

# 1 February 25th, 2021

## 1.1 Gauss-Seidel Iteration

Recall that in Jacobi,  $\xi_i^{(k+1)}$  are updated parallelly in Algorithm ??.

**Gauss-Seidel** is just modifying Jacobi to be updated “successively”. When we modify  $\xi_i^{(k+1)}$  we can use  $\xi_1^{(k+1)}, \dots, \xi_{i-1}^{(k+1)}$ . This gives us Algorithm 1.

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### Algorithm 1 Gauss-Seidel Iteration

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1: for  $k = 0, 1, 2, \dots$  do
2:   for  $i = 1, \dots, n$  do
3:      $\xi_i^{(k+1)} = (\beta_i - \sum_{j>i} a_{ij}\xi_j^{(k)} - \sum_{j<i} a_{ij}\xi_j^{(k+1)})/a_{ii}$ 
4:   end for
5: end for

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In the matrix reformulation, recalling that  $A = D - E - F$ , where  $E$  is lower triangular and  $F$  is upper triangular, we have:

$$\begin{aligned}
 x_{k+1} &= D^{-1}(b + Fx_k + Ex_{k+1}) \\
 \iff (D - E)x_{k+1} &= b + Fx_k \\
 \iff x_{k+1} &= (D - E)^{-1}Fx_k + (D - E)^{-1}b \\
 \iff x_{k+1} &= \underbrace{(I - (D - E)^{-1}A)}_G x_k + \underbrace{(D - E)^{-1}b}_f.
 \end{aligned}$$

which is in the stationary iteration form.

**Remark 1.1** — In the algorithm above, we are updating  $\xi_i^{(k+1)}$  in regular numerical order (from 1 to  $n$ ). However, we can use other ordering, such as going from  $n$  to 1 in reverse order. Gauss-Seidel is sensitive to the ordering of updating of unknowns.

### 1.1.1 Convergence of Gauss-Seidel

Since Gauss-Seidel is a stationary iteration, we can use the same framework as the Jacobi iteration, with:

$$(x_k - x_*) = G^k(x_0 - x_*)$$

#### Theorem 1.2

Gauss-Seidel converges to  $x_*$  for any  $x_0$  if and only if  $\rho(G) < 1$ .

The convergence speed is  $\|G\|$ , as we have:

$$\|x_{k+1} - x_*\| \leq \|G\|\|x_k - x_*\|$$

In particular, we have:

$$\|G\| = \begin{cases} \rho(G) & \text{if } G \text{ is symmetric and } \|\cdot\| = \|\cdot\|_2 \\ \rho(G) + \epsilon & \text{otherwise} \end{cases}$$

with  $\epsilon$  arbitrarily small.

### Theorem 1.3

Let  $A \in \mathbb{R}^{n \times n}$  be SPD. Then  $\rho(G) < 1$ .

*Proof.* Recall  $G = I - (D - E)^{-1}A$ . Let  $(\lambda, u)$  be an eigenpair of  $G$  (note that  $\lambda, u$  may be complex). We have:

$$\begin{aligned} Gu &= \lambda u \\ (I - (D - E)^{-1}Au) &= \lambda u \\ (D - E)^{-1}Fu &= \lambda u \\ Fu &= \lambda(D - E)^{-1}u. \end{aligned}$$

Since  $A$  is symmetric, we have:

$$A = A^T \implies F = E^T.$$

Thus, we have:

$$\begin{aligned} \lambda(D - E)u &= E^T u \\ \lambda u^*(D - E)u &= u^* E^T u. \end{aligned}$$

by left-multiplying both side by  $u^* = \bar{u}^T$

Let us set  $u^* E^T u = \alpha + i\beta \in \mathbb{C}$  and  $u^* Du = \delta \in \mathbb{R}$ , since  $D$  is real and symmetric. Then, we have:

$$u^* E^T u = (u^* E^T u)^T = u^T E (u^*)^T = \overline{u^* E u} = \overline{u^* E u} = \alpha - i\beta$$

Since  $A$  is SPD,  $D$  is also SPD, giving us  $\lambda = u^* Du > 0$ .

Thus

$$\lambda(\delta - (\alpha + i\beta)) = \alpha - i\beta \implies \lambda = \frac{\alpha - i\beta}{(\delta - \alpha) - i\beta}$$

Giving us:

$$|\lambda|^2 = \frac{(\alpha^2 + \beta^2)}{(\delta - \alpha)^2 + \beta^2}$$

Because:

$$(\delta - \alpha)^2 + \beta^2 = \delta^2 + \alpha^2 + \beta^2 - 2\alpha\delta = (\alpha^2 + \beta^2) + \underbrace{\delta}_{>0}(\delta - 2\alpha)$$

we need to check the sign of  $\delta - 2\alpha$ .

Since  $A$  is SPD,  $u^* Au > 0 \implies u^*(D - E - E^T)u = \delta - (\alpha + i\beta) - (\alpha - i\beta) = \delta - 2\alpha > 0$ . As such:

$$(\delta - \alpha)^2 + \beta^2 > \alpha^2 + \beta^2 \implies |\lambda|^2 < 1$$

□

**Remark 1.4** — This is more general than Jacobi, which is convergent if and only if  $2D - A$  is SPD.

### Example 1.5

Recall that the 1D discrete Laplacian is SPD, meaning that Gauss-Seidel converges. Similarly for the 2D case.

### Example 1.6

Gauss-Seidel converges for irreducible diagonally dominant matrices. This is used for laplacian on irregular domains.

## 1.2 Acceleration of G-S / and Jacobi (SOR)

Here we try to accelerate Gauss-Seidel and Jacobi by considering momentum acceleration. We can consider the iteration of  $x_k$  to  $x_{k+1}$  as a movement by the Gauss-Seidel iteration. The idea of SOR is to consider the momentum of the movement until some stage.

For Gauss-Seidel, the components are:

- $x_k : \xi_i^{(k)}$
- $x_{k+1} : \xi_i^{(k+1)} = (\beta_i - \sum_{j>i} a_{ij}\xi_j^{(k)} - \sum_{j<i} a_{ij}\xi_j^{(k+1)})/a_{ii}$

To consider the momentum, we would have:

$$\begin{aligned}\xi_i^{(k+1)} &= \xi_i^{(k)} + \omega(\xi_i^{(k+1)} - \xi_i^{(k)}) \\ &= \xi_i^{(k)} + \omega \left( (\beta_i - \sum_{j>i} a_{ij}\xi_j^{(k)} - \sum_{j<i} a_{ij}\xi_j^{(k+1)})/a_{ii} - \xi_i^{(k)} \right).\end{aligned}$$

for some  $\omega > 1$ , giving us the Algorithm 2.

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### Algorithm 2 SOR Iteration

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1: for  $k = 0, 1, 2, \dots$  do
2:   for  $i = 1, \dots, n$  do
3:      $\xi_i^{(k+1)} = \xi_i^{(k)} + \omega \left[ (\beta_i - \sum_{j>i} a_{ij}\xi_j^{(k)} - \sum_{j<i} a_{ij}\xi_j^{(k+1)})/a_{ii} - \xi_i^{(k)} \right]$ 
4:   end for
5: end for
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In the matrix reformulation, we have:

$$x_{k+1} = x_k + \omega [D^{-1}(b + Fx_k + Ex_{k+1}) - x_k]$$

Expressing this in its stationary form, we have:

$$x_{k+1} = \underbrace{(D - \omega E)^{-1} ((1 - \omega)D + \omega F)}_{G_\omega} x_k + \underbrace{\omega(D - \omega E)^{-1} b}_{f_\omega}$$

Now the question is finding a suitable  $\omega$ .

## 1.2.1 Convergence of SOR

**Theorem 1.7**

If  $\omega \notin (0, 2)$ , then  $\rho(G_\omega) > 1$ .

*Proof.* We have:

$$\det(G_\omega) = \det((D - \omega E)^{-1}) \cdot \det((1 - \omega)D + \omega F)$$

- Since  $D - \omega E$  is lower triangular,  $(D - \omega E)^{-1}$  is also lower triangular. Thus,  $\det((D - \omega E)^{-1})$  is the product of its diagonals, which are the inverse of the diagonal of  $D - \omega E$  which are  $\frac{1}{a_{ii}}$ . As such, we have:

$$\det((D - \omega E)^{-1}) = \prod_{i=1}^n a_{ii}^{-1}$$

- Since  $(1 - \omega)D + \omega F$  is upper triangular, we have:

$$\det((1 - \omega)D + \omega F) = \prod_{i=1}^n (1 - \omega) a_{ii}$$

Thus, we have:

$$\det(G_\omega) = \prod_{i=1}^n (1 - \omega) a_{ii} \cdot a_{ii}^{-1} = (1 - \omega)^n = \prod_{i=1}^n \lambda_i$$

This gives us:

$$\begin{aligned} |1 - \omega|^n &= |\lambda_1| |\lambda_2| \dots |\lambda_n| \leq (\rho(G_\omega))^n \\ \implies |1 - \omega| &\leq \rho(G_\omega). \end{aligned}$$

□

**Corollary 1.8**

In order for SOR to converge, we must have  $\omega < 2$ .

- If  $\omega = 1$ , we have Gauss Seidel
- If  $\omega \in (0, 1)$ , we have successive under relaxation
- If  $\omega \in (1, 2)$  we have successive over relaxation (SOR).

**Theorem 1.9**

If  $A$  is SPD, then  $\rho(G_\omega) < 1$  for all  $\omega \in (0, 2)$ .

*Proof.* Similar to Gauss-Seidel.

□

How do we pick the optimal  $\omega$ ? Roughly speaking, with the optimal  $\omega$ :

- $\rho(G_{\omega_{\text{opt}}}) \sim \rho(G_1)^2$ . Where  $G_1$  is the iteration matrix in G-S.

**Remark 1.10** — This means that SOR is roughly two times faster than G-S, since it is the square.

- $\rho(G_1) \sim (\rho(G_{\text{Jacobi}}))^2$ .