1 Computational Complexity

(ϵ, δ)-DP *partition selection* is considerably more efficient and easier to implement than current state-of-the-art DP prefix-based query methods. To support large scale applications of our *TraVaS* framework, we will briefly discuss the most important complexity considerations of all associated algorithms (see our paper).

In the preprocessing phase of both the *TraVaS* core and the optimizer shell *uTraVaS*, an event log L needs to be transformed into its aggregated simple form, holding only variant-frequency pairs $(\sigma, L(\sigma))$. This operation can be described as a Group-By-Count query that consumes $\mathcal{O}(n)$ steps with n being the number of cases in L. Although large input logs may lead to runtime bottlenecks, the aggregation must only be executed once independent of further privatization processes. The *TraVaS* core then requires i.i.d *k-TSGD* random samples and variant deletion checks for all m unique variants of L. Hence, we face a separate maximum complexity of $\mathcal{O}(m)$ with each *TraVaS* run. Eventually, our extension *uTraVaS* calls u *TraVaS* instances, implicating a worst-case complexity of $\mathcal{O}(u^2 \cdot m)$ according to the *Gaussian* sum formula. In addition, the same number of steps are needed for all merging and mean operations to obtain the respective candidate logs. Together with u utility computations and the final argmax function, Algorithm 2 (*uTraVaS*) therefore runs with a cumulative complexity of $\mathcal{O}(u^2 \cdot m + n)$. Note, that the cost function might possess a larger complexity than the *TraVaS* body.

One compelling advantage of our algorithm design is that only the initial event log transformation represents a sequential task. All subsequent routines such as the sublog construction, the merging or the cost evaluation can be perfectly parallelized on multiple computing units. As a result, industrial applications may partially compensate large inputs by stronger technology backbones.

1.1 Algorithms

In the following, we list both *TraVaS* algorithms (*TraVaS* and *uTraVaS*) for the sake of completeness.

Algorithm 1: Differentially Private Trace Variant Selection (TraVaS)

Input: Event log L, DP-Parameters (ϵ, δ)
Output: (ϵ, δ)-DP log L'

function travas (L, ϵ, δ)
  $p = 1 - e^{-\epsilon}$ `// compute probability`
  $k = \lceil 1/\epsilon \times \ln\left((e^{\epsilon} + 2\delta - 1)/(\delta(e^{\epsilon} + 1))\right) \rceil$ `// compute threshold`
  forall $(\sigma, L(\sigma)) \in L$ **do**
    $x_\sigma = \text{rTSGD}\ (p, k)$ `// generate i.i.d k-TSGD noise`
    if $L(\sigma) + x_\sigma > k$ **then**
      add $(\sigma, L(\sigma) + x_\sigma)$ to L'
  return L'

Algorithm 2: Utility-Aware TraVaS (uTraVaS)

Input: Event log L, DP-Parameters (ϵ, δ), Subquery bound u
Output: (ϵ, δ)-DP log L'

function optimize (L, ϵ, δ, u)
  forall $i \in \{1, 2, \ldots, u\}$ **do**
    create i empty *sublogs*: $\{sL_1, sL_2, \ldots, sL_i\}$ `// initialize`
    forall $j \in \{1, \ldots, i-1, i\}$ **do**
      add travas(L, ϵ/i, δ/i) to sL_j `// run TraVaS`
    add all unique variants σ from *sublogs* to a simple log pL'
    forall $\sigma \in pL'$ **do**
      forall $sL \in sublogs$ **do**
        if $\sigma \in sL$ **then**
          $f_\sigma = sL(\sigma)$ `// get frequency from sublog`
        else
          f_σ = freqEstimate (ϵ, δ) `// estimate frequency`
        update $pL'(\sigma)$ to $pL'(\sigma) + f_\sigma$ in pL'
    update $pL'(\sigma)$ to $pL'(\sigma)$ / i in pL' `// compute mean`
    add pL' to the set *pool*. `// optimization domain`
  $L' = \underset{pL' \in pool}{\text{argmax}}$ utility(L, pL') `// optimize data utility`
  return L'

1.2 Note of Authorship

This document is part of the supplementary material created for the paper "TraVaS: Differentially Private Trace Variant Selection for Process Mining", written by Majid Rafiei, Frederik Wangelik and Wil M.P. Van der Aalst. Please contact *frederik.wangelik@rwth-aachen.de* for further information.