

# 动态规划(Dynamic Programming)





扫码了解极客时间《算法面试通关40讲》视频课程



### 本节内容

- 1. 递归+记忆化 -> 递推
- 2. 状态的定义: opt[n], dp[n], fib[n]
- 3. 状态转移方程: opt[n] = best\_of(opt[n-1], opt[n-2], ...)
- 4. 最优子结构

### 斐波那契数列

0, 1, 1, 2, 3, 5, 8, 13, 21, ...

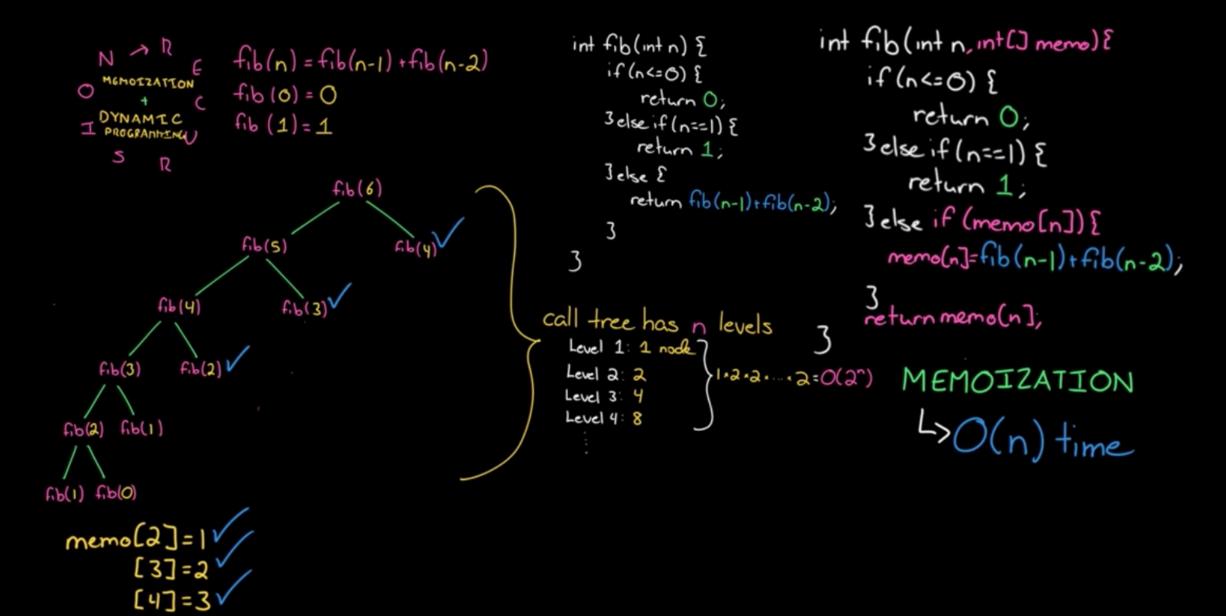
递推公式: F[n] = F[n-1] + F[n-2]

```
int fib (int n) {
                     fib(n) = fib(n-1) + fib(n-2)
                                                                     if (n<=0) {
    MEMOTZATION
 0
                     fib (0) = 0
                                                                         return 0;
    DYNAMIC
                                                                    3 else if (n==1) {
                     fib (1) = 1
 I PROGRAMMINU
                                                                        return 1;
      S
                                                                    Jelse &
                                  f.b(6)
                                                                        return fib (n-1)+fib(n-2);
                      fib(5)
                                         F.P(A)
                                                f.b(2)
           fib (4)
                                        f_{1b}(3)
                           f(P(3))
                                                            call tree has n levels
                                                                Level 1: 1 nock
                         fib(2) fib(1) fib(2)
    fib(3)
              f.b(2)
                                                                Level 2: 2
                                                                                  112 x2 x... x 2 = O(2")
                                   t.P(1) t.P(0)
                                                               Level 3: 4
                                                                Level 4: 8
      t'P(1) t'P(1) t'P(0)
f.p(2)
```



```
int fib(int n) {
  return n <= 1 ? n : fib(n - 1) + fib(n - 2);
}</pre>
```

```
int fib (int n) {
                       fib(n) = fib(n-1) + fib(n-2)
                                                                      if (n<=0) {
      MEMOTZATION
   0
                       f_1b(0) = 0
                                                                          return 0;
     DYNAMIC
                                                                      3 else if (n==1) {
                       fib (1) = 1
   I PROGRAMMINU
                                                                          return 1;
       S
                                                                      Jelse &
                                   f.p(6)
                                                                         return fib (n-1)+fib(n-2);
                        fib(s)
                                           F.P(A)
                                                  f.b(2)
             fib (4)
                                          fib(3)
                             f(P(3))
                                                              call tree has n levels
                                                                  Level 1: 1 node 7
                           fib(2) fib(1) fib(2)
      fib(3)
                f.b(2)
                                                                                   >1*2 *2 * ... * 2 = O(2")
                                                                 Level 2: 2
                                     t.P(1) t.P(0)
                                                                 Level 3: 4
                          £₽(1) £₽(0)
                                                                 Level 4: 8
        4P(1) 4P(1) 4P(0)
(P(1) tP(0)
```



[5]=5



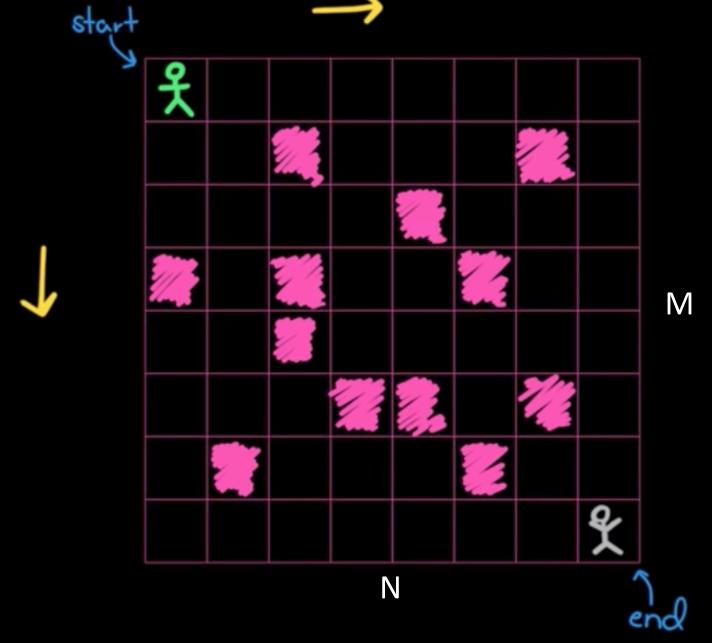
### 递归+记忆化==>递推

递推公式: F[n] = F[n-1] + F[n-2]

```
F[0] = 0, F[1] = 1;
for (int i = 2; i <= n; ++i) {
   F[i] = F[i - 1] + F[i - 2];
}</pre>
```

COUNT THE PATHS





```
COUNT THE PATHS
               paths(start, end) =
                    + paths(B,enQ)
   paths(A,en@)
paths(D,end)+paths(C,end) paths(C,end)+paths(E,end)
        int countPaths(boolean[]] grid, int row, int col) {
          if (!validSquare(grid, row, col)) return 0;
          if (isAtEnd(grid, row, col)) return 1;
          return countPaths(grid, row+1, col) + countPaths(grid, row, col+1);
```



COON	IHE	PATHS	
start	$\rightarrow$	columns	(x)

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						ı	4





COUNT THE PATHS

columns (x) start

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. Q.	l.
	1

웃	Bo <sup>+</sup>	• E					
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D				纜			
		獲			絮		
		**					
				黎		*	
							1
_	1	ı	1	1	1	ı	4



```
opt[i, j] = opt[i - 1, j] + opt[i, j - 1]
```

\_\_\_\_\_\_

```
if a[i, j] = '空地':
    opt[i , j] = opt[i - 1, j] + opt[i, j - 1]
else: // 石头
    opt[i , j] = 0
```

COUNT THE PATHS

columns (x) start

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(4	)		
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			1

웃	Bo <sup>X</sup>	• E					
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#### COUNT THE PATHS



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5	S	2	2	覆	3	3	١
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7	3	茅	1	1	1		ı
4	3	3	7	黎	0	1/4	ı
1	青	3	a	ı		2	41
١	1	1	1	1	1	1	4





## 动态规划 Dynamic Programming

- 1. 递推! (递推 + 记忆化)
- 2. 状态的定义: opt[n], dp[n], fib[n]
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### 动态规划 vs 回溯 vs 贪心算法

•回溯(递归)—重复计算

• 贪心算法 — 永远局部最优

• 动态规划 — 记录局部最优子结构 / 多种记录值



### 实战题目

- 1. <a href="https://leetcode.com/problems/climbing-stairs/description/">https://leetcode.com/problems/climbing-stairs/description/</a>
- 2. <a href="https://leetcode.com/problems/triangle/description/">https://leetcode.com/problems/triangle/description/</a>
- 3. <a href="https://leetcode.com/problems/maximum-product-subarray/description/">https://leetcode.com/problems/maximum-product-subarray/description/</a>
- 4. <a href="https://leetcode.com/problems/best-time-to-buy-and-sell-stock/#/description">https://leetcode.com/problems/best-time-to-buy-and-sell-stock/#/description</a>
- 5. <a href="https://leetcode.com/problems/best-time-to-buy-and-sell-stock-ii/">https://leetcode.com/problems/best-time-to-buy-and-sell-stock-ii/</a>
- 6. <a href="https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iii/">https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iii/</a>
- 7. <a href="https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iv/">https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iv/</a>
- 8. <a href="https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-cooldown/">https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-cooldown/</a>
- 9. <a href="https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-transaction-fee/">https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-transaction-fee/</a>
- 10. <a href="https://leetcode.com/problems/longest-increasing-subsequence">https://leetcode.com/problems/longest-increasing-subsequence</a>
- 11. <a href="https://leetcode.com/problems/coin-change/">https://leetcode.com/problems/coin-change/</a>
- 12.<a href="https://leetcode.com/problems/edit-distance/">https://leetcode.com/problems/edit-distance/</a>





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