

# Model

## Variables

$tasks$	The number of tasks to schedule
$machines$	The number of machines available
$T = \{1 \dots tasks\}$	The tasks
$M = \{1 \dots machines\}$	The machines
$d(task_i) \in \{0 \dots 2^{63} - 1\}, \quad task_i \in T$	The duration of task $i$
$s(task_i) \in \{0 \dots 2^{63} - 1\}, \quad task_i \in T$	The start time of task $i$
$m(task_i) \in M, \quad task_i \in T$	The machine task $i$ uses
$dep(task_i, task_j) \in \{0, 1\}, \quad task_i, task_j \in T$	The dependency between task $i$ and task $j$ . If 1, task $i$ needs to be completed before task $j$
$End = \{0..D\}, \quad D = sum(\{d(task_i)   \forall task_i \in T\})$	Marks the end of the schedule, i.e. makespan. The variable that is minimized.
$same\_machine_k \subset T$	Set of tasks needing the same machine
$concurrent\_tasks_k \subset T$	Set of tasks needing concurrent execution
$occupied_k \subset T$	Set of tasks occupying a machine in the order specified by $dep$
$overlap(task_i, task_j), \quad task_i, task_j \in T$	$task_i$ overlaps $task_j$

## Constraints

$$End = \max(\{s(task_i) + d(task_i) | \forall task_i \in T\})$$

The end is where the last task ends

$$task_i \prec task_j, \quad task_i, task_j \in T \wedge dep(task_i, task_j) = 1$$

$task_i$  has to precede  $task_j$  if specified in  $dep$ .

$$overlap(task_i, task_j) \implies m(task_i) \neq m(task_j),$$

If  $task_i$  and  $task_j$  overlap, they cannot use the same machine.

$$task_i, task_j \in T$$

$$m(task_i) = m(task_j), \quad \forall task_i, task_j \in same\_machine_k$$

Every task in a  $same\_machine$  set should use the same machine.

$$s(task_i) = s(task_j) \wedge m(task_i) \neq m(task_j),$$

Every task in a  $concurrent\_tasks$  set should execute at the same time, but they cannot use the same machine.

$$\forall task_i, task_j \in concurrent\_tasks_k$$

$$m(task_k) = m(task_i) \implies task_k \prec task_i \vee task_k \succ task_j,$$

If two tasks are set to occupy a machine following each other, there cannot be another task between them using the same machine.

$$\forall task_k \in T, \quad task_i, task_j \in occupied_k, \quad dep(task_i, task_j) = 1,$$

$$task_i \neq task_k, \quad task_j \neq task_k$$