# EP2200 Course Project 2021

# Project I - Error Control in Relay Networks

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#### I. AF RELAYING WITH END-TO-END ARQ

Consider the queuing network for AF relaying with end-toend ARQ shown. Give the queuing model for the RS and MS nodes.

For each RS and MS, the model is M/M/1.

1) The end-to-end error probability of packet transmission  $(p_{e,e2e})$  and the packet arrival rate at the first relay  $(\lambda_1)$ . The end-to-end error probability of packet transmission  $p_{e,e2e}$ :

$$p_{e,e2e} = 1 - \prod_{k=1}^{r+1} (1 - p_{(e,k)})$$
 (1)

The packet arrival rate at the first delay:

$$\begin{cases} \lambda_1 = \lambda_{r+1} \cdot p_{e,e2e} + \lambda \\ \lambda_2 = \lambda_1 \\ \lambda_3 = \lambda_2 \\ \dots \\ \lambda_{r+1} = \lambda_r \end{cases}$$

According to equation set above,

$$\lambda_{r+1} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{e,k})}$$

Therefore,

$$\lambda_1 = \lambda_{r+1} \cdot p_{e,e2e} + \lambda = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{e,k})}$$
 (2)

2) The average queuing delay and the average number of packets (waiting or under processing) at every RS and ms  $(\bar{W}_i)$  and  $(\bar{N}_i)$ .

For RS, the server utilization  $\rho_i^{AF}$ :

$$\rho_i^{AF} = \frac{\lambda_i^{AF}}{\mu_i^{AF}} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{e,k}) \cdot \mu_{AF}}$$

Because for M/M/1 system, the waiting time is  $W = \frac{1}{\mu - \lambda} - \frac{1}{\mu}$ . Therefore, the average queuing delay for every RS is:

$$W_{RS} = \frac{\lambda}{\mu_{AF}^2 \prod_{k=1}^{r+1} (1 - p_{e,k}) - \lambda \cdot \mu_{AF}}$$
(3)

Because for M/M/1 system, the average number in the system is  $N = \frac{\rho}{1-\rho}$ . Thus, the average number of packets (waiting or under processing) is:

$$N_{RS} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{e,k}) \cdot \mu_{AF} - \lambda}$$
 (4)

For MS, the server utilization  $\rho^{MS}$ :

$$\rho^{MS} = \frac{\lambda_{r+1}}{\mu_{MS}} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{e,k}) \cdot \mu_{MS}}$$

Therefore, the average queuing delay for MS is:

$$W_{MS} = \frac{\lambda}{\mu_{MS}^2 \cdot \prod_{k=1}^{r+1} (1 - p_{e,k}) - \lambda \cdot \mu_{MS}}$$
 (5)

The average number of packets (waiting or under processing) is:

$$N_{MS} = \frac{\lambda}{\prod_{k=1}^{r+1} (1 - p_{e,k}) \cdot \mu_{MS} - \lambda}$$
 (6)

3) The average end-to-end delay, that is the average time from the generation of a packet until it is received successfully at the MS  $(\bar{T})$ .

Because, the average time from the generation of a packet until it is received successfully at the MS is the sum of all the time spent in the RS + all the time spent in the MS, i.e.

$$\bar{T} = \sum_{i=1}^{r} (W_{RS} + \frac{1}{\mu_i^{AF}}) + W_{MS} + \frac{1}{\mu_{MS}}$$

Furthermore, since a packet is decoded and checked only at the last hop when the MS receives it, the time spent in each RS is the same. Therefore, the average end-to-end delay is:

$$\bar{T} = \frac{r \cdot \prod_{k=1}^{r+1} (1 - p_{e,k})}{\mu_{AF} \cdot \prod_{k=1}^{r+1} (1 - p_{e,k}) - \lambda} + \frac{\prod_{k=1}^{r+1} (1 - p_{e,k})}{\mu_{MS} \cdot \prod_{k=1}^{r+1} (1 - p_{e,k}) - \lambda}$$
(7)

# II. DF RELAYING WITH END-TO-END OR DF RELAYING WITH HOP-BY-HOP ARQ

Consider the queuing network for DF relaying with end-toend ARQ and with hop-by-hop ARQ.

## A. DF Relaying with end-to-end ARQ

1) The packet arrival rate at each RS and MS  $(\lambda_j)$ :

$$\lambda_{1} = \lambda + \lambda_{1} \cdot p_{e,1} + \lambda_{2} \cdot p_{e,2} + \dots + \lambda_{r+1} \cdot p_{e,r+1}$$

$$= \lambda + \lambda_{1} \cdot \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1 - p_{e,k})]$$

$$= \frac{\lambda}{1 - \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{i=1}^{k} (1 - p_{e,i})]}$$

Therefore,

$$\lambda_{i} = \lambda_{1} \prod_{k=0}^{i-1} (1 - p_{e,k})$$

$$= \frac{\lambda \cdot \prod_{k=1}^{i-1} (1 - p_{e,k})}{1 - \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1 - p_{e,j})]}$$
(8)

2) The average queuing delay and the average number of packets (waiting or under processing) at every RS and MS ( $\bar{W}_i$  and  $\bar{N}_i$ )

For RS:

The average queuing delay  $W_{RS}^i = \frac{1}{\mu_{DF} - \lambda_i} - \frac{1}{\mu_{DF}}$ :

$$\frac{\lambda \cdot \prod_{k=1}^{i-1} (1-p_{e,k})}{\mu_{DF}^2 [1-\sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1-p_{e,j})]] - \lambda \mu_{DF} \prod_{k=1}^{i-1} (1-p_{e,k})} \quad ^{(9)}$$

Because the utilization is:

$$\begin{split} \rho_{RS}^i &= \frac{\lambda_i}{\mu_{DF}} \\ &= \frac{\lambda \cdot \prod_{k=1}^{i-1} (1 - p_{e,k})}{\mu_{DF} - \mu_{DF} \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1 - p_{e,j})]} \end{split}$$

Therefore, the average number of packets at every RS is  $N_{RS}^i = \frac{\rho_{RS}^i}{1 - \rho_{DS}^i}$ :

$$\frac{\lambda \cdot \prod_{k=1}^{i-1} (1 - p_{e,k})}{\mu_{DF} - \mu_{DF} \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{i=1}^{k} (1 - p_{e,j})] - \lambda \cdot \prod_{k=1}^{i-1} (1 - p_{e,k})} \quad (10)$$

For MS:

The average queuing delay  $W_{MS}=rac{1}{\mu_{MS}-\lambda_{r+1}}-rac{1}{\mu_{MS}}$ :

$$\frac{\lambda \cdot \prod_{k=1}^{r} (1 - p_{e,k})}{\mu_{MS}^{2} [1 - \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1 - p_{e,j})]] - \lambda \mu_{MS} \prod_{k=1}^{r} (1 - p_{e,k})}$$
(11)

Because the utilization is:

$$\rho_{MS} = \frac{\lambda_{r+1}}{\mu_{MS}}$$

$$= \frac{\lambda \cdot \prod_{k=1}^{r} (1 - p_{e,k})}{\mu_{MS} - \mu_{MS} \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1 - p_{e,j})]}$$

Therefore, the average number of packets at every MS is  $N_{MS} = \frac{\rho_{MS}}{1-\rho_{MS}}$ :

$$\frac{\lambda \cdot \prod_{k=1}^{r} (1 - p_{e,k})}{\mu_{MS} - \mu_{MS} \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{i=1}^{k} (1 - p_{e,j})] - \lambda \cdot \prod_{k=1}^{r} (1 - p_{e,k})}$$
(12)

3) The average end-to-end delay, that is, the average time from the generation of a packet until it is received successfully at the MS  $(\bar{T})$ .

Because, the average time from the generation of a packet until it is received successfully at the MS is the sum of all the time spent in the RS + all the time spent in the MS, i.e.

$$\bar{T} = \sum_{i=1}^{r} (W_{RS}^{i} + \frac{1}{\mu_{i}^{DF}}) + W_{MS} + \frac{1}{\mu_{MS}}$$

Therefore, the average end-to-end delay is  $\bar{T}$ :

$$\begin{split} \sum_{i=1}^{r} \frac{1 - \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1 - p_{e,j})]}{\mu_{DF} - \mu_{DF} \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1 - p_{e,j})] - \lambda \prod_{k=1}^{i-1} (1 - p_{e,k})} \\ + \frac{1 - \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1 - p_{e,j})]}{\mu_{MS} - \mu_{MS} \sum_{k=0}^{r} [p_{e,k+1} \cdot \prod_{j=1}^{k} (1 - p_{e,j})] - \lambda \prod_{k=1}^{r} (1 - p_{e,k})} \end{split}$$

#### B. DF Relaying with hop-by-hop ARQ

1) The packet arrival rate at each RS and MS  $(\lambda_i)$ :

$$\begin{cases} \lambda_1 = \lambda + \lambda_1 \cdot p_{e,1} \\ \lambda_2 = \lambda_1 \cdot (1 - p_{e,1}) + \lambda_2 \cdot p_{e,2} \\ \dots \\ \lambda_r = \lambda_{r-1} \cdot (1 - p_{e,r-1}) + \lambda_r \cdot p_{e,r} \\ \lambda_{r+1} = \lambda_r \cdot (1 - p_{e,r}) + \lambda_{r+1} \cdot p_{e,r+1} \end{cases}$$

Therefore, from the equation set above, we can get the equation of  $\lambda_i$  and  $\lambda_{r+1}$ , which are:

$$\begin{cases} \lambda_i = \frac{\lambda}{1 - p_{e,i}} \\ \lambda_{r+1} = \frac{\lambda}{1 - p_{e,r+1}} \end{cases}$$
 (14)

2) The average queuing delay and the average number of packets (waiting or under processing) at every RS and MS ( $\bar{W}_i$  and  $\bar{N}_i$ ).

For RS:

The average queuing delay is  $W_{RS}^i = \frac{1}{\mu_{RS} - \lambda_i} - \frac{1}{\mu_{RS}}$ :

$$W_{RS}^{i} = \frac{\lambda}{\mu_{DF} \cdot [\mu_{DF} \cdot (1 - p_{e,i}) - \lambda]}$$
(15)

Because the utilization is:

$$\rho_{RS}^i = \frac{\lambda_i}{\mu_{DF}} = \frac{\lambda}{\mu_{DF} \cdot (1 - p_{e,i})}$$

Therefore, the average number of packet at every RS is  $N_{RS}^i=rac{
ho_{RS}^i}{1ho_{DS}^i}$ :

$$N_{RS}^{i} = \frac{\lambda}{\mu_{DF}(1 - p_{e,i}) - \lambda}$$
 (16)

For MS:

The average queuing delay is  $W_{MS} = \frac{1}{\mu_{MS} - \lambda_i} - \frac{1}{\mu_{MS}}$ :

$$W_{MS} = \frac{\lambda}{\mu_{MS} \cdot [\mu_{MS} \cdot (1 - p_{e,r+1}) - \lambda]} \tag{17}$$

Because the utilization is:

$$\rho_{MS} = \frac{\lambda_{r+1}}{\mu_{MS}} = \frac{\lambda}{\mu_{MS} \cdot (1 - p_{e,r+1})}$$

Therefore, the average number of packet at MS is  $N_{MS} = \frac{\rho_{MS}}{1-\rho_{MS}}$  :

$$N_{MS} = \frac{\lambda}{\mu_{MS} \cdot (1 - p_{e,r+1}) - \lambda} \tag{18}$$

3) The average end-to-end delay, that is, the average time from the generation of a packet until it is received

successfully at the MS  $(\bar{T})$ .

Because, the average time from the generation of a packet until it is received successfully at the MS is the sum of all the time spent in the RS + all the time spent in the MS, i.e.

$$\bar{T} = \sum_{i=1}^{r} (W_{RS}^{i} + \frac{1}{\mu_{i}^{DF}}) + W_{MS} + \frac{1}{\mu_{MS}}$$

Therefore, the average end-to-end delay is  $\bar{T}$ :

$$\bar{T} = \sum_{i=1}^{r} \frac{1 - p_{e,i}}{\mu_{DF}(1 - p_{e,i}) - \lambda} + \frac{1 - p_{e,r+1}}{\mu_{MS}(1 - p_{e,r+1}) - \lambda}$$
(19)

#### III. STABILITY REGION

A. Stability Region for AF relaying with end-to-end ARQ

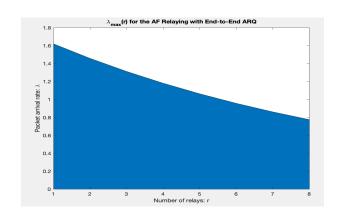


Fig. 1. The maximal allowed arrival rate of packets  $\lambda_{max}(r)$  for the AF relaying with end-to-end ARQ.

As  $\lambda \leq 2 \cdot 0.9^{r+1}$ , when r goes to infinity,  $2 \cdot 0.9^{r+1}$  will go to 0. Therefore,  $\lambda_{max}(r)$  will go to 0 as well.

B. Stability Region for DF relaying with end-to-end ARQ

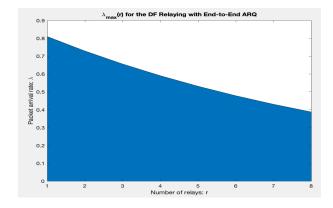


Fig. 2. The maximal allowed arrival rate of packets  $\lambda_{max}(r)$  for the DF relaying with end-to-end ARQ.

As  $(\lambda \leq 0.9^{r+1}) \cap (\lambda \leq 1.8)$ , when r goes to infinity,  $2 \cdot 0.9^{r+1}$  will go to 0. Therefore,  $\lambda_{max}(r)$  will go to 0 as well.

# C. Stability Region for DF relaying with hop-by-hop ARQ

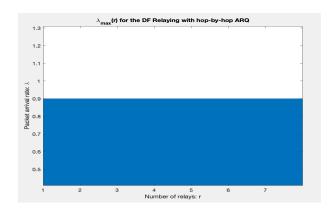


Fig. 3. The maximal allowed arrival rate of packets  $\lambda_{max}(r)$  for the DF relaying with hop-by-hop ARQ.

Since  $(\lambda \le 0.9) \cap (\lambda \le 1.8)$ , when r goes to infinity, there is no influence on  $\lambda_{max}(r)$ .

#### IV. ARRIVAL RATE-DELAY CHARACTERISTICS

## A. Arrival rate-delay for AF relaying with end-to-end ARQ

$$T = \frac{5 \cdot 0.9^5}{2 \cdot 0.9^5 - \lambda}$$

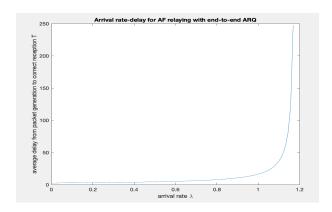


Fig. 4. arrival rate-delay for AF relaying with end-to-end ARQ.

When  $p_{e,j}$  increases, the arrival rate-delay curves moves to the left.

## B. Arrival rate-delay for DF relaying with end-to-end ARQ

$$T = \sum_{i=1}^{4} \frac{0.9^5}{0.9^5 - \lambda \cdot 0.9^{i-1}} + \frac{0.9^5}{2 * 0.9^5 - 0.9^4 \cdot \lambda}$$

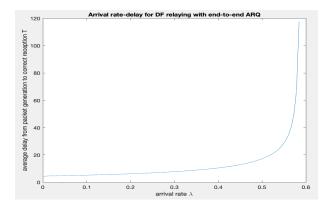


Fig. 5. arrival rate-delay for DF relaying with end-to-end ARQ.

When  $p_{e,j}$  increases, the arrival rate-delay curves moves to the left.

# C. Arrival rate-delay for DF relaying with hop-by-hop ARQ

$$T = \frac{3.6}{0.9 - \lambda} + \frac{0.9}{1.8 - \lambda}$$

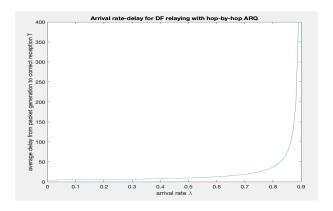


Fig. 6. arrival rate-delay for DF relaying with hop-by-hop ARQ.

When  $p_{e,j}$  increases, the arrival rate-delay curves moves to the left.

#### V. AF or DF?

## A. Condition I

Denote  $p_{e,j}=p, j=1,2,\cdots,r+1$ . In this condition, the average system delay for the three models is shown as

followed:

The average delay for AF relaying with end-to-end ARQ:

$$T_1 = 5$$

The average delay for DF relaying with end-to-end ARQ:

$$T_2 = \sum_{i=1}^{4} \frac{(1-p)^5}{(1-p)^5 - 0.5 \cdot (1-p)^{i+4}} + \frac{(1-p)^5}{2 \cdot (1-p)^5 - 0.5 \cdot (1-p)^9}$$

The average delay for DF relaying with hop-by-hop ARQ:

$$T_3 = \frac{26}{3}$$

Therefore, the figure can be drawn with  $0.05 \le j \le 0.95$ .

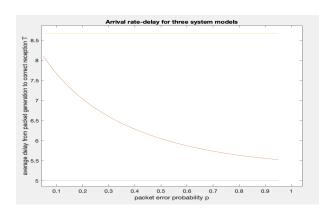


Fig. 7. arrival rate-delay for three system models.

As can be seen from Fig.7 above, choosing AF relaying with end-to-end ARQ will be the optimal choice.

#### B. Condition II

In this scenario, the average delay for AF relaying with end-to-end ARQ:

$$T = \frac{4}{k-1} + 1$$

The average delay for DF relaying with end-to-end ARQ:

$$T = \sum_{i=1}^{4} \frac{0.9^5}{0.9^5 - 0.5 \cdot 0.9^{i+4}} + \frac{0.9^5}{2 \cdot 0.9^5 - 0.5 \cdot 0.9^9}$$

The average delay for DF relaying with hop-by-hop ARQ:

$$T = \frac{26}{3}$$

Therefore, the figure can be drawn when k increases from 1 to 4.

As can be observed from the graph that, when  $1 \leq k \leq 1.6$ , DF relaying with end-to-end ARQ will be the optimal choice. However, when  $1.6 \leq k \leq 4$ , AF relaying with end-to-end will shows its advantage in system average delay.

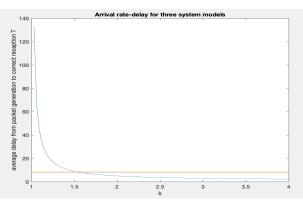


Fig. 8. arrival rate-delay for three system models.

# VI. PROBABILISTIC ARQ - EXTRA EXERCISE FOR THE $$\operatorname{\textbf{Best}}$$

## A. The P-ARQ queuing network mode

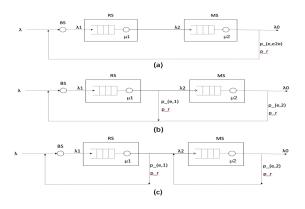


Fig. 9. P-ARQ queuing network model for (a)AF relaying with end-to-end ARQ, (b)DF relaying with end-to-end ARQ, and (c)DF relaying with hop-by-hop ARQ

#### B. Performance evaluation

1) P-ARQ - AF - end-to-end  $T_{AF-ETE}^{P}$ :

$$\begin{split} &\frac{1-(1-p_{e,1})[1-(1-p_{e,1})(1-p_{e,2})]p_r}{\mu_{AF}-\mu_{AF}p_r[1-(1-p_{e,1})(1-p_{e,2})]-\lambda} \\ &+\frac{1-(1-p_{e,1})[1-(1-p_{e,1})(1-p_{e,2})]p_r}{\mu_{MS}-\mu_{MS}p_r[1-(1-p_{e,1})(1-p_{e,2})]-\lambda} \end{split} \tag{20}$$

2) P-ARQ - DF - end-to-end  $T^P_{DF-ETE}$ :

$$\begin{split} &\frac{1-p_{e,1}p_r-(1-p_{e,1})p_{e,2}p_r}{\mu_{DF}-\mu_{DF}p_{e,1}p_r-\mu_{DF}(1-p_{e,1})p_{e,2}p_r-\lambda}\\ &+\frac{(1-p_{e,1})(1-p_{e,1}p_r-(1-p_{e,1})p_{e,2}p_r)}{\mu_{MS}-\mu_{MS}p_{e,1}p_r-\mu_{MS}(1-p_{e,1})p_{e,2}p_r-\lambda(1-p_{e,1})} \end{split}$$

3) P-ARQ - DF - hop-by-hop  $T^P_{DF-HBH}$ :

$$\frac{1 - p_{e,1}p_r}{\mu_{DF} - \mu_{DF}p_{e,1}p_r - \lambda} + \frac{(1 - p_{e,2}p_r)(1 - p_{e,1}p_r)}{\mu_{MS}(1 - p_{e,2}p_r)(1 - p_{e,1}p_r) - (1 - p_{e,1})\lambda}$$
(22)

4) Plotting the figure with the given parameters, the equations above can be switched into the following:

$$\begin{split} T_{AF-ETE}^P &= \frac{2 - 0.342 \cdot p_r}{1.8 - 0.38 \cdot p_r} \\ T_{DF-ETE}^P &= \frac{1 - 0.19 \cdot p_r}{0.8 - 0.19 \cdot p_r} + \frac{0.9 - 0.171 \cdot p_r}{1.82 - 0.38 \cdot p_r} \\ T_{DF-HBH}^P &= \frac{1 - 0.1 \cdot p_r}{0.8 - 0.1 \cdot p_r} + \frac{(1 - 0.1 \cdot p_r)^2}{2 \cdot (1 - 0.1 \cdot p_r)^2 - 0.18} \end{split}$$

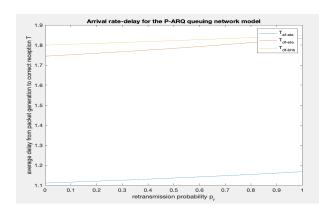


Fig. 10. end-to-end delay for P-ARQ queuing network model.