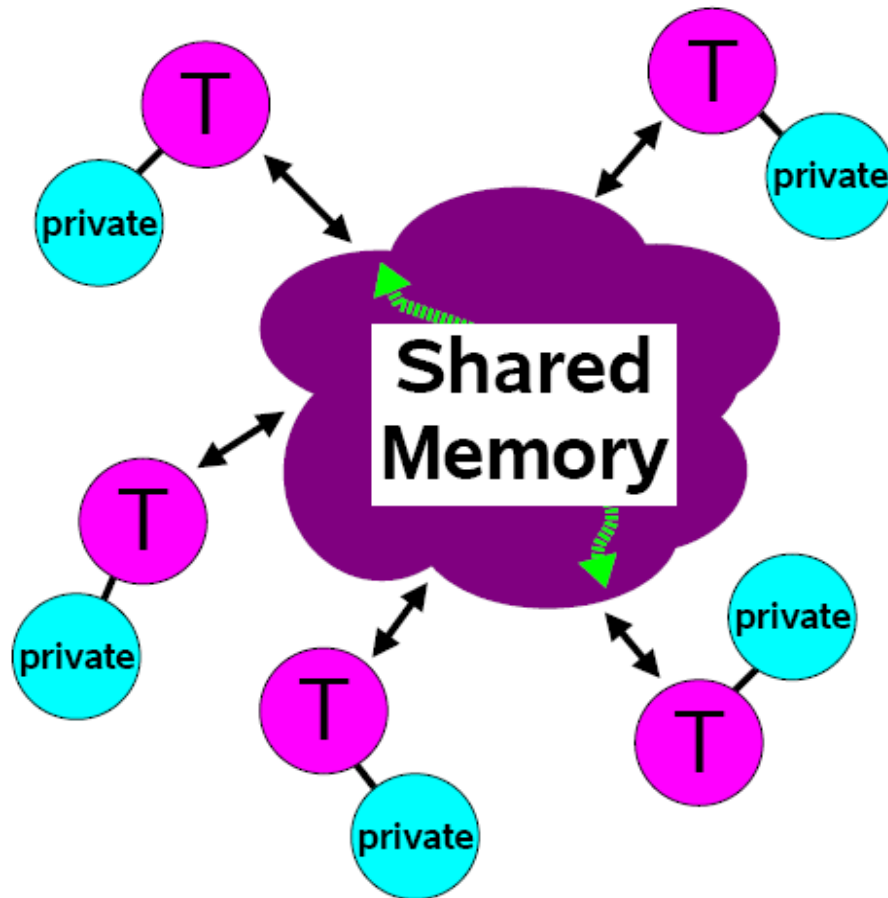


Lecture 5

OpenMP: Contents

- OpenMP's constructs fall into 5 categories:
 - ◆ Parallel Regions
 - ◆ Worksharing
 -  ◆ Data Environment
 - ◆ Synchronization
 - ◆ Runtime functions/environment variables

Review: Shared Memory Model



Programming Model

- ✓ All threads have access to the same, globally shared, memory
- ✓ Data can be shared or private
- ✓ Shared data is accessible by all threads
- ✓ Private data can be accessed only by the threads that owns it
- ✓ Data transfer is transparent to the programmer
- ✓ Synchronization takes place, but it is mostly implicit

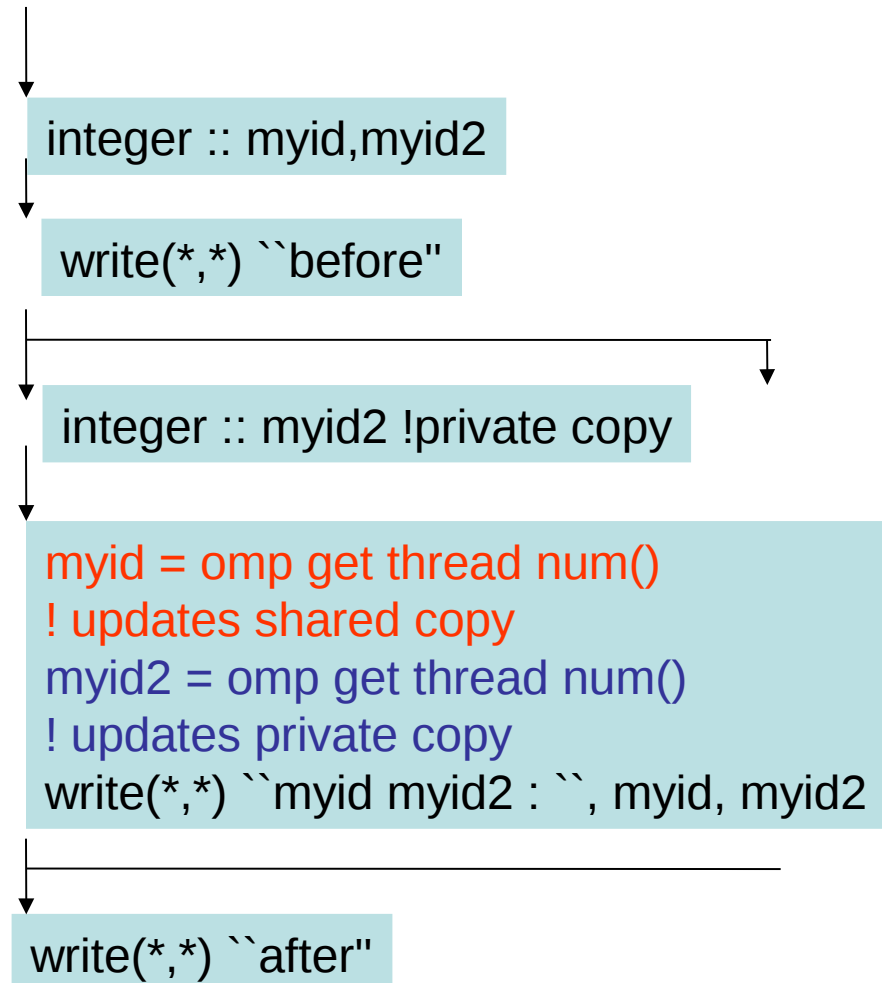
Data Scope Clauses

- **SHARED (list)**
- **PRIVATE (list)**
- **FIRSTPRIVATE (list)**
- **LASTPRIVATE (list)**
- **DEFAULT (list)**
- **THREADPRIVATE (list)** (it is actually a directive, not clause)
- **COPYIN (list)**
- **REDUCTION (operator | intrinsic : list)**

Data Scope Example (shared vs private)

All sample codes are in </home/syam/ces745/openmp/Fortran/data-scope/>

```
program scope
  implicit none
  integer :: myid, myid2
  write(*,*) "before"
  !$omp parallel private(myid2)
    myid = omp_get_thread_num()
    myid2 = omp_get_thread_num()
    write(*,*) "myid myid2 : ", myid, myid2
  !$omp end parallel
  write(*,*) "after"
end program
```



```
[syam@silky:~/ces745/openmp/Fortran/data-scope] ./scope-ifort
```

before

myid myid2 :	50	8
myid myid2 :	32	18
myid myid2 :	62	72
myid myid2 :	79	17
myid myid2 :	124	73
myid myid2 :	35	88
myid myid2 :	35	37

.....

.....

myid myid2 :	35	114
myid myid2 :	35	33
myid myid2 :	35	105
myid myid2 :	35	122
myid myid2 :	35	68
myid myid2 :	35	51
myid myid2 :	35	81

after

```
[syam@silky:~/ces745/openmp/Fortran/data-scope]
```

Data Environment:

Default storage attributes

- Shared Memory programming model:
 - Most variables are shared by default
- Global variables are SHARED among threads
 - Fortran: COMMON blocks, SAVE variables, MODULE variables
 - C: File scope variables, static
- But not everything is shared...
 - Stack variables in sub-programs called from parallel regions are PRIVATE
 - Automatic variables within a statement block are PRIVATE.



Data Environment:

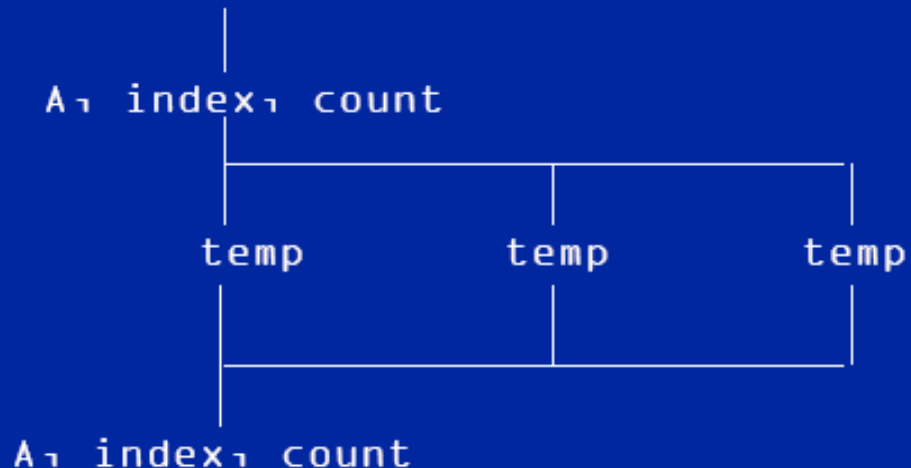
Example storage attributes

```
program sort
common /input/ A(10)
integer index(10)
call input
C$OMP PARALLEL
  call work(index)
C$OMP END PARALLEL
print*, index(1)
```

```
subroutine work
common /input/ A(10)
real temp(10)
integer count
save count
.....
```

A, index and count are shared by all threads.

temp is local to each thread



Data Environment:

Changing storage attributes

- One can selectively change storage attributes constructs using the following clauses*
 - SHARED
 - PRIVATE
 - FIRSTPRIVATE
 - THREADPRIVATE*
- The value of a private inside a parallel loop can be transmitted to a global value outside the loop with:
 - LASTPRIVATE
- The default status can be modified with:
 - DEFAULT (PRIVATE | SHARED | NONE)

All the clauses on this page only apply to the *lexical extent* of the OpenMP construct.

Private Clause

- `private(var)` creates a local copy of `var` for each thread.
 - The value is uninitialized
 - Private copy is *not* storage associated with the original

```
program wrong
  IS = 0
  C$OMP PARALLEL DO PRIVATE(IS)
    DO J=1,1000
      IS = IS + 1
    1000 CONTINUE
  print *, IS
```

IS was not
initialized



</home/syam/ces745/openmp/Fortran/data-scope/private-11.f90>

```
program wrong
use omp_lib
integer :: myid, nthreads
IS=10
!$omp parallel private(IS, myid)
  myid = omp_get_thread_num()
  nthreads = omp_get_num_threads()
  print *, IS
  !$omp do
    do j=1, 10
      IS = IS + j
      print *, j, IS, nthreads, myid
    end do
  !$omp end do
!$omp end parallel
IS = IS + 1
print *, IS
end
```

[syam@saw-login1:~] export OMP_NUM_THREADS=4

[syam@saw-login1:~] ./private-11_pathscale

0

345315584

1 345315585 4 0

2 345315587 4 0

3 345315590 4 0

10 10 4 3

0

7 7 4 2

8 15 4 2

9 24 4 2

0

4 4 4 1

5 9 4 1

6 15 4 1

11

Different IS from threads

[syam@saw-login1:~] ./private-11_intel

0

1 1 4 0

2 3 4 0

3 6 4 0

0

9 9 4 3

10 19 4 3

0

4 4 4 1

5 9 4 1

6 15 4 1

0

7 7 4 2

8 15 4 2

11

Firstprivate Clause

- Firstprivate is a special case of private.
 - Initializes each private copy with the corresponding value from the master thread.

```
program almost_right
  IS = 0
  C$OMP PARALLEL DO FIRSTPRIVATE(IS)
    DO J=1,1000
      IS = IS + 3
    END DO
  print *, IS
```

Each thread gets its own IS
with an initial value of 0

private-3.f90

```
program good
use omp_lib
integer x, y
x = 1
y = 2
!$omp parallel private(x) firstprivate(y)
  x = 3
  y = y+2
print *, omp_get_thread_num(), x, y
!$omp end parallel

print *, x, y
end
```

private-4.f90

```
! without use omp_lib
program wrong

integer x, y
x = 1
y = 2
!$omp parallel private(x) firstprivate(y)
  x = 3
  y = y+2
print *, omp_get_thread_num(), x, y
!$omp end parallel

print *, x, y
end
```

use omp_lib module (head file) is important in intel Fortran compiler

private-3.f90

```
[syam@saw-login1:] ./private-3_ifort
```

0	3	4
3	3	4
1	3	4
2	3	4
1	2	

private-4.f90

```
[syam@saw-login1:~] ./private-4_ifort
```

0.00000000E+00	3	4
0.00000000E+00	3	4
0.00000000E+00	3	4
0.00000000E+00	3	4
1	2	

Lastprivate Clause

- Lastprivate passes the value of a private from the last iteration to a global variable.

```
program closer
  IS = 0
  C$OMP PARALLEL DO FIRSTPRIVATE(IS)
  C$OMP+ LASTPRIVATE(IS)
    DO J=1,1000
      IS = IS + 3
    END DO
  print *, IS
```

Each thread gets its own IS
with an initial value of 0

IS is defined as its value at the last
iteration (i.e. for J=1000)

</home/syam/ces745/openmp/Fortran/data-scope/first-last.f90>

```
use omp_lib
integer :: myid, nthreads
integer :: x, y
x=10
print *, "x, j, nthreads, myid"
!$omp parallel do private(myid) firstprivate(x) lastprivate(j, x)
  do j=1, 10
    myid = omp_get_thread_num()
    nthreads = omp_get_num_threads()
    x = x + j
    print *, x, j, nthreads, myid
  end do
!$omp end parallel do

y = x + 1
print *, "x, j, y"
print *, x, j, y
end
```

OpenMP:

Another data environment example

- Here's an example of PRIVATE and FIRSTPRIVATE

```
variables A,B, and C = 1  
C$OMP PARALLEL PRIVATE(B)  
C$OMP& FIRSTPRIVATE(C)
```

- Inside this parallel region ...
 - “A” is shared by all threads; equals 1
 - “B” and “C” are local to each thread.
 - B’s initial value is undefined
 - C’s initial value equals 1
- Outside this parallel region ...
 - Original values of “B” and “C” are restored

OpenMP: Default Clause

- Note that the default storage attribute is **DEFAULT(SHARED)** (so no need to specify)
- To change default: **DEFAULT(PRIVATE)**
 - ◆ *each* variable in *static* extent of the parallel region is made private as if specified in a private clause
 - ◆ mostly saves typing
- **DEFAULT(NONE)**: *no* default for variables in static extent. Must list storage attribute for each variable in static extent

Only the Fortran API supports default(private).

C/C++ only has default(shared) or default(none).

OpenMP: Default Clause Example

```
    itotal = 1000  
C$OMP PARALLEL PRIVATE(np, each)  
    np = omp_get_num_threads()  
    each = itotal/np  
    .....  
C$OMP END PARALLEL
```

```
    itotal = 1000  
C$OMP PARALLEL DEFAULT(PRIVATE) SHARED(itotal)  
    np = omp_get_num_threads()  
    each = itotal/np  
    .....  
C$OMP END PARALLEL
```

**These two
codes are
equivalent**

Changing default scoping rules: C vs Fortran

- Fortran

default (shared | private | firstprivate | none)

index variables (serial, and *parallel do*) are private

- C/C++

default (shared | none)

- no default (private): many standard C libraries are implemented using macros that reference global variables
- serial loop index variables are shared, *parallel for* index variables are private

Default (none): helps catch scoping errors

Default scoping rules in Fortran

subroutine caller(a, n)

Integer n, a(n), i, j, m
m = 3

!\$omp parallel do

```
do i = 1, n
  do j = 1, 5
    call callee(a(j), m, j)
  end do
end do
end
```

subroutine callee(x, y, z)

common /com/ c
Integer x, y, z, c, ii, cnt
save cnt

```
cnt = cnt + 1
do ii = 1, z
  x = y + z
end do
end
```

Variable	Scope	Is Use Safe?	Reason for Scope
a	shared	yes	declared outside par construct
n	shared	yes	declared outside par construct
i	private	yes	parallel loop index variable
j	private	yes	Fortran seq. loop index var
m	shared	yes	declared outside par construct
x	shared	yes	actual param. is a, which is shared
y	shared	yes	actual param. is m, which is shared
z	private	yes	actual param. is j, which is private
c	shared	yes	in a common block
ii	private	yes	local stack var of called subrout
cnt	shared	no	local var of called subrout with save attribute

Default scoping rules in C

```
void caller(int a[ ], int n)
{
    int i, j, m=3;
```

```
#pragma omp parallel for
    for ( i = 0; i<n; i++){
        int k = m;
        for(j=1; j<=5; j++){
            callee(&a[i], &k, j);
        }
    }
extern int c;
```

```
void callee(int *x, int *y, int z)
{
    int ii;
    static int cnt;

    cnt++;
    for(ii=0; ii<z; ii++){
        *x = *y + c;
    }
```

Variable	Scope	Is Use Safe?	Reason dor Scope
a	shared	yes	declared outside par construct
n	shared	yes	declared outside par construct
i	private	yes	parallel loop index variable
j	shared	no	loop index var, but not in Fortran
m	shared	yes	declared outside par construct
k	private	yes	auto var declared inside par constr.
x	private	yes	Value parameter
*x	shared	yes	actual param. is a, which is shared
y	private	yes	Value parameter
*y	private	yes	actual param. is k, which is private
z	private	yes	Value parameter
c	shared	yes	declared as extern
ii	private	yes	local stack var of called subrout
cnt	shared	no	declared as static

OpenMP: Reduction

- Another clause that effects the way variables are shared:
 - **reduction (op : list)**
- The variables in “list” must be shared in the enclosing parallel region.
- Inside a parallel or a worksharing construct:
 - **A local copy of each list variable is made and initialized depending on the “op” (e.g. 0 for “+”)**
 - **pair wise “op” is updated on the local value**
 - **Local copies are reduced into a single global copy at the end of the construct.**

reduction(operator|intrinsic:var1[,var2])

- Allows safe **global** calculation or comparison.
- A private copy of each listed variable is created and initialized depending on operator or intrinsic (e.g., 0 for +).
- Partial sums, local mins etc. are determined by the threads in parallel.
- Partial sums are added together from one thread at a time to get global sum.
- Local mins are compared from one thread at a time to get gmin.
- At the end of the region for which the reduction clause was specified, the original list item is updated by combining its original value with the final value of each of the private copies, using the operator specified.

```
sum = 0.0
c$omp do shared(x) private(i)
c$omp& reduction(+:sum)
do i = 1, N
sum = sum + x(i)
end do
```

```
gmin = 1e30
c$omp do shared(x) private(i)
c$omp& reduction(min:gmin)
do i = 1,N
gmin = min(gmin,x(i))
end do
```

reduction(operator|intrinsic:var1[,var2])

- Listed variables must be shared in the enclosing parallel context.
- In Fortran
 - operator can be **+**, *****, **-**, **.and.**, **.or.**, **.eqv.**, **.neqv.**
 - intrinsic can be **max**, **min**, **iand**, **ior**, **ieor**
- In C
 - operator can be **+**, *****, **-**, **&**, **^**, **|**, **&&**, **||**
 - pointers and reference variables are not allowed in reductions!

OpenMP:

Reduction example

```
#include <omp.h>
#define NUM_THREADS 2
void main ()
{
    int i;
    double ZZ, func(), res=0.0;
    omp_set_num_threads(NUM_THREADS)
    #pragma omp parallel for reduction(+:res) private(ZZ)
    for (i=0; i< 1000; i++){
        ZZ = func(i);
        res = res + ZZ;
    }
}
```

PROGRAM REDUCTION

```
USE omp_lib
IMPLICIT NONE
INTEGER tnumber
INTEGER I,J,K
I=1
J=1
K=1
```

```
PRINT *, "Before Par Region: I=",I," J=", J," K=",K
PRINT *, ""
```

/home/syam/ces745/openmp/Fortran/data-scope/reduction

```
!$OMP PARALLEL PRIVATE(tnumber) REDUCTION(+:I) REDUCTION(*:J)
REDUCTION(MAX:K)
```

```
tnumber=OMP_GET_THREAD_NUM()
```

```
I = I+tnumber
```

```
J = J*tnumber
```

```
K = MAX(K,tnumber)
```

```
PRINT *, "Thread ",tnumber, "      I=",I," J=", J," K=",K
```

```
!$OMP END PARALLEL
```

```
PRINT *, ""
```

```
print *, "Operator      +      *      MAX"
```

```
PRINT *, "After Par Region: I=",I," J=", J," K=",K
```

```
END PROGRAM REDUCTION
```

```
[syam@nar316 reduction]$ ./para-reduction
Before Par Region: I= 1  J= 1  K= 1
```

```
Thread 0      I= 0  J= 0  K= 0
```

```
Thread 1      I= 1  J= 1  K= 1
```

```
Operator      +      *      MAX
```

```
After Par Region: I= 2  J= 0  K= 1
```

```
[syam@nar316 reduction]$
```

Scope clauses that can appear in a parallel construct

- *shared* and *private* explicitly scope specific variables
- *firstprivate* and *lastprivate* perform initialization and finalizing of privatized variables
- *default* changes the default rules used when variables are not explicitly scoped
- *reduction* explicitly identifies reduction variables

General Properties of Data Scope Clauses

- Directive with the scope clause must be within the lexical extent of the declaration
- A variable in a data scoping clause cannot refer to a portion of an object, but must refer to the entire object (e.g., not an individual array element but the entire array)
- A directive may contain multiple shared and private scope clauses; however, an individual variable can appear on at most a single clause (e.g., a variable cannot be declared as both shared and private). Exception: `firstprivate` & `lastprivate`
- Data references to variables that occur within the lexical extent of the parallel loop are affected by the data scope clauses; however, references from subroutines invoked from within the loop are not affected

Threadprivate directive

- Makes global data private to a thread
 - ◆ Fortran: **COMMON** blocks
 - ◆ C: File scope and static variables
- Different from making them **PRIVATE**
 - ◆ with **PRIVATE** global variables are masked.
 - ◆ **THREADPRIVATE** preserves global scope within each thread
- Threadprivate variables can be initialized using **COPYIN** or by using **DATA** statements.



A threadprivate example

Consider two different routines called within a parallel region.

```
subroutine poo
parameter (N=1000)
common/buf/A(N) , B(N)
C$OMP THREADPRIVATE (/buf/)
do i=1, N
    B(i)= const* A(i)
end do
return
end
```

```
subroutine bar
parameter (N=1000)
common/buf/A(N) , B(N)
C$OMP THREADPRIVATE (/buf/)
do i=1, N
    A(i) = sqrt(B(i))
end do
return
end
```

Because of the **threadprivate** construct, each thread executing these routines has its own copy of the common block /buf/.

Data scoping clause across lexical and dynamic extents

```
program wrong
common /bounds/ istart, iend
integer iarray(10000)

N=10000
!$omp parallel private(iam, nthreads) &
!$omp private (chunk, istart, iend)
! Compute the subset of iterations
! Executed by each thread
nthreads = omp_get_num_threads()
iam = omp_get_thread_num()
chunk = (N+nthreads-1)/nthreads
istart = iam*chunk + 1
iend = min((iam+1)*chunk, N)

  call work(iarray)
!$omp end parallel
end
```

subroutine work(iarray)

```
! Subroutine to operate on a thread's
! Portion of the array "iarray"
common /bounds/ istart, iend
integer iarray(10000)
```

```
do I = istart, iend
  iarray(i) = i*i
end do
return
end
```

Problem:

Private clause applies only within the lexical scope of the parallel region. Reference to istart, iend from within the work subroutine are not affected, and directly access the shared instances from the common block which are undefined.

Fixing data scoping through parameters

```
program correct
common /bounds/ istart, iend
integer iarray(10000)

N=10000
!$omp parallel private(iam, nthreads) &
!$omp private (chunk, istart, iend)
! Compute the subset of iterations
! Executed by each thread
nthreads = omp_get_num_threads()
iam = omp_get_thread_num()
chunk = (N+nthreads-1)/nthreads
istart = iam*chunk + 1
iend = min((iam+1)*chunk, N)

call work(iarray, istart, iend)
!$omp end parallel
end
```

```
subroutine work(iarray, istart, iend)
! Subroutine to operate on a thread's
! Portion of the array "iarray"
```

```
integer iarray(10000)

do I = istart, iend
  isrray(i) = i*i
end do
return
end
```

Solution 1:

Pass the values of istart and iend as parameters to the work subroutine

Fixing data scoping using the threadprivate directive

```
program correct
common /bounds/ istart, iend
!$omp threadprivate(/bounds/)
integer iarray(10000)

N=10000
!$omp parallel private(iam, nthreads, chunk)
! Compute the subset of iterations
! Executed by each thread
nthreads = omp_get_num_threads()
iam = omp_get_thread_num()
chunk = (N+nthreads-1)/nthreads
istart = iam*chunk + 1
iend = min((iam+1)*chunk, N)

  call work(iarray)
!$omp end parallel
end
```

```
subroutine work(iarray)
! Subroutine to operate on a thread's
! Portion of the array "iarray"
common /bounds/ istart, iend
!$omp threadprivate(/bounds/)
integer iarray(10000)

do I = istart, iend
  isrray(i) = i*i
end do
return
end
```

Solution 2:

Using the threadprivate directive. It effectively behaves like a private clause except that it applies to the entire program. Both the main program and the subroutine access the same threadprivate copy of the variables.



OpenMP: Contents

- OpenMP's constructs fall into 5 categories:
 - ◆ Parallel Regions
 - ◆ Worksharing
 - ◆ Data Environment
 -  ◆ Synchronization
 - ◆ Runtime functions/environment variables

OpenMP: Synchronization

- OpenMP has the following constructs to support synchronization:

- atomic
- critical section
- barrier
- flush
- ordered
- single
- master

We discuss this here, but it really isn't a synchronization construct. It's a work-sharing construct that includes synchronization.

We discuss this here, but it really isn't a synchronization construct.

Synchronization categories

- **Mutual Exclusion Synchronization**

critical
atomic

- **Event Synchronization**

barrier
ordered
master

- **Custom Synchronization**

flush
(lock – runtime library)

OpenMP: Synchronization

- Only one thread at a time can enter a **critical** section.

```
C$OMP PARALLEL DO PRIVATE(B)
C$OMP& SHARED(RES)
    DO 100 I=1,NITERS
        B = DOIT(I)
C$OMP CRITICAL
        CALL CONSUME (B, RES)
C$OMP END CRITICAL
100  CONTINUE
```

Named Critical Sections

A named critical section must synchronize with other critical sections of the same name but can execute concurrently with critical sections of a different name.

```
cur_max = min_infinity
cur_min = plus_infinity

!$omp parallel do
  do i = 1, n
    if (a(i).gt. cur_max) then
!$omp critical (MAXLOCK)
      if (a(i).gt. cur_max) then
        cur_max = a(i)
      endif
!$omp end critical (MAXLOCK)
    endif
  enddo
enddo

if (a(i).lt. cur_min) then
!$omp critical (MINLOCK)
  if (a(i).lt. cur_min) then
    cur_min = a(i)
  endif
!$omp end critical (MINLOCK)
endif
enddo
!$omp end parallel do
```

necessary but not sufficient



OpenMP: Synchronization

- **Atomic** is a special case of a critical section that can be used for certain simple statements.
- It applies only to the update of a memory location (the update of X in the following example)

```
C$OMP PARALLEL PRIVATE(B)
    B = DOIT(I)
C$OMP ATOMIC
    X = X + B
C$OMP END PARALLEL
```



```
program sharing_par2
```

```
use omp_lib
```

```
implicit none
```

```
integer, parameter :: N = 50000000
```

```
integer(selected_int_kind(17)) :: x(N)
```

```
integer(selected_int_kind(17)) :: total
```

```
integer :: i
```

```
total = 0
```

```
!$omp parallel
```

```
!$omp do
```

```
do i = 1, N
```

```
x(i) = i
```

```
end do
```

```
!$omp end do
```

```
!$omp do
```

```
do i = 1, N
```

```
!$omp atomic
```

```
total = total + x(i)
```

```
end do
```

```
!$omp end do
```

```
!$omp end parallel
```

```
write(*,*) "total = ", total
```

```
end program
```

! Parallel code with openmp do directives

! Synchronized with atomic directive

! which give correct answer, but cost more

```
[syam@saw-login1:~] ./sharing-atomic-par2
```

```
total =      12500000025000000
```

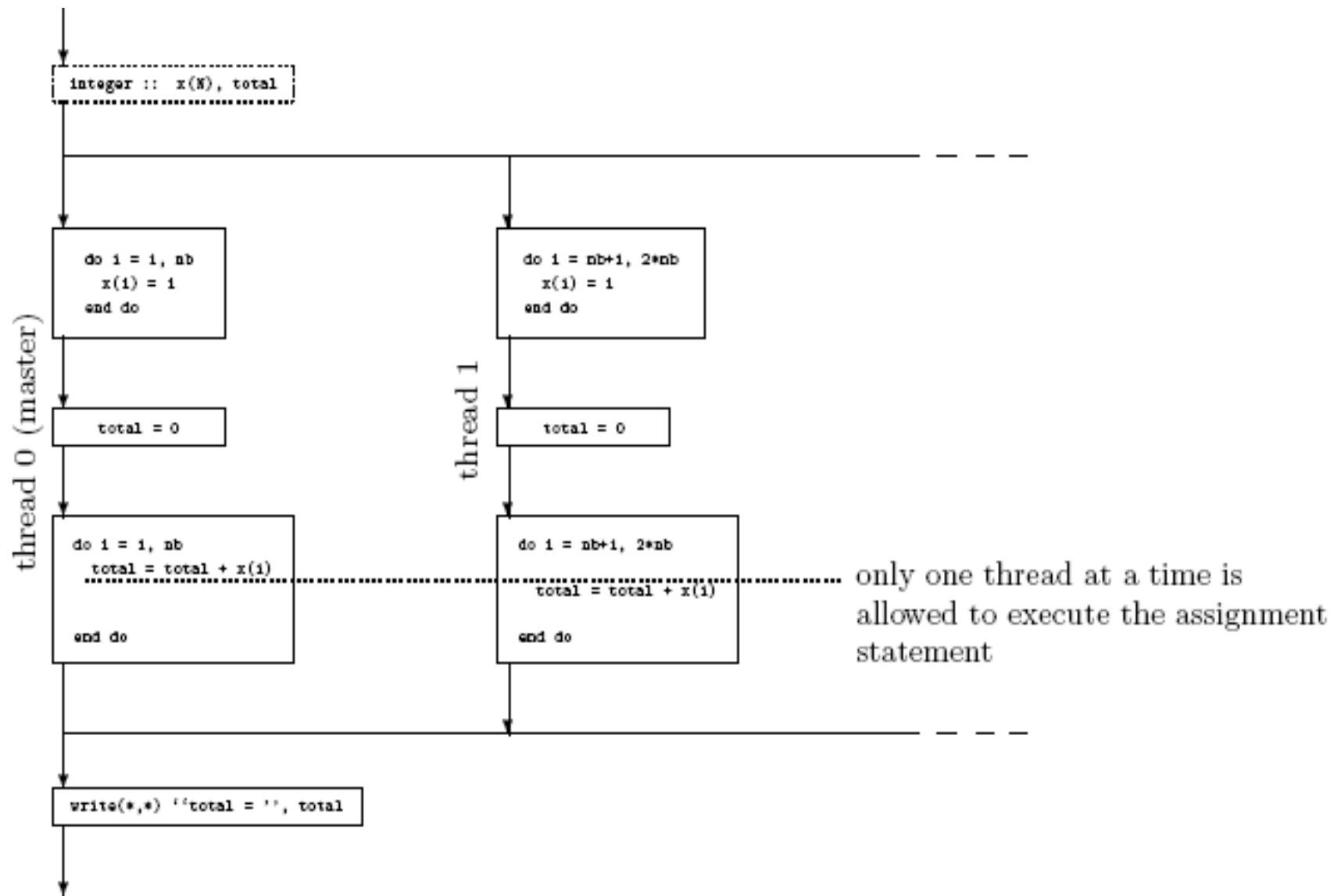
```
[syam@saw-login1:~] ./sharing-atomic-par2
```

```
total =      12500000025000000
```

```
[syam@saw-login1:~] ./sharing-atomic-par2
```

```
total =      12500000025000000
```





OpenMP: Synchronization

- **Barrier**: Each thread waits until all threads arrive.

```
#pragma omp parallel shared (A, B, C) private(id)
{
    id=omp_get_thread_num();
    A[id] = big_calc1(id);
    #pragma omp barrier
    #pragma omp for
        for(i=0;i<N;i++){C[i]=big_calc3(i,A);}
    #pragma omp for nowait
        for(i=0;i<N;i++){ B[i]=big_calc2(C, i); }
    A[id] = big_calc3(id);
}
```

implicit barrier at the
end of a for work-
sharing construct

implicit barrier at the end
of a parallel region

no implicit barrier
due to nowait

Barriers are used to synchronize the execution of multiple threads within a parallel region, not within a work-sharing construct.

Ensure that a piece of work has been completed before moving on to the next phase

```
!$omp parallel private(index)
  index = generate_next_index()
  do while (index .ne. 0)
    call add_index (index)
    index = generate_next_index()
  enddo

  ! Wait for all the indices to be generated
!$omp barrier
  index = get_next_index()
  do while (index .ne. 0)
    call process_index (index)
    index = get_next_index()
  enddo
!omp end parallel
```



OpenMP: Synchronization

- The **single** construct denotes a block of code that is executed by only one thread.
- A barrier and a flush are implied at the end of the single block.

```
#pragma omp parallel private (tmp)
{
    do_many_things();
    #pragma omp single
        { exchange_boundaries(); }
    do_many_other_things();
}
```