



Chapter 10

Signal Approximation and Compressed Sensing

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- Compressed Sensing

- **Equivalence between ℓ_0 and ℓ_1 Recovery**

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- Restricted Isometry Property

1.1. Orthogonal Bases

- (general) Signal:

e.g. data such as sea water levels, audio recordings, photographic images, video data, and financial data

represent the signal by a vector $\theta^* \in R^p$

1.1. Orthogonal Bases

- Orthonormal basis: A basis with finite dimension whose vectors are all unit vectors and orthogonal to each other
- $\{\psi_j\}_{j=1}^p$ orthonormal basis of $R^p \rightarrow \Psi := [\psi_1 \ \psi_2 \dots \psi_p]$ is a $p \times p$ matrix with $\Psi^T \Psi = I_{p \times p}$

1.1. Orthogonal Bases

- Basis coefficient $\beta^* = \Psi^T \theta^* = \begin{pmatrix} \psi_1 \\ \psi_2 \\ \dots \\ \psi_p \end{pmatrix} \theta^* = \begin{pmatrix} \sum_{i=1}^p \theta_i^* \psi_{i1} \\ \sum_{i=1}^p \theta_i^* \psi_{i2} \\ \dots \\ \sum_{i=1}^p \theta_i^* \psi_{ip} \end{pmatrix}$ of a signal

$$\theta^* \in R^p \text{ with } \theta^* := \sum_{j=1}^p \beta_j^* \psi_j$$

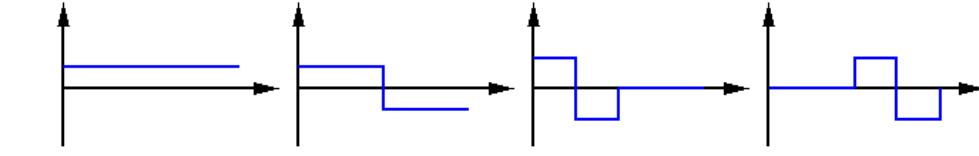
- the j-th basis coefficient can be written as a inner product

$$\beta_j^* := \langle \theta^*, \psi_j \rangle = \sum_{i=1}^p \theta_i^* \psi_{ij}$$

1.1. Orthogonal Bases

- Example (wavelet transform)

$$\Psi = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} & \frac{1}{\sqrt{2}} & 0 \\ \frac{1}{2} & \frac{1}{2} & \frac{-1}{\sqrt{2}} & 0 \\ \frac{1}{2} & \frac{-1}{2} & 0 & \frac{1}{\sqrt{2}} \\ \frac{1}{2} & \frac{-1}{2} & 0 & \frac{-1}{\sqrt{2}} \end{pmatrix}$$



<http://fourier.eng.hmc.edu/e161/lectures/Haar/index.html>

as orthonormal basis matrix

4 column vectors = 4 bases for Haar transform (signal length is 4)

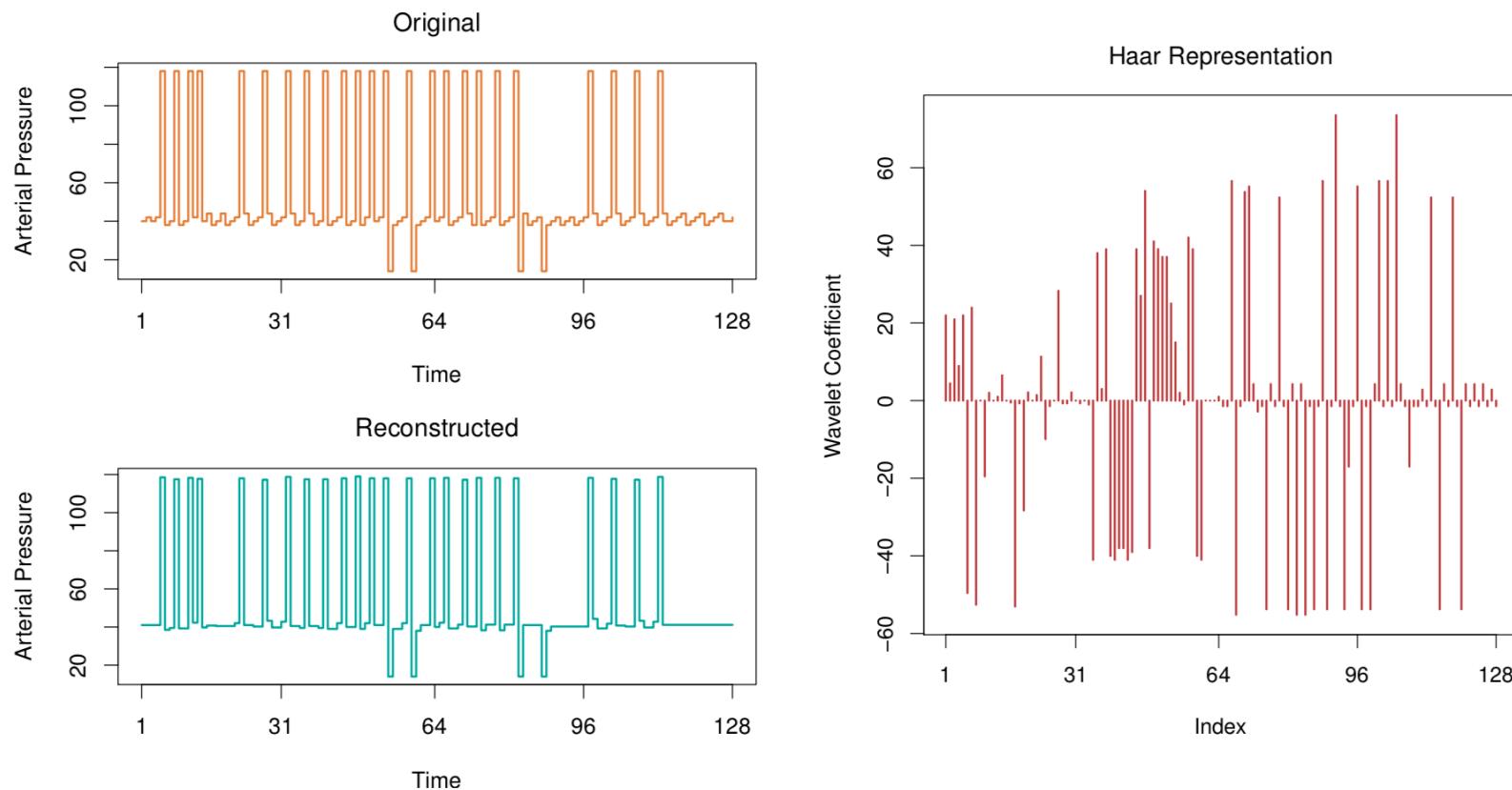


Figure 10.1 Illustration of sparsity in time series data. Left, top panel: Signal $\theta^* \in \mathbb{R}^p$ of arterial pressure versus time over $p = 128$ points. Left, bottom panel: Reconstruction $\hat{\theta}^{128}$ based on retaining the largest (in absolute amplitude) $k = 64$ coefficients from the Haar basis. Right: Haar basis coefficients $\beta^* = \Psi^T \theta^*$ of the signal.

1.2. Approximation in Orthogonal Bases

- Goal of signal compression: represent signal $\theta^* \in R^p$ using $k < p$ coefficients
- How?
- Use only sparse subset of the orthogonal vectors $\{\psi_j\}_{j=1}^p$ for $k \in \{1, \dots, p\}$: consider reconstruction

$$\Psi\beta = \sum_{j=1}^p \beta_j \psi_j, \|\beta\|_0 := \sum_{j=1}^p I[\beta_j \neq 0] \leq k$$

1.2. Approximation in Orthogonal Bases

- Optimal k-sparse approximation:
 - Compute: $\widehat{\beta}^k \in \operatorname{argmin}_{\beta \in R^p} \|\theta^* - \Psi^T \beta\|_2^2$ so that $\|\beta\|_0 \leq k$
 - Reconstruction: $\theta^k := \sum_{j=1}^p \widehat{\beta}_j^k \psi_j$
- best least-squares approximation to the signal $\theta^* \in R^p$ based on k terms
- non-convex and combinatorial problem

1.2. Approximation in Orthogonal Bases

- solve by taking first k coefficients with largest absolute values:
 1. Order vector $\beta^* \in R^p$ of basis coefficients: $|\beta_{(1)}^*| \geq |\beta_{(2)}^*| \geq \dots \geq |\beta_{(p)}^*|$
 2. For given $k \in \{1, \dots, p\}$ choose the first k terms: $\hat{\theta}_k := \sum_{j=1}^k \beta_{(j)}^* \psi_{\sigma(j)}$

1.2. Approximation in Orthogonal Bases

- Complexity of the procedure
 1. Compute basis coefficients $\beta_j^* := \langle \theta^*, \psi_j \rangle = \sum_{i=1}^p \theta_j^* \psi_{ij} \longrightarrow O(p^2)$
 2. Sort coefficients in terms of absolute values $\longrightarrow O(p \log p)$
 3. Extract first k coefficients
 4. Compute the best k-term approximation $\hat{\theta}_k := \sum_{j=1}^k \beta_{(j)}^* \psi_{\sigma(j)}$



(a)



(b)

Figure 10.3 Illustration of image compression based on wavelet thresholding. (a) Zoomed portion of the original “Boats” image from Figure 10.2(a). (b) Reconstruction based on retaining 5% of the wavelet coefficients largest in absolute magnitude. Note that the distortion is quite small, and concentrated mainly on the fine-scale features of the image.

1.3. Reconstruction in Overcomplete Bases

- Problem: Only limited class of signals has sparse representations in ANY orthonormal bases
- Solution: Combine different orthonormal bases
→ use subsets of vectors from both bases simultaneously

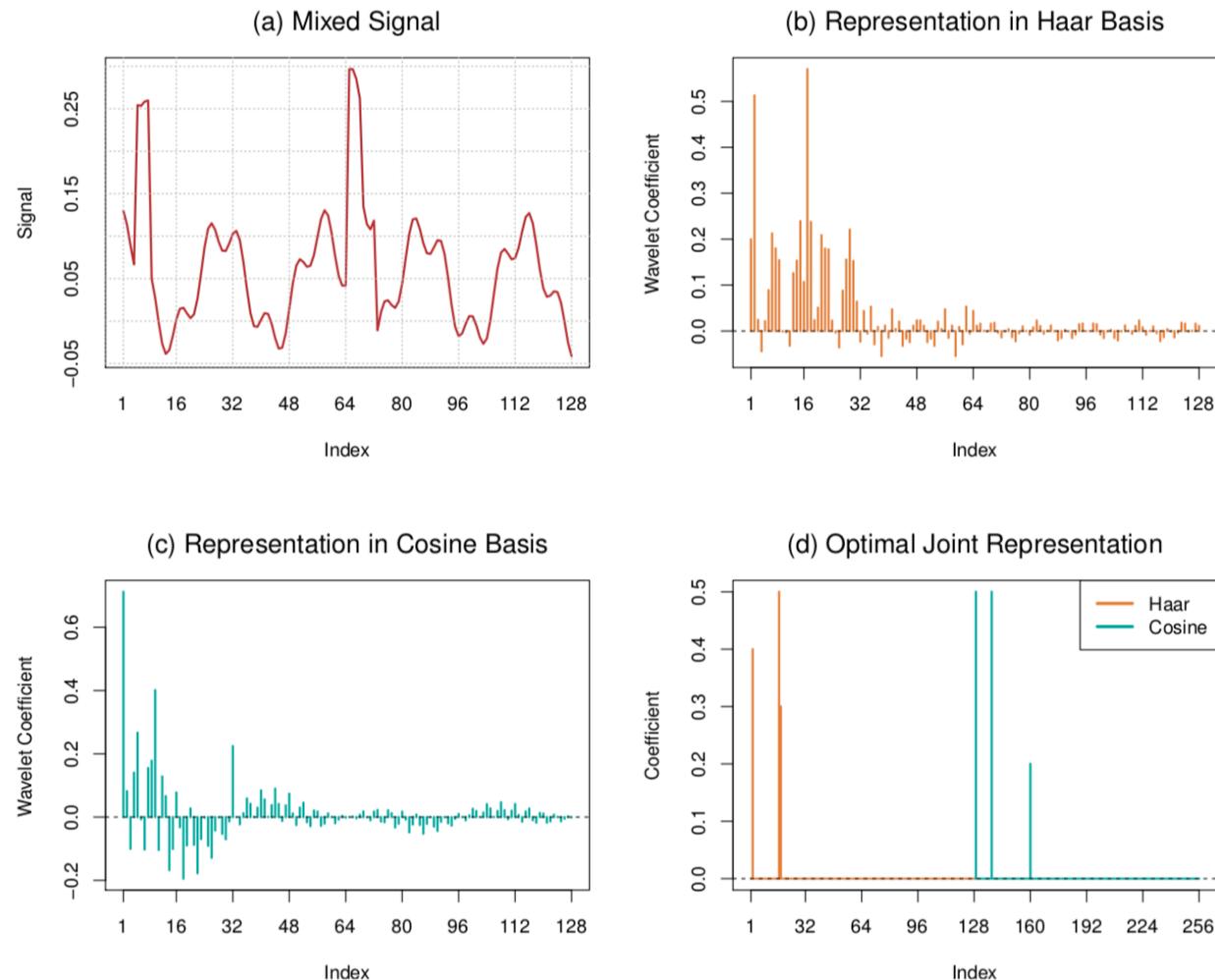


Figure 10.4 (a) Original signal $\theta^* \in \mathbb{R}^p$ with $p = 128$. (b) Representation $\Psi^T \theta^*$ in the Haar basis. (c) Representation $\Phi^T \theta^*$ in the discrete cosine basis. (d) Coefficients $(\hat{\alpha}, \hat{\beta}) \in \mathbb{R}^p \times \mathbb{R}^p$ of the optimally sparse joint representation obtained by solving basis pursuit linear program (10.11).

1.3. Reconstruction in Overcomplete Bases

- two pairs of orthonormal bases: $\{\psi_j\}_{j=1}^p$, $\{\phi_j\}_{j=1}^p$
- reconstruction of the form:

$$\sum_{j=1}^p \alpha_j \phi_j + \sum_{j=1}^p \beta_j \psi_j \text{ so that } \|\alpha\|_0 + \|\beta\|_0 \leq k$$

- optimisation problem:

$$\underset{(\alpha, \beta) \in R^p \times R^p}{\text{minimize}} \|\theta^* - \Phi\alpha - \Psi\beta\|_2^2 \text{ so that } \|\alpha\|_0 + \|\beta\|_0 \leq k$$

1.3. Reconstruction in Overcomplete Bases

- difficult to solve (non convex)
- consider convex program:

$$\text{minimize}_{(\alpha, \beta) \in R^p \times R^p} \|\theta^* - \Phi\alpha - \Psi\beta\|_2^2 \text{ so that } \|\alpha\|_1 + \|\beta\|_1 \leq R$$

→ constrained version of the lasso program/ basis-pursuit program

- consider even simpler problem:

$$\text{minimize}_{(\alpha, \beta) \in R^p \times R^p} \|\alpha\|_1 + \|\beta\|_1 \text{ so that } \theta^* = [\Phi \ \Psi] \begin{bmatrix} \alpha \\ \beta \end{bmatrix}$$

→ linear program (LP)/ basis-pursuit linear program

1.3. Reconstruction in Overcomplete Bases

- Question: When is basis-pursuit linear program equivalent to ℓ_0 constraint problem in general?
- Answer: It depends on the degree of coherence between two bases. (s. 3.2)

2.1. Johnson-Lindenstrauss Approximation

For any $0 < \epsilon < 1$ and any integer n , let k be a positive integer such that

$$k \geq 4(\epsilon^2/2 - \epsilon^3/3)^{-1} \ln n.$$

Then for any set V of n points in R^d , there is a map $f : R^d \rightarrow R^k$ such that for all $u, v \in V$,

$$(1 - \epsilon) \|u - v\|_2^2 \leq \|F(u) - F(v)\|_2^2 \leq (1 + \epsilon) \|u - v\|_2^2.$$

Further this map can be found in polynomial.

2.1. Johnson-Lindenstrauss Approximation

- establish the existence of distance-preserving dimension reduction projection
- provide an explicit bound on the dimension required for approximate distance preserving
- provide an explicit construction of the random projection

2.1. Johnson-Lindenstrauss Approximation

- Definition of Random Projection:

A random projection of a signal θ^* is a measurement of the form

$$y_i = \langle z_i, \theta^* \rangle = \sum_{j=1}^p z_{ij} \theta_j^*$$

with $z_i \in R^p$ random vector.

2.1. Johnson-Lindenstrauss Approximation

■ PART 1

- Given: n data points (d dimension), relative error ϵ
- Compute: minimum dimension k with

$$k = 4(\epsilon^2/2 - \epsilon^3/3)^{-1} \ln n$$

- larger ϵ smaller k
- k only depends on number of points n , not on dimension of points d

2.1. Johnson-Lindenstrauss Approximation

- PART 2
 - Given: input matrix X of n points in d -dimensional space
 - Compute: map F which preserves distance approximately
 - F can be a simple matrix A constructed by sampling $k \times d$ i.i.d. draws from a Gaussian matrix with mean 0 and variance $1/k$

$$F(u) := \frac{1}{\sqrt{k}} Z u, \quad Z \in R^{k \times d} \text{ random matrix with each } Z_{ij} \sim \mathcal{N}(0,1) \text{ i.i.d}$$

Compressed Sensing

- Motivation: In previous algorithm we discard most β_j^* 's, so do we really need to calculate them all?
- Oracle technique: We know which subset of k coefficients will be retained for sparse approximation → only need to compute this subset of basis coefficients
- Compressed sensing:
 - Instead of precomputing all coefficients $\beta^* = \Psi^T \theta^*$, we compute N random projections $y = Z\theta$ (with $N \ll p$)
 - Mimics behaviour of the oracle technique with only little computational overhead.

Compressed Sensing

- Setup of our problem:
 - Given: $y \in \mathbb{R}^N$: Vector of random projections of signal θ^*
 $Z \in \mathbb{R}^{N \times p}$: Design matrix used to compute random projections
 - Goal: Recover signal $\theta^* \in \mathbb{R}^p$
 - Problem: $y = Z\theta$ is highly underdetermined as $N \ll p$
- Example :
 - $y = x_1 + x_2$ if $y=1$:
 - $x_1 = 1$ and $x_2 = 0$
 - $x_1 = 0.5$ and $x_2 = 0.5$
 - etc...

$$\mathbf{y} = \mathbf{Z} \theta$$

\mathbf{Z}
 $N \times p$

$\underset{\theta \in \mathbb{R}^p}{\text{minimize}} \|\Psi^T \theta\|_0$ such that $\mathbf{y} = \mathbf{Z}\theta$

$\downarrow \ell_1 - \text{relaxation}$

$\underset{\theta \in \mathbb{R}^p}{\text{minimize}} \|\Psi^T \theta\|_1$ such that $\mathbf{y} = \mathbf{Z}\theta$

(a)

$$\mathbf{y} = \mathbf{Z} \Psi \beta$$

\mathbf{Z}
 $N \times p$

Ψ

Equivalently we can write:

$\underset{\beta \in \mathbb{R}^p}{\text{minimize}} \|\beta\|_1$ such that $\mathbf{y} = \tilde{\mathbf{Z}}\beta$ where $\tilde{\mathbf{Z}} = \mathbf{Z}\Psi$

Compressed Sensing

1. For given sample size N , compute random projections $y_i = \langle z_i, \theta^* \rangle$
(Or ideally measure the random projections y instead of full signal θ^*)
2. Estimate θ^* by solving linear program:

$$\underset{\theta \in \mathbb{R}^p}{\text{minimize}} \left\| \Psi^T \theta \right\|_1 \text{ such that } y = Z\theta, \text{ to obtain } \hat{\theta}.$$

- Family of procedures depending on random projection vectors

simplest choice: $z_{ij} \sim N(0, \frac{1}{N})$

- Combination of random projection and l_1 – relaxation

Compressed Sensing

- Before: minimize $\left\| \Psi^T \theta \right\|_1$ such that $y = Z\theta$
 - In reality we often have noisy data: $y = Z\theta^* + \varepsilon$ with $\left\| \varepsilon \right\|_2 \leq e$.
- We introduce a bound $e > 0$ on the noise level :

$$\text{minimize}_{\theta \in \mathbb{R}^p} \left\| \Psi^T \theta \right\|_1 \text{ such that } \left\| y - Z\theta \right\|_2 \leq e.$$

Sensing matrices in Compressed Sensing

Random Sensing Matrices $Z \in \mathbb{R}^{N \times p}$:

- Gaussian Random Matrix: $z_{i,j} \sim N(0, 1/N)$

- Bernoulli Random Matrix: $z_{i,j} = \begin{cases} \frac{1}{\sqrt{N}} \\ -\frac{1}{\sqrt{N}} \end{cases}$

- Satisfy RIP with almost optimal order → allows sparse recovery using l_1 – minimization.
- No fast matrix multiplication that may speed up the algorithm is available.
- Storing an unstructured matrix may be difficult

} large scale problems not practicable
with Gaussian/Bernoulli matrices

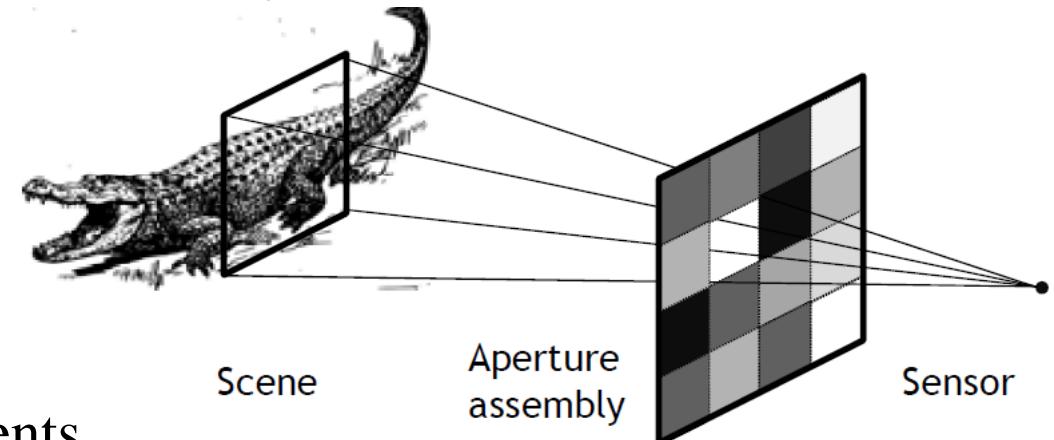
→ Deterministic Sensing Matrices

Application: Lensless Imaging

Architecture:

- Aperture assembly:
 - Consists of 2 dimensional array of aperture elements
 - The transmittance of each aperture element is independently controllable.
- Sensor:
 - Used for taking compressive measurements
- Compressive sensing matrix
 - Implemented by adjusting transmittance of individual aperture elements according to the values of the sensing matrix.

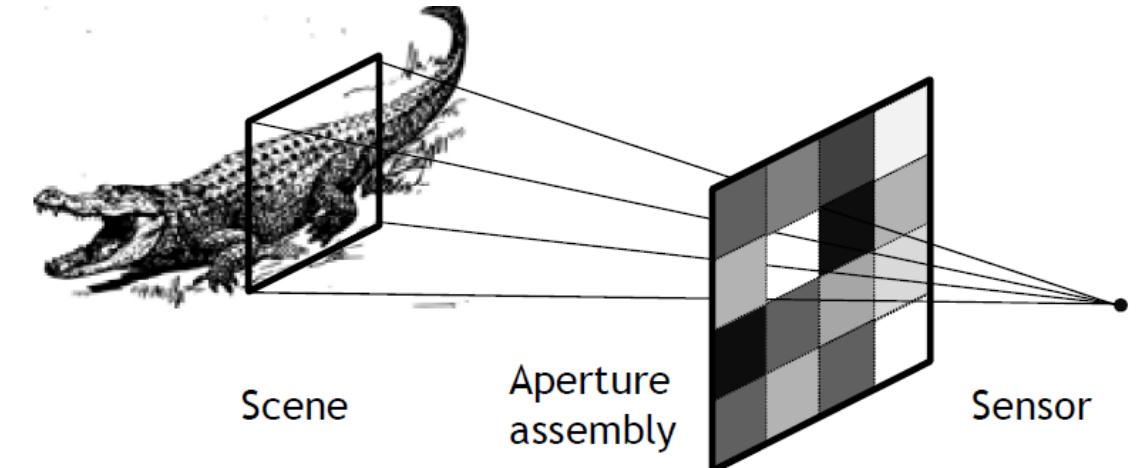
Picture and Information from:
G. Huang, H Jiang, K. Matthews, P. Wilford: Lensless Imaging by Compressive Sensing, arXiv, 2013.



Application: Lensless Imaging

How it works:

- One Element of aperture together with sensor
→ cone of rays.
- Integration of rays in this cone is defined as a pixel value of the image



Example: "traditional way of representing an image, pixel by pixel"

1. Open one aperture element and close all others
 2. Take a measurement, get one pixel value for that aperture element
 3. Repeat that for all aperture elements
- We take as many measurements as number of pixels

Compressive Measurements

- Goal: #Measurements << #Pixels

1. Define a sensing matrix:

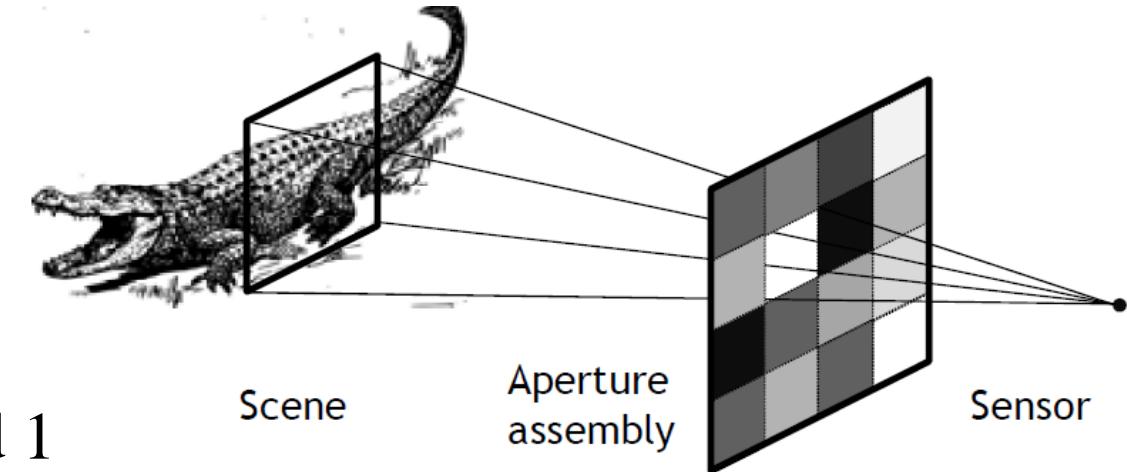
- Entries = random numbers between 0 and 1
- Each row defines a pattern for the aperture elements
- #Columns = #Elements in aperture assembly
→ each element in a row corresponds to one aperture element

2. Sensor takes one measurement for each pattern

(integrates all rays after their intensity is modulated by transmittances)

→ measurement = projection of image onto row of sensing matrix

- We take less measurements than pixels and try to reconstruct the image using CS



Compressive Measurements

- Goal: #Measurements << #Pixels

1. Define a sensing matrix:

- Entries = random numbers between 0 and 1
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- We take less measurements than pixels and try to reconstruct the image

Picture from: G. Huang, H Jiang, K. Matthews, P. Wilford: Lensless Imaging by Compressive Sensing, arXiv, 2013.



Figure 5. Reconstructed images of "Soccer", 12.5%.

Other Applications

- X-Ray (to reduce radiation exposure)
- Single Pixel Camera
- Magnetics Resonance Imaging
- Radar
- Machine learning
- ⋮

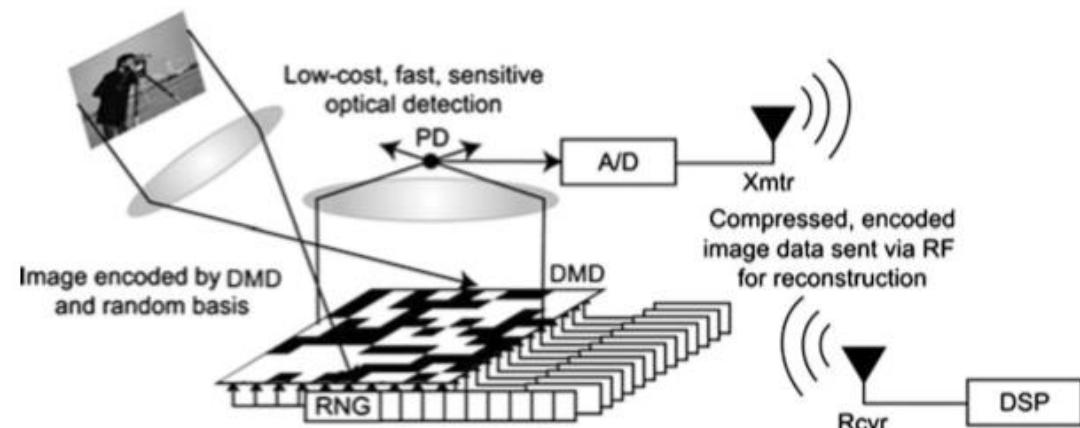


Fig. 1.5 Schematic representation of a single-pixel camera (Image courtesy of Rice University)

https://www.google.ch/search?q=single+pixel+camera&source=lnms&tbo=isch&sa=X&ved=0ahUKEwjWzob_-I3fAhXWFsAKHQlfDvwQ_AUDigB&biw=1536&bih=747#imgrc=-4OBKTuGUiQpM:



Figure 3. Prototype device. Top: lab setup. Bottom left: the LCD screen as the aperture assembly. Bottom right: the sensor board with two sensors, indicated by the red circle.

Picture:
G. Huang, H Jiang, K. Matthews, P. Wilford: Lensless Imaging by Compressive Sensing, arXiv, 2013.

Equivalence between l_0 and l_1 Recovery

Question

When is solving the l_1 -relaxation equivalent to solving the original l_0 -problem?

Equivalence between l_0 and l_1 Recovery

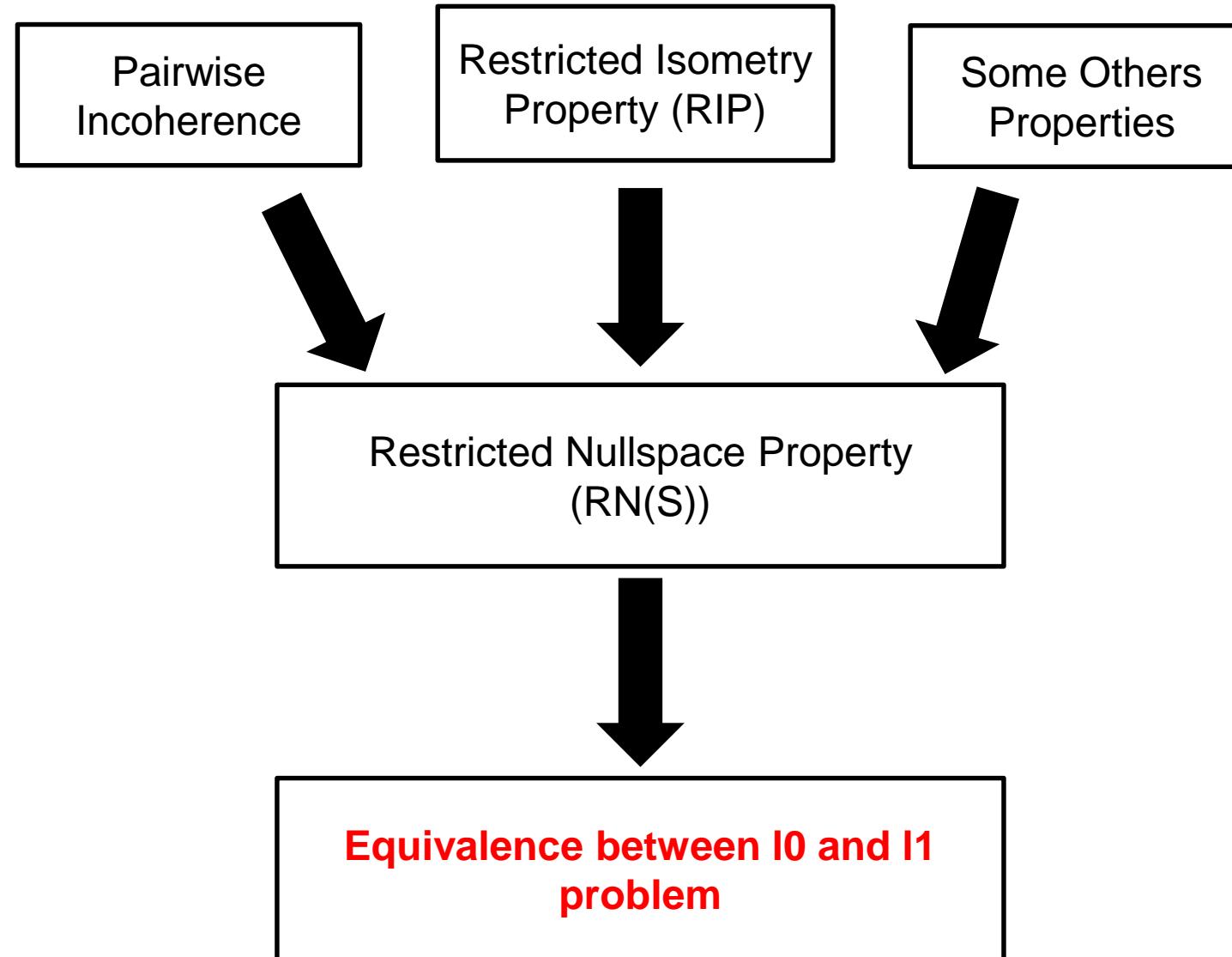
Given: $y \in \mathbb{R}^N$ and $X \in \mathbb{R}^{N \times p}$

l_0 problem: minimize $\|\beta\|_0$ such that $X\beta = y$
 $\beta \in \mathbb{R}^p$

l_1 problem: minimize $\|\beta\|_1$ such that $X\beta = y$
 $\beta \in \mathbb{R}^p$

Example: Compressed Sensing

$$X = \tilde{Z} = Z\Psi$$



Restricted Nullspace Property

An $N \times p$ matrix X satisfies the *restricted nullspace property* for a set $S \subseteq \{1, \dots, p\}$ if

$$\|\beta_S\|_1 < \|\beta_{S^c}\|_1 \text{ for all } \beta \in \ker(X) \setminus \{0\}$$

It is said to satisfy the null space property of order k if it satisfies the Nullspace property for any set S with $\text{card}(S) \leq k$.

Theorem : Equivalence between l_0 and l_1 Recovery

If a given matrix $X \in \mathbb{R}^{N \times p}$ satisfies the null space property for a set S , every vector $\beta^* \in \mathbb{R}^p$ supported on this set S is the unique solution of the l_1 – problem with $y = X\beta^*$.

- l_1 problem: minimize $\|\beta\|_1$ such that $y = X\beta$

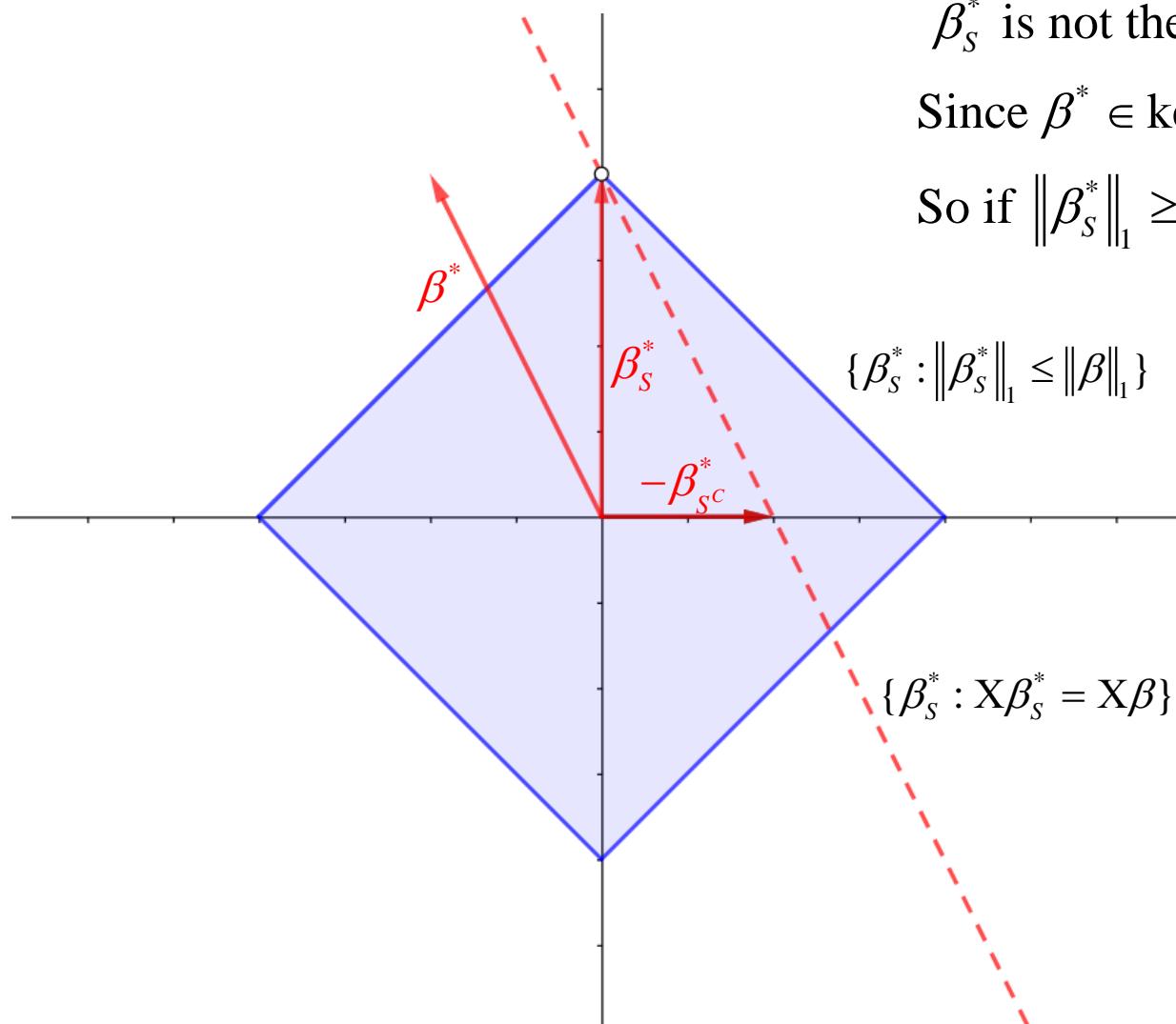
What goes wrong if $\|\beta_s\|_1 \geq \|\beta_{s^c}\|_1$?

Chose $\beta^* \in \ker(X) \setminus \{0\}$ and show:

β_s^* is not the unique minimizer of $\|\beta\|_1$ subject to $X\beta = X\beta_s^*$

Since $\beta^* \in \ker(X)$ we have: $X(\beta_s^* + \beta_{s^c}^*) = 0 \rightarrow X\beta_s^* = -X\beta_{s^c}^*$

So if $\|\beta_s^*\|_1 \geq \|\beta_{s^c}^*\|_1$ then β_s^* is not a unique solution anymore.



Verifying the restricted nullspace property

Restricted Isometry Property (RIP) :

For tolerance $\delta \in (0,1)$ and $2k \in \{1,\dots,p\}$ we say that $RIP(2k, \delta)$ holds if

$$\frac{\|X_S u\|_2^2}{\|u\|_2^2} \in [1 - \delta, 1 + \delta] \text{ for all } u \in \mathbb{R}^k \setminus \{0\}$$

for all subsets $S \subset \{1,\dots,p\}$ of kardinality $2k$.

- Intuition:
 - RIP holds if X_S changes length of vectors very little, Eigenvalues close to 1
 - Every set of columns of size at most $2k$ approximatlly behaves like orthonormal system

RIP implies restricted nullspace property

RIP($2k, \delta$) holds with $\delta < \frac{1}{3} \rightarrow$ Nullspace property of order k holds

\rightarrow the l_1 – relaxation is exact for all vectors supported on at most k elements.

- Advantage: RIP constant δ does not depend on k .
- Problem: constraint on a huge number of submatrices, $\binom{p}{2k}$ in total.
- Various choices of random projection matrix X satisfy RIP with high probability

as long as $N \gtrsim k \log \frac{ep}{k}$. (only small multiplicative overhead compared to oracle)

Proof of Theorem:

- Suppose X satisfies RN(S) property.
- β_0, β_1 : optimal solution to l_0, l_1 – problem.
- Define error vector $\Delta := \beta_0 - \beta_1$
- Goal: Show that $\Delta = 0$, for that we try to show: $\Delta \in \ker(X)$ and $\|\Delta_S\|_1 \geq \|\Delta_{S^c}\|_1$
- We have: $X\beta_0 = y = X\beta_1 \rightarrow X(\beta_0 - \beta_1) = 0 \rightarrow \Delta \in \ker(X)$
- Writing $\beta_0 = \beta_1 + \Delta$, we get:
$$\|\beta_{0,S}\|_1 \geq \|\beta_1\|_1 = \|\beta_{0,S} + \Delta_S\|_1 + \|\Delta_{S^c}\|_1 \geq \|\beta_{0,S}\|_1 - \|\Delta_S\|_1 + \|\Delta_{S^c}\|_1$$
- Rearranging terms $\rightarrow \|\Delta_S\|_1 \geq \|\Delta_{S^c}\|_1$

Summary

- Sparse Representation of Signals
 - Sparsity of a Signal depends on the Choice of the Basis.
 - Reconstruction in Overcomplete Bases
- Random Projection
 - Johnson-Lindenstrauss Approximation
 - R-Code showing distance preserving property
- Compressed Sensing
 - R-Code showing recovery of an Image
 - Applications
- Equivalence between ℓ_0 and ℓ_1 Recovery
 - Restricted Nullspace Property and Geometric Intuition
 - Restricted Isometry Proeprty

Reference

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