

# Iterative Hessian Sketch for Constrained Least-Squares

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- Consider the **constrained least-squares**

$$x_{ls} := \arg \min_{x \in C} \frac{1}{2} \|Ax - y\|_2^2$$

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- $S \in \mathbb{R}^{m \times n}$ : sketching matrix;  $ES^T S = I_n$ .

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- random row sampling: randomly select  $m$  rows from  $I_n$  (rescaled by  $\sqrt{\frac{n}{m}}$ );
- randomized orthogonal system: randomly select  $m$  rows from an orthogonal matrix  $H$  (rescaled by  $\sqrt{\frac{n}{m}}$ );



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- We expect the optimal method satisfies  $m \approx O(p)$  **Think about this part more carefully?**

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- **Sub-optimal** in the same sense;
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# Iterative Hessian Sketch

- Idea: construct a new least-squares problem for which the optimal solution is  $x^{LS} - x^t$ . Then applying HS to the problem will produce a new  $x^{t+1}$  whose distance to  $x^{LS}$  has been reduced by a factor of  $\rho$ .
- $\hat{u} := \operatorname{argmin}_{u \in \{C - x^t\}} \left\{ \frac{1}{2} \|Au\|_2^2 - \langle A^T(y - Ax^t), u \rangle \right\}$ .  
By construction,  $\hat{u} = x^{LS} - x^t$ .
- Advantages: sample size efficiency and computational time saving

# Iterative Hessian Sketch Algorithm

## Algorithm

- 1 Initialize  $x^0 = 0$ .
- 2 For iterations  $t = 0, 1, 2, \dots, N - 1$ , generate independent sketches  $S^{t+1} \in \mathbb{R}^{m \times n}$ , and perform the updates
$$x^{t+1} = \operatorname{argmin}_{x \in C} \left\{ \frac{1}{2m} \|S^{t+1} A(x - x^t)\|_2^2 - \langle A^T(y - Ax^t), x \rangle \right\}.$$
- 3 Return  $\hat{x} = x^N$ .







Sarlos (2006)

Title of the publication

*Journal Name* 12(3), 45 – 678.

# Evaluation for Sketch Performance

## Cost Approximation

In terms of  $f$ -cost, the approximated  $\tilde{x}$  is said to be  $\epsilon$ -optimal if

$$f(x^{LS}) \leq f(\tilde{x}) \leq (1 + \epsilon)^2 f(x^{LS}).$$

## Solution Approximation

In terms of  $f$ -cost, the prediction norm defined as

$$\|\tilde{x} - x^{LS}\|_A := \frac{1}{\sqrt{n}} \|A(\tilde{x} - x^{LS})\|_2.$$

Thank you!