2.

Symbol	Address	External
Main	00400000	No
In_string	n/a	Yes
Out_string	n/a	Yes
Fib	n/a	Yes

3.

Address	Entry
0x00400004	lui \$a0, upper(in_string)
0x00400008	ori \$a0, lower(in_string)
0x00400024	jal fib
0x00400034	lui \$a0, upper(out_string)
0x00400038	ori \$a0, lower(out_string)

5.

Symbol	Address	External
epilogue	0x00400064	No
do_recurse	0x0040002c	No
in_string	0x10010000	Yes
out_string	0x1001000c	Yes
fib	0x00400000	Yes

6.

Address	Entry
0x00400038	Jal Fib
0x0040004c	Jal Fib

Bkpt	Address	Code	Basic			Source
	0x00400000	0x24020004	addiu \$2,\$0,4	8:	li \$v0, 4	# system call code for printing string = 4
	0x00400004	0x3c011001	lui \$1,4097	9:	la \$a0, in string	# load address of string to be printed into \$a0
	0x00400008	0x34240000	ori \$4,\$1,0			
	0x0040000c	0x0000000c	syscall	10:	syscall	# call operating system to perform print operation
	0x00400010	0x24020005	addiu \$2,\$0,5	13:	li \$v0, 5	# system call code for read integer = 5
	0x00400014	0x0000000c	svscall	14:	svscall	# call operating system
			addu \$16,\$0,\$2	15:	move \$s0, \$v0	# value read from keyboard returned in register \$v0; transfer to \$s0
			sw \$16,0(\$29)	17:	sw \$s0, (\$sp)	# push argument for Fib on stack
			addi \$29,\$29,-4	18:	addi \$sp,\$sp,-4	# and decrement stack pointer
			jal 0x00400050	19:	jal Fib	# jump to subroutine
			addi \$29,\$29,4	20:	addi \$sp,\$sp,4	# increment stack pointer
			lw \$17,0(\$29)	21:	lw \$s1, (\$sp)	# and pop result from stack
			addiu \$2,\$0,4	24:	li \$v0, 4	# system call code for printing string = 4
			lui \$1,4097	25:	la \$a0, out string	# load address of string to be printed into \$a0
			ori \$4,\$1,28	20.		and danced of boring of he printed and the
	0x0040003c			26:	syscall	# call operating system
			addiu \$2,\$0,1	29:	li \$v0, 1	# system call code for printing integer = 1
			addu \$4,\$0,\$17	30:	move \$a0, \$s1	# move integer to be printed into \$a0: \$a0 = \$s1
	0x00400048			31:	syscall	# call operating system to perform print
	0x0040004c			32.	ir Sra	" our operating byboth to perform print
			sw \$31,0(\$29)	24:	sw \$ra,(\$sp)	# save return address on stack, since recursive,
			addi \$29,\$29,-4	25:	addi \$sp.\$sp4	# and decrement stack pointer
			sw \$30,0(\$29)	26:	sw \$fp, (\$sp)	# save previous frame pointer on stack
			addi \$29.\$294	27:	addi \$sp.\$sp4	# and decrement stack pointer
			addi \$30,\$29,12	28:	add \$fp,\$sp,12	# set frame pointer to point at base of stack frame
			lw \$8.0(\$30)	30:	lw \$t0,(\$fp)	# copy argument to \$t0: \$t0 = n
			addiu \$9,\$0,2	31:	li \$t1, 2	# copy argument to etc. etc - n
			slt \$1,\$9,\$8	32:		# if argument n >= 2, branch to recursive sequence
			bne \$1,\$0,2	52.	bgc vco,vci,do_iecaise	# II digament in >= 2, branch to recursive sequence
			addiu \$8,\$0,1	33:	li \$t0, 1	<pre># else set result to 1 (base cases n = 1 and n = 2)</pre>
	0x00400074			34:	b epiloque	# branch to end
			addi \$8,\$8,-1	37:	addi \$t0,\$t0,-1	# \$t0 = n-1
			sw \$8,0(\$29)	38:	sw \$t0, (\$sp)	# push argument n-1 on stack
			addi \$29,\$29,-4	39:	addi \$sp,\$sp,-4	# and decrement stack pointer
			jal 0x00400050	40:	jal Fib	# call Fibonacci with argument n-1
			lw \$8,0(\$30)	42:	lw \$t0, (\$fp)	# re-copy argument to \$t0: \$t0 = n
			addi \$8,\$8,-2	43:	addi \$t0,\$t0,-2	# \$t0 = n-2
			sw \$8,0(\$29)	44:	sw \$t0, (\$sp)	# push argument n-2 on stack
			addi \$29,\$29,-4	45:	addi \$sp,\$sp,-4	# and decrement stack pointer
			ial 0x00400050	46:	ial Fib	# call Fibonacci with argument n-2
			addi \$29.\$29.4	47:	addi \$sp,\$sp,4	# increment stack pointer
			lw \$8,0(\$29)	48:	lw \$t0, (\$sp)	# and pop result of Fib(n-2) from stack into \$t0
			addi \$29,\$29,4	49:	addi \$sp,\$sp,4	# increment stack pointer
			lw \$9,0(\$29)	50:	lw \$t1, (\$sp)	# and pop result of Fib(n-1) from stack into \$t1
			add \$8,\$8,\$9	51:	add \$t0,\$t0,\$t1	# \$t0 = Fib(n-2) + Fib(n-1); have result
			addi \$29,\$29,4	54:	addi \$sp,\$sp,4	# increment stack pointer
			lw \$30,0(\$29)	55:	lw \$fp, (\$sp)	# and pop saved frame pointer into \$fp
			addi \$29,\$29,4	56:	addi \$sp,\$sp,4	# increment stack pointer
			lw \$31.0(\$29)	57:	lw \$ra, (\$sp)	# and pop return address into \$ra
			addi \$29,\$29,4	58:	addi \$sp,\$sp,4	# increment stack pointer
			sw \$8,0(\$29)	60:	add1 \$5p,\$5p,4 sw \$t0,(\$5p)	# push result onto stack
			3W \$0,0(\$29)	61:	addi can can _/	# push result onto stack
			274 274 27		The same of the sa	N. M.

8. Those lines were the only thing that changed. What the assembler does is it concatenates the two programs together to create a new program. The first part comes from lab5part1 and the second part goes after that and thus all the addresses of the lab5-part2's code gets offset by the length(part1)*4. The assembler also does this with the data section by combining the two data sections, but in lab5-part1.asm there's nothing in there so the data section remains the same.

9.

Label	Address ▲
(global)	
Fib	0x00400050
in_string	0x10010000
out_string	0x1001001c
lab5-par	t1.asm
main	0x00400000
lab5-par	t2.asm
do_recurse	0x0040007c
epilogue	0x004000b4

10.

Address	Entry
00400004	Lui \$1, 4097
00400008	Ori \$4,\$1,0
00400024	Jal Fib
00400088	Jal Fib
0040009c	Jal Fib
00400034	Lui \$1, 4097
00400038	Ori \$4,\$1,28

11.

The two files could've been combined into one file. The makes it so that the linker doesn't have to do any work to link the files together and make it so that the two files are not interdependent.

12.

This is the same program that you designed in lab 4. Compile the two files, and examine their object files. Look at the symbol tables for both files and fill in the following table. Write 'UND' for undefined symbols, and write 'N/A' for symbols not present in a particular file. Also include the section (.data, .text, or another section) for each symbol.

Symbol	Address in file 1	Address in file 2		2	Address in linked file Section
print_int	00000000	Γ	UND		00400570 text
print_string	00000010		UND		00400580 text
read_int	00000020	Γ	UND		00400590 text
prodMessage	UND	Γ	UND		00410800 data
my_mul	N/A	Γ	00000000		004005a0 text
main	N/A		00000018		004005b8 text

1.

The only instruction that changed was the j loop instruction. The native instruction changed because the nop instructions at the beginning offset the rest of the program because they took up addresses and thus the jump loop had to change to reflect the change in address of the jump. The la \$t0, n didn't change because the nop instructions didn't affect the static data of the program.