

# NCKU Programming Contest Training Course Math 2016/03/09

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#### **Greatest Common Divisor**



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• 輾轉相除法 (Euclid's Algorithm)

```
int gcd(int a, int b)
{
   if (a == 0) return b;
   return gcd(b % a, a);
}
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```

- a x b = GCD(a, b) x LCM(a, b)
- GCB(a, b) = GCD(b, a % b)



# Practice

### UVA 408 Uniform Generator



- aX + bY = GCD(a, b), find a particular solution of (X, Y).
- Ex: 12X + 9Y = 3, (X, Y) = (1, -1)



- aX + bY = GCD(a, b), find a particular solution of (X, Y).
- Ex: 12X + 9Y = 3, (X, Y) = (1, -1)
- Theorem: GCD(a, b) = GCD(b, a % b)

• 
$$aX + bY = GCD(a, b) = GCD(b, a\%b) = bX' + (a\%b)Y'$$
  
 $aX + bY = bX' + [a - (a/b)b]Y' = aY' + b(X' - (a/b)Y')$   
 $X = Y'$   
 $Y = X' - (a/b)Y'$ 



```
int exGCD(int a, int b, int &X, int &Y)
   if (b == 0) {
       X = 1;
       Y = 0;
        return a;
    } else {
        int gcd = exGCD(b, a % b, X, Y);
        int tmp = X;
        X = Y;
        Y = tmp - (a/b)*Y;
        return gcd;
```



找出等式 AX + BY = N 的所有解:

- 1. 確認 N 是否能整除 GCD(A, B) - 若不能 => 無解
- 2. 等式兩邊同除以GCD(A, B),變成 aX + bY = n,其中 GCD(a, b) = 1
- 3. 使用 exGCD 找出特定解(x, y)
- 4. 則可以得到 X = n \* x + b \* t Y = n \* y - a \* t 其中 t 為整數

ex: 
$$12X + 9Y = 9$$
  
 $4X + 3Y = 3$   
exGCD(4, 3) =>  $(x, y) = (1, -1)$   
 $(X, Y) = (3 + 3t, -3 - 4t)$ 

NCKU Programming Contest Training Course



# Practice

### UVA 10104 Euclid Problem







- 用篩選方式建質數表
  - 由小到大選擇質數,並刪除其倍數



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```
#include <cstring>
#include <cmath>
#define MAX 1000000
bool isprime[MAX];
void Sieve()
    memset(isprime, true, sizeof(isprime));
    isprime[0] = isprime[1] = false;
    for (int i = 2; i \leftarrow sqrt(MAX); ++i)
        if (isprime[i])
            for (int j = i + i; j < MAX; j += i)
                isprime[j] = false;
```





質數從 5 開始後,只會出現在+2,+4,+2,+4…的數字上



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```
#define MAX 1000000
vector<int> prime;
bool isPrime(int n)
    for (int i = 0; prime[i] * prime[i] <= n; ++i)</pre>
        if (n % prime[i] == 0)
            return false;
    return true;
void MakePrime()
    prime.push_back(2);
    prime.push_back(3);
    for (int i=5, gap=2; i < MAX; i+=gap, gap=6-gap)</pre>
        if (isPrime(i)) prime.push back(i);
```



- 方法二比方法一慢,但較省空間
- 其他方法: 演算法筆記 - Prime
- 但是前頁的方法已經可以解幾乎所有的題目



### Practice

### UVA 10392 Factoring Large Numbers









#### Float:

- 數值範圍: -3.4e-38~3.4e38
- 十進位精確度位數: 6~7

#### Double:

- 數值範圍: -1.7e-308~1.7e308
- 十進位精確度位數: 14~15



### **Epsilon**

#### Example:

```
#include <cstdio>
#include <cmath>
int main()
    double a = asin(sqrt(2.0) / 2) * 4.0;
    double b = acos(-1.0);
    printf("a = \%.201f\n", a);
    printf("b = \%.20lf(n", b);
    printf("a - b = \%.201f\n", a - b);
    printf("a == b? %s\n", a == b ? "True" : "False");
```





#### Result:

```
a = 3.14159265358979360000
b = 3.14159265358979310000
a - b = 0.00000000000000044409
a == b? False
```



# **Epsilon**

引入eps,來判斷浮點數是否相等

eps = 1e-8



# **Epsilon**

#### 引入eps,來判斷浮點數是否相等

$$eps = 1e-8$$

整數	浮點數
a == b	a - b  < eps
a != b	a - b  > eps
a < b	a – b < -eps
a > b	a – b > eps



### Practice

### UVA 906 Rational Neighbor







- Array
- 習慣將低位數放在 index 比較小的位置



- Array
- 習慣將低位數放在 index 比較小的位置
- ex: 680468975231245

• 右方補0



- 加法:位數各自相加後,由低至高位依序進位
- 減法: 位數各自相減後,由低至高位依序借位
- 乘法: 直式乘法
- 除法: 長除法



Ex:加法

```
void add(int a[100], int b[100], int c[100])
    for (int i = 0; i < 100; ++i)
        c[i] = a[i] + b[i];
    for (int i = 0; i < 100-1; ++i) {
        c[i+1] += c[i] / 10;
        c[i] %= 10;
```



# Practice

### UVA 10106 Product



# **Problem List**

#### **UVA**

113, 11538, 10035, 338, 375, 10494, 746, 12404, 906, 10867, 10311, 11408, 10407, 10606, 10090