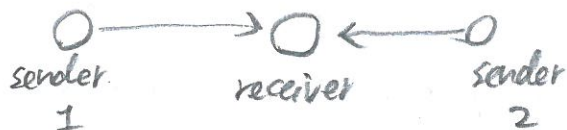
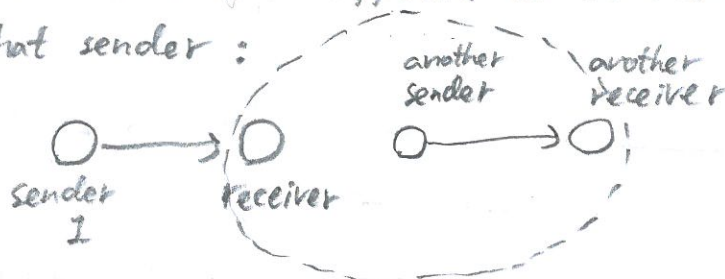


- (1.1) Primary interference refers to the case where there are multiple senders trying to send data to the same receiver at the same time:



- (1.2) Secondary interference refers to the case where a receiver was interfered by a sender for someone else because the receiver happened to be within the transmission range of that sender:



- (1.3) Graph coloring algorithms try to assign mutual colors to adjacent entities (depending on the context, the entities may be vertices, links, geographic regions, etc.). The requirement of a collision-free data communication schedule is to assign mutual time slots to wireless links that may suffer from primary/secondary interferences (because they are adjacent in the sense that at least one wireless node from each link is within the transmission range of another).

P2 (2.2) The contributions are summarized in the fourth paragraph:

- ① a holistic analysis for real-time reliable IoT edge computing;
- ② the ARREC architecture;
- ③ an implementation of ARREC and its empirical performance evaluation.

We may think in this way: an architecture needs to be based on some theoretical analysis that addresses the challenges we are to tackle; the analysis offers guidelines for us to design an architecture. An architecture needs to be implemented in a way that demonstrates the promised performance; the implementation and evaluation offer evidences for us to see the architecture's capability in dealing with the challenges mentioned.

(3.2.1) Here's a reason for the misleading result:

The data in `e2e.delay.misleading` contains some latency results obtained while the broker was still busying accepting new connections. Those results have larger latency. You may verify this by examining the first few lines in `e2e.delay.misleading`. Now, with a longer sampling duration we essentially average out these abnormal results, and therefore the averaged end-to-end latency is shorter.

An important lesson that we may learn here:

A software system often needs some time to "warm up", and thus we should avoid mixing the measurement obtained in the warm-up phase with that obtained in the later phase, to prevent some misinterpretation of the result.

(2.1) This is described in the related work section.

With a higher workload it is likely that $T \uparrow$,
and therefore from $N = \lambda T$ we see that $N \uparrow$,
which implies ^{that} in average there will be more

data in the system at any given point of time.

So the system is more likely to have data losses should
the primary IoT gateway fail to work.

(2.2) In scenarios such as temperature monitoring,
provided that the sampling rate is high, we may
use interpolation (內插法) to have a reasonable
estimation of the lost data.

(2.3) In this case we have the assurance that the system
can always recover some lost data so that there will be
no two consecutive data missings, which by definition
meets the application requirement " $L_i \geq 1$ ".

(3.2) @s prints the number of seconds since the Epoch.

@N prints the number of nanoseconds that have passed in
the current second.

For the last question, the discrepancy could result from
the fact that this public broker is very busy (people around
the world are using it) and thus a message may need to
wait for a while for the broker to complete processing those
messages that have arrived earlier. #