CMPS 102 — Fall 2018 — Homework 4

"I have read and agree to the collaboration policy." - Kevin Wang

Collaborators: None

Solution to Problem 2: DNA

2 DNA sequences $S=\{s_1,s_2,\cdots,s_i,\cdots,s_m\}$ and $T=\{t_1,t_2,\cdots,t_j,\cdots,t_n\}$ have lengths m and n respectively.

Let LCS be the longest common sub-sequence.

LCS[i, j] is the length of the largest common sub-sequence, LCS, shared by two sequences of sizes i and j.

Base Case: An empty sequence does not share a common sub-sequence.

$$LCS[i, 0] = 0$$
 and $LCS[0, j] = 0$

Case 1: $s_i = t_j$; the element is a part of the LCS.

$$LCS[i, j] = 1 + LCS[i - 1, j - 1]$$

Case 2: $s_i \neq t_j$; we ignore one of the elements.

$$LCS[i, j] = \max(LCS[i-1, j], LCS[i, j-1])$$

Therefore, the LCS of S and T has length LCS[m][n]. We can then walk back through the matrix to find an LCS.

Continued...

Algorithm 1 Finds the LCS of of two sequences S and T with sizes m and n, respectively

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LCS (m, n):
Let LCS[m][n] be the sub-sequence size matrix
\mathbf{for}\ i=0\ \mathsf{to}\ m\ \mathbf{do}
  for j = 0 to n do
    if i = 0 or j = 0 then
       LCS[i][j] = 0
     else if s_i = t_j then
       LCS[i][j] = LCS[i-1][j-1] + 1
    else
       LCS[i][j] = \max(LCS[i-1][j], LCS[i][j-1])
     end if
  end for
end for
LCS[m][n] is the length of the LCS
Let LCS be an empty sub-sequence
Let i = m and j = n
while i > 0 and j > 0 do
  if s_i = t_j then
    s_i :: LCS
    i - -, j - -
  else if LCS[i-1][j] \ge LCS[i][j-1] then
     ***Using\ LCS[i-1][j] > LCS[i][j-1] can return another different LCS***
    i - -
  else
    j - -
  end if
end while
```

Time Complexity: O(mn)

The 2 for-loops iterate m and n times, respectively, taking time O(mn). Iterating back through the matrix to find the elements of the LCS takes O(m+n). Total time used is O(mn+m+n).

Space Complexity: O(mn)

The LCS matrix takes size O(mn). The LCS sequence, takes (at most) $O(\max(m, n))$. Total space used is $O(mn + \max(m, n))$.