, the rotations of *acab* are *acab*, *caba*, *abac*, and *baca*.

booth.cpp

**Description:** Use a modified version of KMP to find the lexicographically minimal string rotation

**Time:**  $\mathcal{O}(n)$  64184b, 30 lines

```
// Booth Algorithm
ll least_rotation(string &s) { // O(n)
    ll n = s.length();
    vector<ll> f(2*n, -1);
    ll k = 0;
    for(ll j=1; j<2*n; j++) {
        ll i = f[j-k-1];
        while(i != -1 and s[j % n] != s[(k+i+1) % n] ) {
            if (s[j % n] < s[(k+i+1) % n])
                k = j - i - 1;
            i = f[i];
        }
        if (i == -1 and s[j % n] != s[(k+i+1) % n] ) {
            if (s[j % n] < s[(k+i+1) % n])
                k = j;
            f[j - k] = -1;
        }
        else
            f[j - k] = i + 1;
    }
    return k;
}

int32_t main(){ sws;
    string s; cin >> s;
    ll n = s.length();
    ll ans_idx = least_rotation(s);
    string tmp = s + s;
    cout << tmp.substr(ans_idx, n) << endl;
}
```

## 9.7 Suffix Automaton

The goat!!!

The implementation can be changed to use a fixed vector for adjacent edges instead of a map. This will reduce the time complexity to  $\mathcal{O}(n)$  and the memory to  $\mathcal{O}(nk)$ .

suffix-automaton.cpp

**Description:** Suffix automaton, each node represents a set of end-pos equivalent substrings. Solves A LOT of tasks!

**Time:**  $\mathcal{O}(n \log k)$ ,  $n = s.size()$ ,  $k = \text{characters used}$  22f143, 247 lines

```
struct Automaton {
    struct State {
        ll len = 0, link = 1;
        map<char, ll> down = {}; // 0 => non existent edge
        ll cnt = 0, fpos = -1;

        ll& operator [] (const char c) {
            return down[c];
        }
    };
};
```

```
};

ll n = 2; // number of states
vector<State> ton; // short for automaton :D
vector<ll> order_len;
string s;

Automaton(string ss) : s(ss) {
    // root = 1, root.link = 0 (0 is a dummy node)
    ton.assign(2, {0, 0});

    for(auto c : s) add(c);

    for(ll i=1; i<n; i++) order_len.pb(i);
    sort(order_len.begin(), order_len.end(), [&](auto l,
        auto r) {
            return ton[l].len > ton[r].len;
        });
    for(auto i : order_len) {
        ll j = ton[i].link;
        if (j != 1) ton[j].cnt += ton[i].cnt;
    }
}

ll last = 1;
void add(char c) {
    ll cur = n++;
    ll p = last;
    last = cur;

    State node; // state[cur]
    node.len = ton[p].len + 1;
    node.cnt = 1;
    node.fpos = node.len - 1;
    ton.pb(node);

    for (; p and !ton[p].down.count(c); p = ton[p].link)
        ton[p][c] = cur;

    if (p == 0) return;

    ll q = ton[p][c];
    if (ton[p].len + 1 == ton[q].len) {
        ton[cur].link = q;
        return;
    }

    ll clone = n++;
    State node2 = ton[q]; // state[clone]
    node2.cnt = 0;
    node2.len = ton[p].len + 1;
    ton.pb(node2);

    ton[cur].link = ton[q].link = clone;

    for (; ton[p][c] == q; p = ton[p].link)
        ton[p][c] = clone;
}

// ----- //
// Tasks //
// ----- //

// s1. Number of distinct substrings
// separated in a vector by their lengths
// using the property that: minlen(cur) = 1 + len(link(cur))
// and that state[cur] cover all the substrings (suffixes)
// of size [minlen, len] represented by this state
```

```
vector<ll> distinctSubstrings() { // O(n)
    ll sz = s.size();
    vector<ll> ans(sz+1, 0);

    for(ll i=2; i<n; i++) {
        ll j = ton[i].link;
        ll minlen = 1 + ton[j].len;
        ll len = ton[i].len;

        ans[minlen] += 1;
        if (len + 1 <= sz)
            ans[len + 1] -= 1;
    }

    // delta encoding
    for(ll len=1; len<=sz; len++) {
        ans[len] += ans[len-1];
    }
    // ans[0] = 0, because the empty string is not
    // considered as a substring
    return ans;
}

// s2. Find the lexicographically k-th substring (one can
// consider only the distincts or not)
// The k-th substring corresponds to the lexicographically
// smallest one,
// which is also the k-th path in the suffix automaton

// Additionally, by creating the automaton on the
// duplicated string (S+S),
// the k-th substring with k = s.size(), will give us the
// Smallest cyclic shift (Minimal Rotation)
// This may have a big complexity nevertheless, need to
// optimize

// number of substrings below node (including node)
// ps[0] -> include repeated substring, ps[1] -> consider
// only distinct
vector<ll> ps[2];

void buildPS() {
    ps[0].assign(n, 0), ps[1].assign(n, 0);

    for(ll k : {0, 1}) {
        for(auto i : order_len) {
            if (i != 1)
                ps[k][i] = (k ? 1 : ton[i].cnt);

            for(auto [c, j] : ton[i].down) {
                ps[k][i] += ps[k][j];
            }
        }
    }
}

string substring(ll k, bool distinct = true) { // O(V+E) =
    O(2sz + 3sz) = O(5sz), sz = s.size()
    assert(!ps[0].empty());

    string ans = ""; // {k = 0} will return the empty
    string ""

    function<void (ll)> dfs = [&](ll u) {
        if (k <= 0) return;
        for(auto [c, v] : ton[u].down) {
            ll sum = ps[distinct][v];
            if (k <= sum) {
```

```

        ans += c;
        k -= (distinct ? 1 : ton[v].cnt);
        dfs(v);
        if (k <= 0) return;
    }
    else k -= sum; // optimization
};

dfs(1);
return ans;
}

// ===== //
// Patterns //
// ===== //

// p1. Check for occurrence of a pattern P
// by returning the length of the longest prefix of P in S
// A match occurs when len(prefix-pattern) == len(pattern)

ll prefixPattern(string &p) { // O( p.size() )
    ll ans = 0, cur = 1;
    for(auto c : p) {
        if (ton[cur].down.count(c)) {
            cur = ton[cur][c];
            ans += 1;
        }
        else break;
    }
    return ans;
}

// p2. Count the numbers of occurrences of a pattern P

ll countPattern(string &p) { // O( p.size() )
    ll u = 1;
    for(auto c : p) {
        if (ton[u].down.count(c)) {
            u = ton[u].down[c];
        }
        else return 0; // no match
    }
    return ton[u].cnt;
}

// p3. Find the first position in which occurred the
// pattern (0-idx)

ll firstPattern(string &p) { // O( p.size() )
    ll u = 1;
    for(auto c : p) {
        if (ton[u].down.count(c)) {
            u = ton[u].down[c];
        }
        else return -1; // no match
    }
    ll sz = p.size();
    return ton[u].fpos - sz + 1;
}

// p4. Longest Common Substring of P and S

string lcs(string &p) { // O( p.size() )

    ll cur = 1, match = 0, best = 0, pos = 0;

    for(ll i=0; i<(ll)p.size(); i++) {
        auto c = p[i];

```

```

        while(cur > 1 and !ton[cur].down.count(c)) {
            cur = ton[cur].link;
            match = ton[cur].len;
        }
        if (ton[cur].down.count(c)) {
            cur = ton[cur][c];
            match++;
        }
        if (match > best) {
            best = match;
            pos = i;
        }
    }

    return p.substr(pos - best + 1, best);
}

// ===== //
// Debugging //
// ===== //

using T = tuple<ll, ll, char>;
vector<T> edges() {
    vector<T> ans;
    for(ll i=1; i<n; i++) {
        for(auto [c, j] : ton[i].down) {
            ans.pb({i, j, c});
        }
    }
    return ans;
}

vector<pll> links() {
    vector<pll> ans;
    for(ll i=1; i<n; i++) {
        ans.pb({i, ton[i].link});
    }
    return ans;
}

}
};

```

## Miscellaneous (10)

### 10.1 Random Generator

random.cpp

**Description:** Good randomizer to generate int in a range or shuffle vectors  
**Time:**  $\mathcal{O}(1)$  for randint,  $\mathcal{O}(n \log(n))$  for shuffle

55c5b9, 13 lines

```

mt19937 rng(chrono::steady_clock::now().time_since_epoch().
    count());

// or for 64 bits
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().
    count());

// to shuffle a vector
vector<int> vec;
shuffle(vec.begin(), vec.end(), rng);

// to limit the number to the range [l, r]
int randint(int l, int r) {
    return (rng() % (r-l+1)) + l;
}

```

### 10.2 Read an Fraction Input

char c;

```

ll num, den;
cin >> num >> c >> den;

```

## 10.3 Ternary Search

ternary-search.cpp

**Description:** Computes the min/max for a function that is monotonically increasing then decreasing or decreasing then increasing.

**Time:**  $\mathcal{O}(N \log N_3)$

c3a5d7, 48 lines

```

/*
Float and Min Version: Requires EPS (precision usually defined
in the question text)
*/

ld f(ld d){
    // function here
}

// for min value
ld ternary_search(ld l, ld r){
    while(r - l > EPS){
        // divide into 3 equal parts and eliminate one side
        ld m1 = l + (r - l) / 3;
        ld m2 = r - (r - l) / 3;
        if (f(m1) < f(m2)) {
            r = m2;
        }
        else {
            l = m1;
        }
    }
    return f(l); // check here for min/max
}

/*
Integer and Max Version:
*/

ll f(ll idx) {
    // function here
}

// for max value, using integer idx
ll ternary_search(ll l, ll r) {
    while(l <= r) {
        // divide into 3 equal parts and eliminate one side
        ll m1 = l + (r-l)/3;
        ll m2 = r - (r-l)/3;
        if (f(m1) < f(m2)) {
            l = m1+1;
        }
        else {
            r = m2-1;
        }
    }
    return f(l); // check here for min/max
}

```

# Techniques (A)

techniques.txt	160 lines
Recursion	
Divide and conquer	
Finding interesting points in N log N	
Algorithm analysis	
Master theorem	
Amortized time complexity	
Greedy algorithm	
Scheduling	
Max contiguous subvector sum	
Invariants	
Huffman encoding	
Graph theory	
Transform edges into vertices, duplicating the nodes of the graph	
Dynamic graphs (extra book-keeping)	
Breadth first search	
Depth first search	
* Normal trees / DFS trees	
Dijkstra's algorithm	
MST: Prim's algorithm	
Bellman-Ford	
Konig's theorem and vertex cover	
Min-cost max flow	
Lovasz toggle	
Matrix tree theorem	
Maximal matching, general graphs	
Hopcroft-Karp	
Hall's marriage theorem	
Graphical sequences	
Floyd-Warshall	
Euler cycles	
Flow networks	
* Augmenting paths	
* Edmonds-Karp	
Bipartite matching	
Min. path cover	
Topological sorting	
Strongly connected components	
2-SAT	
Cut vertices, cut-edges and biconnected components	
Edge coloring	
* Trees	
Vertex coloring	
* Bipartite graphs (=> trees)	
* 3^n (special case of set cover)	
Diameter and centroid	
K'th shortest path	
Shortest cycle	
Dynamic programming	
Knapsack	
Coin change	
Longest common subsequence	
Longest increasing subsequence	
Number of paths in a dag	
Shortest path in a dag	
Dynprog over intervals	
Dynprog over subsets	
Dynprog over probabilities	
Dynprog over trees	
3^n set cover	
Divide and conquer	
Knuth optimization	
Convex hull optimizations	
RMQ (sparse table a.k.a 2^k-jumps)	
Bitonic cycle	

Log partitioning (loop over most restricted)
Combinatorics
Computation of binomial coefficients
Pigeon-hole principle
Inclusion/exclusion
Catalan number
Pick's theorem
Number theory
Integer parts
Divisibility
Euclidean algorithm
Modular arithmetic
* Modular multiplication
* Modular inverses
* Modular exponentiation by squaring
Chinese remainder theorem
Fermat's little theorem
Euler's theorem
Phi function
Frobenius number
Quadratic reciprocity
Pollard-Rho
Miller-Rabin
Hensel lifting
Vieta root jumping
Game theory
Combinatorial games
Game trees
Mini-max
Nim
Games on graphs
Games on graphs with loops
Grundy numbers
Bipartite games without repetition
General games without repetition
Alpha-beta pruning
Probability theory
Optimization
Binary search
Ternary search
Unimodality and convex functions
Binary search on derivative
Numerical methods
Numeric integration
Newton's method
Root-finding with binary/ternary search
Golden section search
Matrices
Gaussian elimination
Exponentiation by squaring
Sorting
Radix sort
Geometry
Coordinates and vectors
* Cross product
* Scalar product
Convex hull
Polygon cut
Closest pair
Coordinate-compression
Quadtrees
KD-trees
All segment-segment intersection
Sweeping
Discretization (convert to events and sweep)
Angle sweeping
Line sweeping
Discrete second derivatives
Strings

Longest common substring
Palindrome subsequences
Knuth-Morris-Pratt
Tries
Rolling polynomial hashes
Suffix array
Suffix tree
Aho-Corasick
Manacher's algorithm
Letter position lists
Combinatorial search
Meet in the middle
Brute-force with pruning
Best-first (A*)
Bidirectional search
Iterative deepening DFS / A*
Data structures
LCA (2^k-jumps in trees in general)
Pull/push-technique on trees
Heavy-light decomposition
Centroid decomposition
Lazy propagation
Self-balancing trees
Convex hull trick (wcipeg.com/wiki/Convex_hull_trick)
Monotone queues / monotone stacks / sliding queues
Sliding queue using 2 stacks
Persistent segment tree