Bigtable: A Distributed Storage System for Structured Data

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One-line Summary

BigTable is a distributed storage system for managing structured data with a simple data model and providing flexible controls over data layout, locality, storage medium, and format, but has poor performance and consistency.

Overview/Main Points

- Data model
 - o a sparse, distributed, persistent multi-dimensional sorted map
 - indexed by a row key, column key, and a timestamp (a 64-bit integer, "real time" in ms)
 - o row
 - row key is 64KB in size, but typical 10-100 bytes
 - every r/w for a single row key is atomic
 - data sorted in lexicographically by row key
 - each tablet contains all data associated with a row range by dynamically partition
 - reversed URL: com.cnn.www
 - column
 - keys group as column families as the basic unit of access control and both disk and memory accounting, and they are always the same type and compress stored.
 - family:qualifier
 - Each tablet is represented by several SSTables
 - 1. key value pairs
 - 2. key offset index
- APIs
 - o create
 - lookup
 - update
 - o delete
 - o scan
- Arch
 - master
 - assign tablets to tablet servers
 - tablet server management: add or expire
 - load balance of tablet-server
 - garbage collection of files in GFS
 - handle schema changes including table and column family creation
 - o tablet server
 - Only one tablet is assigned to one server

- handle r/w to the tablets
- split tablets if too large
- Chubby server
 - a high-available and persistent distributed lock service based on Paxos.
 - 5 active replicas, one of which is master to serve requests if the majority live and can talk to each other.
 - namespace consists of directories and small files, which act as locks
 - Each r/w to files are atomic
 - Chubby client registers callbacks for notification of changes or session expirations.

- o client
 - cache tablet locations

- GFS server for storing logs and data files in SSTable file format
 - SSTable: a persistent, ordered immutable key-value map
 - Each SSTable is represented by a set of blocks associated with indices (load into memory when opened for future lookup in a single disk seek)
- o cluster management system
 - job schedule
 - resource management on shared machiens
 - machine failure recovery, including master restart
 - machine status monitor

Operations

- Lookup
 - performs a three-level search starting from location of the root tablet, which contains the location of all tablets in a METADATA table, which in turn contains the location of user tablets
 - The METADATA table stores the location tablets whose row keys are an encoding of the tablet's table identifier and its end row.

- Read
 - merge data from in-memory Memtables and on-disk SSTables
 - Bloom filters for only necessary reads
- Write
 - write both in-memory Memtables (in a sorted manner) and on-disk commit logs
- Compaction in tablet servers
 - minor compaction: flush in-memory Memtables to on-disk SSTables (GFS)
 - major compaction: merge on-disk SSTables

Recovery

- o Tablet server failure: the same to assign a tablet to a new tablet server
- Master failure: restart by the cluster management system
 - 1. the master acquires a unique master lock in Chubby
 - 2. scan Chubby's servers directory for live servers
 - 3. talks to live servers about their assigned tablets
 - 4. scan the METADATA table for the set of tablets, and processes a tablet that has no associated server
- User's dynamic control
 - Data layout: segregate several column families to one SSTable
 - Locality: leverage row-key range in tablets
 - Storage medium: load all SSTables into memory

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Relevance

Flaws