

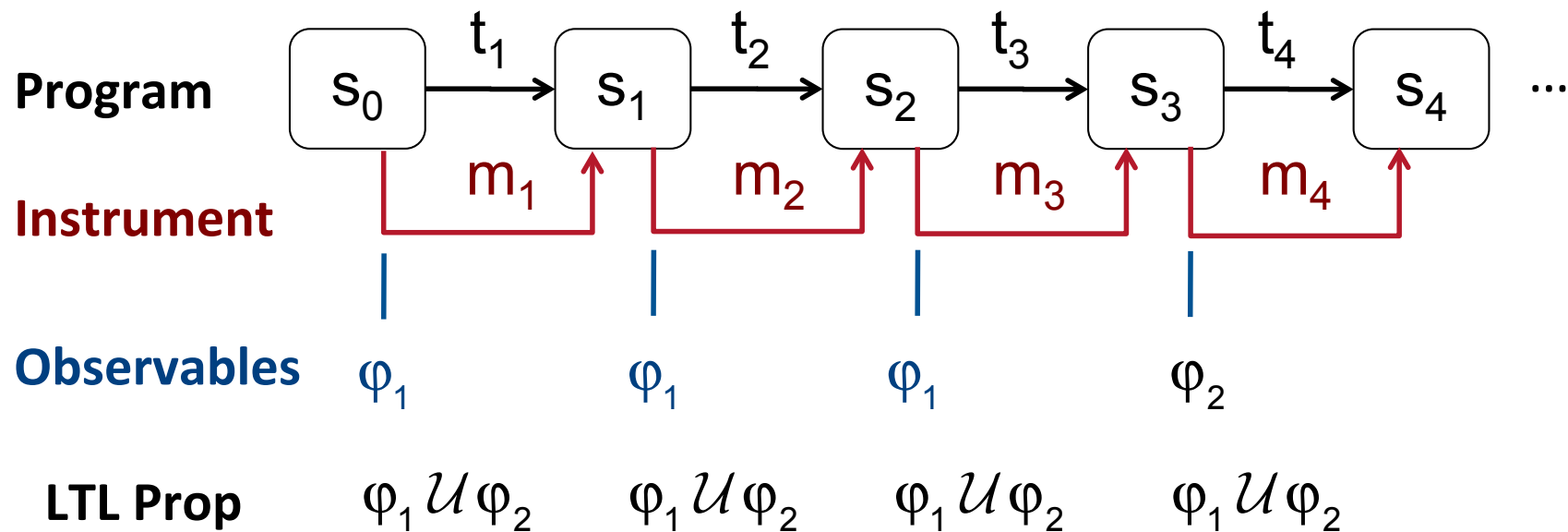
Adaptive Runtime Verification

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Joint work with:

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S.A. Smolka, S.D. Stoller, J. Seyster, E. Zadok**

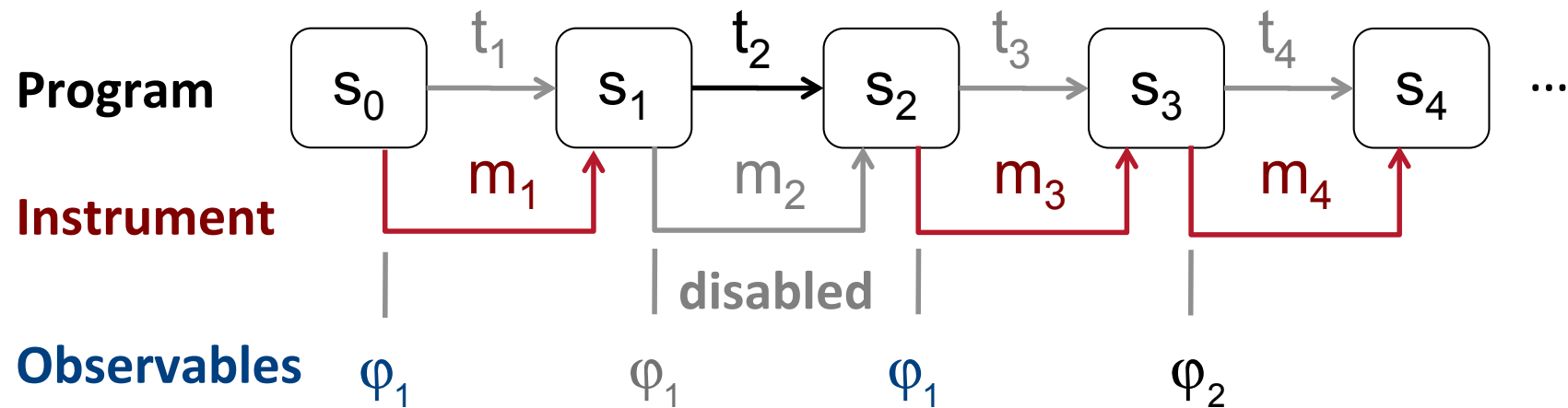
Runtime Verification Problem



Why RV? Most properties cannot be checked statically

Problem: It introduces overhead!

Overhead Control Problem



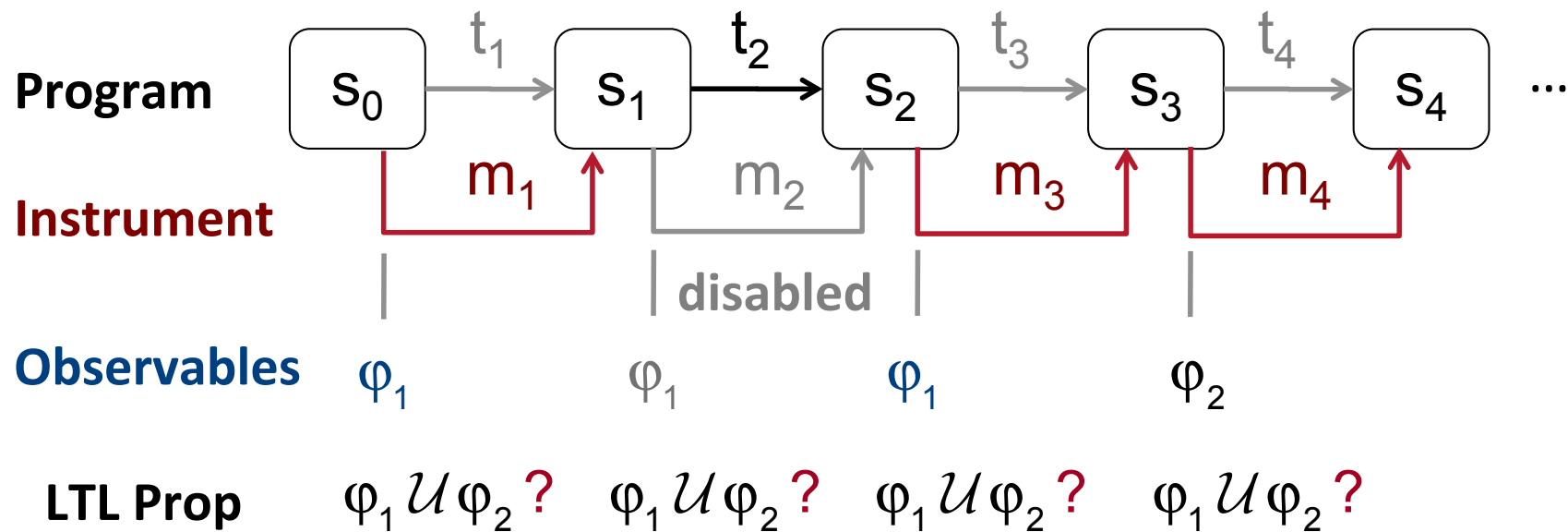
Solution: Use a PID controller

LTL Prop $\varphi_1 \mathcal{U} \varphi_2 ?$ $\varphi_1 \mathcal{U} \varphi_2 ?$ $\varphi_1 \mathcal{U} \varphi_2 ?$ $\varphi_1 \mathcal{U} \varphi_2 ?$

Problem: We do not know anymore if the LTL property holds!

Fix: Learn and Estimate state with a Hidden Markov Model

Stateful Overhead Control



Adaptive RV: Knowing the state one can tune budget towards Instances that are most likely to violate the LTL property next.

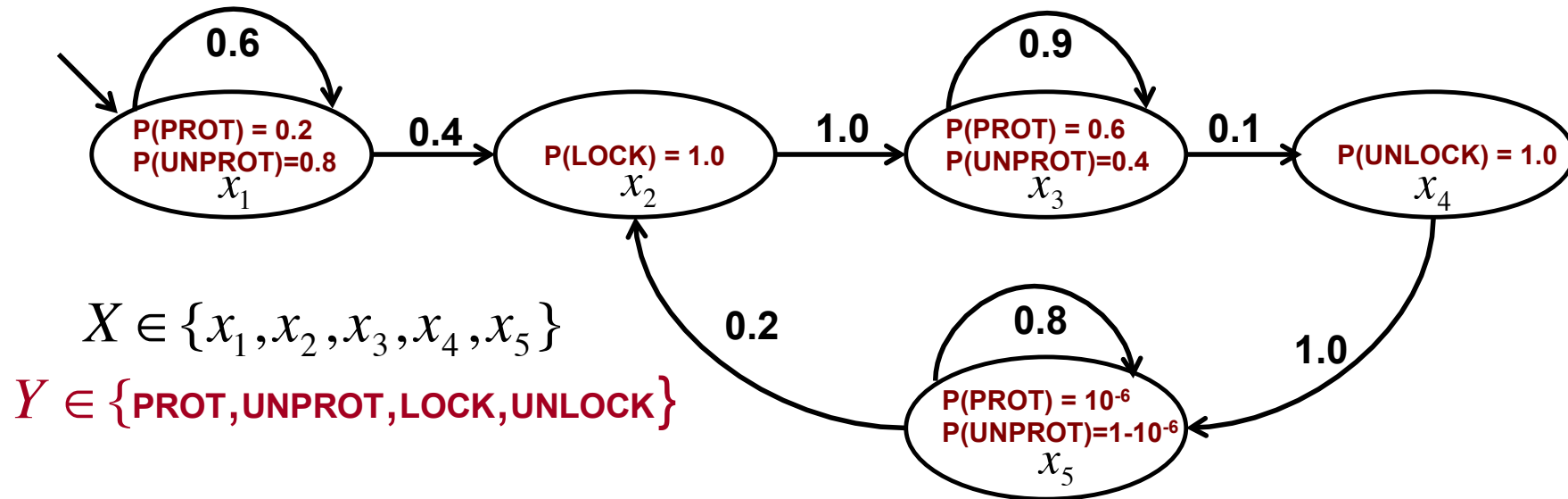
Outline

- **Program Model: Hidden Markov Model**
- **Property: Deterministic Finite Automaton**
- **Corresponding Dynamic Bayesian Network**
- **State Estimation with Particle Filters**
- **Comparison of State-Estimation Techniques**
- **Adaptive Runtime Verification**

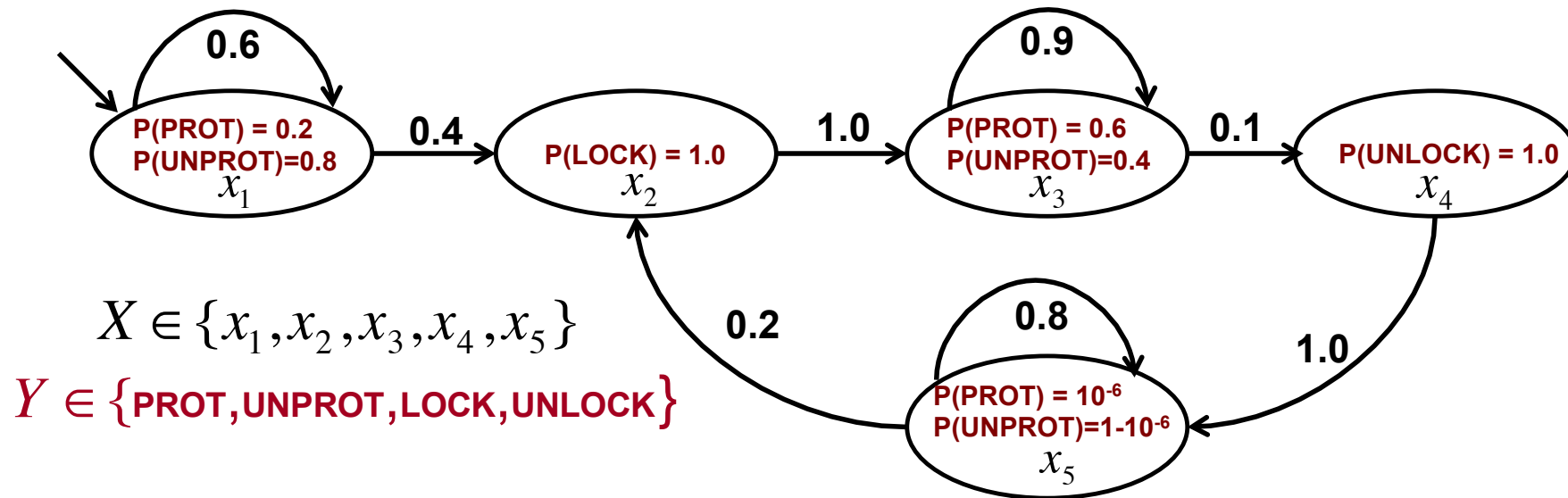
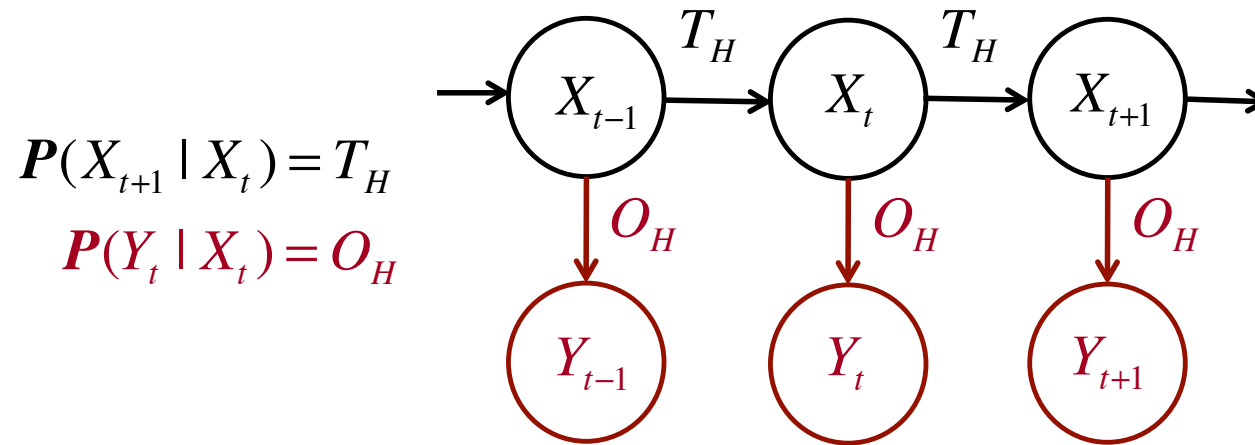
Program Model: Hidden Markov Model

T_H	x_1	x_2	x_3	x_4	x_5
x_1	0.6	0.4			
x_2			1		
x_3			0.9	0.1	$P(x_5 x_3)$
x_4					1
x_5		0.2			0.8

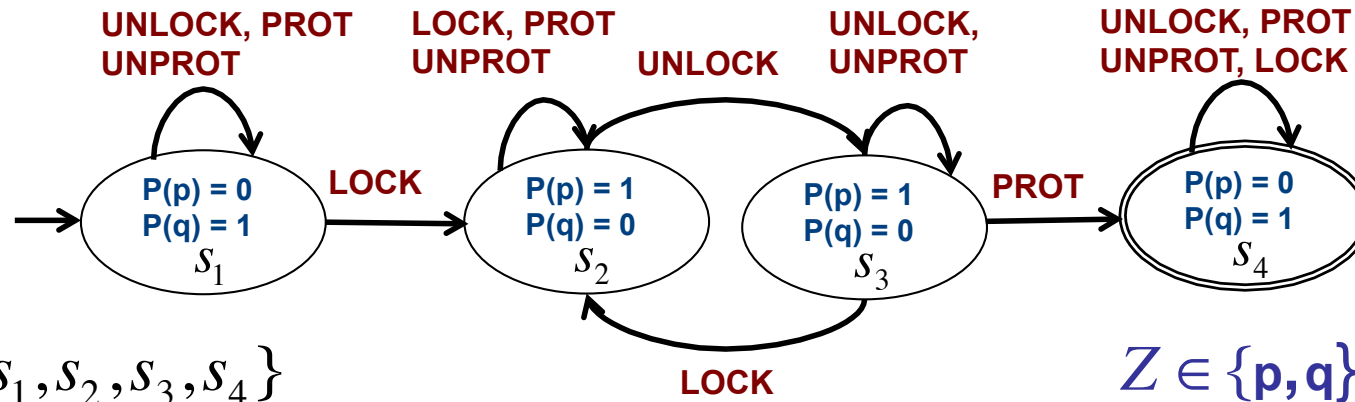
O_H	PROT	UNPROT	LOCK	UNLOCK
x_1	0.2	0.8		
x_2			1	
x_3	0.6	0.4		$P(\text{UNLOCK} x_3)$
x_4				1
x_5	10^{-6}	$1 - 10^{-6}$		



Program Model: Hidden Markov Model



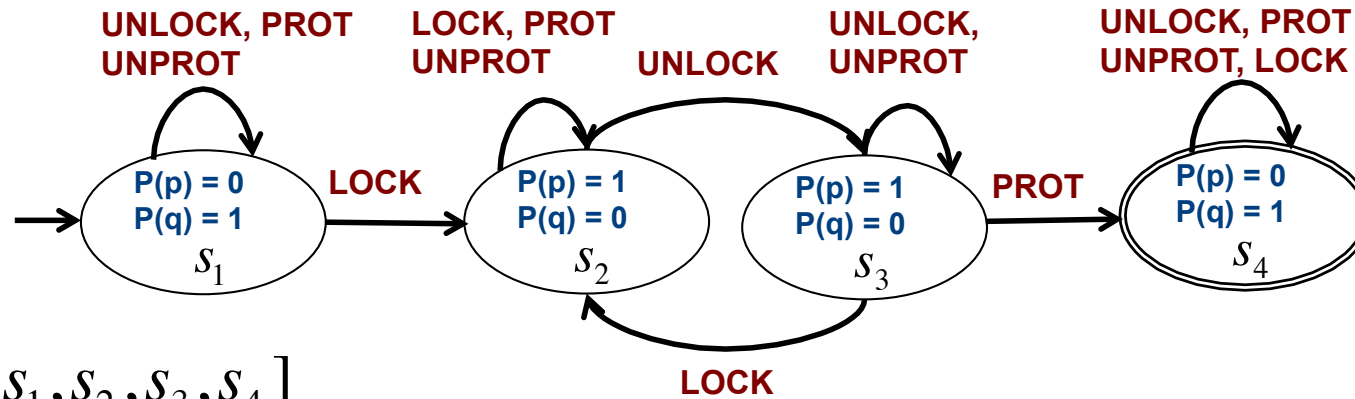
Property: DFA



T_D	s_1	s_2	s_3	s_4
s_1	PROT UNPROT UNLOCK	LOCK		
s_2		PROT UNPROT LOCK	UNLOCK	
s_3		LOCK	UNPROT UNLOCK	$P(s_4 s_3, \text{PROT})$
s_4				PROT LOCK UNPROT UNLOCK

O_D	p	q
s_1	0	1
s_2	1	0
s_3	1	$P(q s_3)$
s_4	0	1

Property: DFA

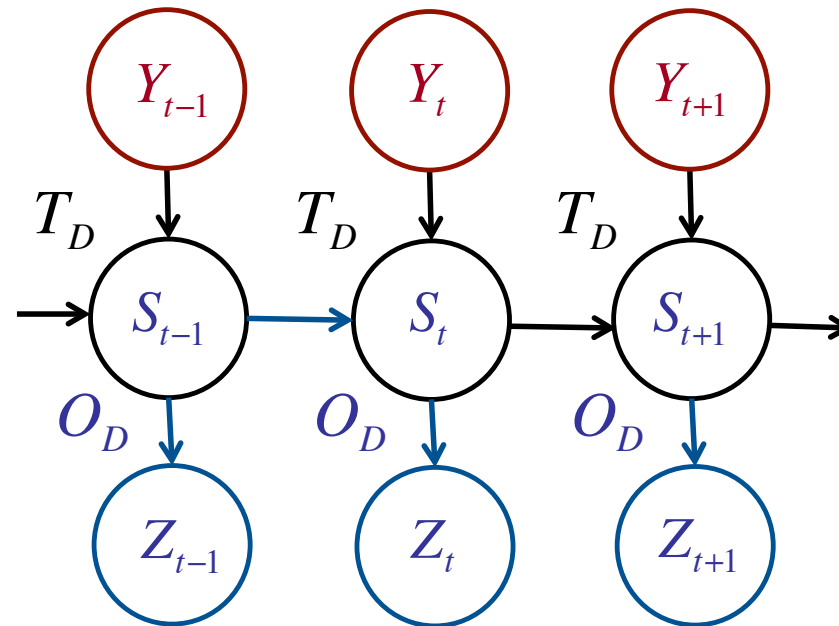


$$S = [s_1, s_2, s_3, s_4]$$

$$Z \in \{p, q\}$$

$$P(S_{t+1} | S_t, Y_{t+1}) = T_D$$

$$P(Z_t | S_t) = O_D$$



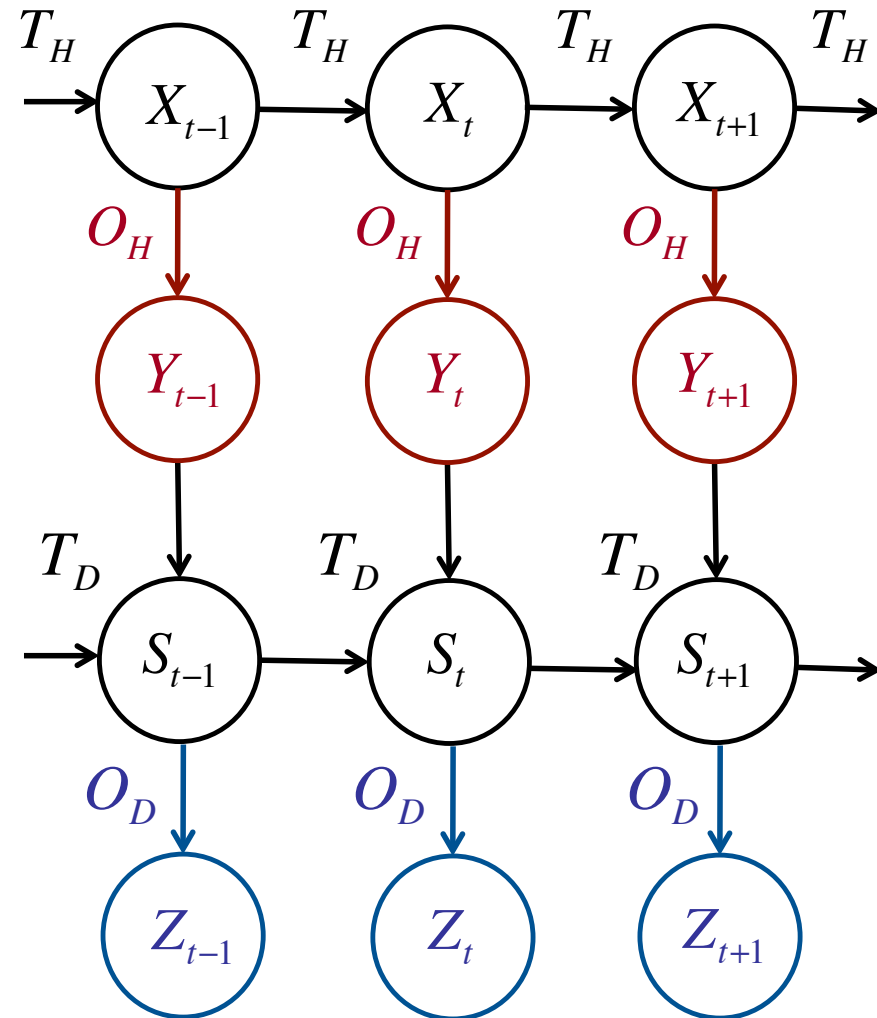
Filtering in DBNs

Given

- $P(X_{t+1} | X_t) = T_H, P(X_1)$
- $P(Y_t | X_t) = O_H$
- $P(S_{t+1} | S_t, Y_{t+1}) = T_D, P(S_1)$
- $P(Z_t | S_t) = O_D$
- $y_1 \dots y_{t+1}, z_1 \dots z_{t+1}$

Find

- $P(S_{t+1} | y_1 \dots y_{t+1}, z_1 \dots z_{t+1})$



Computing $P(X_{t+1} | y_{1:t+1})$

Exact computation (RVSE):

- $P(X_{t+1} | y_{1:t+1})^t = \alpha O_{y_{t+1}}^d T^t P(X_t | y_{1:t})^t$
- **Expensive**: Matrix multiplication consumes 80% of time

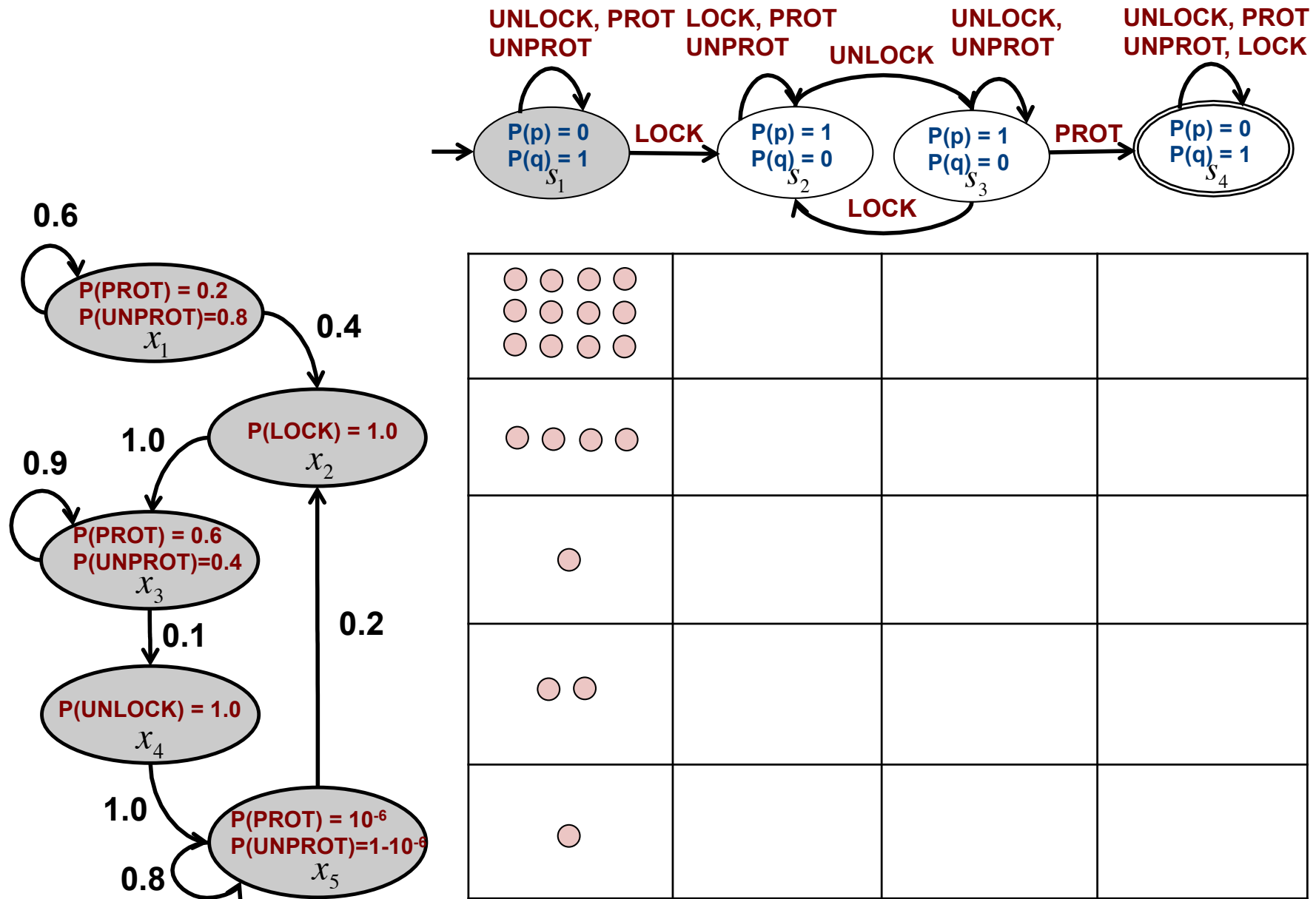
Offline Tabling (AP - RVSE):

- Approximate computation of $P(X | y_{1:*})^t, \forall y_{1:*}$ for error ε
- **Fast**: Memory requirements are very high

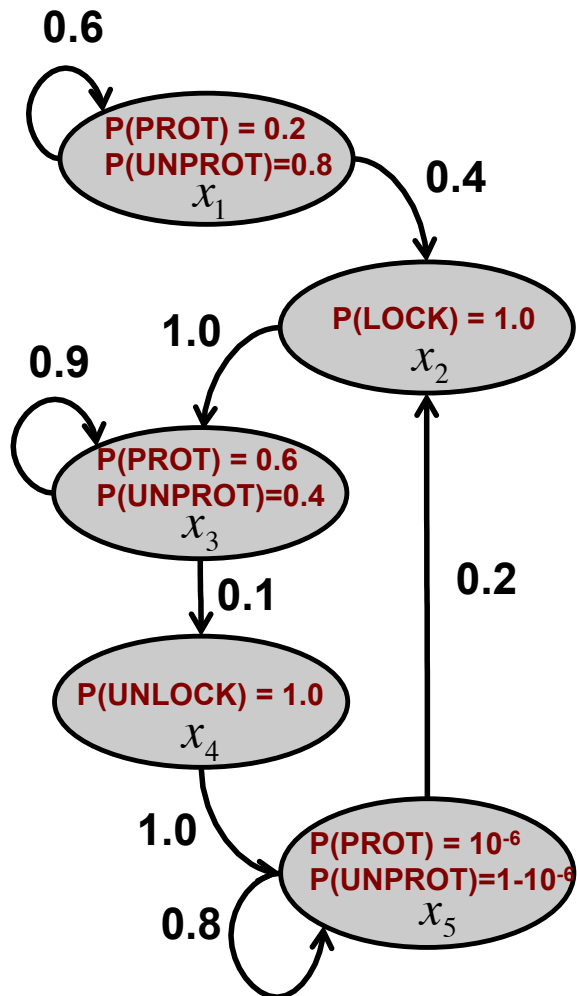
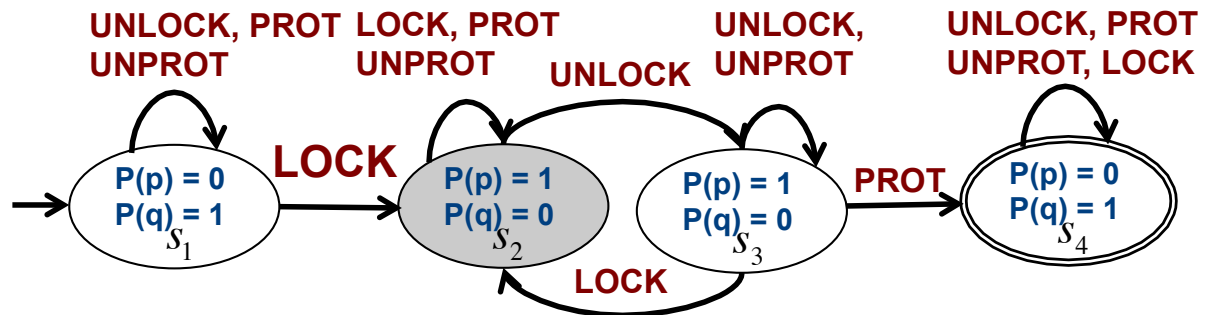
Online Particle Filtering (RVPF):

- Approximate computation of $P(X_{t+1} | y_{1:t+1})^t$ for n particles
- **Adaptive**: Number n of particles balance overhead / accuracy

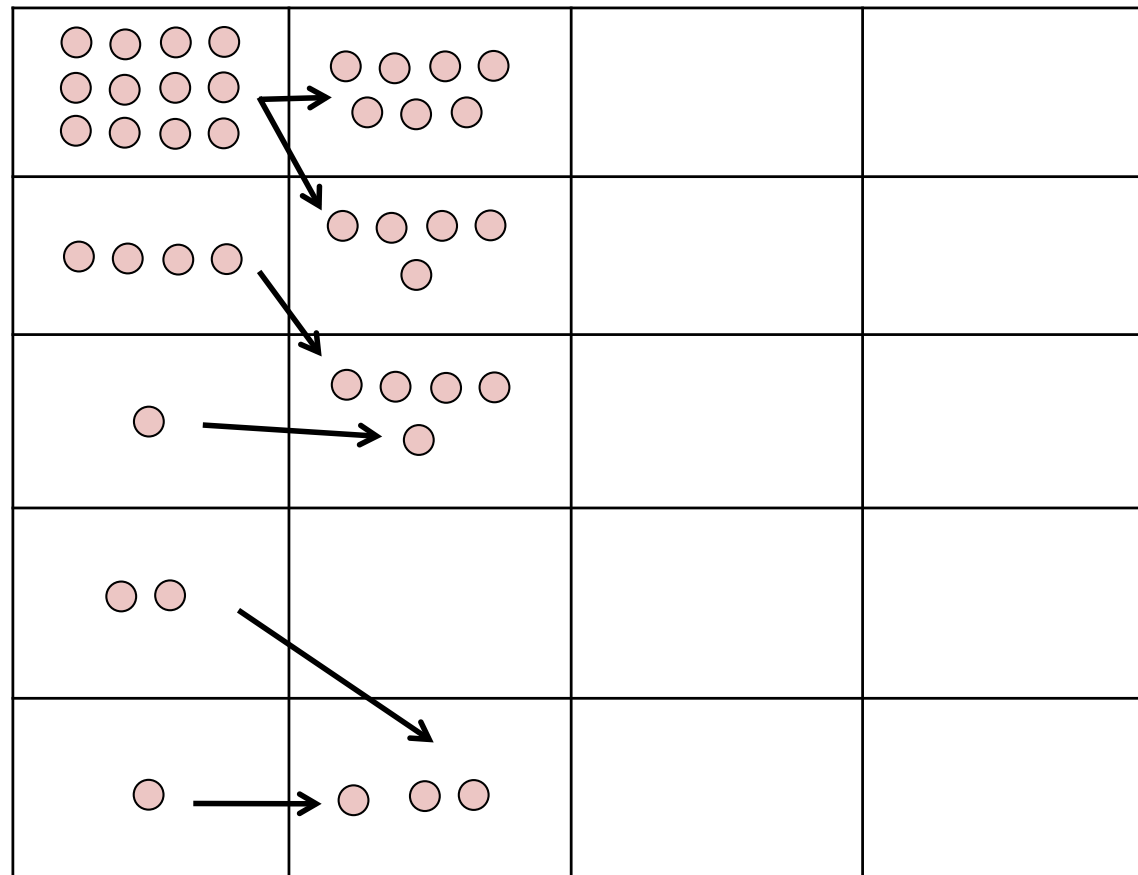
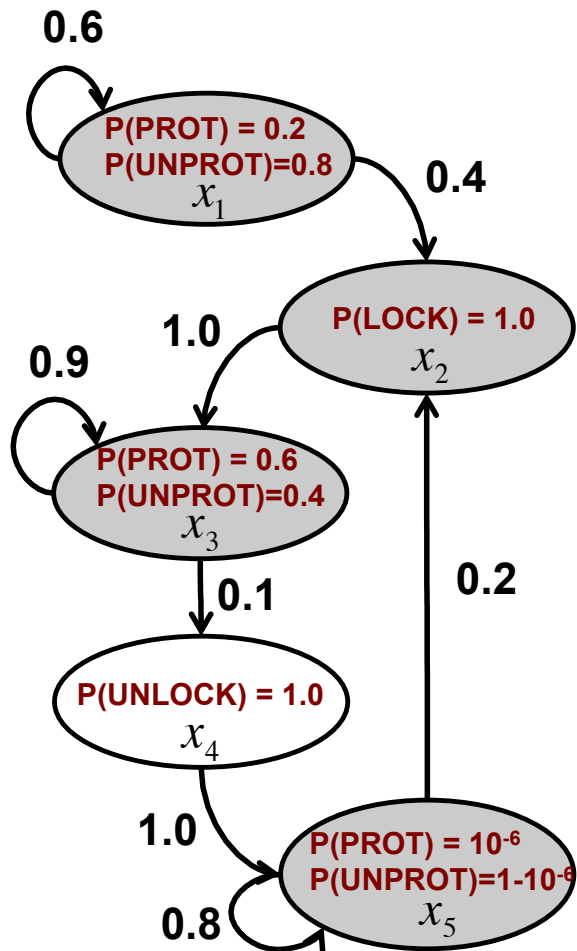
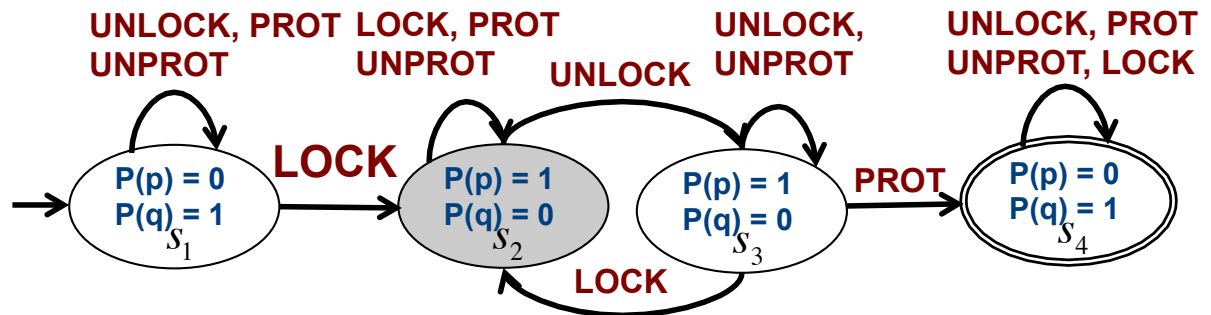
RVPF: Initial Distribution of Particles



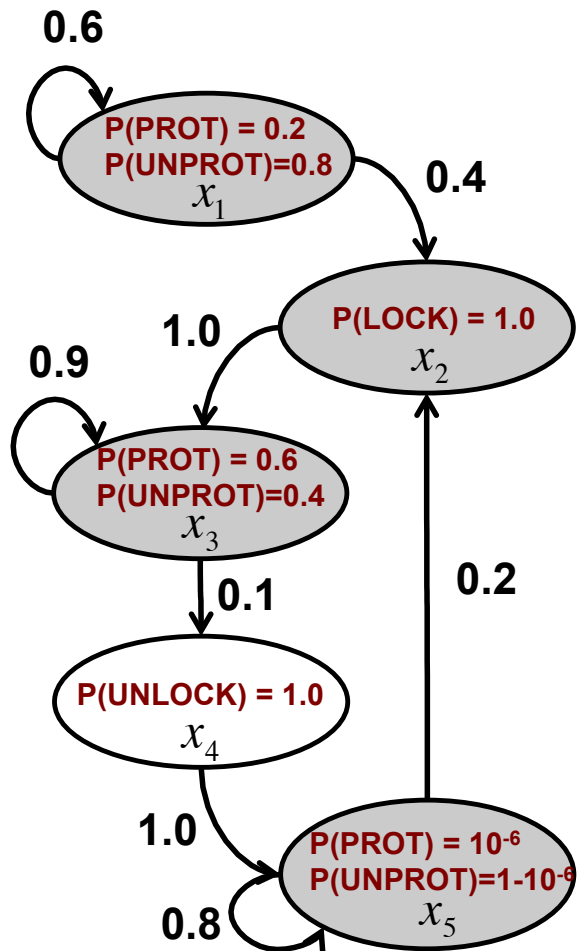
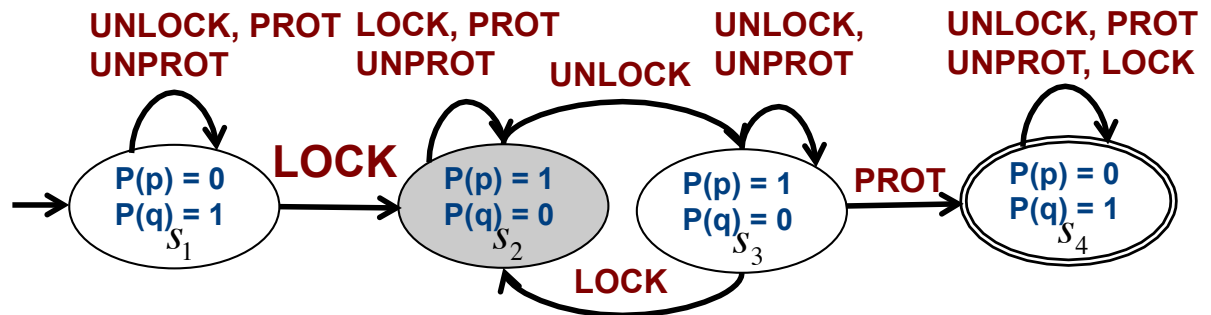
RVPF: Observe LOCK



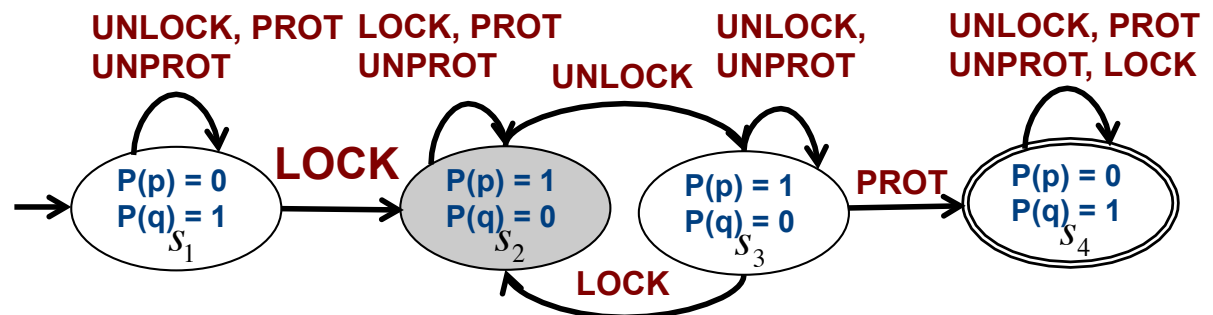
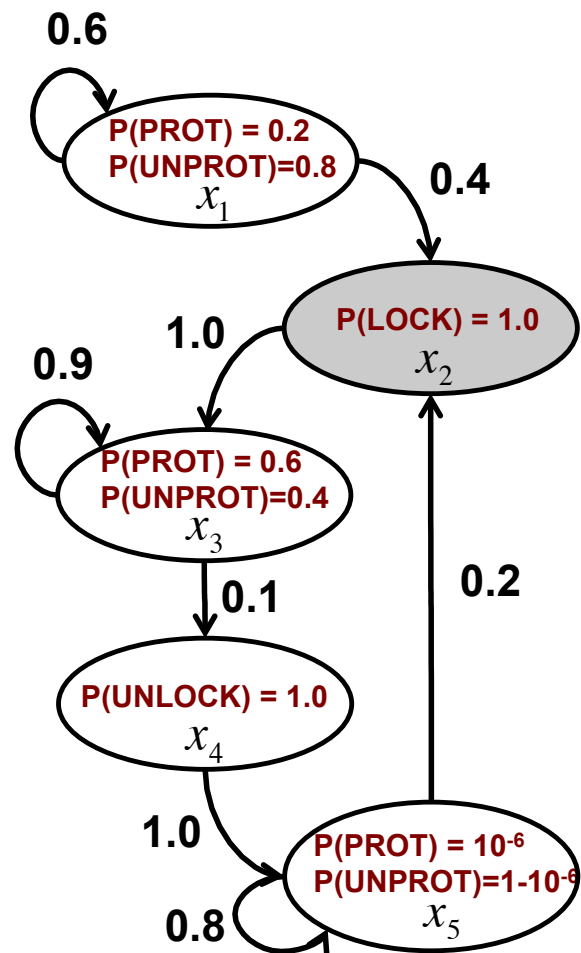
RVPF: Distribute Particles



RVPF: New Configuration of Particles

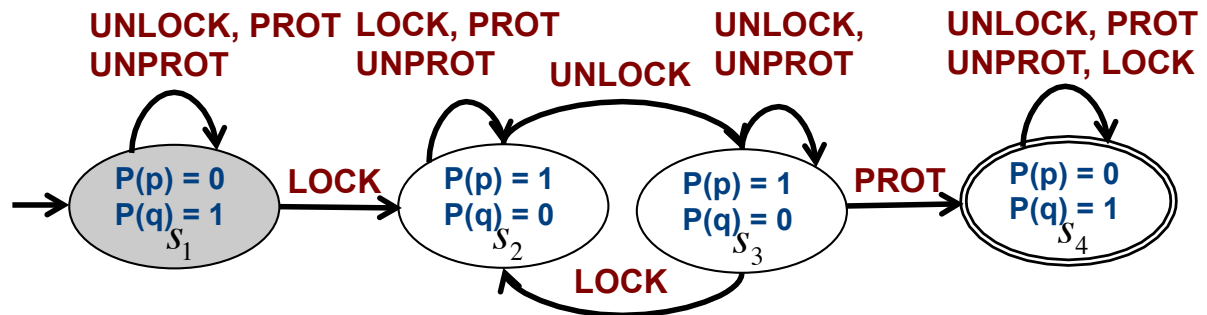
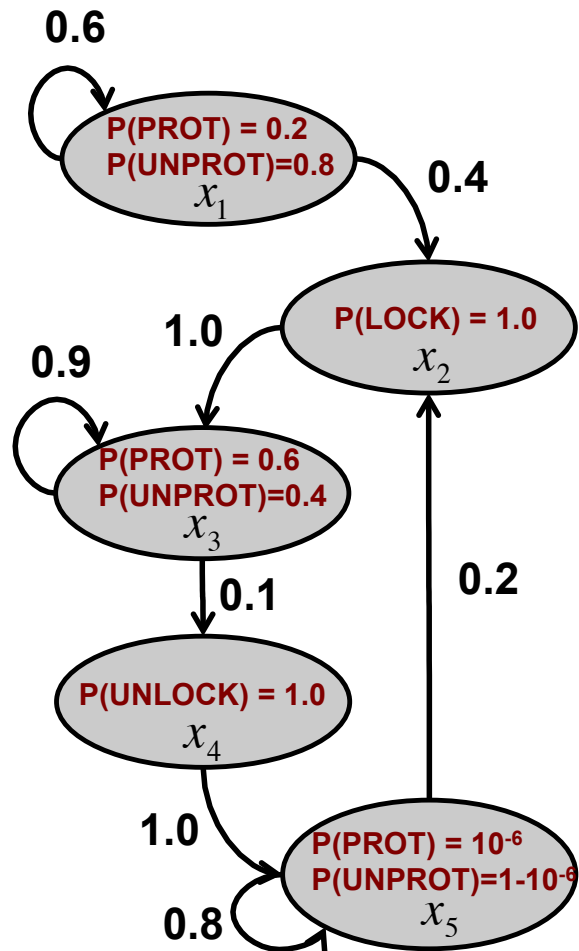


RVPF: Resample Particles for LOCK



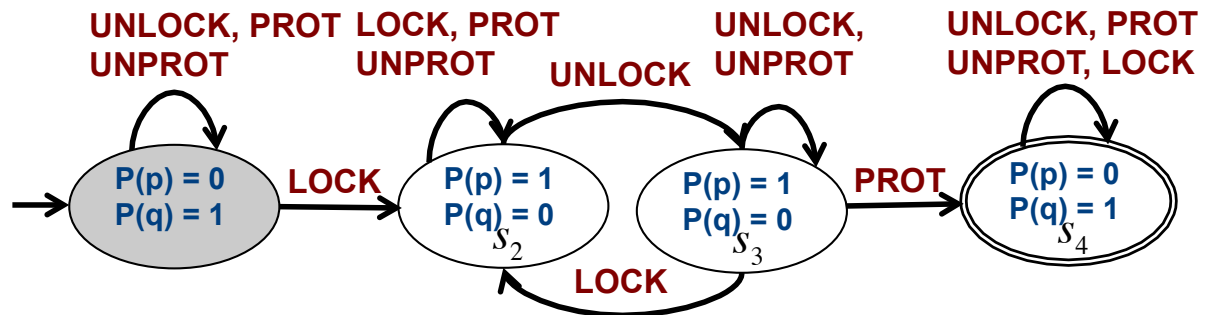
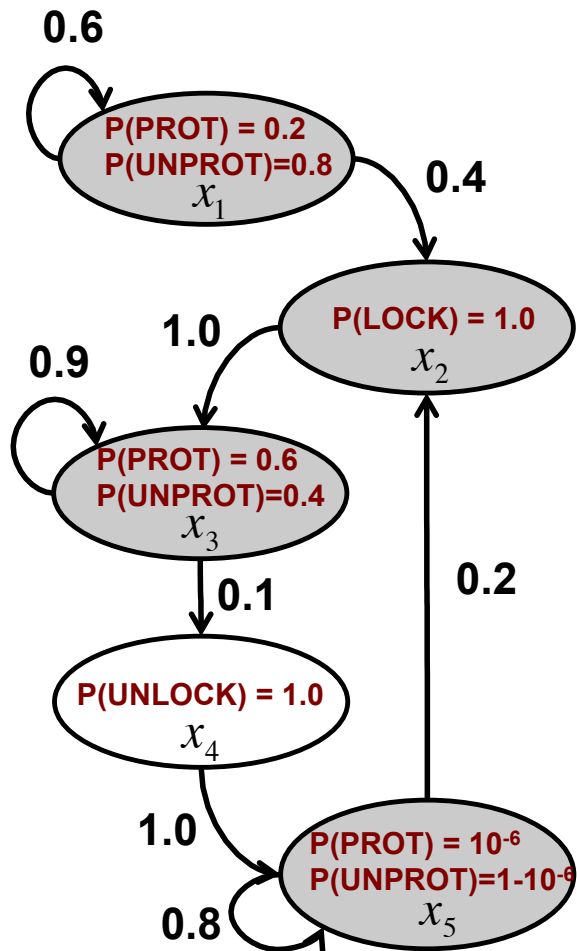
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RVPF: Initial distribution of particles



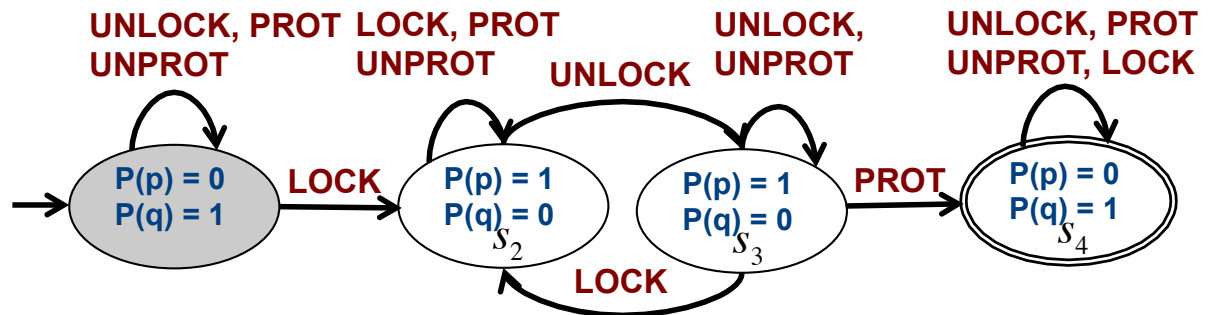
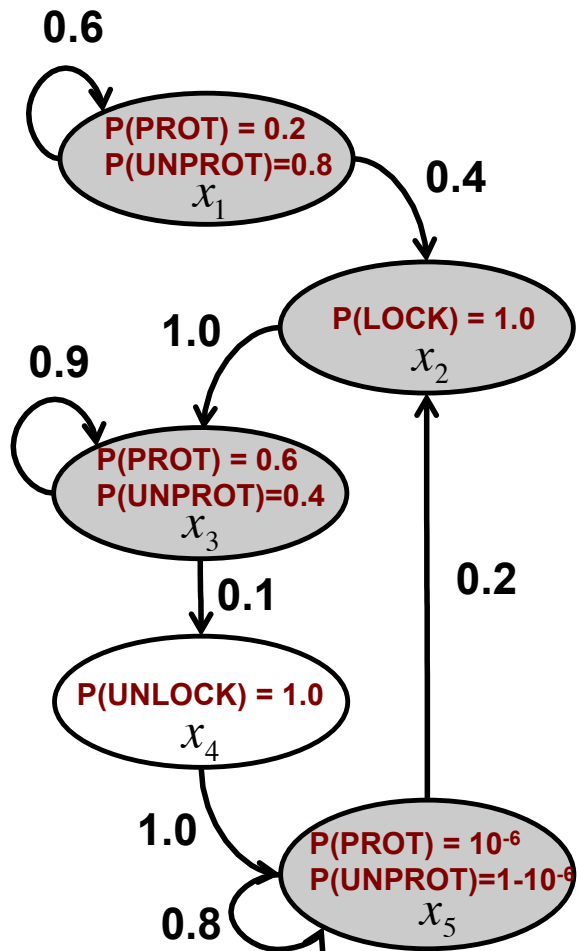
RVPF: Observe GAP

Compute
next X in s_1



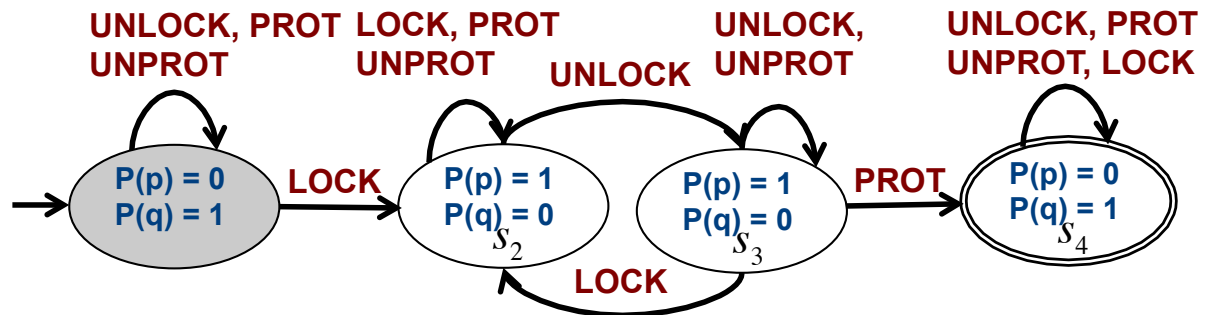
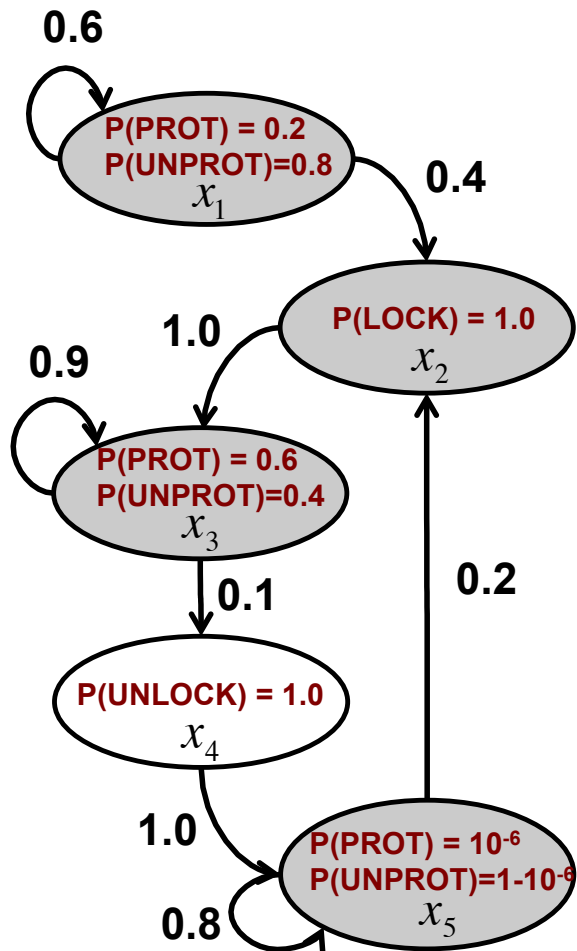
RVPF: Observe GAP

Compute
next X in s_1



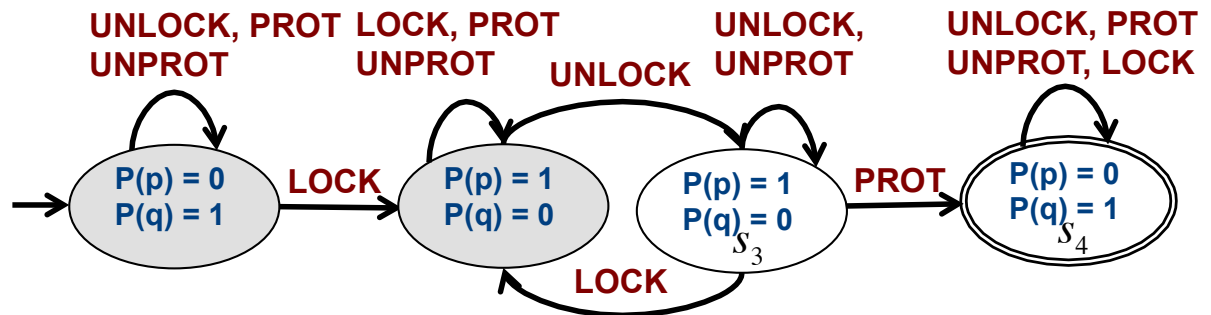
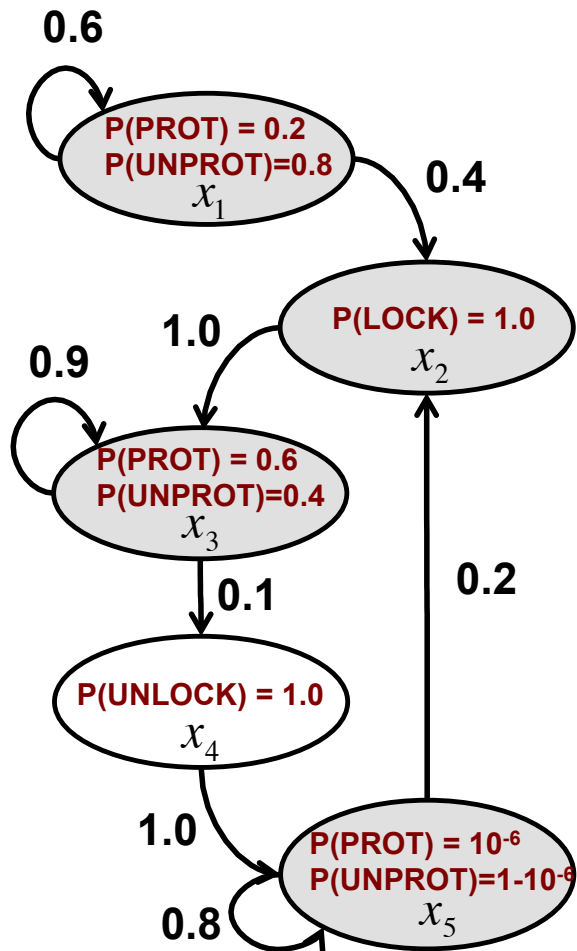
RVPF: Observe GAP

Check O
In every x_i



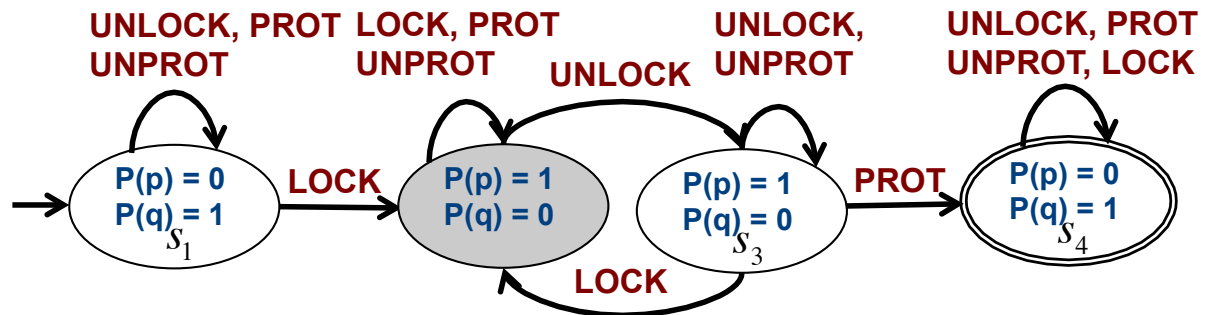
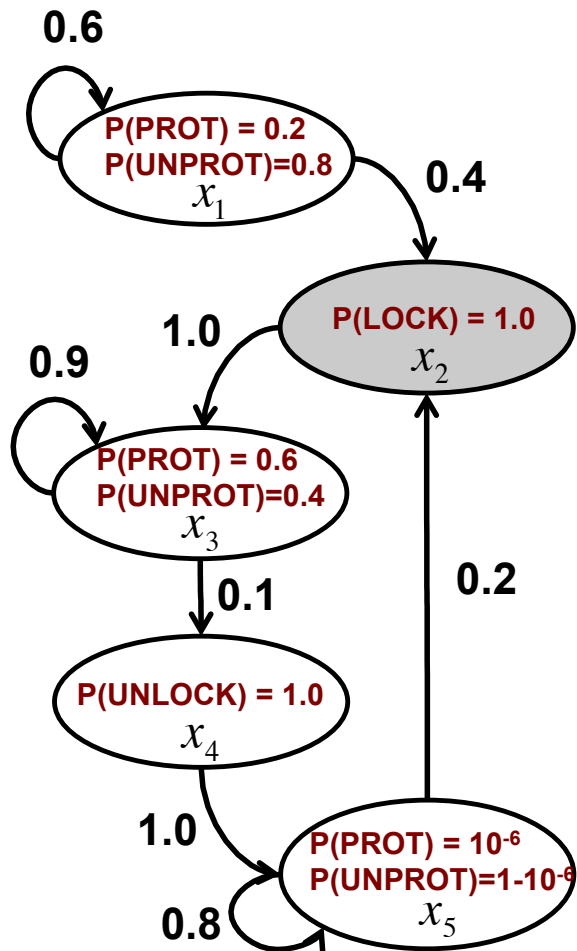
RVPF: Observe GAP

Check O
In every x_i



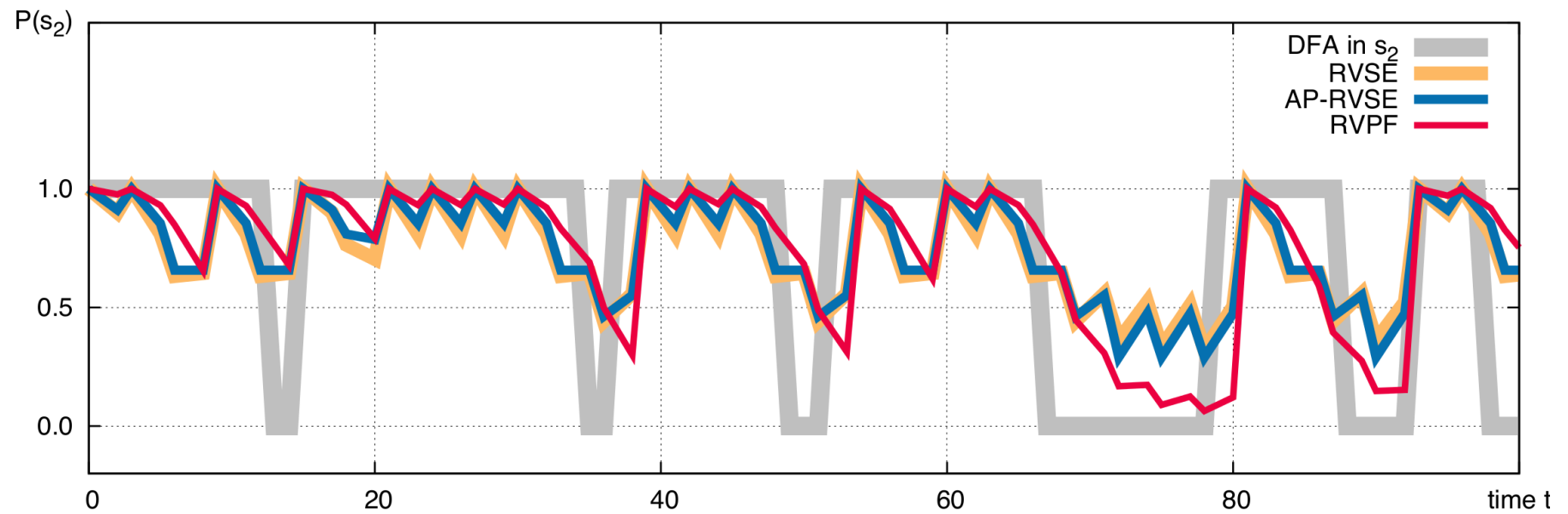
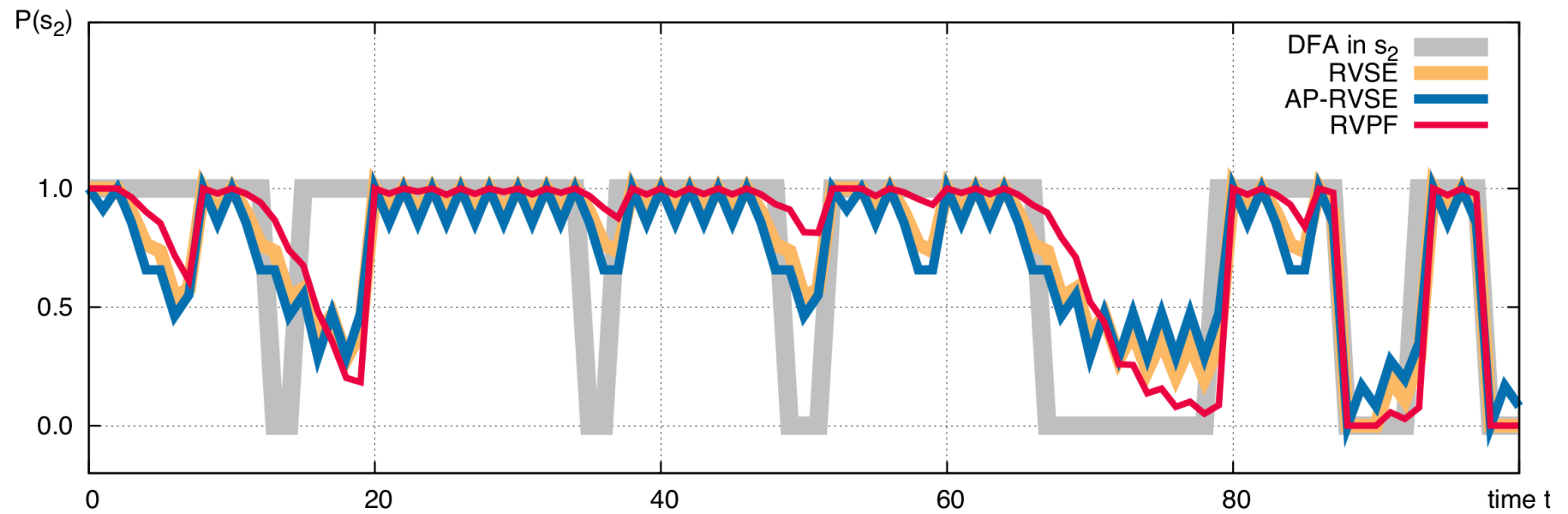
RVPF: Observe a Peek p

Check O
In every x_i

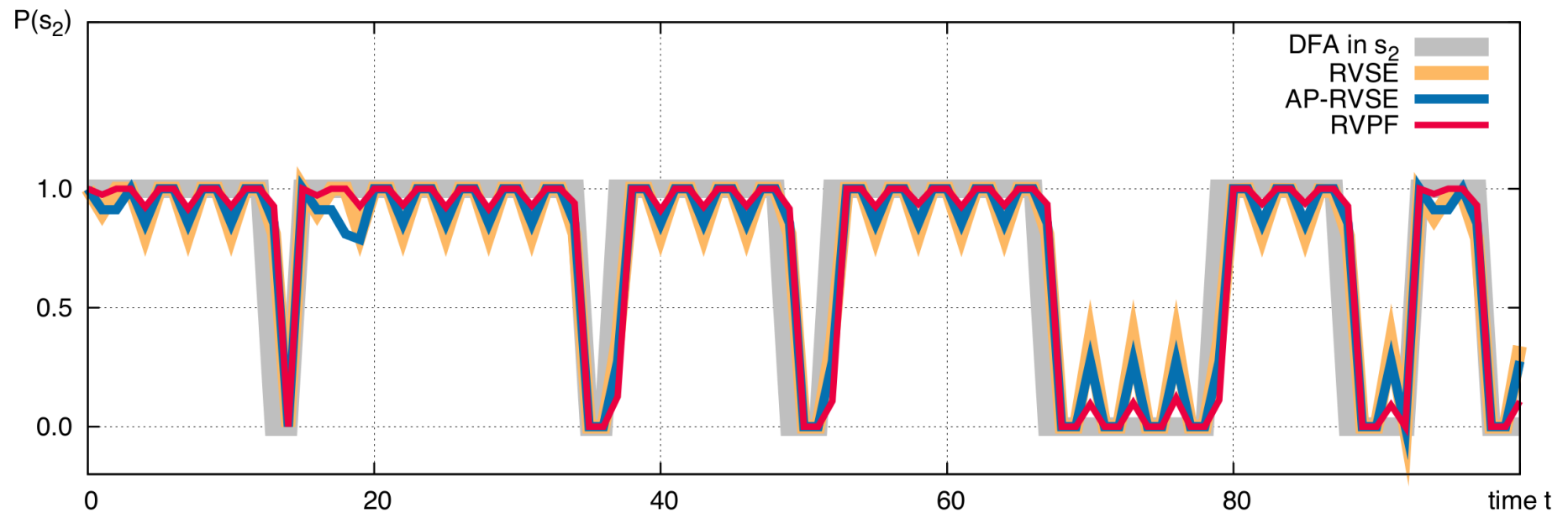


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Accuracy Evaluation

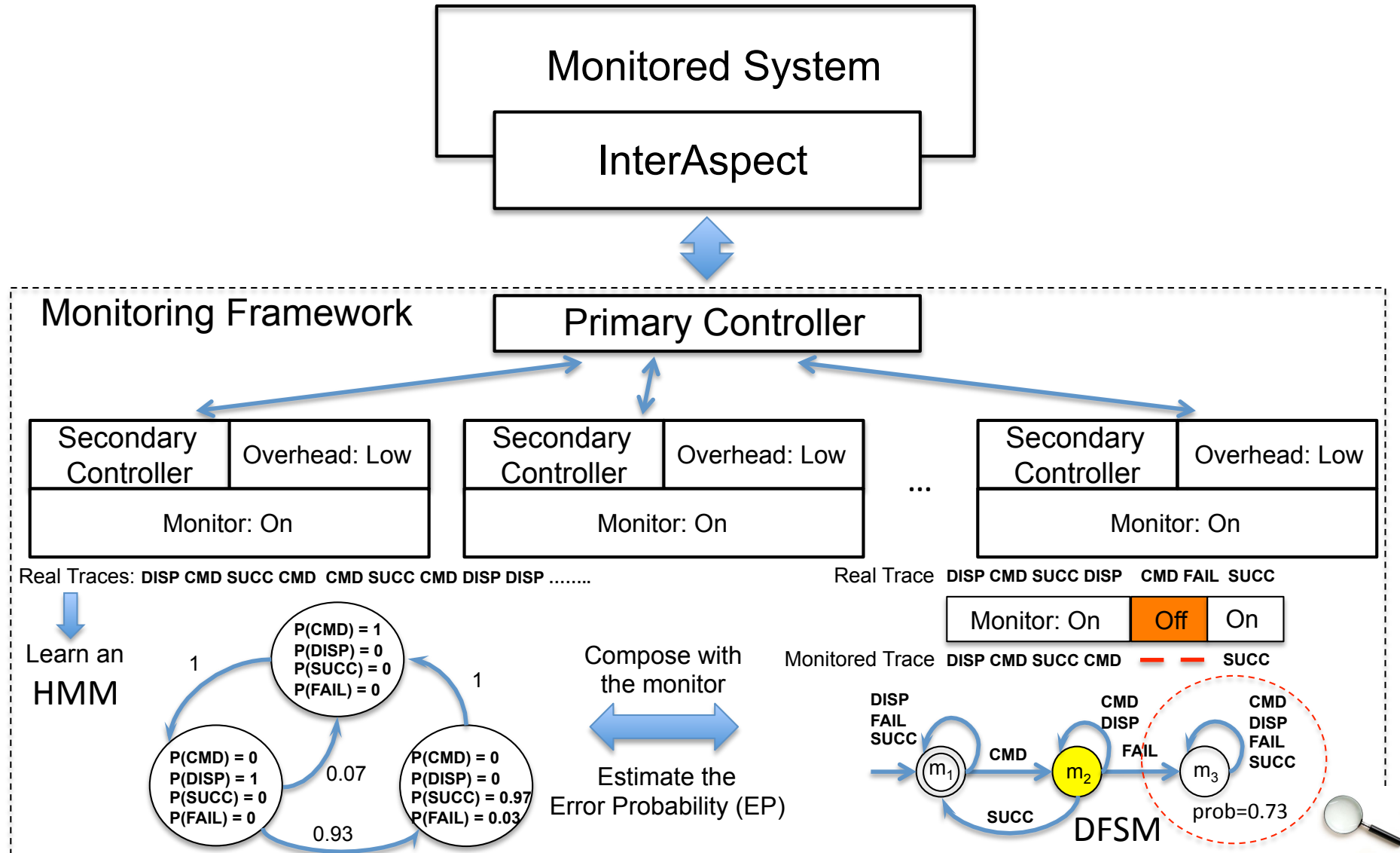


Accuracy Evaluation: Peeking



Peek after gap: significantly improves accuracy

Adaptive Runtime Verification Framework



Adaptive Runtime Verification Framework

