

Arrays (and Hashmaps)

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Chapter 1: Overview

1.1 Arrays

1.1.1 Big-O Notation

- Space: $O(n)$
- Search: $O(n)$
- Access: $O(1)$
- Insertion: $O(n)$
- Deletion: $O(n)$

1.1.2 Creation

Listing 1.1: Arrays in Python

```
# Create empty list
my_nums = []

# Create list with values
my_nums = [5,6,7]

# Create list from string
chars = list('Hello')
# chars => ['h', 'e', 'l', 'l', 'o']

# List from set
unique_set = set(1,2,3)
unique_list = list(unique_set)
# unique_list => [1, 2, 3]

// Size of unique_list
arrLen = len(unique_list)
```


1.2 Hashmaps

1.2.1 Big-O Notation

- Space: $O(n)$
- Get (Access): $O(1)$
- Contains Key (Access): $O(1)$
- Insertion: $O(1)$
- Delete Key: $O(1)$

1.2.2 Creation

Listing 1.4: Dicts (maps) in Python

```
# Create empty dict
fruits = {}

# Add item to dict
fruits['apple'] = 7

# Create dict with items in it
fruits = {
    'apple': 10,
    'banana': 2
}

# Create dict with 'dict' builtin func
cat_names = dict(one='Ham', two='Beanbag')
print(cat_names['two']) # => Beanbag

# You can set a default value for keys
# using a default dict so you don't
# get an error when trying to access
# an undefined key.
#
# This is very helpful for
# frequency lists.
from collections import defaultdict

values = defaultdict(int)
values['a'] = 10

print(values['a']) # => 10
print(values['b']) # => 0
```

Chapter 2: Algorithms

2.1 Binary Search

Binary Search is a popular algorithm often used to find a target in an already-sorted list. The methodology goes as such:

- Create two variables that will act as a **left pointer** and a **right pointer**, with left initialized to 0 to represent the beginning of the array and right initialized to the last index of the array (size - 1)
- Start a while-loop for condition `left <= right`
- Calculate the middle index between left and right:
`middle = left + (right - left) / 2`
- Check if we found target, if so return
- Check if current middle item is smaller than target. If so, discard the smaller (left-hand) half, and set our left pointer to 1 after the middle:
`left = middle + 1`
- Otherwise, our current item is larger than the target. So we discard the larger half, and set our right pointer to 1 less than the middle:
`right = middle - 1`

The time complexity is $O(\log(n))$. This is because at every iteration of the loop, we are halving what part of the array we're looking at.

Listing 2.1: Binary Search In Python

```
def binary_search(items, target):
    left = 0
    right = len(items) - 1

    while left <= right:
        # // in Python rounds to the nearest integer
        mid = (left + right) // 2

        if items[mid] == target:
            return mid
        elif items[mid] < target:
            left = mid + 1
        else items[mid] > target:
            right = mid - 1

    return -1

binary_search([2,3,4,5,6,7,8,9], 3)
```


2.2 Sliding Window

Sliding Window is a technique used to search or traverse an array while looking at a subsection of the array (a "window") at a time. The goal of sliding window algorithms is to minimize nested loops and reduce their time complexities. These algorithms are used often for the following types of problems:

- Minimum / Maximum Sum Array
- Longest Sequence / Substring
- K Closest Elements

Assume we have a problem where we want to find the largest subarray in a given array:

Listing 2.3: Sliding Window In Python

```
# k is the target size of the subarray
def maximum_subarray(nums, k):
    # Set the current maximum to lowest number
    max_sum = float('-inf')

    # Keep track of local sums
    current_sum = 0

    # Calculate the initial max sum
    # in the first window (first k items)
    for i in range(k):
        current_sum += nums[i]

    # Loop through the rest of array
    for i in range(k, len(nums)):
        # Update the largest subarray that we've seen
        max_sum = max(max_sum, current_sum)

        # Add the current num to the current_sum
        # Subtracting the first item from the last window
        current_sum += nums[i] - nums[i - k]

    return max_sum
```

```
maximum_subarray([5,8,3,6,9,1,0,9], 3)
```

```
"""
```

```
So for the above example, we effectively end up seeing:
```

```
Step 1:
```

```
Our window is [5,8,3], the first k elements.
```

```
The pointer is at index 3 (value of k), but we only look
at the values before that (exclusive).
```


2.3 Two Sum

The two sum algorithm is a simple one used when needing to find if two numbers in a list can be added to equate to a given target.

For the simplest form of two sum (exemplified below), the algorithm goes as such:

- Loop through input
- Calculate the difference between the current item and the target
- If the complement is in the hashmap, return the current index and the index of the complement
- Add the current number to a hashmap as a key, whose value is the index

This can be done in $O(n)$ time. The reason this works is because at every iteration of our loop, we're seeing if we already came across a number (an appropriate complement of the current number) that will make the current number equal our target when added together.

Listing 2.5: Two Sum in Python

```
def two_sum(nums, target):
    complements = {}
    for i in range(len(arr)):
        diff = target - nums[i]

        if diff in complements:
            return [i, complements[diff]]
        complements[nums[i]] = i

two_sum([9,6,11,15], 17)
```

Listing 2.6: Two Sum in C++

```
vector<int> twoSum(vector<int>& nums, int target) {
    map<int, int> complements;
    vector<int> result;

    for(int i = 0; i < nums.size(); i++) {
        if (complements.count(target - nums[i]) > 0) {
            return {i, complements[target - nums[i]]};
        }
        complements.insert(pair<int,int>(nums[i], i));
    }

    return {-1, -1};
}
```

