1 November 12th, 2020

1.1 Iterative Methods to Solve Linear Systems

Last time, we introduced three methods, Jacobi, Gauss-Seidel, and SOR. Let us consider the Jacobi iteration. We want to show that it is irreducibly diagonally dominant matrix.

Proof. Assume A is irreducibly diagonally dominant. Recall that the Jacobi iteration is of the form

$$G = D^{-1}(L+U).$$

We want to show that $\rho(G) < 1$. Assume λ is an eigenvalue of G. We need to now show that $|\lambda| < 1$. We have:

$$Gv = \lambda v$$

$$\implies D^{-1}(L+U)v = \lambda v$$

$$\implies (L+U)v = \lambda Dv$$

$$\implies (\lambda D - L - U)v = 0.$$

Writing down each row explicitly, we have:

$$\lambda A_{kk}v_k + \sum_{\ell \neq k} A_{k\ell}v_\ell = 0 \quad \forall k.$$

$$\lambda A_{kk} v_k = -\sum_{\ell \neq k} A_{k\ell} v_\ell \quad \forall k.$$

Taking absolute value on both sides, we have:

$$|\lambda||A_{kk}||v_k| = |\sum_{\ell \neq k} A_{k\ell} v_{\ell}| \le \sum_{\ell \neq k} |A_{k\ell}||v_{\ell}|.$$

If we choose $|v_k| = \max_i |v_i| \neq 0$, we would have:

$$|\lambda| \le \sum_{\ell \ne k} \frac{|A_{k\ell}|}{|A_{kk}|} \frac{|v_{\ell}|}{|v_{k}|} \le \frac{1}{|A_{kk}|} \sum_{\ell \ne k} |A_{k\ell}| \le 1.$$

The last inequality is because we assume that A is diagonally dominant.

If $|\lambda| = 1$, we would have:

$$1 \le \sum_{\ell \ne k} \frac{|A_{k\ell}|}{|A_{kk}|} \frac{|v_{\ell}|}{|v_k|} \le 1 \implies |v_{\ell}| = |v_k|, \forall \ell.$$

Now for all k, we have:

$$\sum_{\ell \neq k} |A_{k\ell}| = |A_{kk}|.$$

which is a contradiction to the irreducibility of A. As such $|\lambda| < 1$ for all eigenvalues of A, meaning that the Jacobi iteration converges.

As shown last time, if we consider the central difference for -u'' = f, then the matrix is irreducibly diagonally dominant. We can compute G, with:

In addition, we can directly compute the eigenvalues of the original differential operator, as we have:

$$-u'' = \lambda u, \quad u(0) = u(1) = 0.$$

Since the general solution to $-u'' = \lambda u$ is:

$$u = \begin{cases} A\sin(\sqrt{\lambda}x) + B\cos(\sqrt{\lambda}x) & \lambda > 0\\ Ae^{\sqrt{-\lambda}x} + Be^{-\sqrt{-\lambda}x} & \lambda < 0 \end{cases}.$$

Note that the second case does not satisfy the boundary condition, as:

$$u(0) = A + B = 0$$
 $u(1) = Ae^{\sqrt{-\lambda}} + Be^{-\sqrt{-\lambda}} = 0.$
 $\implies Ae^{\sqrt{-\lambda}}(1 - e^{-2\sqrt{-\lambda}}) = 0.$

which is untrue, since $e^{-2\sqrt{-\lambda}} < 1$. As such, $\lambda > 0$. Considering the boundary conditions, we have:

$$u(0) = B = 0.$$

 $u(1) = A\sin(\sqrt{\lambda}) = 0 \implies \lambda_n = (n\pi)^2.$

since $A \neq 0$, otherwise we would get the trivial solution. The eigenvectors are thus:

$$u_n(x) = \sin(n\pi x).$$

Now consider the eigenvalues and eigenvectors of G, which are discretized. We predict that the eigenvectors are of form:

$$u_n = \sin(n\pi x_j) = \begin{bmatrix} \sin(n\pi x_1) \\ \sin(n\pi x_2) \\ \vdots \\ \sin(n\pi x_{N-1}) \end{bmatrix}.$$

where $x_i = jh = \frac{j}{N}$. We want to show that:

$$Gu_n = \lambda_n u_n$$
.

Consider the j-th row of Gu_n , we have:

$$\frac{1}{2} \begin{bmatrix} 0 & 0 & \dots & 1 & 0 & 1 & 0 \dots \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \\ \vdots \\ u_j \\ u_{j+1} \\ \vdots \\ u_{N-1} \end{bmatrix} = \frac{1}{2} (u_{j-1} + u_{j+1}) = \lambda_n u_j.$$

Thus, we have:

$$\frac{1}{2}(\sin(n\pi x_{j-1}) + \sin(n\pi x_{j+1})) = \lambda_n \sin(n\pi x_j).$$

$$\implies \frac{1}{2}(\sin(n\pi(j-1)h) + \sin(n\pi(j+1)h)) = \lambda_n \sin(n\pi jh).$$

Recalling that $\sin(x+y) = \sin x \cos y + \cos x \sin y$, giving us:

 $(\sin(n\pi jh)\cos(n\pi h) - \cos(n\pi jh)\sin(n\pi h)) + (\sin(n\pi jh)\cos(n\pi h) + \cos(n\pi jh)\sin(n\pi h)).$

$$\implies \cos(n\pi h)\sin(n\pi jh) = \lambda_n\sin(n\pi jh).$$

This means that:

$$\lambda_n = \cos(n\pi h).$$

with eigen vectors:

$$u_n(x) = \sin(n\pi jh).$$

Thus, we have:

$$|\lambda_n| = |\cos(n\pi h)| < 1, \quad n = 1, 2, \dots, N - 1.$$

We now want to investigate the rate of iteration, which indicates the number of iterations required to obtain a required accuracy.

Definition 1.1. The rate of convergence is

$$Rln - \ln(\rho)$$
.

Example 1.2

To reduce the error by a factor of 10^p , the number of iterations required is $\frac{P}{R}\ln(10)$.

Let's assume that $G = S\Lambda S^{-1}$, i.e. G is diagonalizable. Recall, we have:

$$x_{n+1} = Gx_n + C.$$

$$x_* = Gx_* + C.$$

We have:

$$x_{n+1} - x_* = G(x_n - x_*)$$

 $\implies e_{n+1} = Ge_n = G^2 e_{n-1} = \dots = G^{n+1} e_0.$

This gives us:

$$||e_n|| \le ||G^n|| ||e_0|| \le ||S\Lambda^n S^{-1}|| ||e_0|| \le ||S|| ||S^{-1}|| ||\Lambda^n|| ||e_0||.$$

As such:

$$\frac{\|e_n\|}{\|e_0\|} \le \|S\| \|S^{-1}\| \|\Lambda\|^n \le \|S\| \|S^{-1}\| \rho^n.$$

Say we want to reduce the error by a factor of 10^p , we have:

$$||S|| ||S^{-1}|| \rho^{n} \le 10^{-p}$$

$$\implies (||S|| ||S^{-1}||)^{\frac{1}{n}} \rho \le 10^{-\frac{p}{n}}$$

$$\implies \frac{1}{n} \ln(||S|| ||S^{-1}||) + \ln \rho \le -\frac{p}{n} \ln 10$$

$$\implies \frac{\ln(||S|| ||S^{-1}||)}{\ln \rho} + n \ge -\frac{p}{\ln \rho} \ln 10$$

$$\implies n \ge \frac{P}{R} \ln 10 - C.$$

where $R = -\ln \rho$ for some constant C.

Remark 1.3 — If ρ is close to 1, then $n \sim \infty$.

Remark 1.4 — If N is very large, then the first eigenvalue will be close to 1, meaning that ρ is close to 1. This is not good, since for accuracy reason we want N to be large, but convergence would be slow.